Equality is typable in Semi-Full Pure Type Systems

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Definitions

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  - *Ax* is used to type sorts .
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- Reduction :

$$(\lambda x^A.M)\ N \stackrel{\beta}{\to} M[N/x] + \text{congruences}$$

PTS typing rules

$$\frac{}{\emptyset_{wf}} \qquad \frac{\Gamma \vdash A : s \quad x \notin Dom(\Gamma)}{(\Gamma, x : A)_{wf}} \qquad \frac{\Gamma_{wf} \quad (s,t) \in \mathcal{A}x}{\Gamma \vdash s : t} \qquad \frac{\Gamma_{wf} \quad \Gamma(x) = A}{\Gamma \vdash x : A}$$

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$$\frac{\Gamma \vdash A : s \qquad \Gamma, x : A \vdash B : t \qquad (s,t,u) \in \mathcal{R}el}{\Gamma \vdash \Pi x^A.B : u}$$

$$\frac{\Gamma \vdash M : \Pi x^A.B \qquad \Gamma \vdash N : A}{\Gamma \vdash MN : B[x/N]} \qquad \frac{\Gamma \vdash M : A \qquad A \stackrel{\beta}{\equiv} B \qquad \Gamma \vdash B : s}{\Gamma \vdash M : B}$$

Some known Type Systems

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- Calculus of Constructions:
  $S = \{Prop, Type\} \quad Ax = \{(Prop, Type)\}$
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Some special classes of PTS

- Functional: If $(s, t) \in \mathcal{A}x$ and $(s, t') \in \mathcal{A}x$ then $t = t'$.
  If $(s, t, u) \in \mathcal{R}el$ and $(s, t, u') \in \mathcal{R}el$ then $u = u'$.

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  $\hookrightarrow$ In those PTS, "any" products is typable.
- Semi-full PTS: If $(s,t,u) \in \mathcal{R}el$ then for all t', there is u' such that $(s,t',u') \in \mathcal{R}el$.
  $\hookrightarrow$ If the product $\Pi x^A.B$ is typable, then for any B' well-typed, $\Pi x^A.B'$ is also well-typed (or Π-functionality).

Facts about PTS

- Inversion lemmas :

e.g. if $\Gamma \vdash \lambda x^A.M : T$ then there are s, t, u and B such that

- $(s, t, u) \in \mathcal{R}el,\ T \overset{\beta}{\equiv} \Pi x^A.B$
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- Correctness of types :

If $\Gamma \vdash M : T$ then there is $s \in S$ such that $T = s$ or $\Gamma \vdash T : s$.

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If $\Pi x^A.B \stackrel{\beta}{\equiv} \Pi x^C.D$ then $A \stackrel{\beta}{\equiv} C$ and $B \stackrel{\beta}{\equiv} D$.

- Subject Reduction:

If $\Gamma \vdash M : T$ and $M \stackrel{\beta}{\rightarrow} M'$ then $\Gamma \vdash M' : T$.

Shape of types in PTS

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  - $\forall A, B, \Pi x^A.B \in Ts$
  - if $M \in Ts$, $MN \in Ts$ and $\lambda x^A.M \in Ts$
  - if $M \in Tv$,$\Gamma \vdash M : A$ and $\Gamma \vdash M : B$, then $A \stackrel{\beta}{\equiv} B$.
  - if $M \in Ts$,$\Gamma \vdash M : A$ and $\Gamma \vdash M : B$, then $A \stackrel{\beta}{\twoheadrightarrow} \Pi x_1^{U_1} ... x_n^{U_n}.s$ and $B \stackrel{\beta}{\twoheadrightarrow} \Pi x_1^{U_1} ... x_n^{U_n}.t$.

Why do we want a typed equality ?

- In the conversion rules the intermediate steps are not checked.

$$\frac{\Gamma \vdash M : A \qquad A \overset{\beta}{\equiv} B \qquad \Gamma \vdash B : s}{\Gamma \vdash M : B}$$

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$\hookrightarrow$ all this lead to the definition of PTS *with Judgmental Equality*.

PTSe typing rules (1)

$$\frac{}{\emptyset_{wf_e}} \qquad \frac{\Gamma \vdash_e A : s \quad x \notin Dom(\Gamma)}{(\Gamma, x : A)_{wf_e}} \qquad \frac{\Gamma_{wf_e} \quad (s,t) \in \mathcal{A}x}{\Gamma \vdash_e s : t} \qquad \frac{\Gamma_{wf_e} \quad \Gamma(x) = A}{\Gamma \vdash_e x : A}$$

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$$\frac{\Gamma \vdash_e A = A' : s \qquad \Gamma, x : A \vdash_e B = B' : t \qquad (s,t,u) \in \mathcal{R}el}{\Gamma \vdash_e \Pi x^A.B = \Pi x^{A'}.B' : u}$$

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PTSe typing rules (3)

$$\frac{\Gamma \vdash_e M = M' : A \qquad \Gamma \vdash_e A = B : s}{\Gamma \vdash_e M = M' : B}$$

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The Big Question

Are both systems the same ?

