

M2 BIM/STRUCT - Lecture 1

Folding RNA *in silico*

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1

Introduction

- Dynamic programming 101
- Why RNA?
- RNA folding
- RNA Structure(s)
- Some representations of RNA structure

2

Some flavours of folding prediction

- Thermodynamics vs Kinetics
- Dynamic programming: Reminder

3

Free-energy minimization

- Nussinov-style RNA folding
- Turner energy model
- MFold/Unafold
- Performances and the comparative approach
- Towards a 3D ab-initio prediction

Foreword ...

... or how to make a million bucks by giving change parsimoniously!!

Problem: You have access to unlimited amount of **1, 20** and **50** cents coins.
A client prefers to travel light, i.e. to **minimize the #coins**.
How to give **N** cents back in change without losing a customer?

Strategy #1: Start with *heaviest* coins, and then complete/fill-up with coins of *decreasing* value.

$$21 = ??$$

$$55$$

$$60$$

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$$= \text{£}20 + \text{£}20 + \text{£}20 !$$

Problem *a priori* (?) non-solvable using such a *greedy* approach, as a (simpler) problem is already NP-complete (thus Efficient solution \Rightarrow 1M\$).

Strategy #2: Brute force enumeration $\rightarrow \#Coins^N$ (Ouch!)

Strategy #3: The following recurrence gives the minimal number of coins:

$$\text{Min}\#\text{Coins}(N) = \text{Min} \left\{ \begin{array}{l} \text{1 cent coin} \rightarrow 1 + \text{Min}\#\text{Coins}(N - 1) \\ \text{2 cent coin} \rightarrow 1 + \text{Min}\#\text{Coins}(N - 20) \\ \text{5 cent coin} \rightarrow 1 + \text{Min}\#\text{Coins}(N - 50) \end{array} \right.$$

With some memory (N intermediate computations), the minimum number of coins can be obtained after $N \times \#Coins$ operations. An actual set of coins can be reconstructing by **tracing back** the choices performed at each stage, leading to the minimum.

Remark: We still haven't won the million, as N has **exponential value compared to the length of its encoding**, so the algorithm does not qualify as *efficient* (i.e. polynomial).

Still, this approach is much more efficient than a brute-force enumeration:
 \Rightarrow Dynamic programming.

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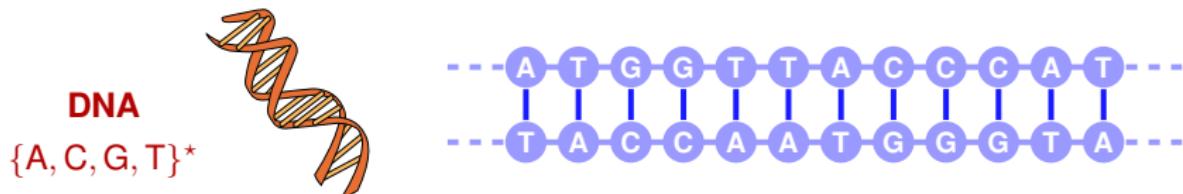
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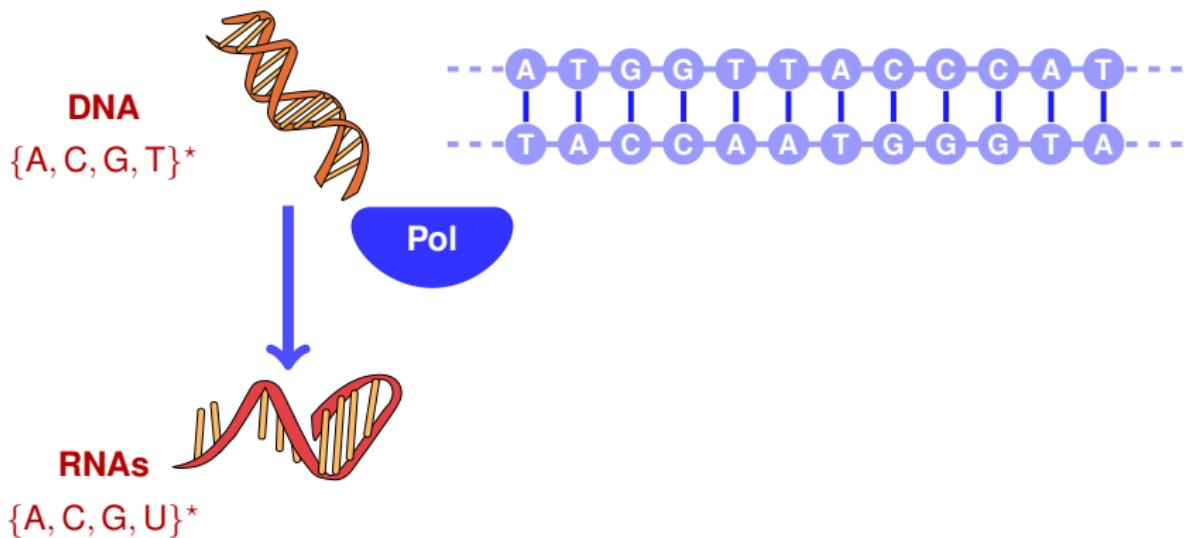
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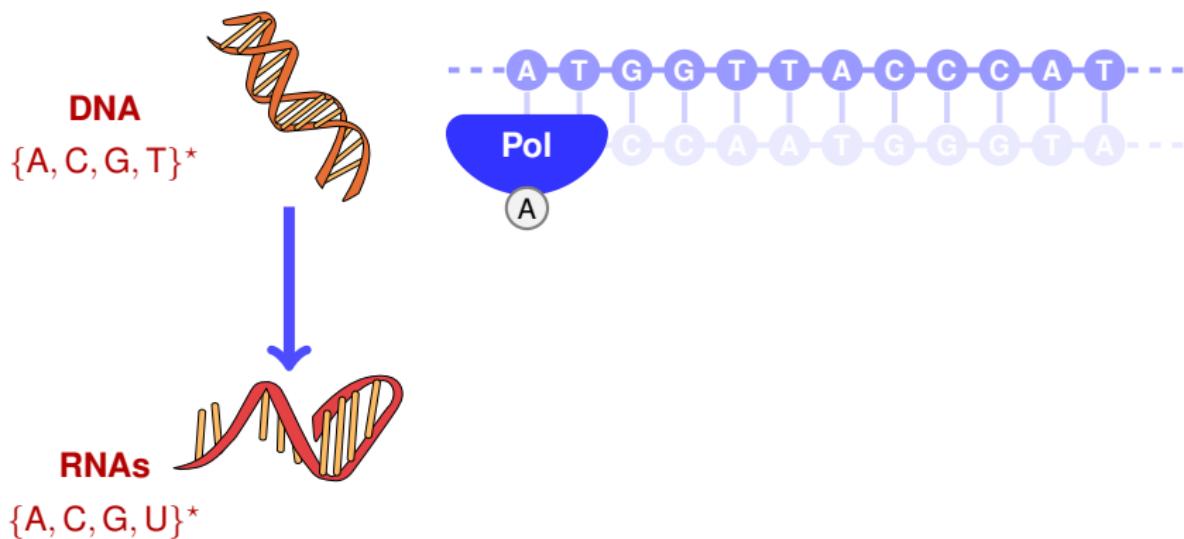
Fundamental dogma of molecular biology



