#### Arbres, cartes et nombres de Hurwitz

GILLES SCHAEFFER CNRS & École Polytechnique

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Plan de l'exposé

Revêtements ramifiés et cartes

Cartes et arbres

Énumération d'arbres et formule d'Hurwitz

Revêtements et cartes aléatoires

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Revêtements et cartes aléatoires

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A mapping  $\phi : \mathcal{D} \to \mathcal{I}$  is a covering if, for all x in  $\mathcal{I}$  there exists  $n \geq 1$  and a neighborhood V of x such that  $\phi^{-1}(V) \sim B \times \{1, \ldots, n\}$ ,

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What is we try to extend from  $A_r$  to B?



Recall  $\phi_k : A_r \to A_{r^k}$  with  $\phi_k(z) = z^k$ . Extend from  $A_r$  to B? The mapping  $\phi_k : B^* \to B^*$  is a covering, but not  $\phi_k : B \to B$ .



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A mapping  $\phi$  is a ramified covering of S by S if there exists a finite subset  $X = \{x_1, \ldots, x_p\}$  such that:

- $\phi_{\mathbb{S} \setminus \phi^{-1}(X)}$  is a covering, and
- $\phi$  is ramified over each  $x_i$



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 $\lambda^{(1)} = 1^5$   $\lambda^{(2)} = 1, 2^2$   $\lambda^{(2)} = 2, 3$ 



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# Ramified coverings of the sphere by itself (Cont'd)



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  - $\Rightarrow \text{The partitions } \lambda^{(i)}$ are partitions of n, degree of the covering.

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coverings with only simple branch points

Today's topic



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**Theorem.** Simple ramified covers of S by itself with m ramifications points are in bijection with increasing labelled quadrangulations with m faces.
### Résumé du 1er épisode

# Compter des classes d'équivalence de revêtements ramifiés compter certaines plongements de graphes

Plan de l'exposé

# Revêtements ramifiés et cartes

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Proof?



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The dual map of a map is the map of incidence between faces.

The dual map of a tree-rooted map is a tree-rooted map: it is naturally endowed with a dual spanning tree.

Proof?

Euler's relation: (#vertices-1)+(#faces-1) = #edges

A planar map is a proper embedding of a connected graph on the sphere (considered up to homeomorphisms).

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The number of tree rooted planar maps with n edges is  $\sum_{i=0}^{n} {\binom{2n}{i}} C_i C_{n-i}$  where  $C_n = \frac{1}{n+1} {\binom{2n}{n}}$  denotes Catalan numbers, counting balanced parenthesis words.



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Let us recycle the idea used for tree-rooted maps, using a canonical spanning tree



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Then write the code of the primal tree on the chosen canonical tree

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The map is recovered from the code by *closure*.

Our code of the map will be a canonical decorated tree

Question is How do we choose the canonical spanning tree so that the resulting decorated trees can be described and counted ?

Orient the tree edges away from the root



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Orient the other edges couterclockwise around the tree



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The resulting orientation has no clockwise circuit.



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The tree is recovered by reconstructing its contour .

Idea:

Choose a minimal accessible orientation to get a spanning tree

Our pb becomes:

How to choose a canonical accessible minimal orientation?

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**Fact:** For many subclasses  $\mathcal{F}$  of planar maps, there exists an  $\alpha_{\mathcal{F}}$  s.t.:

A planar map is in  $\mathcal{F}$  if and only if it admits an  $\alpha_{\mathcal{F}}$ -orientation.

#### $\alpha\text{-orientations}$ for increasing quadrangulations

Recall increasing quadrangulations are planar maps with faces of degree 4 such that:

- faces have labels in  $\{1, \ldots, 2n-2\}$
- around labeled vertices, face labels increase in ccw order
- around white vertices, face labels increase in cw order


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This is our choice of canonical  $\alpha$  to decompose increasing quadrangulations.

















**Proposition.** The resulting simple Hurwitz trees has n unlabelled vertices, n-1 labeled vertices of degree 2, 2n-2 edges that increase ccw around labeled vertices.

