The combinatorics of free bifibrations

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What is a bifibration?

One category living over another category, such that *objects* of the category above may be pushed and pulled along *arrows* of the category below.

Formally:

$$\begin{array}{c} \mathcal{D} \\ {\scriptstyle p}\big\downarrow \\ \mathcal{C} \end{array}$$

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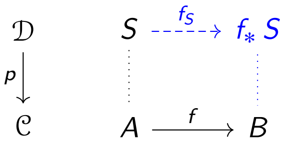

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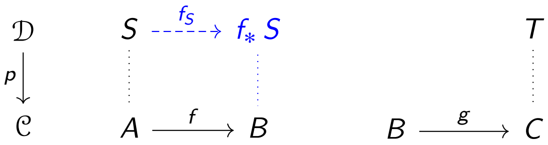

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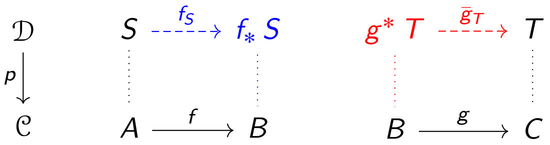

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$$
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\mathcal{D} \\
{\scriptstyle p}\big\downarrow \\
\mathcal{C}
\end{array}
\qquad
\begin{array}{ccc}
S & \overset{f_S}{\dashrightarrow} & f_* \, S \\
\vdots & & \vdots \\
A & \xrightarrow{\ f\ } & B
\end{array}
\qquad
\begin{array}{ccc}
g^* \, T & \overset{\bar{g}_T}{\dashrightarrow} & T \\
\vdots & & \vdots \\
B & \xrightarrow{\ g\ } & C
\end{array}
$$

...and these liftings should be "universal" in an appropriate sense.

What is a bifibration? (cont.)

Pushing and pulling along an arrow $f : A \to B$ of $\mathcal{C}$ induces an *adjunction*

$$\mathcal{D}_A \underset{f^*}{\overset{f_*}{\rightleftarrows}} \mathcal{D}_B \quad (f_* \dashv f^*)$$

between the *fibers* of A and B.

This leads to an equivalent way of seeing bifibrations $\mathcal{D} \to \mathcal{C}$, as pseudofunctors $\mathcal{C} \to \mathcal{A}\mathrm{dj}$ into the category of small categories and adjunctions.

A simple example

Let $\mathcal{S}\mathrm{et}$ be the category of sets and functions.

Let $\mathcal{S}\mathrm{ubset}$ be the category whose objects are subsets, and whose arrows $(S \subseteq A) \longrightarrow (T \subseteq B)$ are functions $f : A \to B$ such that $a \in S$ implies $f(a) \in T$.

The evident forgetful functor $\mathcal{S}\mathrm{ubset} \to \mathcal{S}\mathrm{et}$ is a bifibration:

$$
\begin{array}{c}
\mathcal{S}\mathrm{ubset} \\
\big\downarrow \\
\mathcal{S}\mathrm{et}
\end{array}
\qquad
\begin{array}{ccc}
S & \xrightarrow{f_S} & f(S) \\
\rotatebox{90}{$\subseteq$} & & \rotatebox{90}{$\subseteq$} \\
A & \xrightarrow{f} & B
\end{array}
\qquad
\begin{array}{ccc}
g^{-1}(T) & \xrightarrow{\bar{g}_T} & T \\
\rotatebox{90}{$\subseteq$} & & \rotatebox{90}{$\subseteq$} \\
B & \xrightarrow{g} & C
\end{array}
$$

(Adjunction property: $f(S) \subseteq R \iff S \subseteq f^{-1}(R)$.)

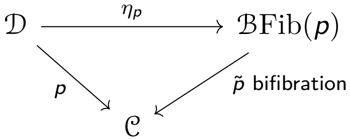
Other motivating examples, from logic

Pushforward and pullback may be used to express:

- strongest postconditions and weakest preconditions in program logic
- existential and universal quantification in predicate logic
- diamond and box in modal logic
- $\otimes$ and $\parr$ in linear logic

Our problem

Most functors are not bifibrations.
Given a functor $p : \mathcal{D} \to \mathcal{C}$, how do we construct the **free bifibration** over p?



A relatively little-studied problem:

- Robert Dawson, Robert Paré, and Dorette Pronk. Adjoining adjoints. *Advances in Mathematics*, 178(1):99–140, 2003.
- François Lamarche. Path functors in Cat. Unpublished, 2010. `https://hal.inria.fr/hal-00831430`.