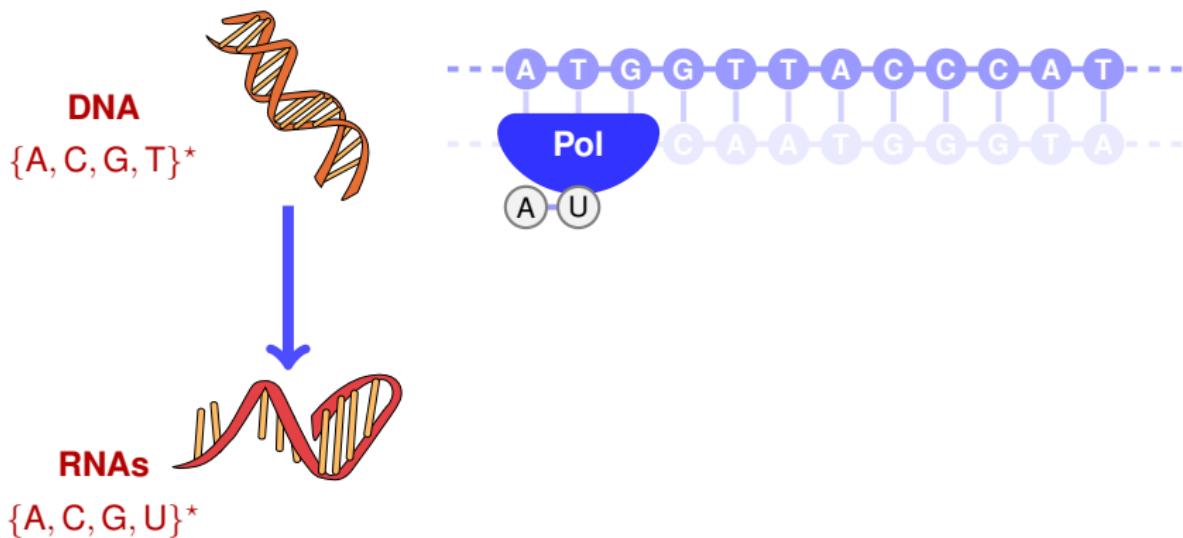
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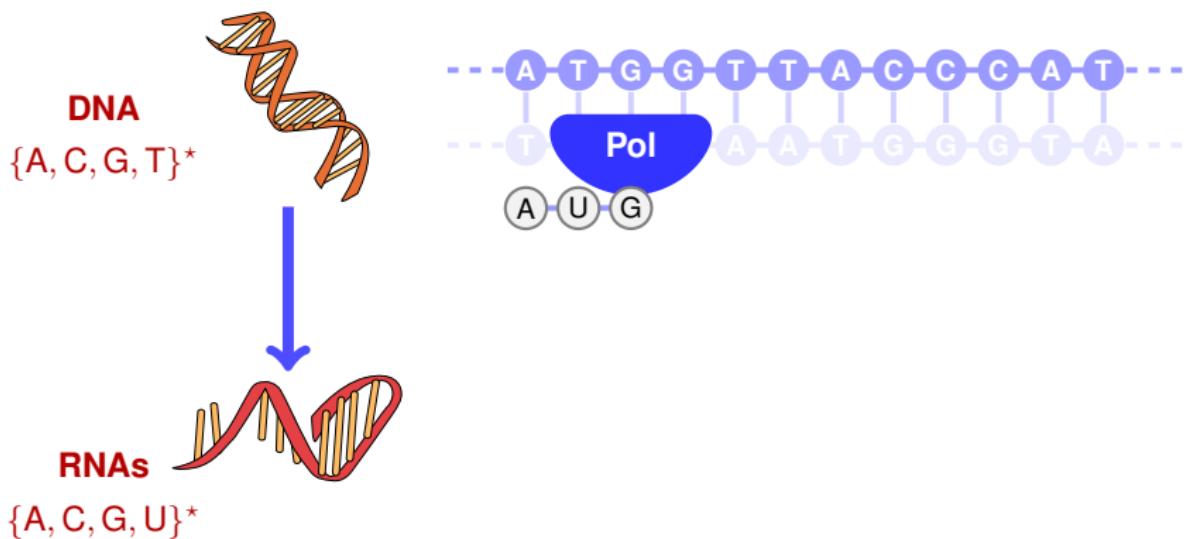
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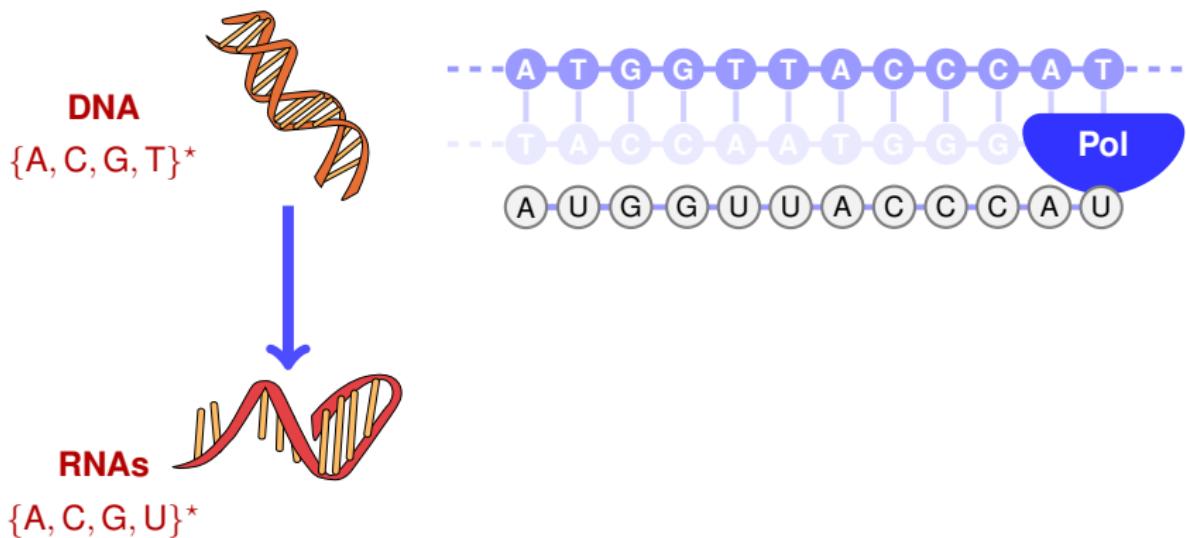
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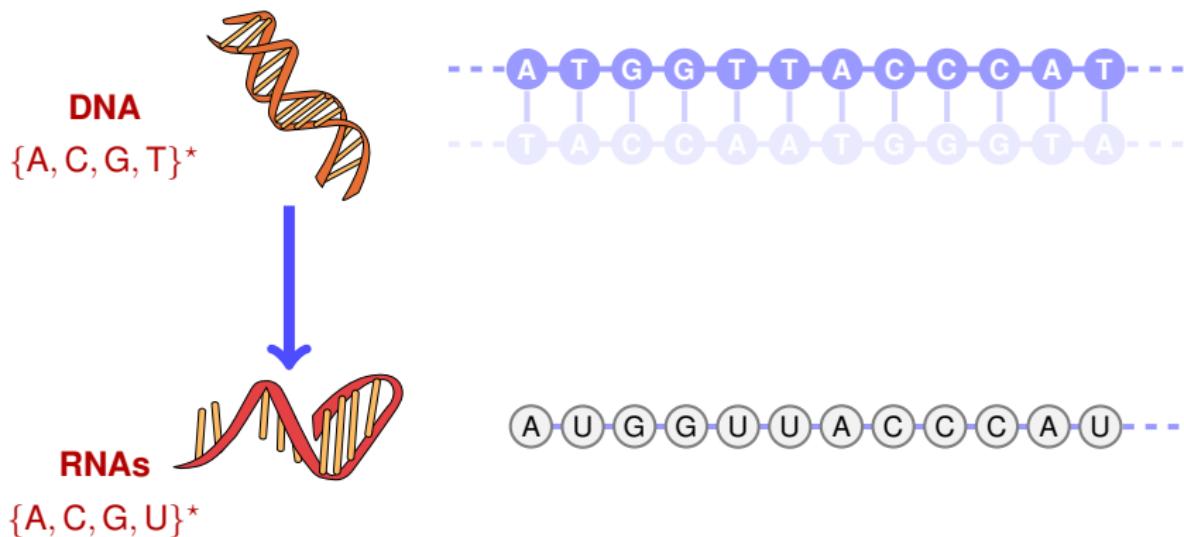
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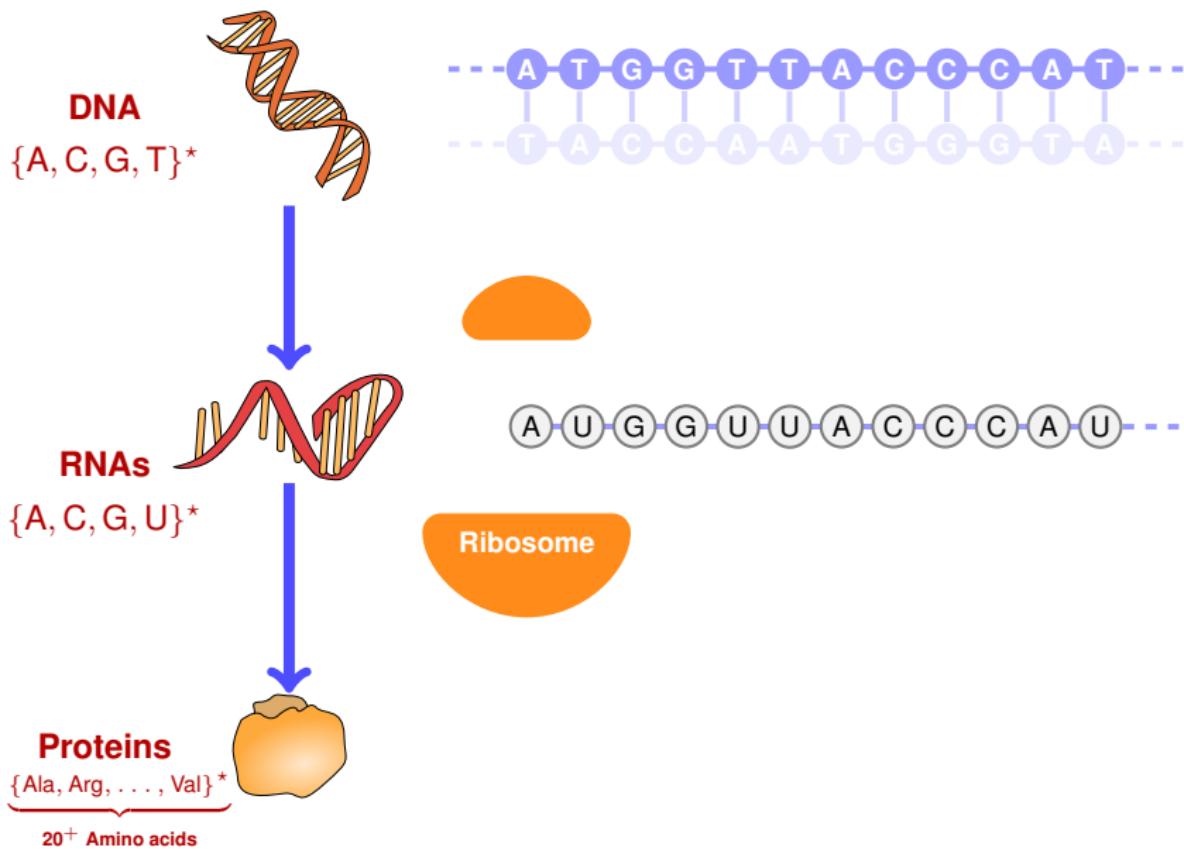
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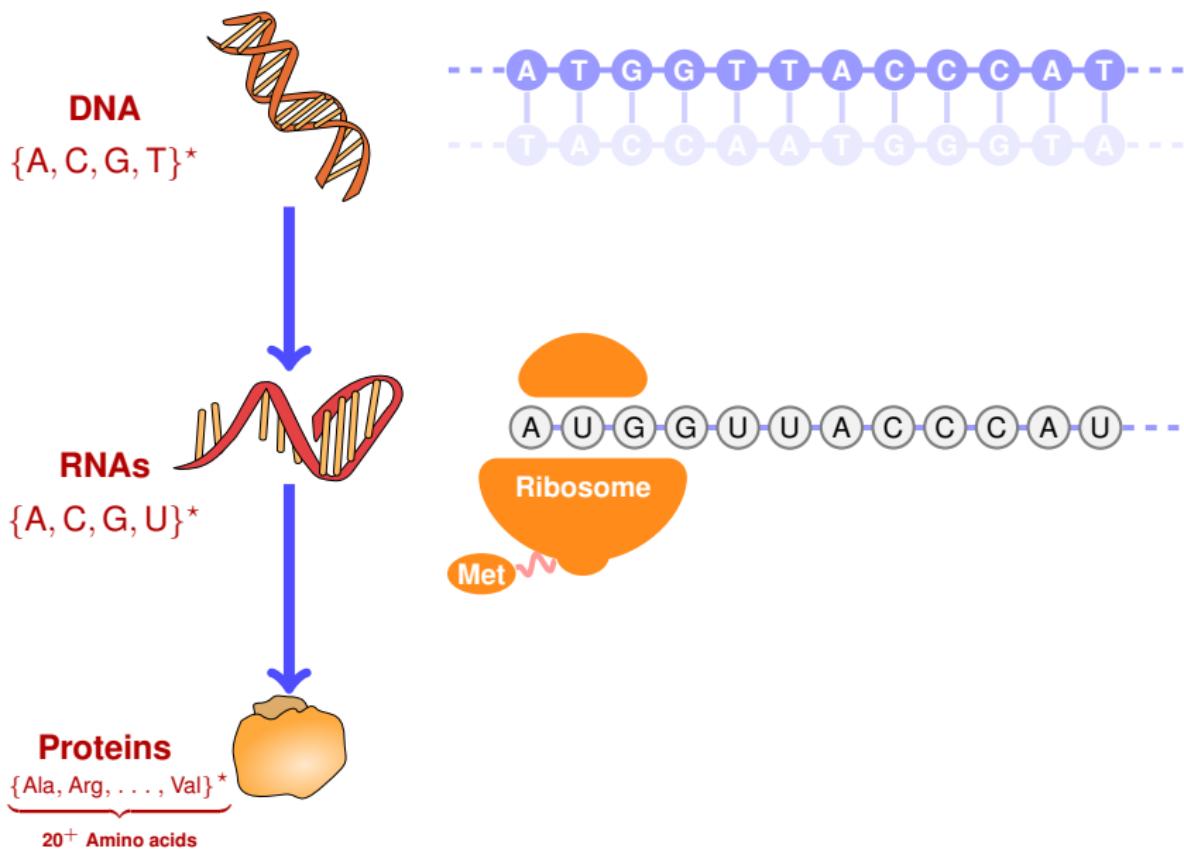
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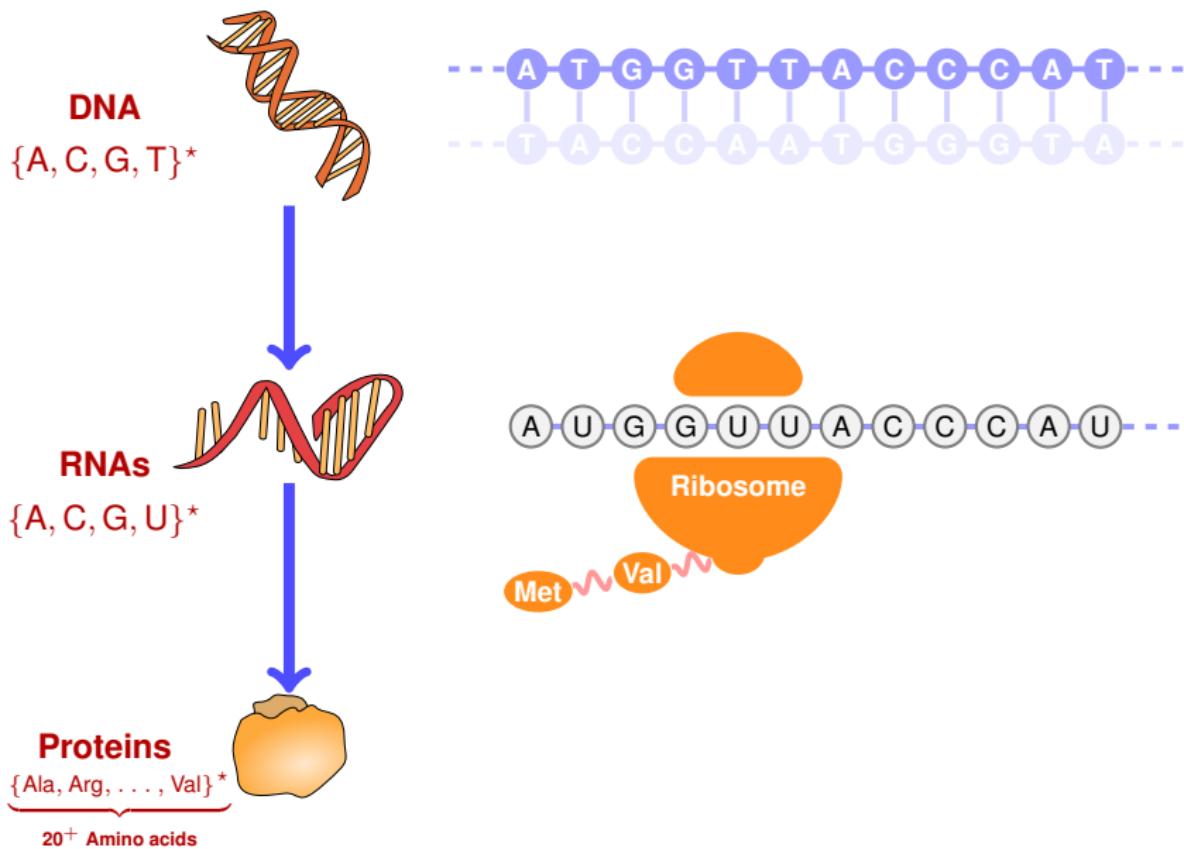
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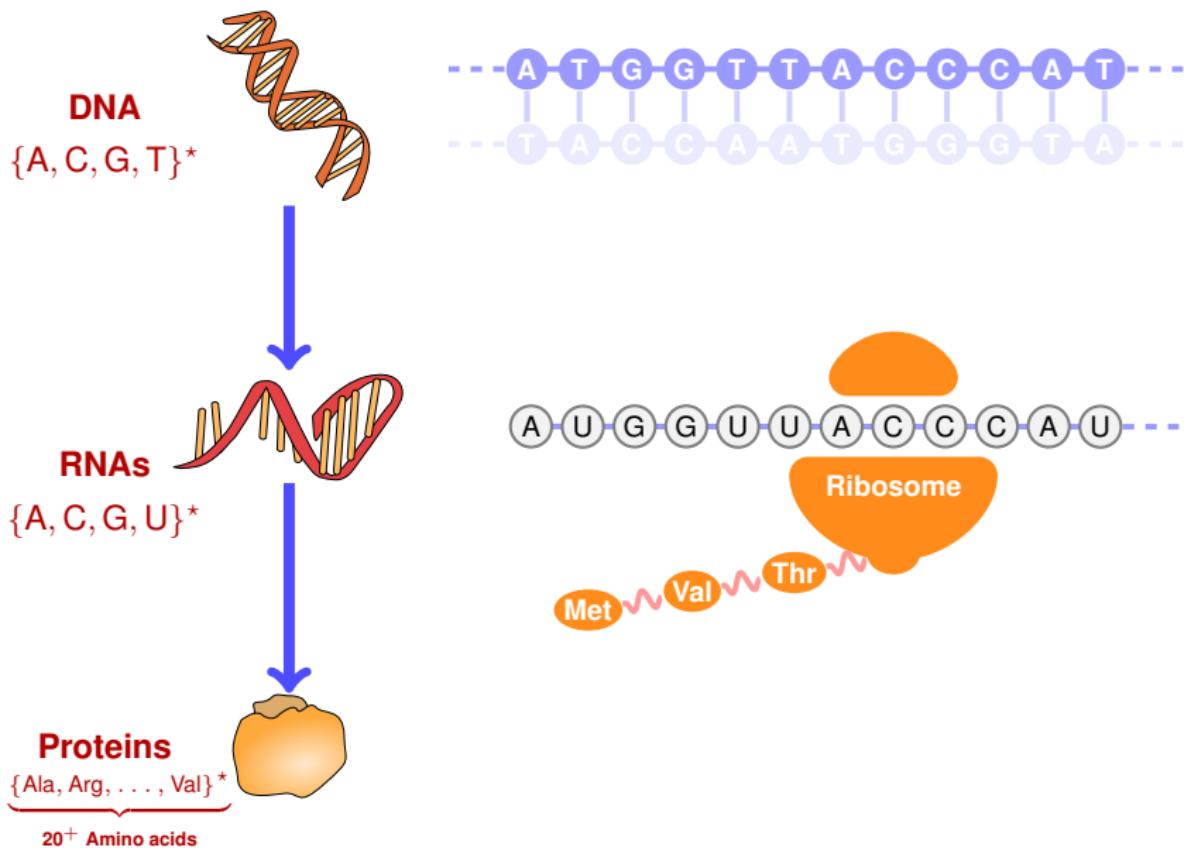
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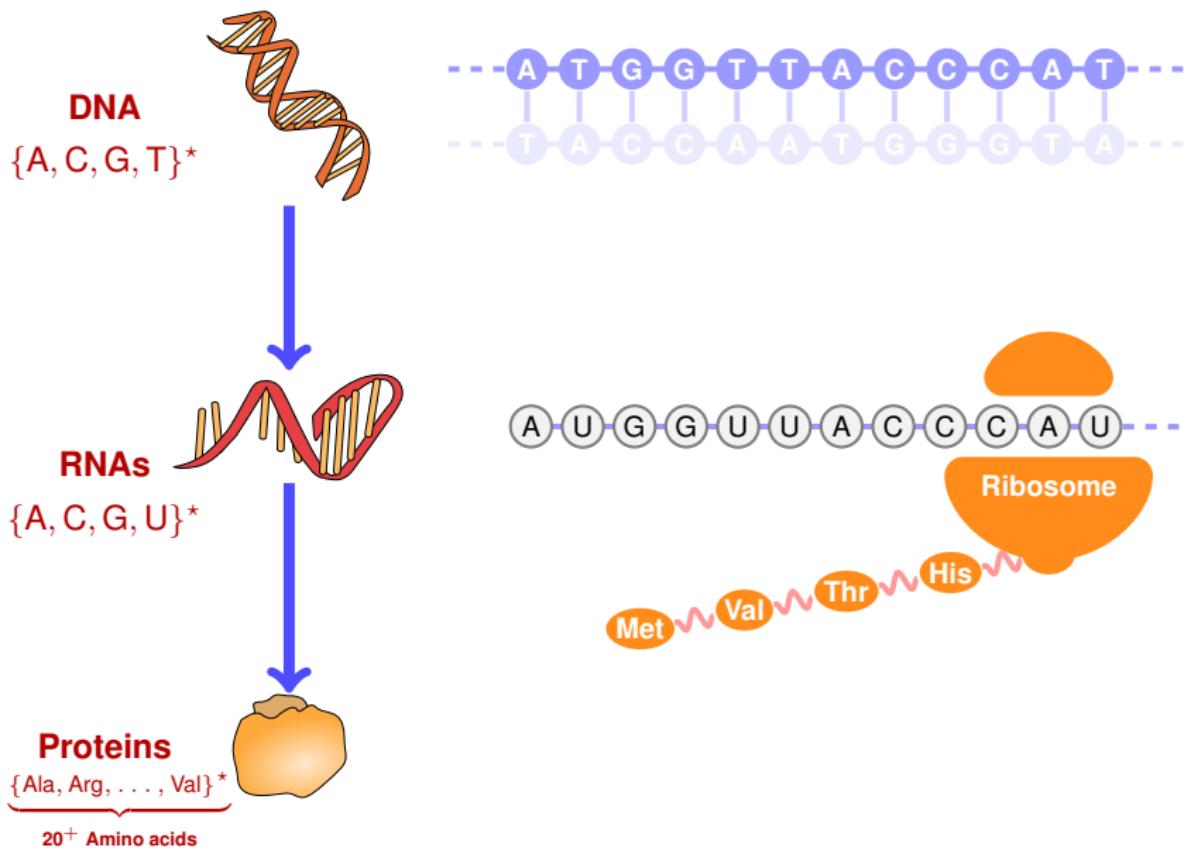
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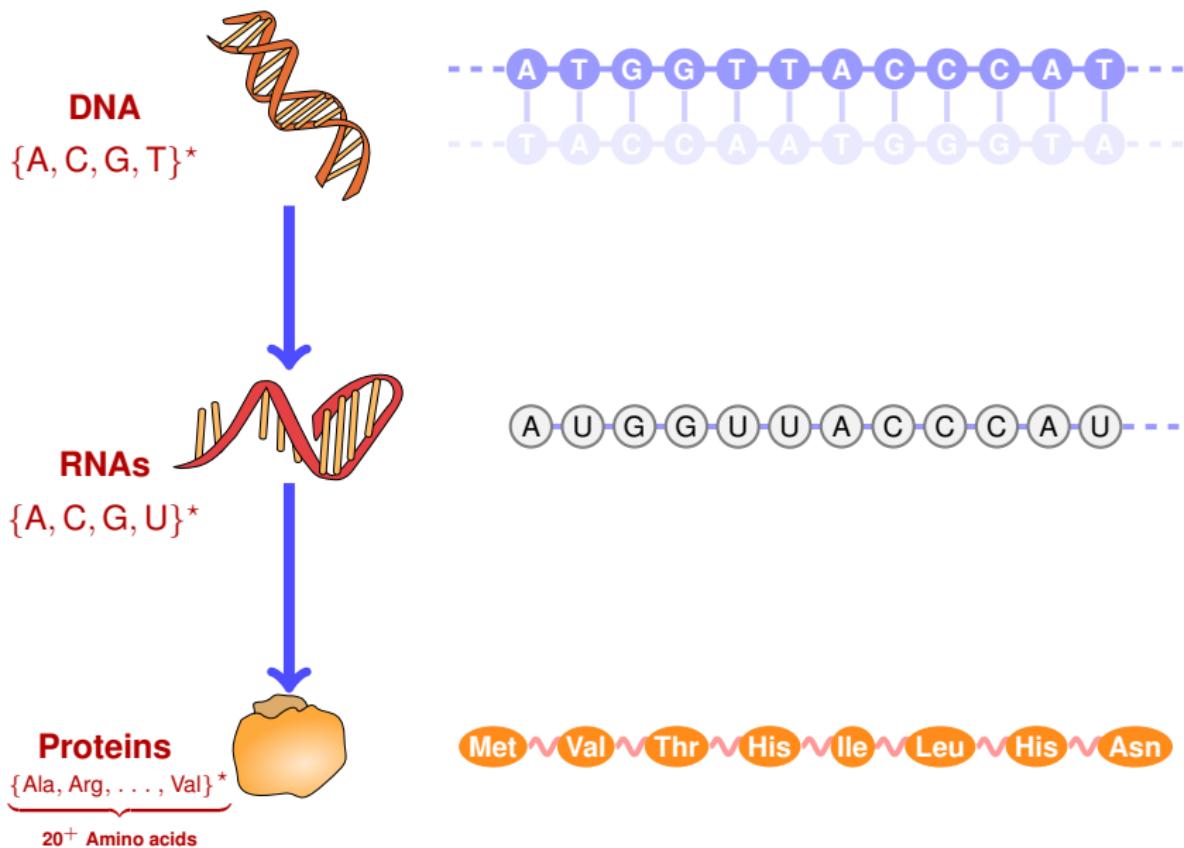
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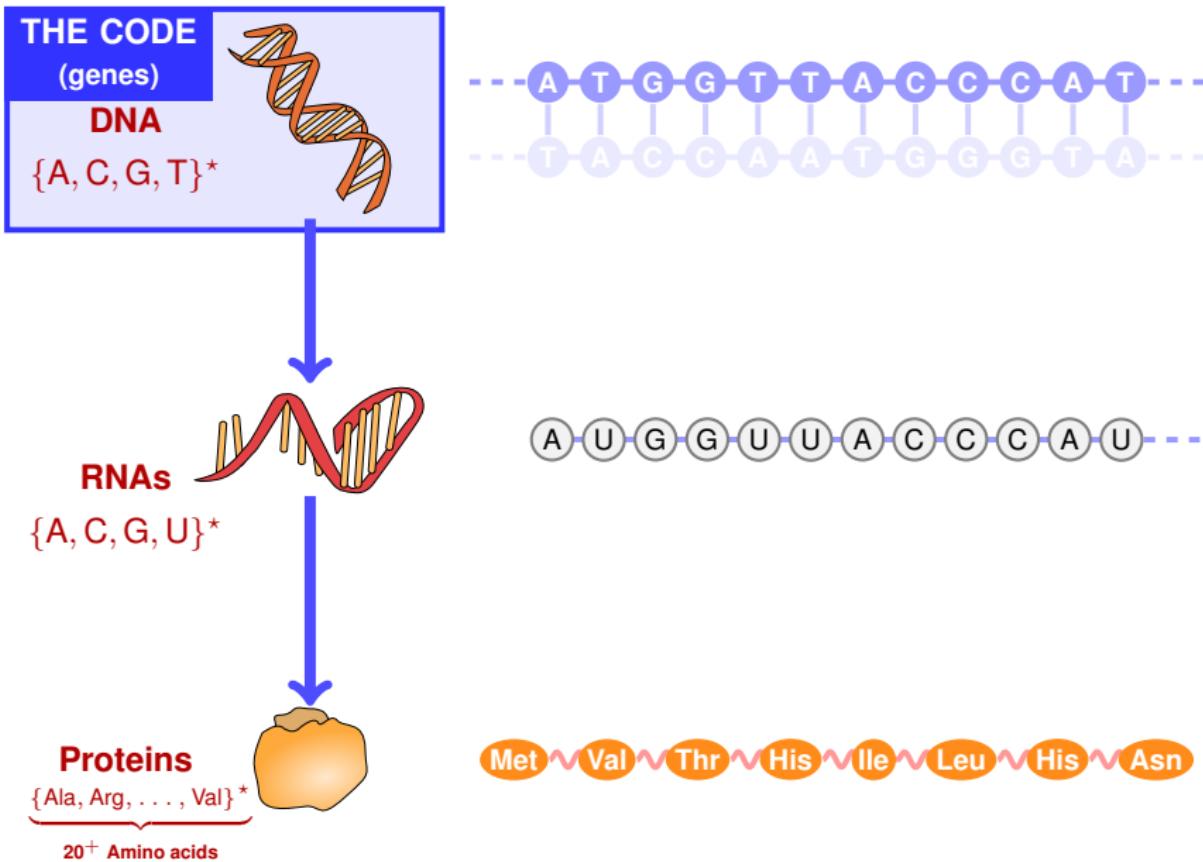
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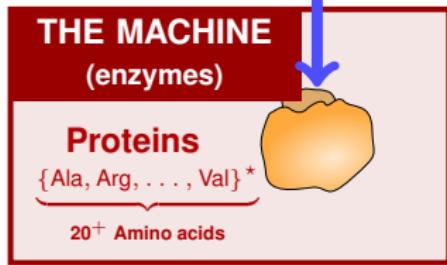
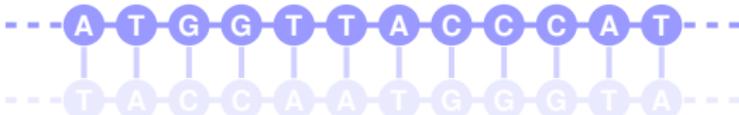
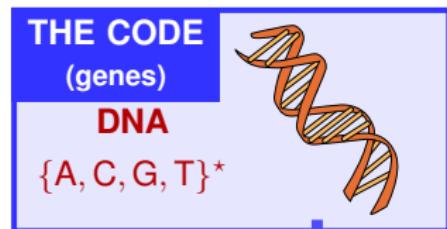
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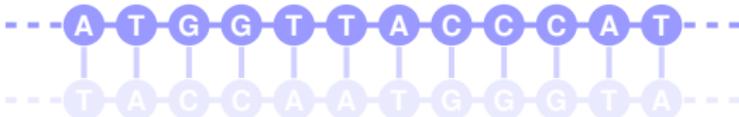
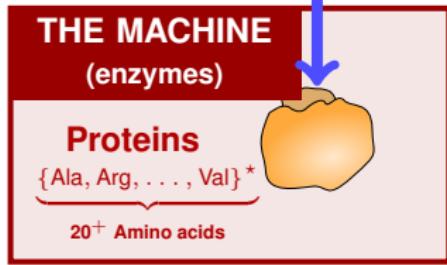
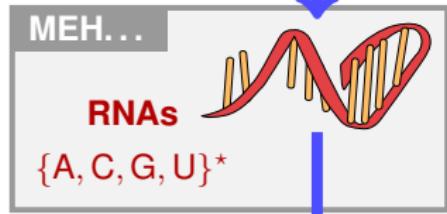
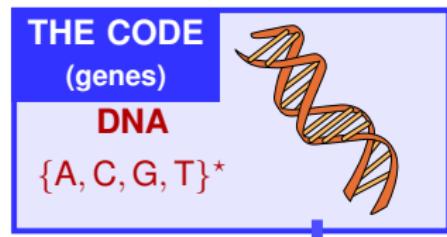
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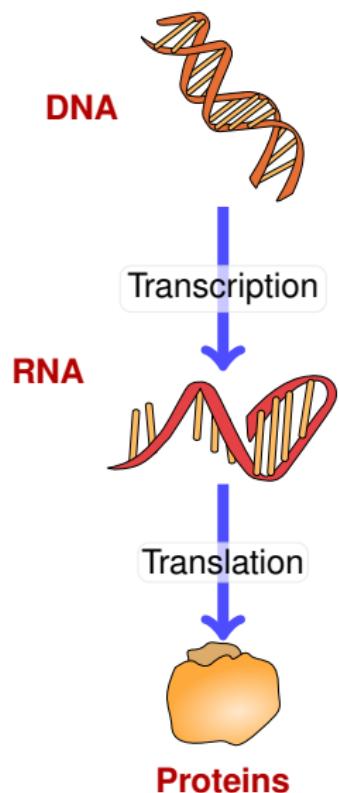
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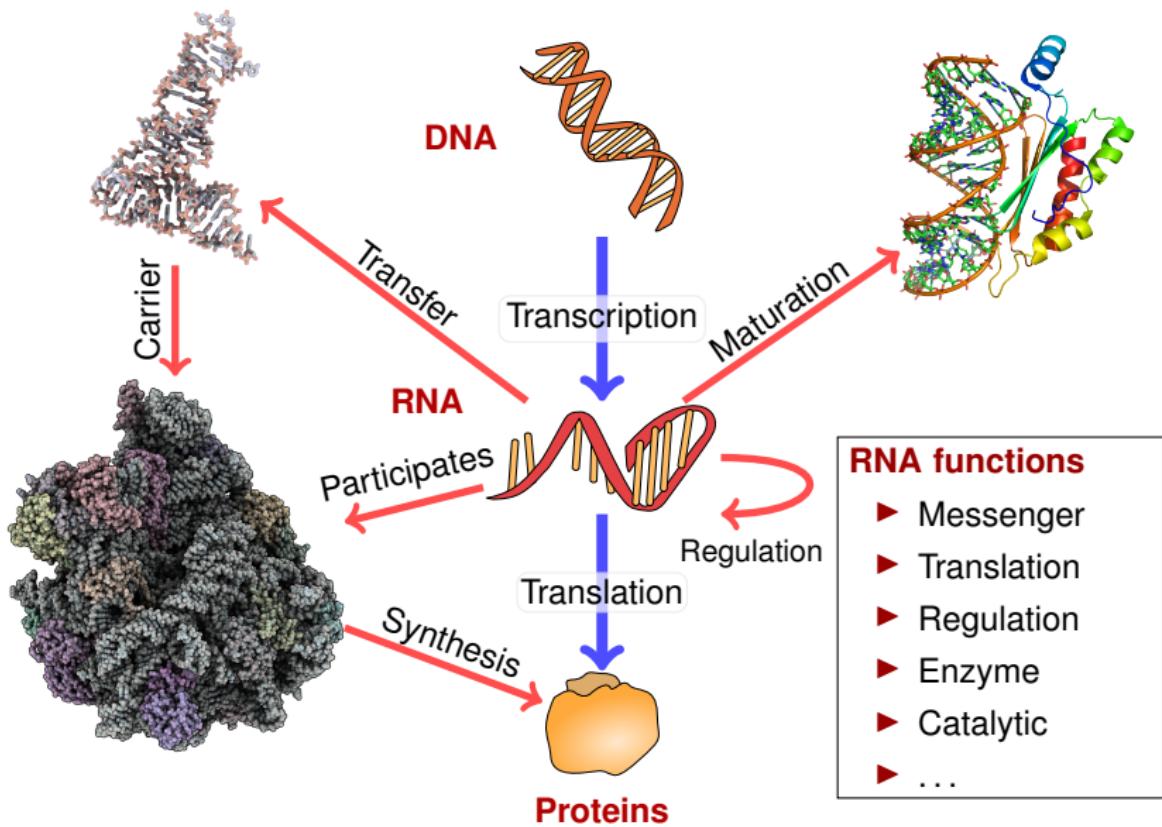
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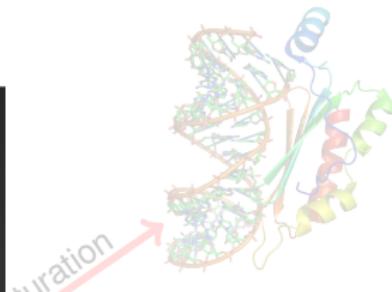
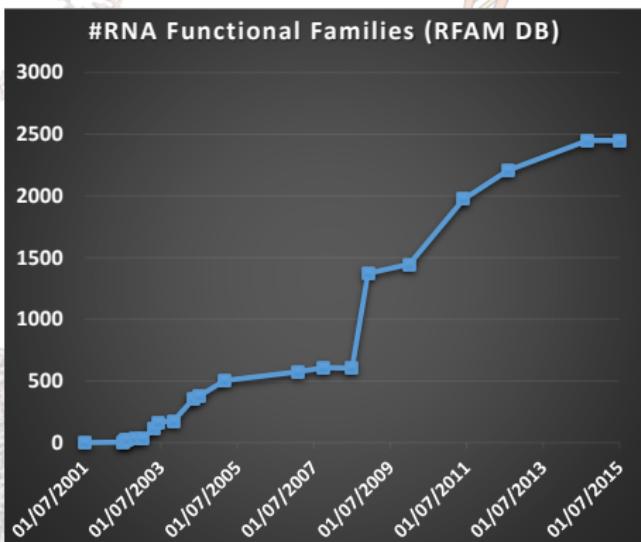
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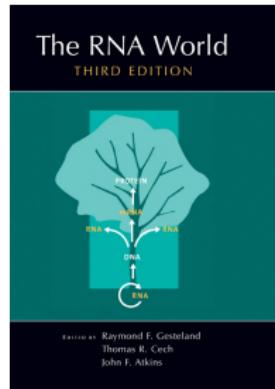
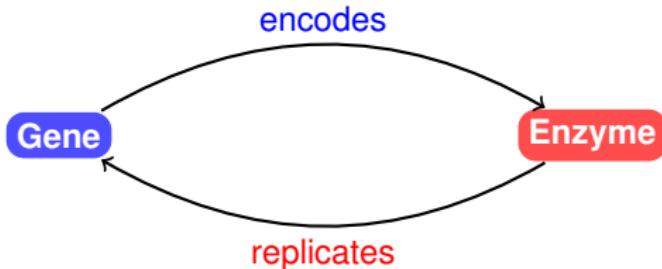