Easy part of the equivalence

We prove by mutual induction that

- If $\Gamma \vdash_e M : T$ then $\Gamma \vdash M : T$.
- If $\Gamma \vdash_e M = N : T$ then $\Gamma \vdash M : T$, $\Gamma \vdash N : T$ and $M \overset{\beta}{\equiv} N$.
- If Γ_{wf_e} then Γ_{wf}.

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Here we just “lose” some information, nothing complicated.

Hard Part

The other way around needs a way to "type" a β-equivalence into a judgmental equality:

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But can we ?

How do we do this ?

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- P is welltyped in PTS by *Subject Reduction.*

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$$
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 & \searrow & & \swarrow & \\
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\end{array}
$$

- P is welltyped in PTS by *Subject Reduction.*
- Is P welltyped in PTSe ?
- How do we type $M = P$ and $N = P$ in PTSe ?

The need of Subject Reduction

To do so, we need to prove that PTSe have the *Subject Reduction* property:

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- for semi-full and full PTS : V. Siles and H. Herbelin [10] “Equality is typable in Semi-Full Pure Type Systems”.
- But the question is still open for general PTS !

Adams' approch

- In order to break the loop, Adams defined a typed version of the usual parallel β-reduction, called *Typed Parallel One Step Reduction* (TPOSR).

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  - Prove that TPOSR is *Church-Rosser*.
  - Prove that TPOSR has *Subject-Reduction*.
  - Prove that TPOSR is equivalent to PTS and PTSe.

TPOSR typing rules (1)

$$\frac{}{\emptyset_{wf}} \qquad \frac{\Gamma \vdash A \rhd A' : s \quad x \notin Dom(\Gamma)}{(\Gamma, x : A)_{wf}} \qquad \frac{\Gamma_{wf} \quad (s,t) \in \mathcal{A}x}{\Gamma \vdash s \rhd s : t} \qquad \frac{\Gamma_{wf} \quad \Gamma(x) = A}{\Gamma \vdash x \rhd x : A}$$

$$\frac{\Gamma \vdash A \rhd A' : s \qquad \Gamma, x : A \vdash B \rhd B' : t \qquad (s,t,u) \in \mathcal{R}el}{\Gamma \vdash \Pi x^A.B \rhd \Pi x^{A'}.B' : u}$$

$$\frac{\begin{array}{c}\Gamma \vdash A \rhd A' : s \\ \Gamma, x : A \vdash B \rhd B' : t \qquad \Gamma, x : A \vdash M \rhd M' : B \qquad (s,t,u) \in \mathcal{R}el\end{array}}{\Gamma \vdash \lambda x^A.M \rhd \lambda x^{A'}.M' : \Pi x^A.B}$$

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TPOSR typing rules (2)

$$\frac{\Gamma \vdash A \rhd A' : s \qquad \Gamma, x : A \vdash B \rhd B' : t \qquad \Gamma, x : A \vdash M \rhd M' : B \qquad \Gamma \vdash N \rhd N' : A \qquad (s, t, u) \in \mathcal{R}el)}{\Gamma \vdash (\lambda x^A . M)_{(x)B} N \rhd M'[x/N'] : B[x/N]}$$

$$\frac{\Gamma \vdash M \rhd N : A \qquad \Gamma \vdash A \rhd B : s}{\Gamma \vdash M \rhd N : B}$$

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We do not keep track of the sort (it requires *Type Uniqueness*).

From TPOSR to PTS and PTSe

Let's consider the $|\ |$ function that removes all annotations on applications, we can easily prove the following lemmas:

From TPOSR to PTS

If $\Gamma \vdash M \rhd N : T$ then $|\Gamma| \vdash |M| : |T|$ $|\Gamma| \vdash |N| : |T|$ and $|M| \stackrel{\beta_{//}}{\rightarrow} |N|$.

From TPOSR to PTSe

If $\Gamma \vdash M \rhd N : T$ then $|\Gamma| \vdash |M| = |N| : |T|$.

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From TPOSR to PTSe

If $\Gamma \vdash M \rhd N : T$ then $|\Gamma| \vdash |M| = |N| : |T|$.

As easy as before by induction, we just remove some information in the derivations.

First step: Church-Rosser

To prove the TPOSR is *Church-Rosser*, we will prove that the *Diamond Property* holds for TPOSR.

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- we only have an hypothesis "$\Gamma, x : C \vdash B \rhd B' : s$"
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$\hookrightarrow$ So we need a way to equal A and C.

Functional vs Semi-Full

- For any functional TPOSR system, *Uniqueness of Types* holds, so we can prove that $\Gamma \vdash A \equiv C$.