Overview of our work

Developed alternative constructions of the free bifibration over a functor $p : \mathcal{D} \to \mathcal{C}$

- a *proof-theoretic* construction, using sequent calculus
- an *algebraic* construction, using double categories
- a *topological* construction, using string diagrams

(These provide three different perspectives, but all closely related.)

We also discovered examples of specific functors $p : \mathcal{D} \to \mathcal{C}$, such that the free bifibration over p has some surprisingly nice combinatorics.

A sequent calculus for the free bifibration over $p : \mathcal{D} \to \mathcal{C}$

Formulas ($S \sqsubset A$):

$$\frac{X \in \mathcal{D} \quad p(X) = A}{X \sqsubset A} \qquad \frac{S \sqsubset A \quad f : A \to B}{f_* S \sqsubset B} \qquad \frac{f : A \to B \quad T \sqsubset B}{f^* T \sqsubset A}$$

Proofs ($S \underset{h}{\Longrightarrow} T$):

$$\frac{S \underset{fg}{\Longrightarrow} T}{f_* S \underset{g}{\Longrightarrow} T} L_f \qquad \frac{S \underset{g}{\Longrightarrow} T}{S \underset{gf}{\Longrightarrow} f_* T} R_f \qquad \frac{S \underset{g}{\Longrightarrow} T}{f^* S \underset{fg}{\Longrightarrow} T} L_{\bar{f}} \qquad \frac{S \underset{gf}{\Longrightarrow} T}{S \underset{g}{\Longrightarrow} f^* T} R_{\bar{f}}$$

$$\frac{\alpha : X \longrightarrow Y \in \mathcal{D} \quad p(\alpha) = g}{X \underset{g}{\Longrightarrow} Y} \alpha$$

Equational theory on derivations

Need to impose four **permutation equivalences** on derivations, including

$$
\cfrac{\cfrac{S \underset{fg}{\Longrightarrow} T}{S \underset{fgh}{\Longrightarrow} h_* \, T}\; R_h}{f_* \, S \underset{gh}{\Longrightarrow} h_* \, T}\; L_f
\sim
\cfrac{\cfrac{S \underset{fg}{\Longrightarrow} T}{f_* \, S \underset{g}{\Longrightarrow} T}\; L_f}{f_* \, S \underset{gh}{\Longrightarrow} h_* \, T}\; R_h
\qquad
\cfrac{\cfrac{S \underset{g}{\Longrightarrow} T}{S \underset{gh}{\Longrightarrow} h_* \, T}\; R_h}{f^* \, S \underset{fgh}{\Longrightarrow} h_* \, T}\; L_{\bar{f}}
\sim
\cfrac{\cfrac{S \underset{g}{\Longrightarrow} T}{S^* \underset{fg}{\Longrightarrow} T}\; L_{\bar{f}}}{f^* \, S \underset{fgh}{\Longrightarrow} h_* \, T}\; R_h
$$

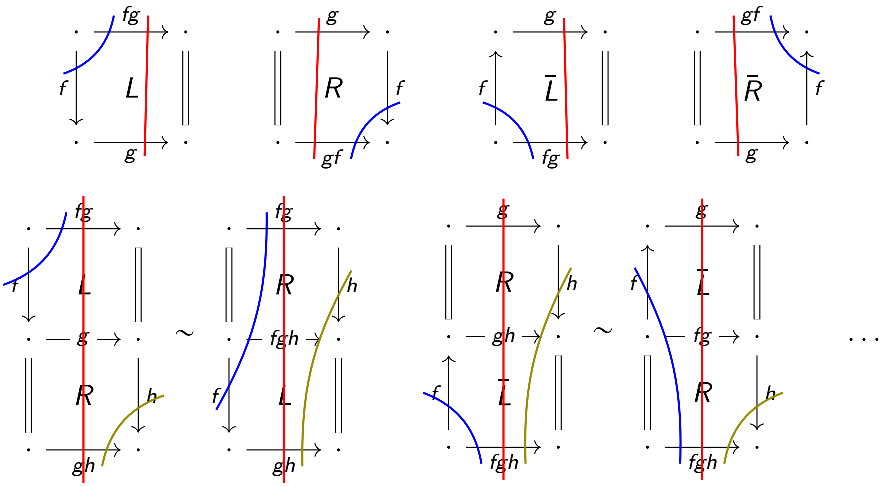
plus their symmetric versions with pushforward and pullback swapped.

Arrows of $\mathcal{B}\mathrm{Fib}(p)$ are equivalence classes of proofs. Composition is by cut-elimination.