RNA functions

- ▶ Messenger
- ▶ Translation
- ▶ Regulation
- ▶ Enzyme
- ▶ Catalytic
- ▶ ...

Proteins

RNA world: Resolving the *chicken vs egg* paradox at the origin of life...

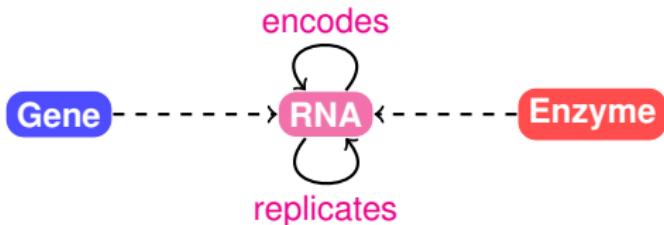


A gene big enough to specify an enzyme would be too big to replicate accurately without the aid of an enzyme of the very kind that it is trying to specify. So the system *apparently cannot get started*.

[...] This is the *RNA World*. To see how plausible it is, we need to look at why proteins are good at being enzymes but bad at being replicators; at why DNA is good at replicating but bad at being an enzyme; and finally why *RNA might just be good enough at both roles to break out of the Catch-22*.

R. Dawkins. *The Ancestor's Tale: A Pilgrimage to the Dawn of Evolution*

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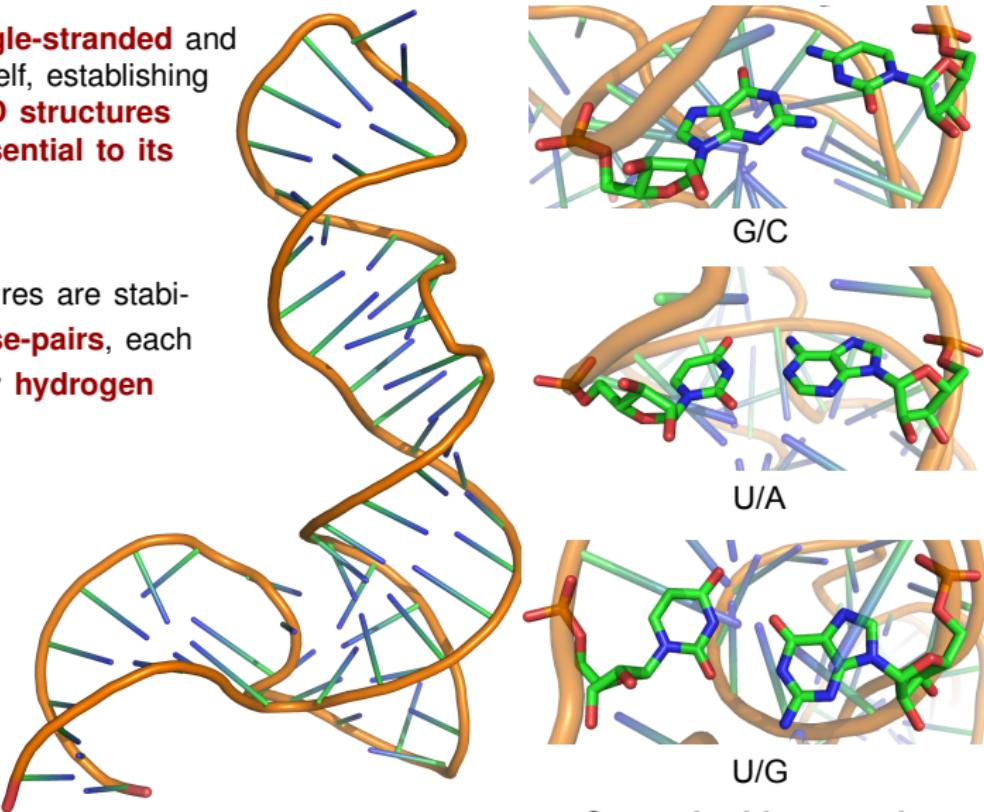
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RNA folding

RNA is **single-stranded** and **folds** on itself, establishing **complex 3D structures** that are **essential to its function(s)**.

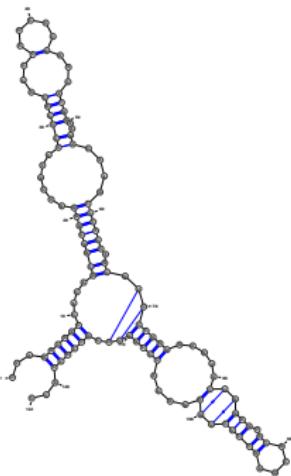
RNA structures are stabilized by **base-pairs**, each mediated by **hydrogen bonds**.