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By applying the previous lemma to N''':

- first case : $\Gamma \vdash A \equiv C$
- second case : A and C only differ by their last sort s and t

Back to the Untyped World, second case

If we erase all the equalities we have so far, by *untyped Confluence* we can conclude that:

- $\Pi x^{|A|}.|B| \stackrel{\beta}{\equiv} \Pi x^{|C|}.|B|$.

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With this, we can now finish to prove everything up to *Subject Reduction*

Validity of Annotations

To close the equivalence, we need to prove that the additional annotations on applications did not change the typing system, that is:

Validity of Annotations

If $\Gamma \vdash M : T$, then $\Gamma^* \vdash M^* \triangleright M^* : T^*$
(for all Γ^*, M^*, T^* such than $|\Gamma^*| = \Gamma$, $|M^*| = M$ and $|T^*| = T$).

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Since there are several ways to annotate a term, the induction can be quite tricky without the following lemma:

Erased Context Conversion

If $\Gamma_1 \vdash M \rhd N : A$, $|\Gamma_1| = |\Gamma_2|$ and $\Gamma_2\ _{wf}$, then $\Gamma_2 \vdash M \rhd N : A$.

Erased Conversion: the second pitfall

To prove this conversion lemma, we need a more general lemma which is easily done for functional PTS, but strangely hard for semi-full:

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If $\Gamma \vdash M \rhd ? : S$, $\Gamma \vdash N \rhd ? : T$ and $|M| = |N|$, then there is P such that:

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$$\begin{array}{ll} |M| = |M'| & |N| = |N'| \\ \Gamma \vdash M \rhd^+ M_0 : \Pi x^A.B & \Gamma \vdash M' \rhd^+ M_0 : \Pi x^{A'}.B' \\ \Gamma \vdash N \rhd^+ N_0 : A & \Gamma \vdash N' \rhd^+ N_0 : A' \\ \Gamma \vdash M_{(x)B} N \rhd ? : B[x/N] & \Gamma \vdash M'_{(x)B'} N' \rhd ? : B'[x/N'] \end{array}$$

Shape of Terms

Shape of Types in TPOSR

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What does it means to be typed by a telescope ?

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(Very Simplified) Shape of Terms in TPOSR in Ts

If $\Gamma \vdash M \rhd ? : \Pi x_1^{U_1} ... x_n^{U_n}.s$ then $\Gamma \vdash M \rhd^+ \lambda x_1^{U_1} ... x_n^{U_n}.P : \Pi x_1^{U_1} ... x_n^{U_n}.s$ and $\Gamma, x_1 : U_1, ..., x_n : U_n \vdash P \rhd P : s$

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By combining Shape of Types and Terms, we can prove that $n > 1$ in our problematic case, thus we can erase the troublesome annotation by performing a β-reduction step first.

Solution to the pitfall

$$
\begin{array}{llll}
\Gamma \vdash M_{(x)B} N & \rhd^+ & M_0\ {}_{(x)B} N & : B[x/N] \\
 & \rhd^+ & (\lambda x^C \lambda\Delta.K)_{(x)B} N & : B[x/N] \\
 & \rhd^+ & \lambda\Delta[x/N].K[x/N] & : B[x/N] \\
 & \rhd^+ & \lambda\Delta[x/N_0].K[x/N_0] & : B[x/N] \\
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\end{array}$$

With *Subject Reduction* and *Validity of Annotations*, we are know able to prove that PTS $\Rightarrow$ TPOSR, and so:

$$\text{PTSe} \Rightarrow \text{PTS} \Rightarrow \text{TPOSR} \Rightarrow \text{PTSe}$$

Possible Extensions of the Proof

They are several ways to enhance the system:

- Change the conversion rule (with η for example).
- Extend the conversion rule with cumulativity : the road to subtyping.

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Possible solutions: adding *Strenghthening* as a primitive rule as in *ICC*, restrict to normalizing systems, only add η-expansion...

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Adding cumulativity for Π-types and sorts requires an odd lemma before being able to prove the *Shape of Types* property (so even far before Π-*injectivity*) which has resisted all attempts until now:

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If $\Gamma \vdash \Pi x_1^{U_1} \ldots x_n^{U_n}.s \equiv \Pi x_1^{V_1} \ldots x_n^{V_n}.s$ then for all t,
$\Gamma \vdash \Pi x_1^{U_1} \ldots x_n^{U_n}.t \equiv \Pi x_1^{V_1} \ldots x_n^{V_n}.t$.

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$\Gamma \vdash \Pi x_1^{U_1} \ldots x_n^{U_n}.t \equiv \Pi x_1^{V_1} \ldots x_n^{V_n}.t$.

However, even if we manage to prove this, the approach used to proof *Validity of Annotations* do not scale to subtyping, so a new way to prove it still needs to be found.

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+ A more precise proof of *Church-Rosser* for TPOSR which works for all useful PTS.

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- Dealing with η-conversion is still the same nightmare
- Subtyping forces us to throw away the *Shape of Types* approach to *Validity of Annotations* and redo it from scratch.

Conclusion: Next step ?

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Thank you for your time. Any questions ?