Example derivations

$$\dfrac{\dfrac{\dfrac{}{S \underset{A}{\Longrightarrow} S}\ \mathrm{id}_S}{S \underset{f}{\Longrightarrow} f_* S}\ R_f}{S \underset{A}{\Longrightarrow} f^* f_* S}\ R_{\bar{f}}
\qquad
\dfrac{\dfrac{S \underset{A}{\Longrightarrow} S'}{S \underset{f}{\Longrightarrow} f_* S'}\ R_f}{f_* S \underset{B}{\Longrightarrow} f_* S'}\ L_f
\qquad
\dfrac{\dfrac{T \underset{B}{\Longrightarrow} T'}{f^* T \underset{f}{\Longrightarrow} T'}\ L_{\bar{f}}}{f^* T \underset{A}{\Longrightarrow} f^* T'}\ R_{\bar{f}}
\qquad
\dfrac{\dfrac{\dfrac{}{T \underset{B}{\Longrightarrow} T}\ \mathrm{id}_T}{f^* T \underset{f}{\Longrightarrow} T}\ L_{\bar{f}}}{f_* f^* T \underset{B}{\Longrightarrow} T}\ L_f$$

Construction via the double category of zigzags, and via string diagrams



Canonical forms in general

A challenge in understanding free bifibrations is getting a handle on the equivalence classes (of proofs/double cells/string diagrams) induced by the permutation relations.

Note that equivalence is in general undecidable![1]

Nevertheless, we (believe we) have a normal form based on *maximal multifocusing*...

$$\frac{N \underset{\pi f}{\Longrightarrow} P}{\pi_* N \underset{f}{\Longrightarrow} P}\ L_\pi \qquad \frac{N \underset{\pi f \rho}{\Longrightarrow} P}{\pi_* N \underset{f}{\Longrightarrow} \rho^* P}\ L_\pi R_{\bar\rho} \qquad \frac{N \underset{f \rho}{\Longrightarrow} P}{N \underset{f}{\Longrightarrow} \rho^* P}\ R_{\bar\rho}$$

$$\frac{P \underset{f}{\Longrightarrow} Q}{\pi^* P \underset{\pi f}{\Longrightarrow} Q}\ L_{\bar\pi} \qquad \frac{P \underset{f}{\Longrightarrow} N}{\pi^* P \underset{\pi f \rho}{\Longrightarrow} \rho_* N}\ L_{\bar\pi} R_\rho \qquad \frac{M \underset{f}{\Longrightarrow} N}{M \underset{f \rho}{\Longrightarrow} \rho_* N}\ R_\rho$$

[1] By adapting a construction in: Robert Dawson, Robert Paré, and Dorette Pronk. Undecidability of the free adjoint construction. *Applied Categorical Structures*, 11:403–419, 2003.

Now for some examples!

Example #1

Consider the following functor:

$$\begin{array}{ccccc} 1 & & 0 & & \\ {\scriptstyle p_0}\big\downarrow & & \vdots & & \\ 2 & & 0 & \xrightarrow{\;f\;} & 1 \end{array}$$

Build the free bifibration $\mathcal{B}\mathrm{Fib}(p_0) \to 2$, and look at the fiber of 0.

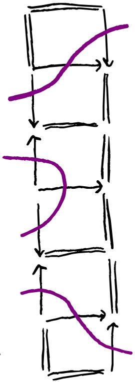
Objects are isomorphic to even-length alternating push/pull sequences $f^* f_* \cdots f^* f_* 0$

Let $d_{m,n}$ be the number of arrows $(f^* f_*)^m 0 \longrightarrow (f^* f_*)^n 0$?

Puzzle: what is $d_{m,n}$?

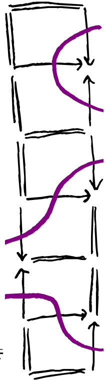
$d_{2,1} = 1$

$$
\begin{array}{cl}
\hline
0 \underset{0}{\Longrightarrow} 0 & id \\
\hline
0 \underset{f}{\Longrightarrow} f_* \, 0 & R_f \\
\hline
f_* \, 0 \underset{1}{\Longrightarrow} f_* \, 0 & L_f \\
\hline
f^* \, f_* \, 0 \underset{f}{\Longrightarrow} f_* \, 0 & L_{\bar{f}} \\
\hline
f_* \, f^* \, f_* \, 0 \underset{1}{\Longrightarrow} f_* \, 0 & L_f \\
\hline
f^* \, f_* \, f^* \, f_* \, 0 \underset{f}{\Longrightarrow} f_* \, 0 & L_{\bar{f}} \\
\hline
f^* \, f_* \, f^* \, f_* \, 0 \underset{0}{\Longrightarrow} f^* \, f_* \, 0 & R_{\bar{f}}
\end{array}
$$