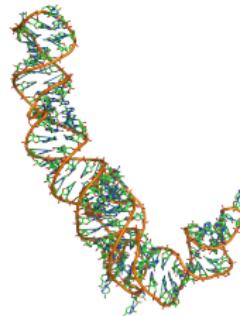
RNA Structure(s)

Three¹ levels of representation:

UUAGGCGGCCACAGC
GGUGGGGUUGCCUC
CGUACCCAUCGGAA
CACGGAAGAUAGCC
CACCAAGCGUUCGGG
GAGUACUGGAGUGCC
CGAGCCUCUGGGAAA
CCCGGUUCGCCGCCA
CC



Primary structure



Secondary structure

Tertiary structure

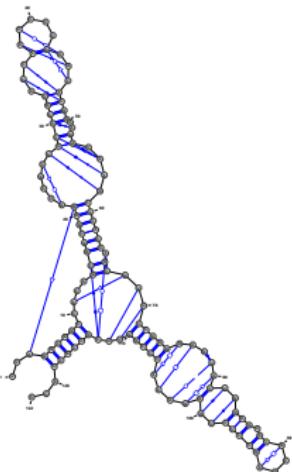
Source: 5s rRNA (PDB 1K73:B)

¹Well, mostly...

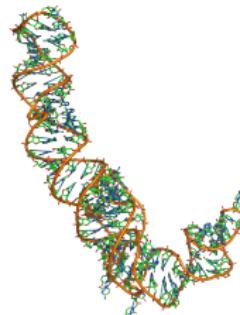
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Primary structure



Tertiary structure

Secondary⁺ structure

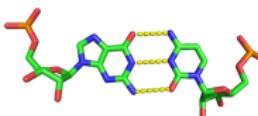
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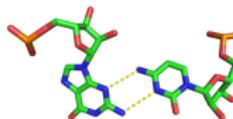
► Non-canonical base-pairs

Any base-pair **other than** {(A-U), (C-G), (G-U)}

Or interacting on non-standard edge (\neq WC/WC-Cis) [LW01].

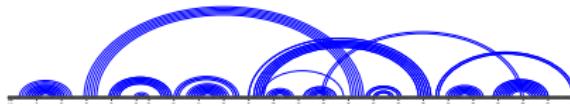


Canonique CG pair(WC/WC-Cis)



Non-canonical CG pair (Sugar/WC-Trans)

► Pseudoknots (PKs)



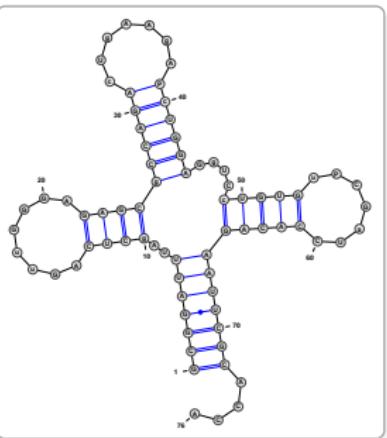
Pseudoknotted structure of group I ribozyme (PDBID: 1Y0Q:A)

Considering PKs may lead to better predictions, **but**:

- Some PK conformations are simply unfeasible;
- Folding *in silico* with general pseudoknots is NP-complete [LP00];

Still, folding on restricted classes of conformations seems promising [CDR⁺04].

Various representations for a versatile biomolecule



Outer-planar graphs

Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*

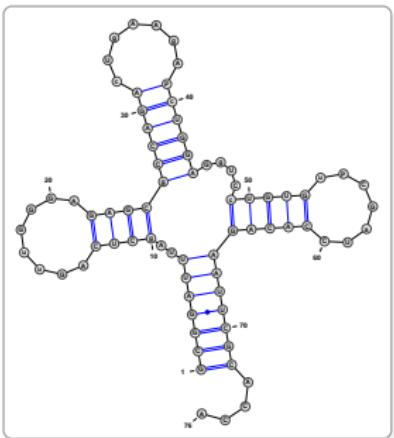
Supporting intuitions

Different representations

Common combinatorial structure

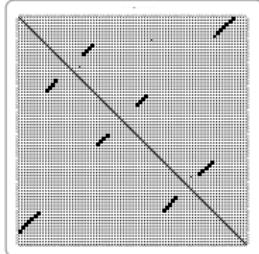
*Additional steric constraints

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Dot plots

Adjacency matrices*

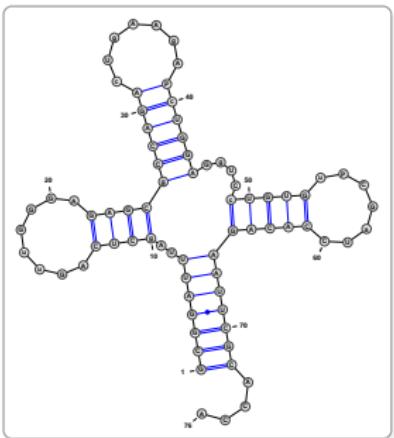
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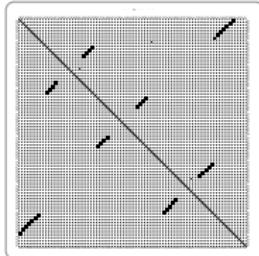
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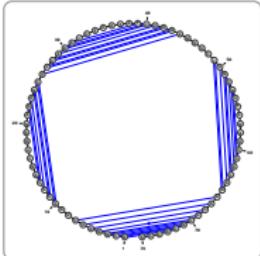


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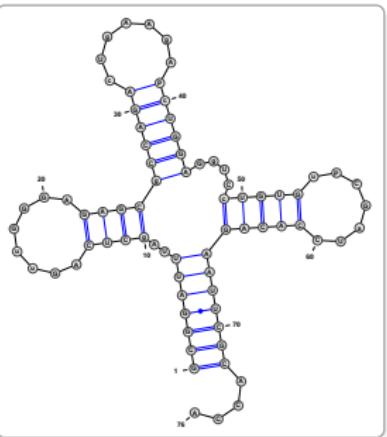
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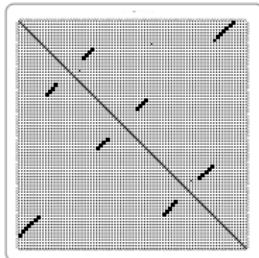


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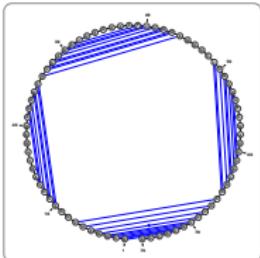
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Motzkin words*



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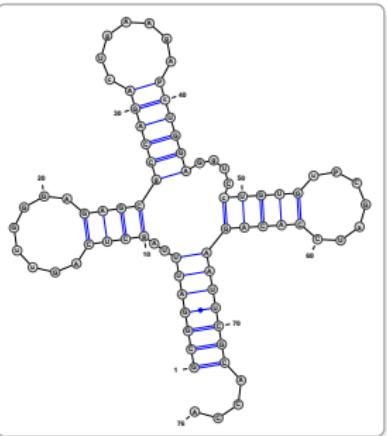


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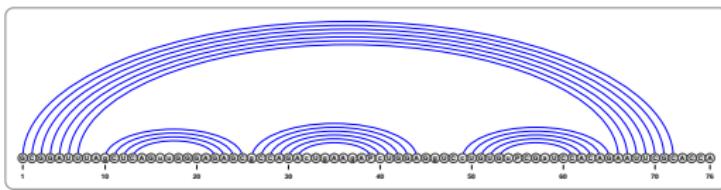


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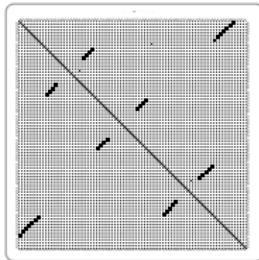
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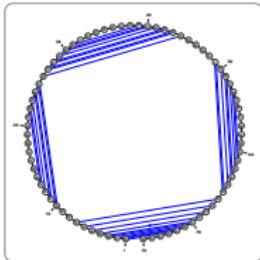
Motzkin words*



Non-crossing arc-annotated sequences*



Dot plots
Adjacency matrices*



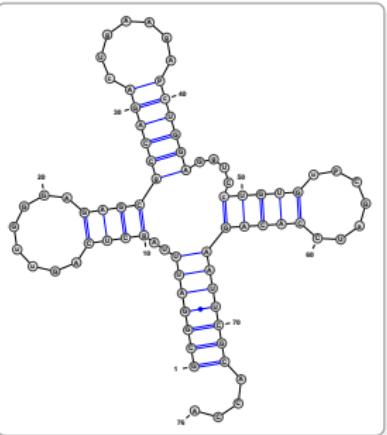
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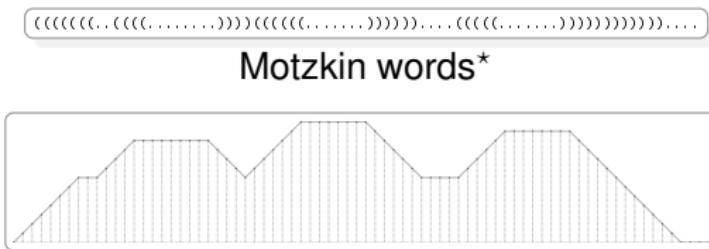
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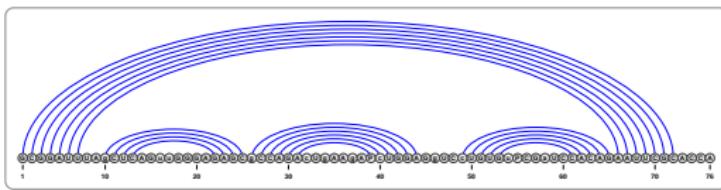


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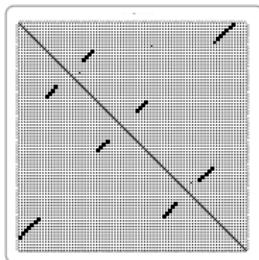
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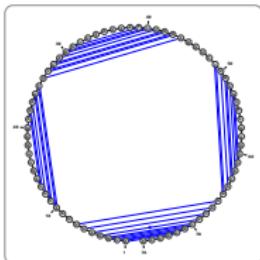
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Non-crossing arc diagrams*

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- RNA Structure(s)
- Some representations of RNA structure

2 Some flavours of folding prediction

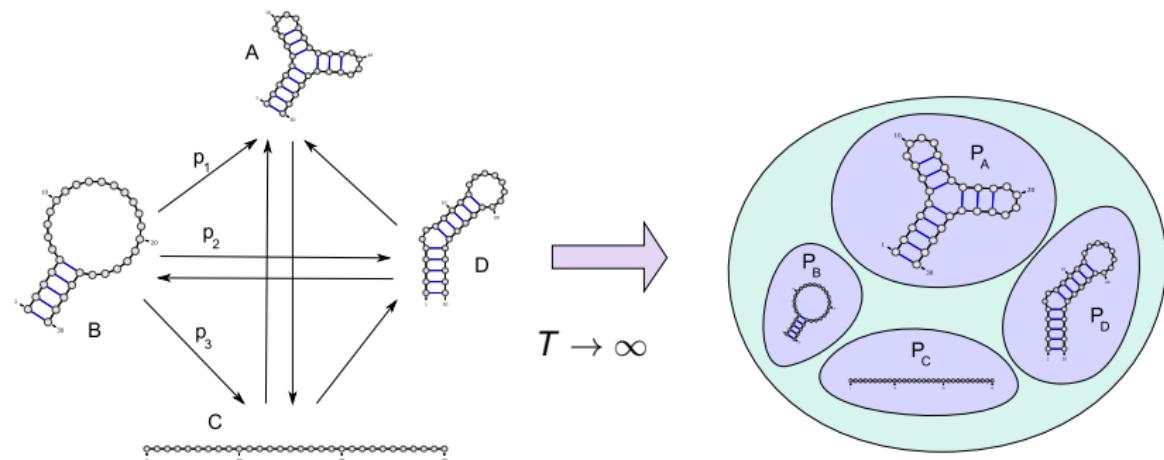
- Thermodynamics vs Kinetics
- Dynamic programming: Reminder

3 Free-energy minimization

- Nussinov-style RNA folding
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- Performances and the comparative approach
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Thermodynamics aparté

At the nanoscopic scale, RNA structure *fluctuates* (\approx Markov process).



Convergence towards a **stationary distribution** at the Boltzmann equilibrium, where the probability of a conformation only depends on its **free-energy**.

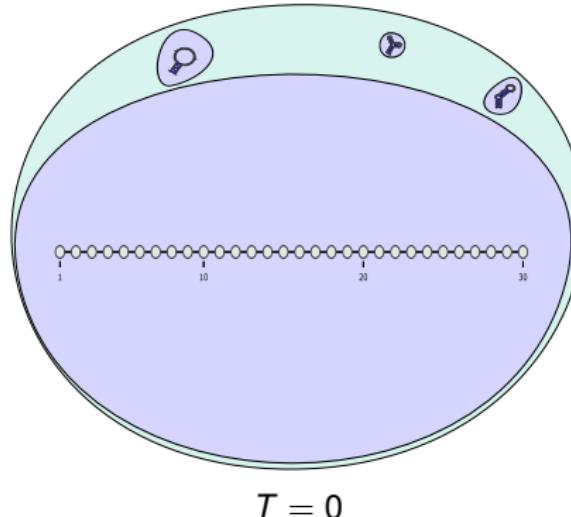
Corollary: Initial conformation does not matter.

Questions: For a given **conformation space** and **free-energy** model:

- Determine most stable (Minimum Free-Energy) structure at equilibrium;
- Compute average properties of Boltzmann ensemble;

Away from equilibrium

Transcription: RNA synthesized, supposedly without structure²

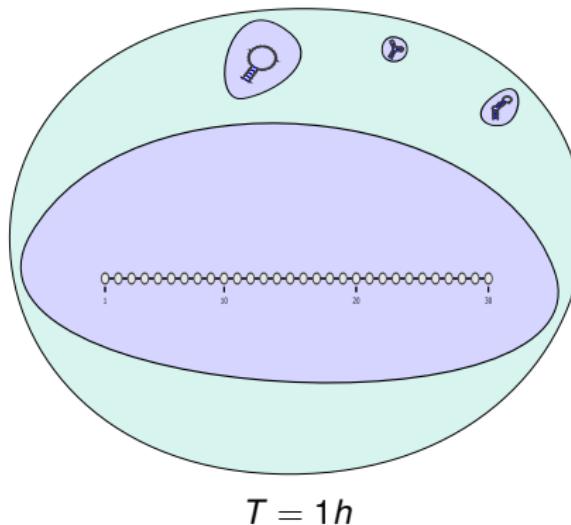


But most mRNAs are degrade before 7h (Org.: Souris [SSN⁺09]).

²Except for co-transcriptional folding...

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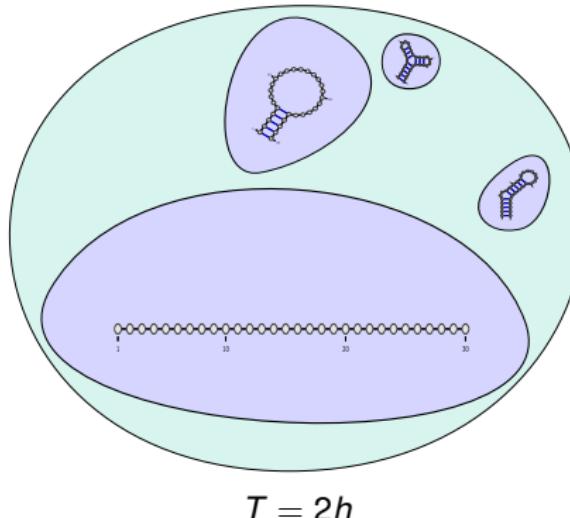


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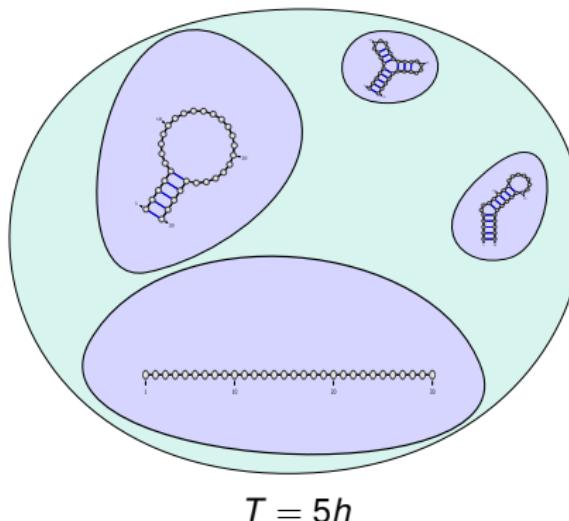


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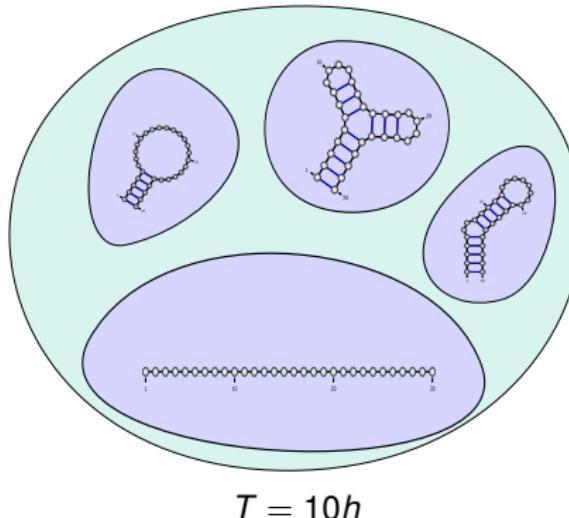


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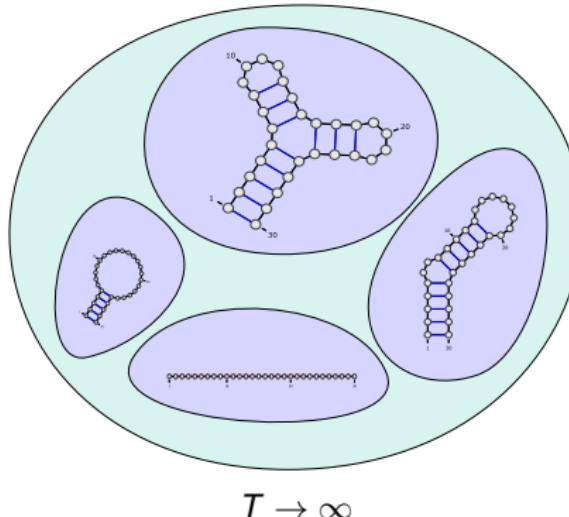


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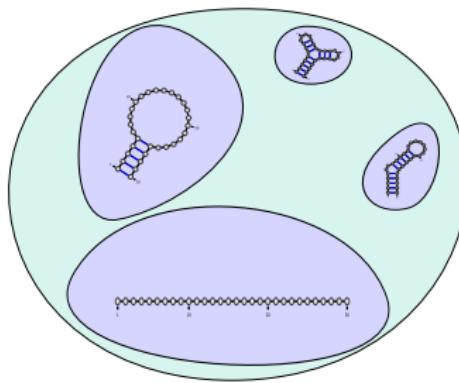


But most mRNAs are degrade before 7h (Org.: Souris [SSN⁺09]).