#### From simple Hurwitz trees to increasing quadrangulations

A local rule to create increasing half edges



Two half-edges with same label  $\Rightarrow$  edge and face of degree 4



Iterate the local rules as long as possible...









for the bijection







vertex label are useless for the bijection

adding buds

Parings and adding buds again



8

12

3

again





adding buds

Parings and adding buds again





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Theorem[Duchi-Poulalhon-S. 2012] Closure is the reverse bijection between

- simple Hurwitz trees of size n, and
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#### Résumé des 2 premiers épisodes

# Compter des classes d'équivalence de revêtements ramifiés compter certaines plongements de graphes compter certains arbres

Plan de l'exposé

Plan de l'exposé

Revêtements ramifiés et cartes

Cartes et arbres

Énumération d'arbres et formule d'Hurwitz

Revêtements et cartes aléatoires











**Theorem**[Duchi-Poulalhon-S. 2012] Increasing quadrangulations (size n) are in bijection with simple Hurwitz trees having n unlabelled vertices, n-1 labeled vertices of degree 2, 2n-2 edges that increase ccw around labeled vertices.



The number of simple ramified cover of S by itself with m = 2n - 2 critical points is  $n^{n-3}(2n-2)!$ .

#### Hurwitz formula for factorizations in transpositions

**Theorem.** Let  $\lambda = 1^{\ell_1}, \ldots, n^{\ell_n}$  be a partition n, and  $\ell = \sum_i \ell_i$ . The number of *m*-uples of transpositions  $(\tau_1, \ldots, \tau_m)$  such that

- (product cycle type)  $\tau_1 \cdots \tau_m = \sigma$  has cycle type  $\lambda$
- (transitivity) the associated graph is connected
- (minimality) the number of factors is  $m = n + \ell 2$

IS 
$$n^{\ell-3} \cdot m! \cdot n! \cdot \prod_{i \ge 1} \frac{1}{\ell_i!} \left(\frac{i^i}{i!}\right)^{\ell_i}$$

#### **Proofs:**

(Hurwitz 1891, Strehl 1996) (Goulden–Jackson 1997) (Lando–Zvonkine 1999) (Bousquet-Mélou–Schaeffer 2000) (recurrences, Abel identities) (gfs and differential eqns) (geometry of LL mapping) (bijection + inclusion/exclusion)

$$\lambda = n$$
, factorizations of *n*-cycles:  $n^{n-2} \cdot (n-1)!$   
 $\lambda = 1^n$ , factorizations of the identity:  $n^{n-3} \cdot (2n-2)$ 

#### A formula for general factorizations [BMS00]

**Theorem.** Let  $\lambda = 1^{\ell_1}, \ldots, n^{\ell_n}$  be a partition of n, and  $\ell = \sum_i \ell_i$ . The number of *m*-uple of permutations  $(\sigma_1, \ldots, \sigma_m)$  such that

- (factorization)  $\sigma_1 \cdots \sigma_m = \sigma$  with cycle type  $\lambda$
- (transitivity)  $\langle \sigma_1, \ldots, \sigma_m \rangle$  acts transitively on  $\{1, \ldots, n\}$
- (minimality) the total rank of factors is  $\sum_i r(\sigma_i) = n + \ell 2$

$$m\frac{((m-1)n-1)!}{(mn-(n+\ell-2))!}\cdot n!\cdot \prod_{i}\frac{1}{\ell_{i}!}\binom{mi-1}{i}$$

 $\ell_i$ 

#### **Proofs:**

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(Bousquet-Mélou–Schaeffer 2000) (Goulden–Serrano 2009) (bijection + inclusion/exclusion)(gfs and differential eqns)

 $\lambda = n$ , factorizations of *n*-cycles:  $\frac{1}{(mn+1)} \binom{mn+1}{n} \cdot (n-1)!$ 

$$\lambda = 1^n$$
, identity factorizations:  $\frac{m}{(m-2)n+2} \frac{(m-1)^{n-1}}{(m-2)n+1} \binom{(m-1)n}{n} \cdot (n-1)!$ 

#### Résumé des 3 premiers épisodes

# Compter des classes d'équivalence de revêtements ramifiés compter certaines plongements de graphes compter certains arbres

les formules simples appellent des preuves constructives
Plan de l'exposé

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#### Quadrangulations croissantes aléatoires uniformes

 $\bar{Q}_n = \{$ quadrangulations croissantes à n faces $\}$ .