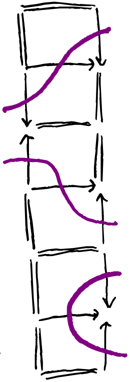


$d_{1,2} = 2$

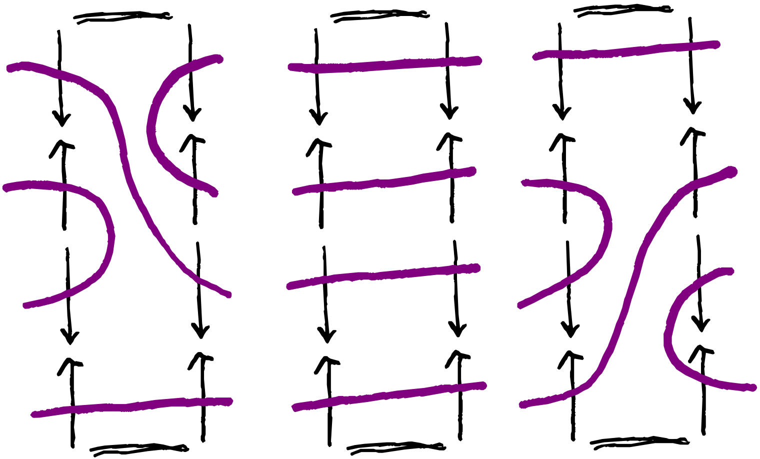
$$
\begin{array}{cl}
\hline
0 \underset{0}{\Longrightarrow} 0 & \mathrm{id}_0 \\
\hline
0 \underset{f}{\Longrightarrow} f_* \, 0 & R_f \\
\hline
0 \underset{0}{\Longrightarrow} f^* \, f_* \, 0 & R_{\bar{f}} \\
\hline
0 \underset{f}{\Longrightarrow} f_* \, f^* \, f_* \, 0 & R_f \\
\hline
f_* \, 0 \underset{1}{\Longrightarrow} f_* \, f^* \, f_* \, 0 & L_f \\
\hline
f^* \, f_* \, 0 \underset{f}{\Longrightarrow} f_* \, f^* \, f_* \, 0 & L_{\bar{f}} \\
\hline
f^* \, f_* \, 0 \underset{0}{\Longrightarrow} f^* \, f_* \, f^* \, f_* \, 0 & R_{\bar{f}}
\end{array}
$$



$$
\begin{array}{cl}
\hline
0 \underset{0}{\Longrightarrow} 0 & \mathrm{id}_0 \\
\hline
0 \underset{f}{\Longrightarrow} f_* \, 0 & R_f \\
\hline
f_* \, 0 \underset{1}{\Longrightarrow} f_* \, 0 & L_f \\
\hline
f^* \, f_* \, 0 \underset{f}{\Longrightarrow} f_* \, 0 & L_{\bar{f}} \\
\hline
f^* \, f_* \, 0 \underset{0}{\Longrightarrow} f^* \, f_* \, 0 & R_{\bar{f}} \\
\hline
f^* \, f_* \, 0 \underset{f}{\Longrightarrow} f_* \, f^* \, f_* \, 0 & R_f \\
\hline
f^* \, f_* \, 0 \underset{0}{\Longrightarrow} f^* \, f_* \, f^* \, f_* \, 0 & R_{\bar{f}}
\end{array}
$$



$d_{2,2} = 3$



$$f^* \, f_* \, f^* \, f_* \, 0 \underset{0}{\Longrightarrow} f^* \, f_* \, f^* \, f_* \, 0$$

Example #1 continued

Arrows $(f^* f_*)^m 0 \longrightarrow (f^* f_*)^n 0$ correspond to monotone maps m $\rightarrow$ n!

Indeed, the free bifibration over $p_0 : 1 \rightarrow 2$ captures the adjunction

$$\Delta \underset{f^*}{\overset{f_*}{\underset{\perp}{\rightleftarrows}}} \Delta_\perp$$

between the category Δ of finite ordinals and order-preserving maps, and the category $\Delta_\perp$ of non-empty finite ordinals and order-and-least-element-preserving maps.

. . . So what's the answer to the puzzle?

Example #1 continued

Arrows $(f^* f_*)^m 0 \longrightarrow (f^* f_*)^n 0$ correspond to monotone maps $\mathrm{m} \to \mathrm{n}$!