²Except for co-transcriptional folding...

Away from equilibrium

Transcription: RNA synthesized, supposedly without structure²



$$T = 10h$$

But most mRNAs are degrade before 7h (Org.: Souris [SSN⁺09]).

- A.** Determine most stable (Minimum Free-Energy) structure at equilibrium;
- B.** Compute average properties of Boltzmann ensemble;
- C. Determine most likely structure at finite time T .**
(c.f. H. Isambert through simulation, NP-complete deterministically [MTSC09])

²Except for co-transcriptional folding...

Dynamic programming: General principle

Dynamic programming = General optimization technique.

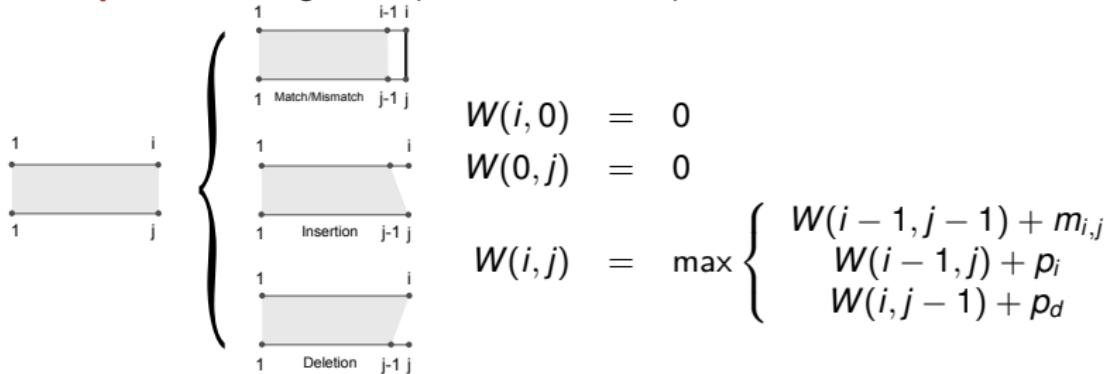
Prerequisite: Optimal solution for problem P can be derived from solutions to strict sub-problems of P .

Bioinformatics :

Discrete solution space (alignments, structures...)

- + Additively-inherited objective function (cost, log-odd score, energy...)
- ⇒ Efficient dynamic programming scheme

Example: Local Alignment(Smith/Waterman)



Algorithmic details

Dynamic programming scheme defines a space of (sub)problems and a recurrence that relates the score of a problem to that of smaller problems.

Given a scheme, two steps :

- ▶ **Matrix filling:** Computation and tabulation of best scores (Computed from smaller problems to larger ones).
- ▶ **Traceback:** Reconstruct best solution from contributing subproblems.

Complexity of algorithm depends on:

- ▶ **Cardinality** of sub-problem space
- ▶ **Number of alternatives** considers at each step (#Terms in recurrence)

Smith&Waterman example:

- ▶ $i: 1 \rightarrow n + 1 \Rightarrow \Theta(n)$
- ▶ $j: 1 \rightarrow m + 1 \Rightarrow \Theta(m)$
- ▶ 3 operations at each step
 $\Rightarrow \Theta(m.n)$ time/memory

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

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	A	C	A	C	A	C	T	A
0	0	0	0	0	0	0	0	0
A	0							
G	0							
C	0							
A	0							
C	0							
A	0							
C	0							
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	A	C	A	C	A	C	T	A
0	0	0	0	0	0	0	0	0
A	0 → 2							
G	0							
C	0							
A	0							
C	0							
A	0							
C	0							
A	0							

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	A	C	A	C	A	C	T	A
0	0	0	0	0	0	0	0	0
A	0	2 → 1						
G	0							
C	0							
A	0							
C	0							
A	0							
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	0	0	0	0	0	0	0	0
A	0	2 → 1 → 2						
G	0							
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0	0	0	0	0	0	0	0	0
A	0	2 → 1	2 → 1					
G	0							
C	0							
A	0							
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	A	C	A	C	A	C	T	A
0	0	0	0	0	0	0	0	0
A	0	2 → 1	2 → 1 → 2	2 → 1 → 0	0	0	0	0
G	0							
C	0							
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A	0							
C	0							
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Complete example

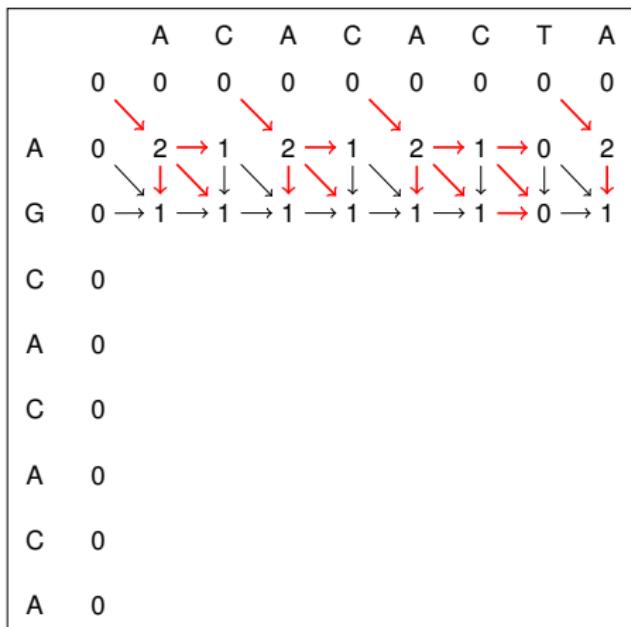
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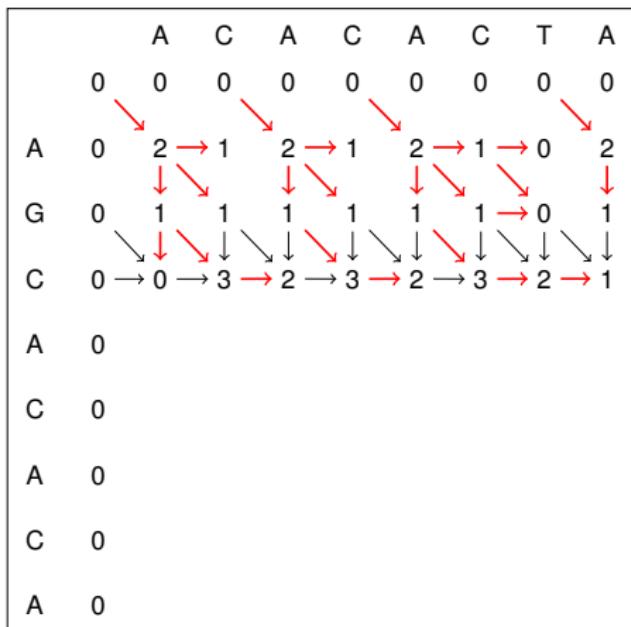
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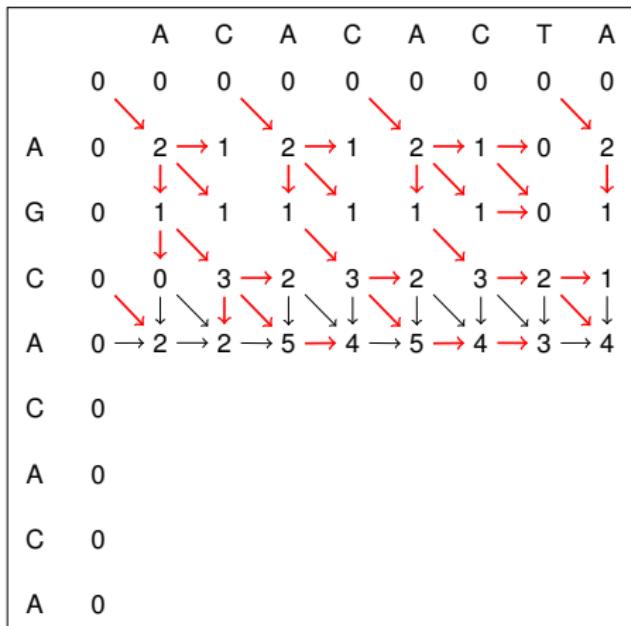
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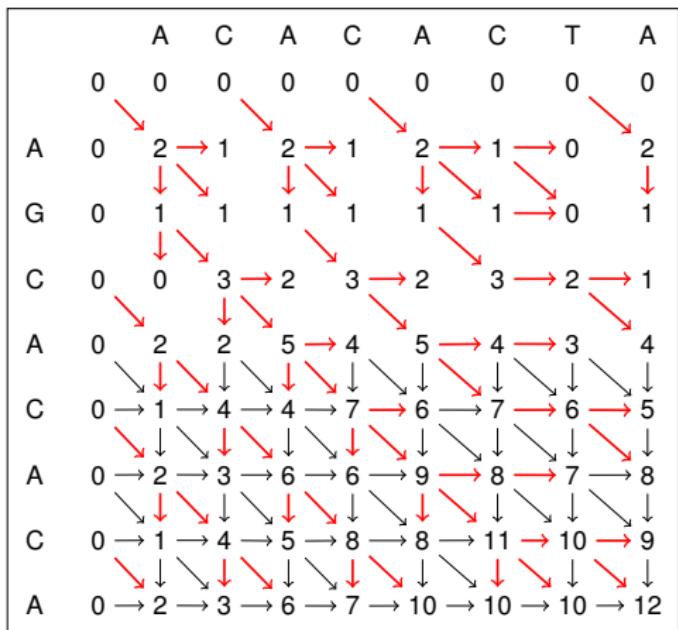
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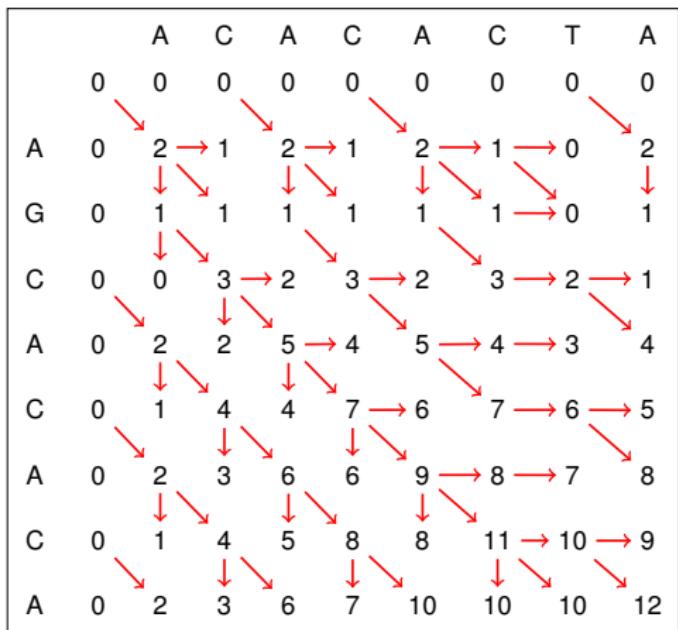
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Complete example

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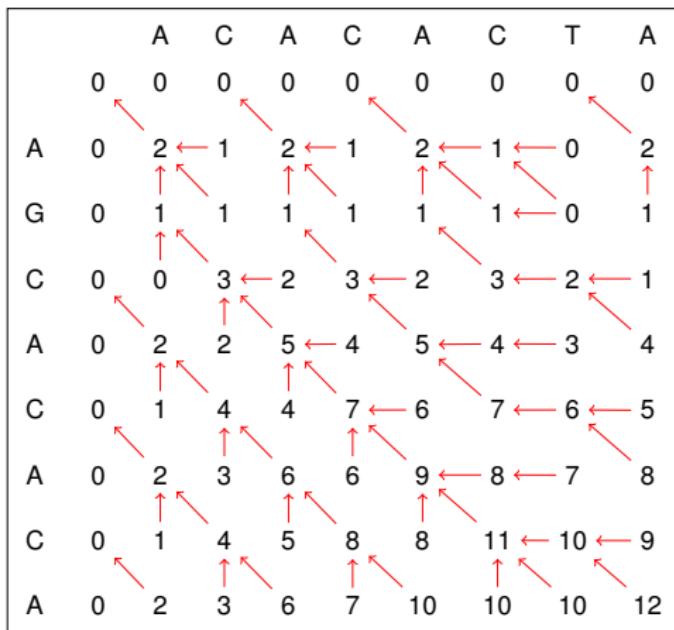
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Best alignment



Complete example

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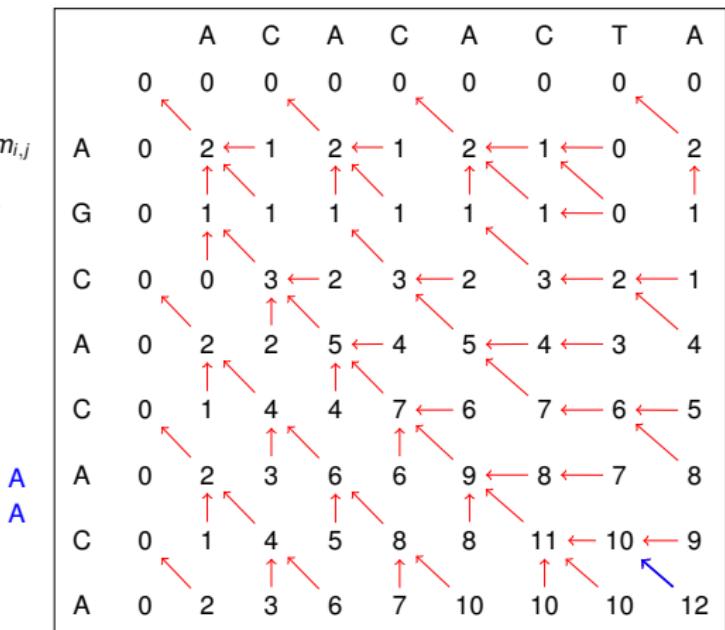
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Complete example

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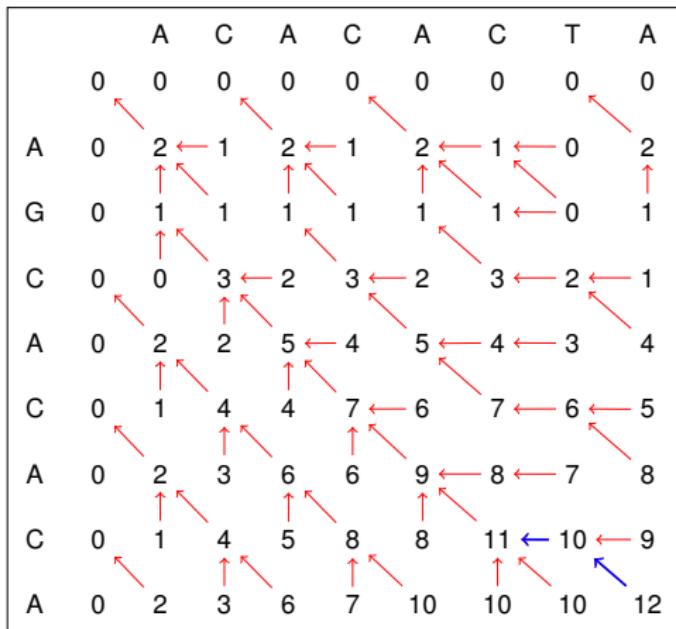
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Best alignment