Quadrangulation croissante uniforme = variable aléatoire  $Q_n$  à valeur dans  $\overline{Q}_n$  avec

$$\Pr(Q_n = q) = \frac{1}{|\bar{\mathcal{Q}}_n|} = \frac{1}{n^{n-3}(2n-2)!} \qquad \text{pour tout } q \in \bar{\mathcal{Q}}_n$$

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• le choix de la distribution uniforme combinatoire est le plus immédiat

Parallèle naturel avec la distribution uniforme sur les quadrangulations enracinées:

$$\Pr(\vec{Q}_n = q) = \frac{1}{|\vec{Q}_n|} = \frac{1}{\frac{2 \cdot 3^n (2n)!}{(n+2)!n!}}$$

pour tout  $q \in ec{\mathcal{Q}}_n$ 

Comment étudier  $Q_n$  ?

# Propriétés des cartes aléatoires uniformes ?



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on est loin d'une discrétisation aléatoire d'une géométrie euclidienne

en physique on lie cela à la modélisation discrète de la gravité quantique

## Quadrangulations uniformes comme surfaces aléatoires

L'allure d'une sphère aléatoire dépend un peu de qui dessine...

**Objectif:** Choisir une métrique intrinsèque et décrire les surfaces ainsi obtenues





Distribution uniforme sur les quadrangulations à n faces, pour n grand

1ère approche: Étudier le comportement asymptotique de paramètres:

- degré d'un sommet aléatoire distance entre 2 sommets aléatoires
- loi 0-1 pour les propriétés locales longueur d'un plus petit cycle diviseur
- $\Rightarrow$  espérance, moments, lois limites discrètes ou continues, qd  $n \rightarrow \infty$

Bender, Canfield *et al* (90's  $\rightarrow$ ) en combinatoire

Ambjørn, Watabiki et al (90's  $\rightarrow$ ) en physique

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Exemple:  $\Delta_n$  = distance entre 2 sommets aléatoires uniformes de  $Q_n$ Théorème (Chassaing-S. 2004)  $\mathbb{E}(\Delta_n) \sim c \cdot n^{1/4}$ 

 $(n^{-1/4}\Delta_n) \xrightarrow{d} \max$  (serpent Brownien)

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2ème approche: Définir des surfaces aléatoires limites

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2ème approche: Définir des surfaces aléatoires limites

- convergence vers une limite d'échelle

(Pb posé au séminaire Hypathie en 2002 à Lyon)

 $\Rightarrow$  la carte Brownienne

Marckert, Mokkadem, Le Gall, Miermont, ...

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 $\Rightarrow$  la carte Brownienne

Marckert, Mokkadem, Le Gall, Miermont, ...

puis Weill, Curien, Benjamini,...

- convergence vers une limite infinie discrète

 $\Rightarrow$  la quadrangulation infinie uniforme (UIPQ)

Angel, Schramm, ...

puis Durhus, Chassaing, Krikun, Bettinelli,...

- L'excursion Brownienne décrit la limite d'échelle de toute sorte d'excursions aléatoires discrètes plus ou moins complexes.

- L'arbre continu aléatoire est limite d'échelle de toute sorte d'arbres aléatoires discrets plus ou moins complexes.

 $\Rightarrow$  On pense qu'il en est de même de la carte Brownienne.

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On dispose d'un cadre bijectif très général pour la construction de cartes par recollements d'arbres (Bernardi-Chapuy-Fusy 2011, Albenque-Poulalhon 2012) On obtient ainsi en particulier un codage d'arbres pour les revêtements... Il reste à utiliser ces constructions pour passer à la limite...