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. . . So what's the answer to the puzzle? $d_{m,n} = \binom{n+m-1}{m}$

Example #2

Now consider the following functor:

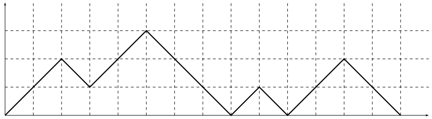
$$
\begin{array}{ccccccccc}
1 & & 0 & & & & & & \\
{\scriptstyle p_0}\big\downarrow & & \vdots & & & & & & \\
\mathbb{N} & & 0 & \longrightarrow & 1 & \longrightarrow & 2 & \longrightarrow & \dots
\end{array}
$$

Build the free bifibration $\mathcal{B}\mathrm{Fib}(p_0) \to \mathbb{N}$, and look at the fiber of 0.

Puzzle: what are its objects?

A category with Dyck walks as objects!

$$f^* f^* f_* f_* f^* f_* f^* f^* f^* f_* f_* f^* f_* f_* 0 \quad = $$

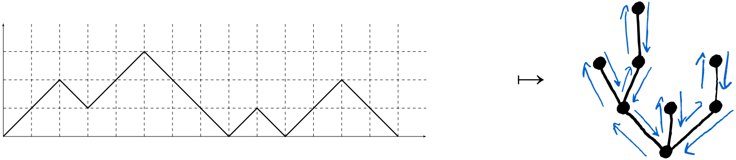


But what is a *morphism* of Dyck walks??

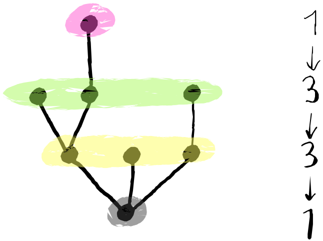
The $\mathcal{B}\mathrm{Fib}(-)$ construction gives an answer. Is it something natural/known?

Reconstructing the Batanin-Joyal category of trees

Dyck paths have a well-known, canonical bijection with (finite rooted plane) trees.



Trees may also be encoded as *functors* $T : \mathbb{N}^{\mathrm{op}} \to \Delta$.



Reconstructing the Batanin-Joyal category of trees

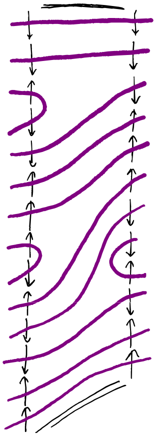
Consider *natural transformations* $\theta : S \Rightarrow T$.

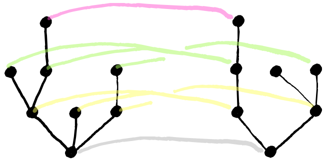
$$\begin{array}{ccc} \vdots & & \vdots \\ \downarrow & & \downarrow \\ S(2) & \xrightarrow{\theta_2} & T(2) \\ \downarrow & & \downarrow \\ S(1) & \xrightarrow{\theta_1} & T(1) \\ \downarrow & & \downarrow \\ S(0) = 1 & = & 1 = T(0) \end{array}$$

In other words, map nodes to nodes of the same height, respecting parents.

Reconstructing the Batanin-Joyal category of trees

Theorem: $\mathcal{B}\mathrm{Fib}(p_0 : 1 \to \mathbb{N})_0 \cong \mathsf{PTree}$.





(More generally, $\mathcal{B}\mathrm{Fib}(p_0)_k \cong \mathsf{PTree}_k$ = category of finite rooted plane trees whose rightmost branch is pointed by a node of height k.)

Example #2 continued

Fix a walk W, and consider the following pair of sequences:

$$in[W]_n = \#\{\, \theta : S \Rightarrow W \mid |S| = n \,\} \qquad out[W]_n = \#\{\, \theta : W \Rightarrow T \mid |T| = n \,\}$$

These seem to be always nice!

W	$out[W]$	$in[W]$
ϵ	A000108	A000007
UD	A000245	A000012
$UUDD$	A000344	A011782
$UDUD$	A099376	A000027

W	$out[W]$	$in[W]$
$UUUDDD$	A000588	A001519
$UUDUDD$	A003517	A001792
$UUDDUD$	A003517	A000079
$UDUUDD$	A003517	A000079
$UDUDUD$	A000344	A000217
$UUDUUDDD$	A003518	A061667

Conclusion

We have a clean and simple construction of the free bifibration over a functor.

An application of proof theory, w/complementary algebraic & topological perspectives.

Some surprisingly rich combinatorics emerges as if out of thin air.

Dziękuję!