- A
T A



Complete example

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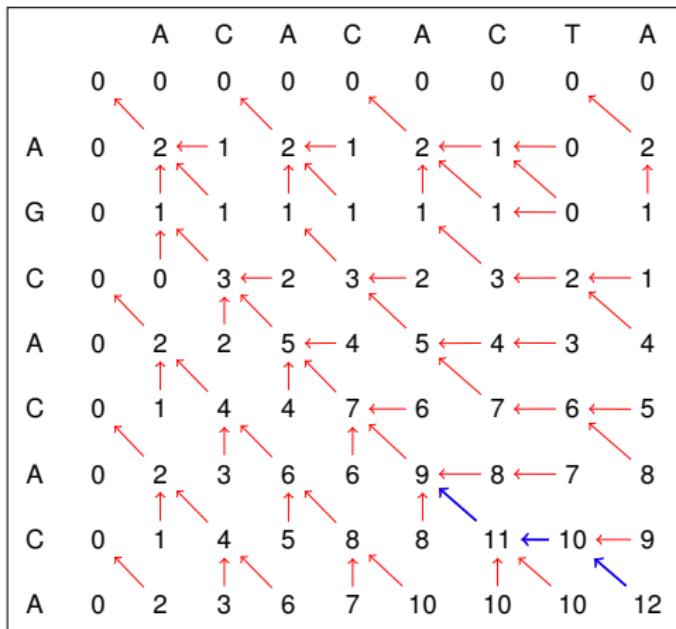
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Best alignment

C - A
C T A



Complete example

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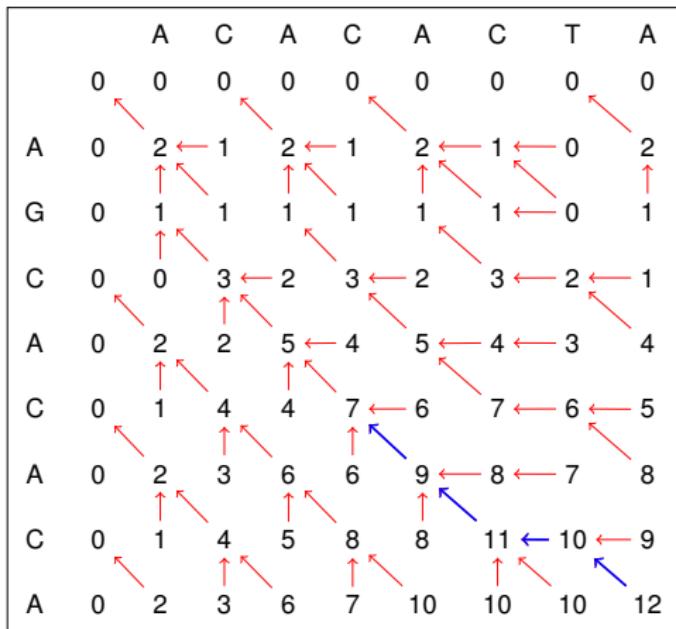
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Best alignment

A	C	-	A
A	C	T	A



Complete example

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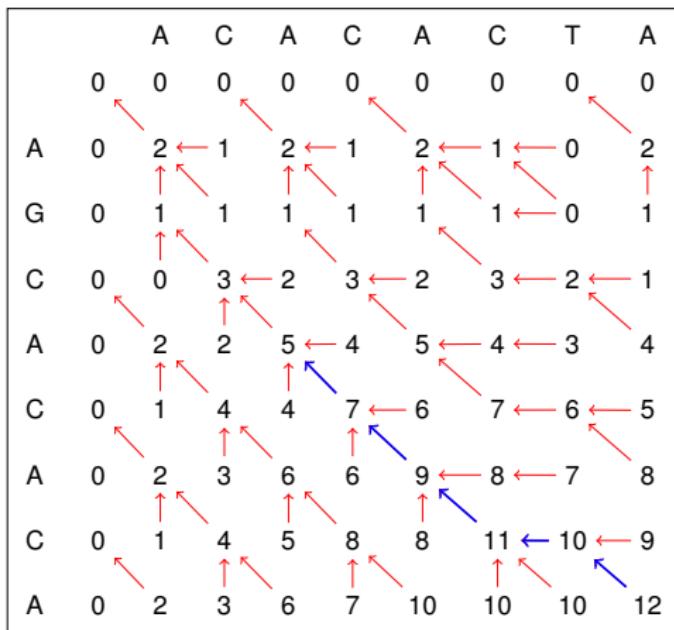
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Best alignment

C	A	C	-	A
C	A	C	T	A



Complete example

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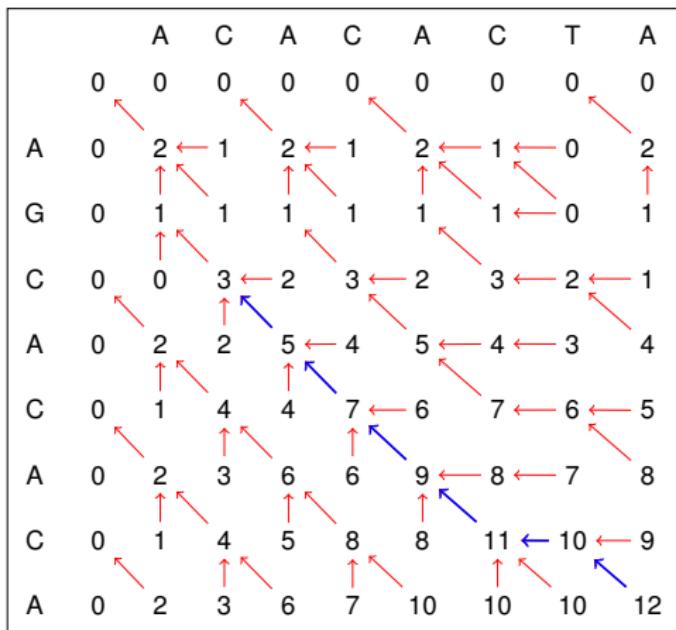
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Best alignment

A	C	A	C	-	A
A	C	A	C	T	A



Complete example

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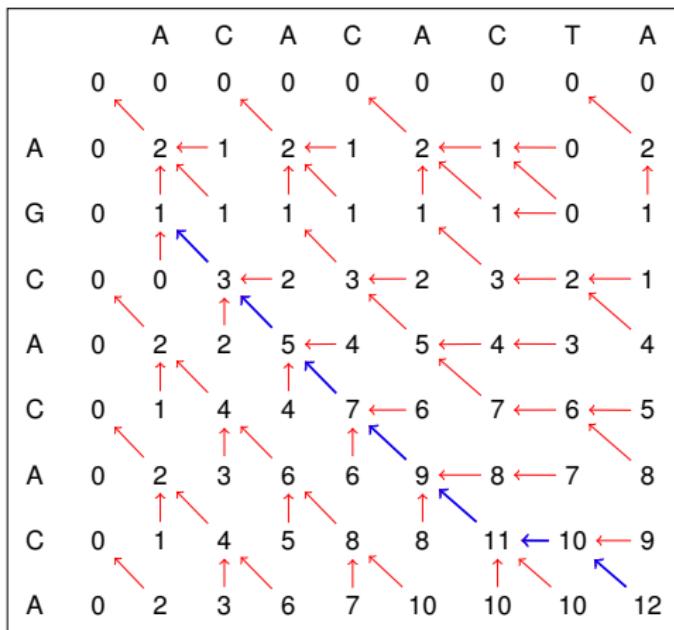
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Best alignment

C A C A C - A
C A C A C T A



Complete example

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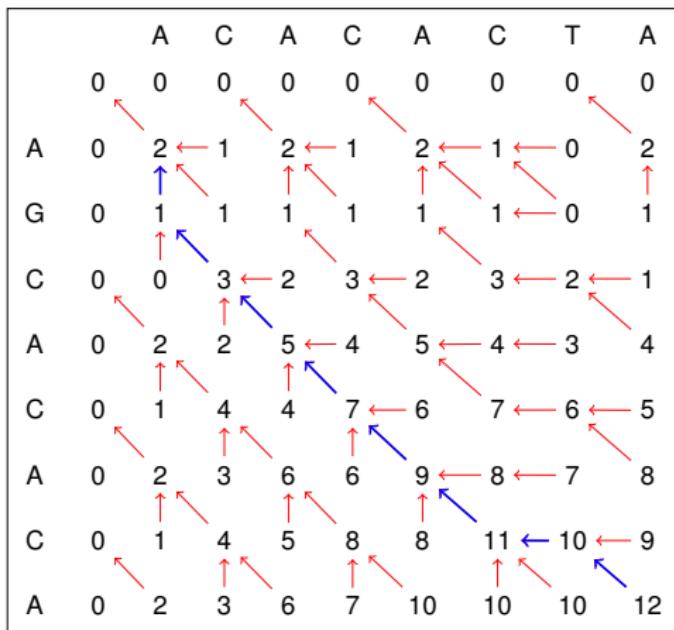
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Best alignment

G C A C A C - A
- C A C A C T A



Complete example

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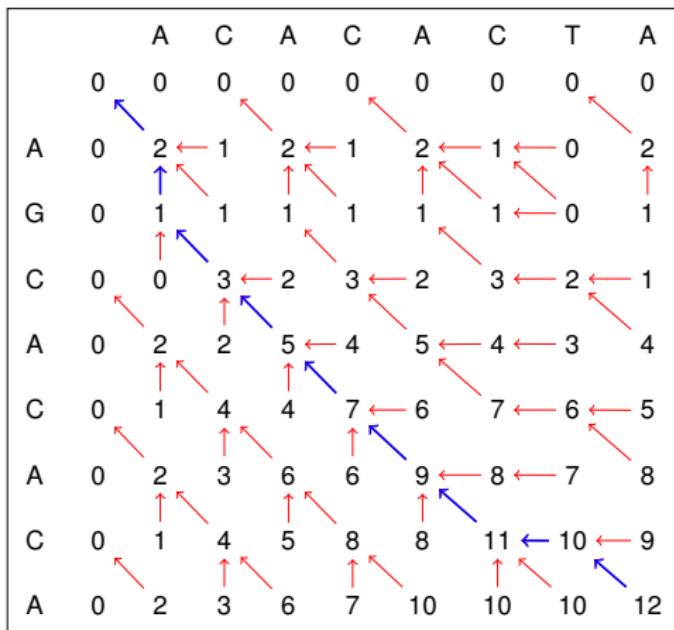
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Best alignment

A G C A C A C - A
A - C A C A C T A



Necessary properties:

- ▶ **Correctness:** \forall sub-problem, the computed value must indeed maximize the objective function .

Proofs usually inductive, and quite technical, but very systematic.

Desirable properties of DP schemes:

- ▶ **Completeness** of space of solutions generated by decomposition.
Algorithmic tricks, by *cutting branches*, may violate this property.
- ▶ **Unambiguity:** Each solution is generated at most once.
⇒ Under these properties, one can enumerate solution space.

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- Dynamic programming 101
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- RNA Structure(s)
- Some representations of RNA structure

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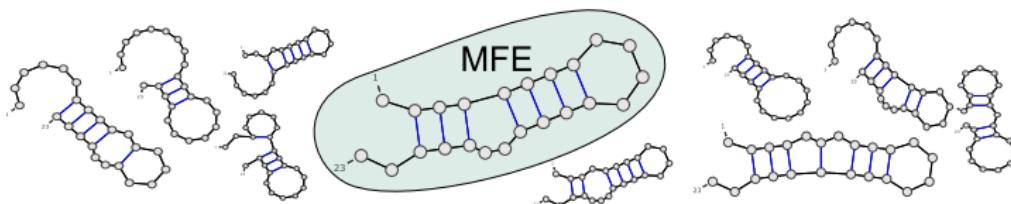
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Problem A: Determine Minimum Free-Energy structure (MFE).

Ab initio folding prediction =

Predict RNA structure from its sequence ω only.



- ▶ **Conformations:** Set S_ω of secondary structures compatible (w.r.t. base-pairing constraints) with primary structure ω .
- ▶ **Free-Energy:** Function $E_{\omega, S}$ (KCal.mol⁻¹), additive on motifs occurring in any sequence/conformation couple (ω, S) .
- ▶ **Native structure:** Functional conformation of the biomolecule.

Remarks:

- ▶ Not necessarily unique (Kinetics, or bi-stable structures);
- ▶ In presence of PKs → Ambiguous: Which is the native conformation?

Nussinov/Jacobson energy model (NJ)

Base-pair maximization (with a twist):

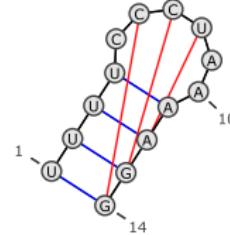
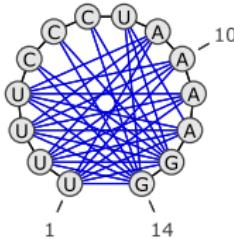
- ▶ Additive model on independently contributing base-pairs;
- ▶ Canonical base-pairs only: Watson/Crick (A/U,C/G) and Wobble (G/U)

$$\Rightarrow E_{\omega,S} = -\# \text{Paires}(S)$$

Folding in NJ model \Leftrightarrow Base-pair (weight) maximization

Example:

UUUUCCCUAAAAGG



Variant: Weight each pair with $-\#$ Hydrogen bonds

$$\Delta G(G \equiv C) = -3$$

$$\Delta G(A = U) = -2$$

$$\Delta G(G - U) = -1$$

Nussinov/Jacobson energy model (NJ)

Base-pair maximization (with a twist):

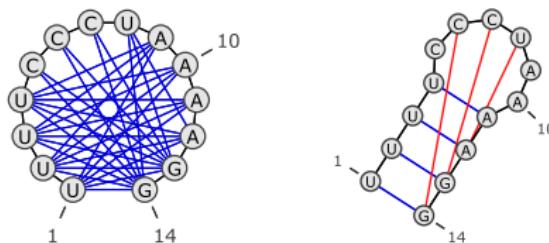
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Folding in NJ model \Leftrightarrow Base-pair (**weight**) maximization

Example:

UUUUUCCCUAAAAGG



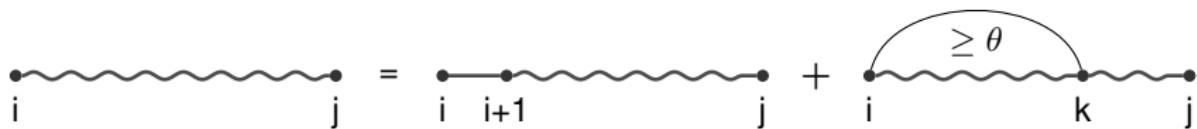
Variant: Weight each pair with $-\#$ Hydrogen bonds

$$\Delta G(G \equiv C) = -3$$

$$\Delta G(A = U) = -2$$

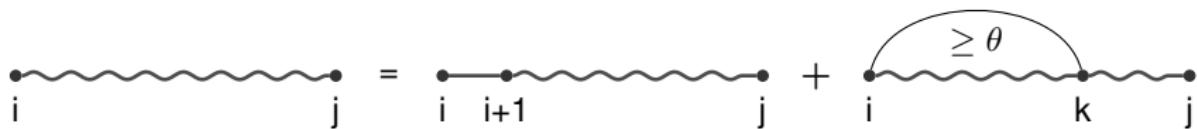
$$\Delta G(G - U) = -1$$

Nussinov/Jacobson DP scheme



$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \left\{ \begin{array}{ll} N_{i+1,j} & i \text{ unpaired} \\ \min_{k=i+\theta+1}^j \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{array} \right.$$



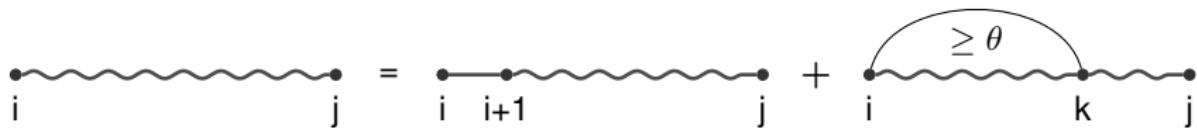
$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \left\{ \begin{array}{ll} N_{i+1,j} & i \text{ unpaired} \\ \min_{k=i+\theta+1}^j \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{array} \right.$$

Correctness. Goal = Show that MFE over interval $[i, j]$ is indeed found in $N_{i,j}$ after completing the computation. Proceed by induction:

- ▶ Assume that property holds for any $[i', j']$ such that $j' - i' < n$.
- ▶ Consider $[i, j], j - i = n$. Let $MFE_{i,j} :=$ Base-pairs of best struct. on $[i, j]$. Then first position i in $MFE_{i,j}$ is either:
 - ▶ **Unpaired:** $MFE_{i,j} = MFE_{i+1,j}$ \rightarrow free-energy = $N_{i+1,j}$
 - ▶ **Paired to k :** $MFE_{i,j} = \{(i, k)\} \cup MFE_{i+1,k-1} \cup MFE_{k+1,j}$
 (Indeed, any BP between $[i+1, k-1]$ and $[k+1, j]$ would violate (i, k))
 \rightarrow free-energy = $\Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j}$

Nussinov/Jacobson DP scheme



$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \left\{ \begin{array}{ll} N_{i+1,j} & i \text{ unpaired} \\ \min_{k=i+\theta+1}^j \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{array} \right.$$

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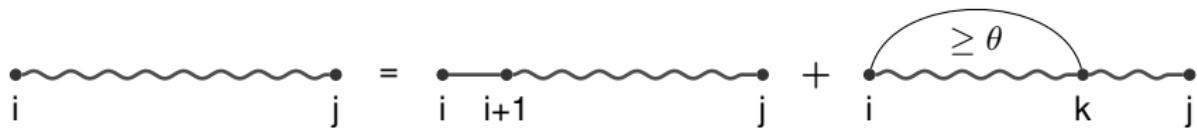
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(Indeed, any BP between $[i+1, k-1]$ and $[k+1, j]$ would cross (i, k))

\rightarrow free-energy = $\Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j}$

Nussinov/Jacobson DP scheme



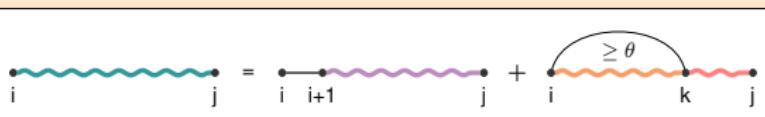
$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \left\{ \begin{array}{ll} N_{i+1,j} & i \text{ unpaired} \\ \min_{k=i+\theta+1}^j \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{array} \right.$$

Correctness. Goal = Show that MFE over interval $[i, j]$ is indeed found in $N_{i,j}$ after completing the computation. Proceed by induction:

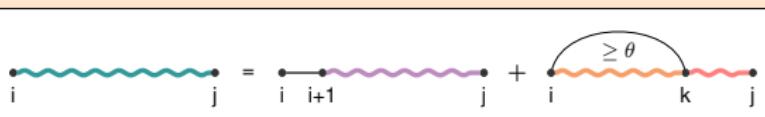
- ▶ Assume that property holds for any $[i', j']$ such that $j' - i' < n$.
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(Indeed, any BP between $[i+1, k-1]$ and $[k+1, j]$ would cross (i, k))
 $\rightarrow \text{free-energy} = \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j}$

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	10
A	0	0	0	0	2	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7	7	7	7
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7	7	7
C	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5	5	5	5
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	2	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0



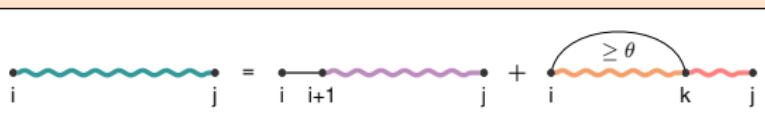
Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	5	5	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7	7	7	7
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7	7	7
C	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5	5	5	5
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	2	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0



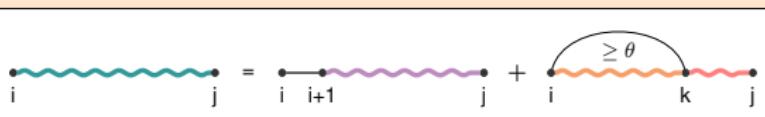
Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	6	7	9	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	5	6	8	10	10	10
A	0	0	0	0	2	2	2	2	4	4	4	5	5	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	5	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7	7	7	7
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7	7	7
C	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5	5	5	5
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	2	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0



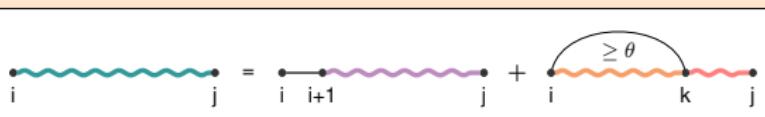
Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10
A	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	5	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5	5		
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0



Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	10
A	0	0	0	0	2	2	2	2	4	4	5	7	7	7	8	10	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5			
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0



Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	5	7	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	2	5	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	2	3	5	5	5	6	7	7	7	8	7
U	0	0	0	0	0	0	2	3	5	5	5	5	5	5	5	7	7	7
C	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	5	5	5
U	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	2	2	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{arc from } i \text{ to } k \text{ with radius } \geq \theta & \text{if } k \text{ is a base pair} \\ 0 & \text{otherwise} \end{cases}$												0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	5	7	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	2	5	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	6	7	7	7
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	5	7	7
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	3	5	5	5
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	2	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \theta \geq \theta \\ \text{---} & \end{cases} k \text{---} j$												0	0	0	0	0	0
C													0	0	0	0	0	0
G													0	0	0	0	0	0
A													0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10
A	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	5	5
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	3	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \theta \geq \theta \\ \text{arc} & \end{cases} k \text{---} j$										0	0	0	0	0	0	0	0
C											0	0	0	0	0	0	0	0
G											0	0	0	0	0	0	0	0
A											0	0	0	0	0	0	0	0

Nussinov/Jacobson

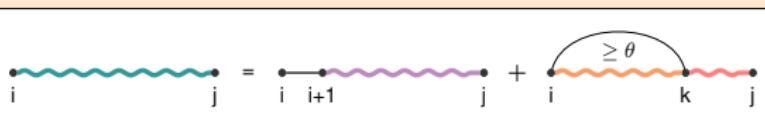
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
(.).	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10
A	0	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	3	5	5	5
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	2	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \theta \geq \theta \\ \text{---} & \end{cases} k \text{---} j$										0	0	0	0	0	0	0	0
C											0	0	0	0	0	0	0	0
G											0	0	0	0	0	0	0	0
A											0	0	0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10
A	0	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5	5	5	5
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	2	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	i	j	=	i	i+1	j	+ 	i	k	j	0	0	0	0	
A																0		

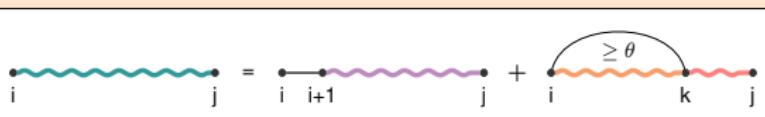
Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	10
A	0	0	0	0	2	2	2	2	2	4	4	5	7	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8	8
C	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	2	3	5	5	5	6	7				
U	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7			
C	0	0	0	0	0	0	3	3	3	3	3	3	3	5	5	5	5	5
U	0	0	0	0	0	0	2	2	2	2	2	2	2	2	2	2	2	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0



Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	10
A	0	0	0	0	2	2	2	2	2	4	4	5	7	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	2	5	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0



	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	10
A	0	0	0	0	2	2	2	2	4	4	5	5	7	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7	7	7	7
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7	7	7
C	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5	5	5	5
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	2	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \theta \geq \theta \\ \text{arc} & \end{cases} k \text{---} j$												0	0	0	0	0	0
C													0	0	0	0	0	0
G													0	0	0	0	0	0
A													0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	((.))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	10
A	0	0	0	0	2	2	2	2	4	4	5	5	7	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5			
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \theta \geq \theta \\ \text{arc} & \end{cases} k \text{---} j$												0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	((.))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	10
A	0	0	0	0	2	2	2	2	4	4	5	5	7	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7	7	7	7
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7	7	7
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	5	5
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	3	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \geq \theta \\ \text{---} & \end{cases} k \text{---} j$												0	0	0	0	0	0
C													0	0	0	0	0	0
G													0	0	0	0	0	0
A													0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	((.))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	5	5	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8	8
C	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7	7	7	7
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7	7	7
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	5	5
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	3	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A													0	0	0	0	0	0
C													0	0	0	0	0	0
G	i		j	=	i	i+1		j		i	k	j						
A																		0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	5	7	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7	7	7	7
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7	7	7
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	5	5
U	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	2	3	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{arc from } i \text{ to } k \geq \theta & \text{if } k \text{ is a base pair} \\ 0 & \text{otherwise} \end{cases}$												0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	5	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10	10	10
A	0	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7	7	7
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7	7	7
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	5
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	3	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \geq \theta \\ \text{arc} & \end{cases} k \text{---} j$												0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10		
A	0	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	5	5
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{arc from } i \text{ to } k \text{ if } \theta \geq \theta \\ 0 \text{ otherwise} \end{cases}$												0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10	10	10
A	0	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	3	5	5	5
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	3	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{arc from } i \text{ to } k \text{ if } \theta \geq \theta \\ 0 \text{ otherwise} \end{cases}$												0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	7	8	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	
A	0	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7		
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	i	j	=	i	i+1	j	i	k	j	≥ θ	0	0	0	0
A																		0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	7	8	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	
A	0	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \theta \\ \text{---} & \geq \theta \end{cases} k \text{---} j$												0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	2	2	2	2	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10	10
A	0	0	0	0	0	0	0	0	0	2	2	2	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7		
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	5
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \theta \geq \theta \\ \text{---} & \end{cases} k \text{---} j$												0	0	0	0	0	0
C													0	0	0	0	0	0
G													0	0	0	0	0	0
A													0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)	.	(.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10	10
A	0	0	0	0	0	0	0	0	0	2	2	2	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7		
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	5
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	$i \text{ --- } j = i \text{ --- } i+1 \text{ --- } j + \begin{cases} \text{arc from } i \text{ to } k \text{ if } \theta \geq \theta \\ 0 \text{ otherwise} \end{cases}$										0	0	0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)	.	(.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10
A	0	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7		
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \geq \theta \\ \text{arc} & \end{cases} k \text{---} j$												0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)	.	(.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10	10
A	0	0	0	0	0	0	0	0	0	2	2	2	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7		
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \theta \geq \theta \\ \text{---} & \end{cases} k \text{---} j$												0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)	.	((.	.	.))))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10	10
A	0	0	0	0	0	0	0	0	0	2	2	2	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7		
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	0	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	0	0	0	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \theta \geq \theta \\ \text{---} & \end{cases} k \text{---} j$												0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

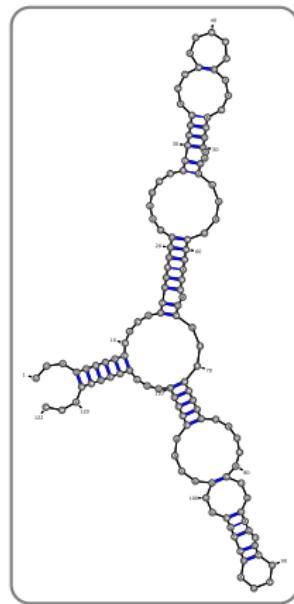
Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)	.	((.	.	.))))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10	10
A	0	0	0	0	0	0	0	0	0	2	2	2	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7		
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	$i \text{---} j = i \text{---} i+1 \text{---} j + \begin{cases} \text{---} & \theta \geq \theta \\ \text{---} & \end{cases} k \text{---} j$												0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Turner energy model

Based on **unambiguous** decomposition of 2^{ary} structure into **loops**:

- ▶ Internal loops
- ▶ Bulges
- ▶ Terminal loops
- ▶ Multi loops
- ▶ Stackings



Free-energy ΔG of a loop depend on
bases, assymmetry, dangles ...

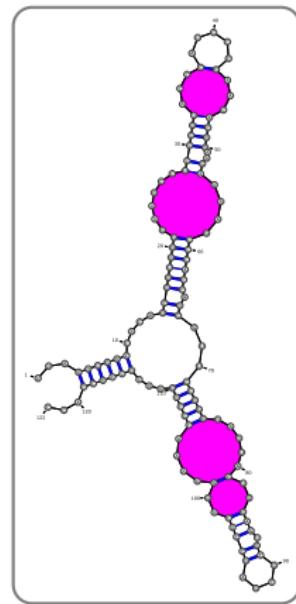
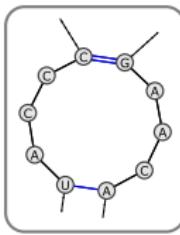
Experimentally determined
+ Interpolated for larger loops.

Improved results by taking stacking into account.

Turner energy model

Based on **unambiguous** decomposition of 2^{ary} structure into **loops**:

- ▶ Internal loops
- ▶ Bulges
- ▶ Terminal loops
- ▶ Multi loops
- ▶ Stackings



Free-energy ΔG of a loop depend on bases, assymmetry, dangles ...

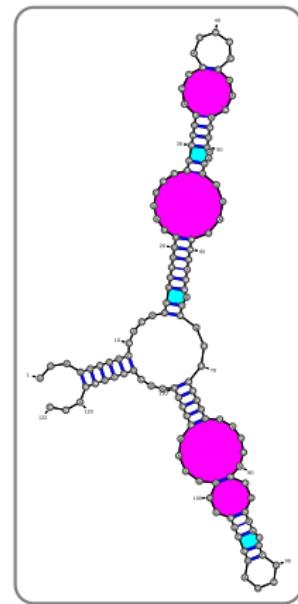
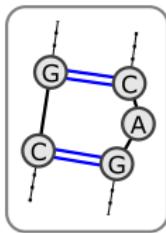
Experimentally determined
+ Interpolated for larger loops.

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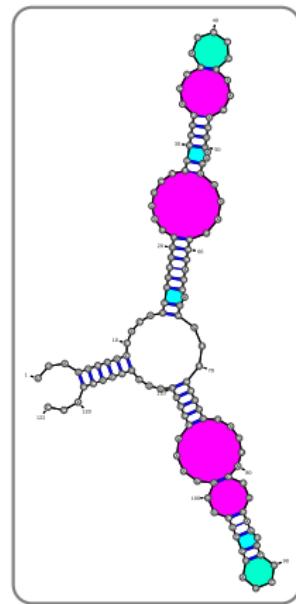
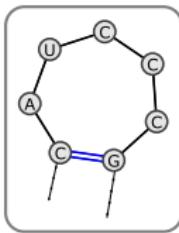
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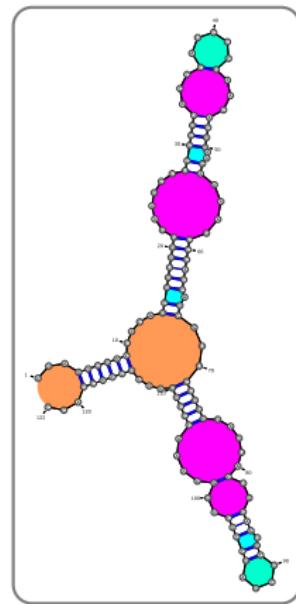
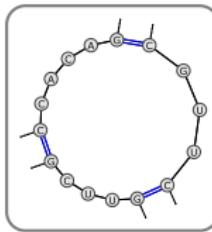
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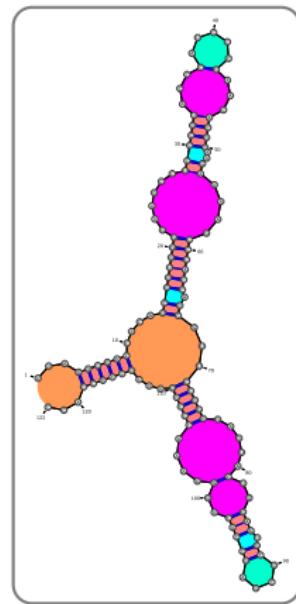
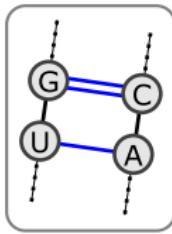
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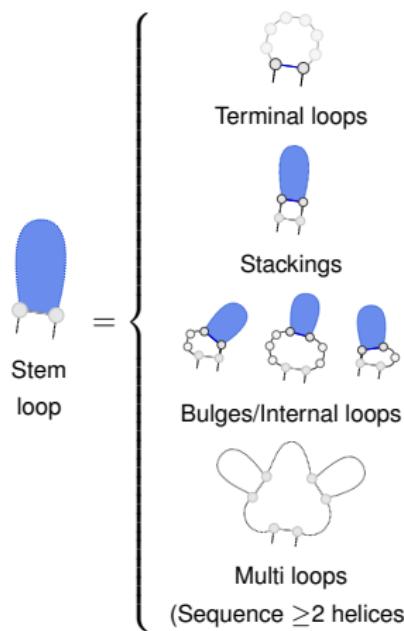


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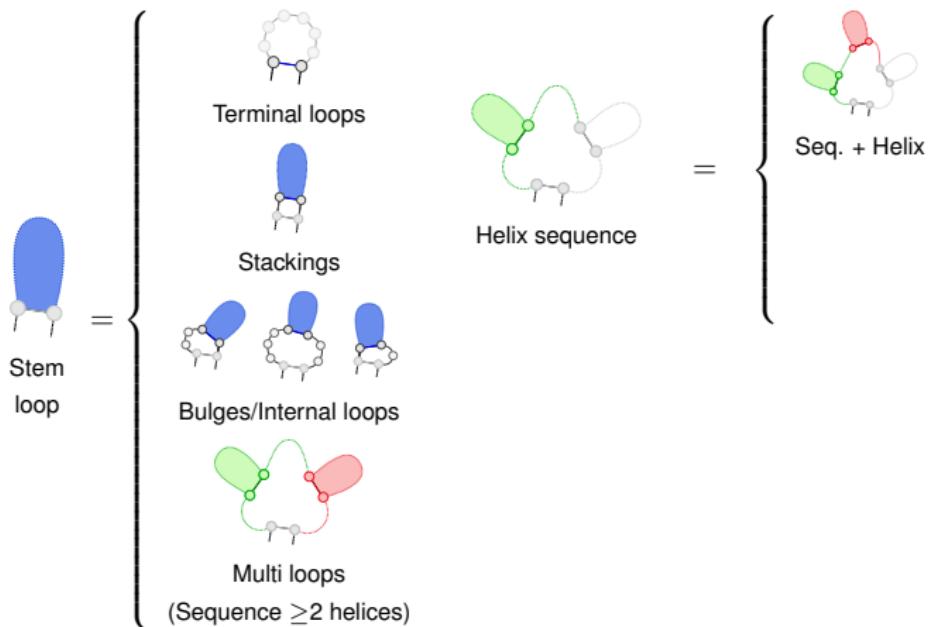
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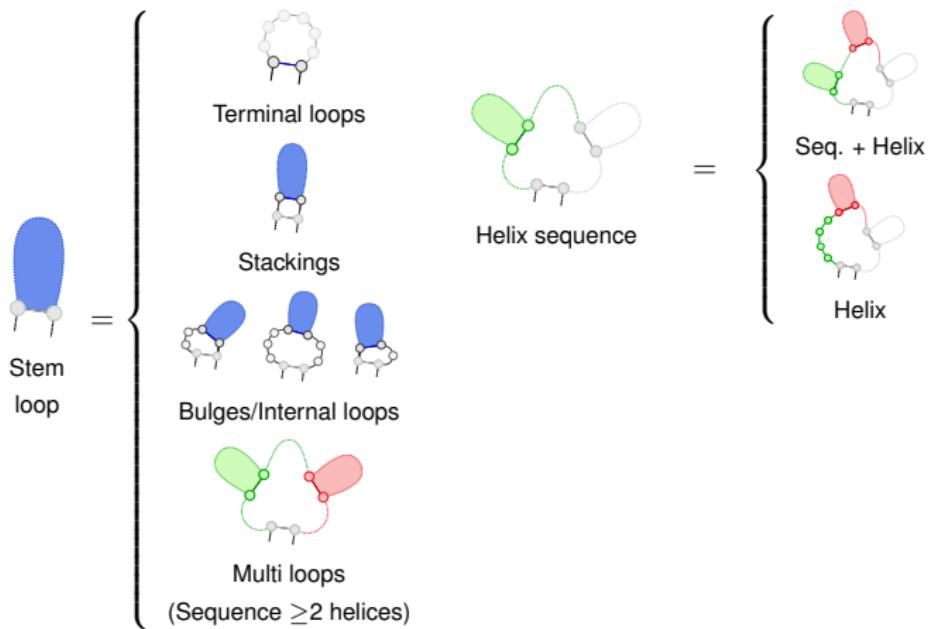
MFE DP equations



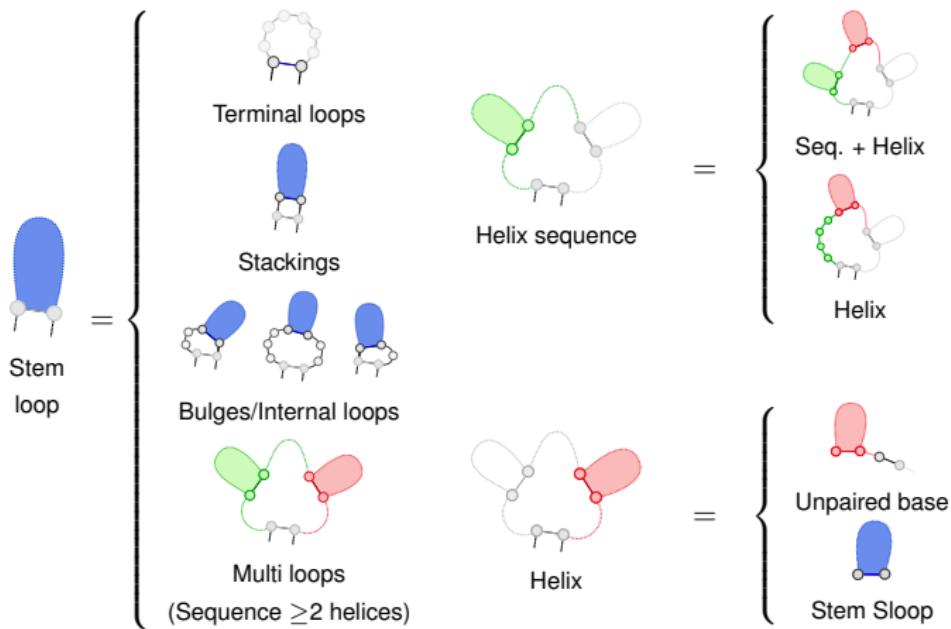
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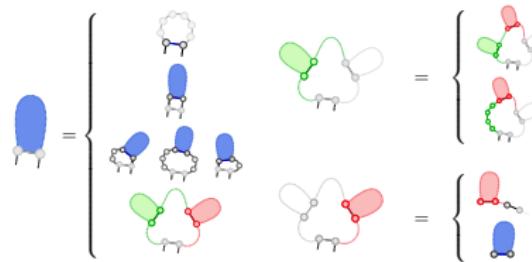


MFE DP equations



MFold Unifold

- ▶ $E_H(i, j)$: Energy of terminal loop *enclosed by* (i, j) pair
- ▶ $E_{BI}(i, j)$: Energy of bulge or internal loop *enclosed by* (i, j) pair
- ▶ $E_S(i, j)$: Energy of stacking $(i, j)/(i + 1, j - 1)$
- ▶ Penalty for multi loop (a), and occurrences of unpaired base (b) and helix (c) in multi loops.



DP recurrence

$$\begin{aligned}\mathcal{M}'_{i,j} &= \min \left\{ \begin{array}{l} E_H(i, j) \\ E_S(i, j) + \mathcal{M}'_{i+1, j-1} \\ \text{Min}_{i', j'} (E_{BI}(i, i', j', j) + \mathcal{M}'_{i', j'}) \\ a + c + \text{Min}_k (\mathcal{M}_{i+1, k-1} + \mathcal{M}^1_{k, j-1}) \end{array} \right\} \\ \mathcal{M}_{i,j} &= \text{Min}_k \left\{ \min (\mathcal{M}_{i, k-1}, b(k-1)) + \mathcal{M}^1_{k, j} \right\} \\ \mathcal{M}^1_{i,j} &= \text{Min}_k \left\{ b + \mathcal{M}^1_{i, j-1}, c + \mathcal{M}'_{i, j} \right\}\end{aligned}$$

Remontée (Backtracking)

Backtracking to reconstruct MFE structure:

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Complexity:

For each min, $\mathcal{O}(n)$ potential contributors

⇒ **Worst-case** complexity in $\mathcal{O}(n^2)$ for **naive backtrack**.

Keep best contributor for each Min ⇒ **Backtracking in $\mathcal{O}(n)$**

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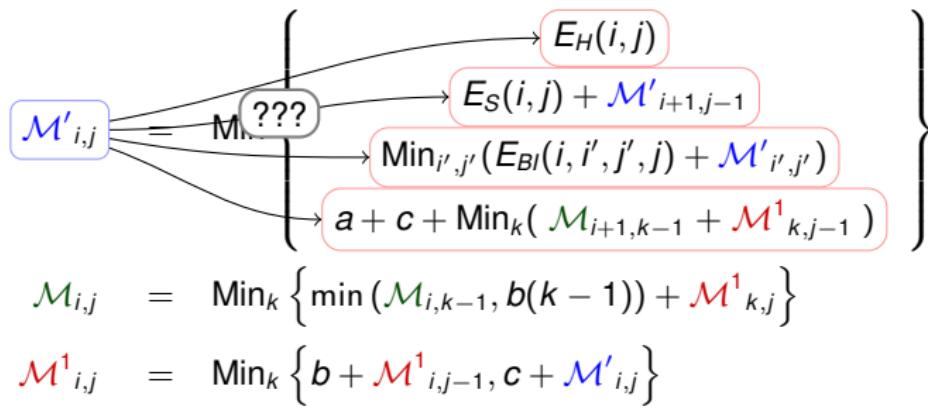
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Definition (Ab initio folding)

Starting from sequence, find conformation that minimizes free-energy.

Advantages:

- ▶ Mechanical nature allows the (in)validation of models
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Limitations:

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- ▶ No cooperativity
- ▶ Limited performances

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Starting from homologous sequences, postulate common structure and find best possible tradeoff between folding & alignment.

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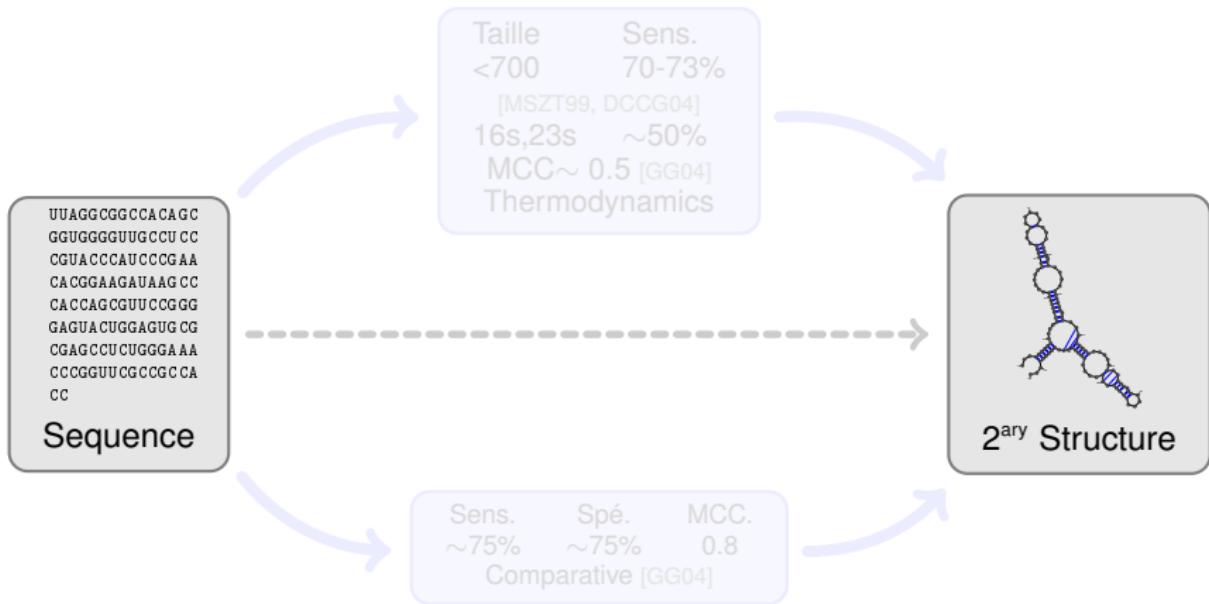
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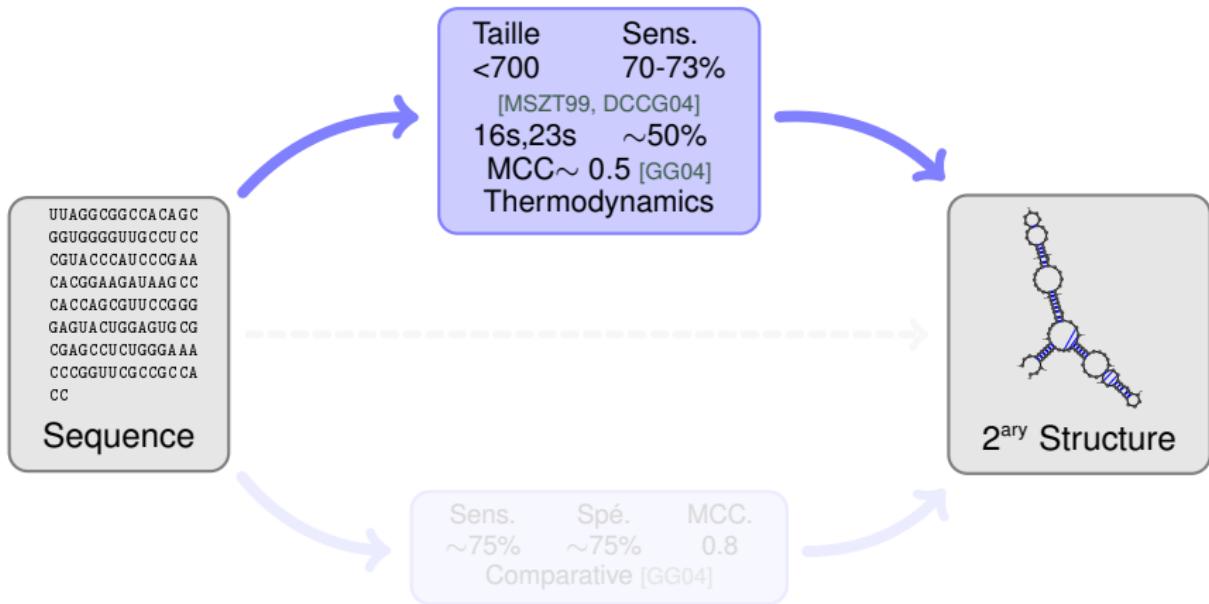
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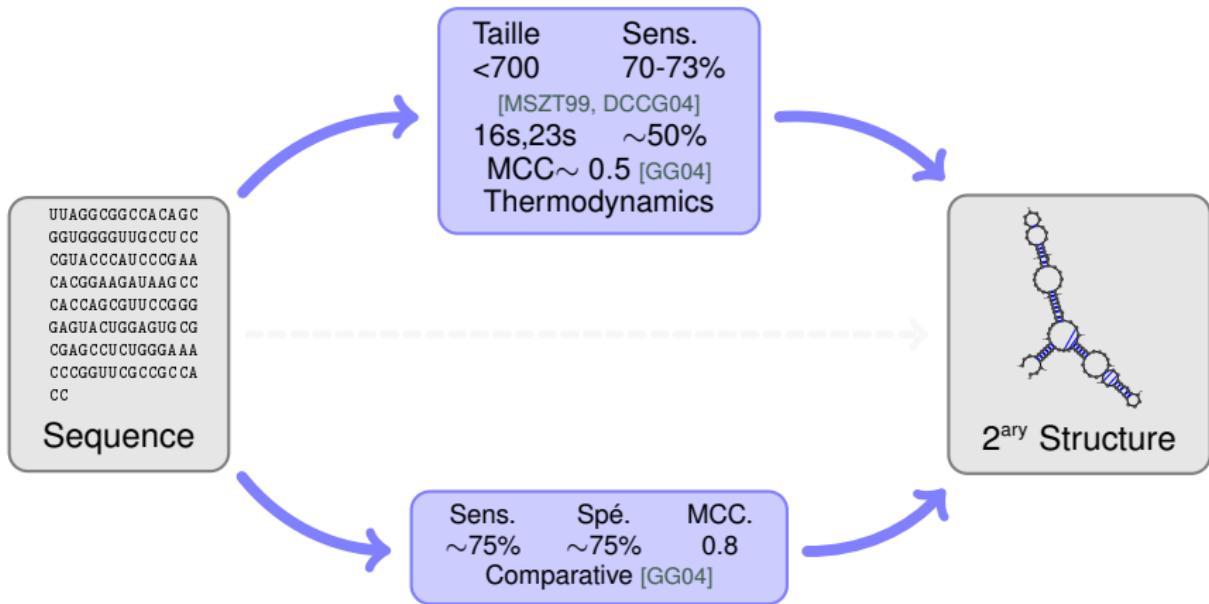
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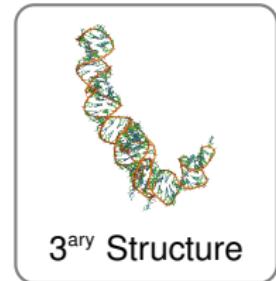
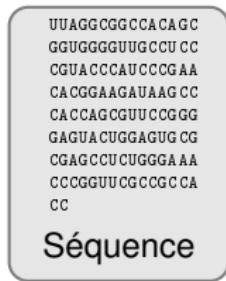


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Goal: From sequence to all-atom/coarse grain 3D models!!!

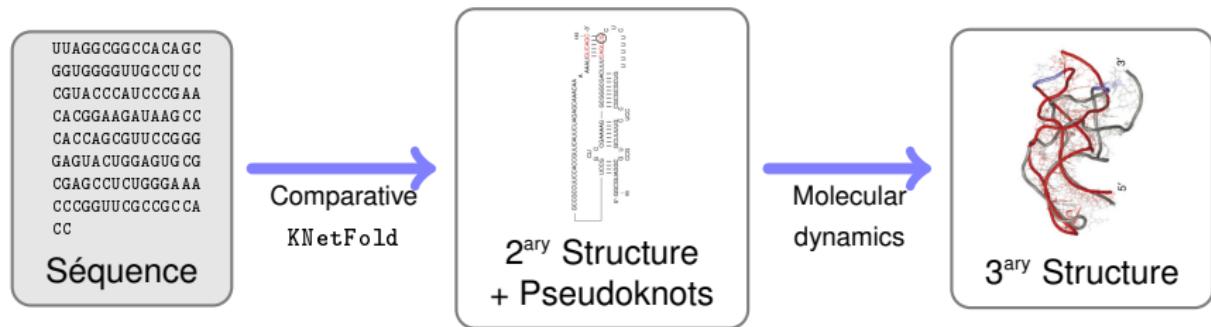
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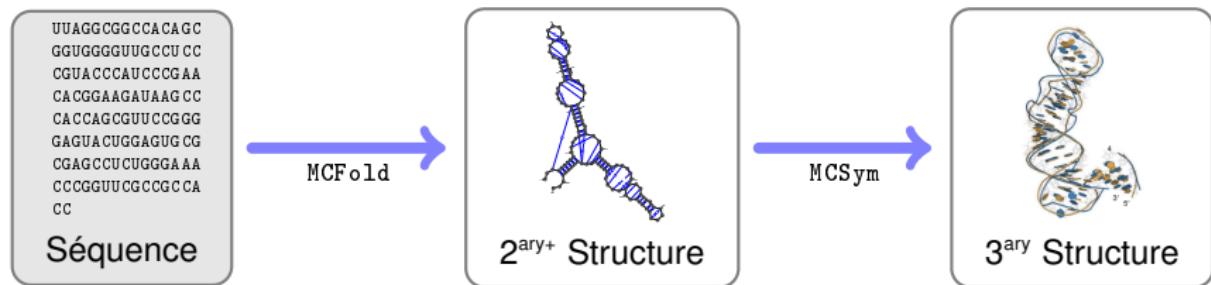
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Exercise: Parsing/#Structs/Folding RNAs (Python)

<http://www.lix.polytechnique.fr/~ponty/#teaching>