

PROBLEM FORMULATION

We consider the following version of the dynamic lot-sizing problem:

$$\text{minimize } \mathbb{E} \left[\sum_{t=1}^T P_t x_t + \psi(s_T) \right]$$

$$\text{s.t. } s_t = s_{t-1} + x_t - D_t \quad \forall t$$

$$s_t \stackrel{\text{a.s.}}{\in} [0, S] \quad \forall t$$

$$|x_t - D_t| \stackrel{\text{a.s.}}{\in} [-C^{\max}, C^{\max}] \quad \forall t$$

$$\text{non-anticipativity } \begin{cases} s_t \text{ a function of } (P_{[t-1]}, D_{[t]}) \\ x_t \text{ a function of } (P_{[t-1]}, D_{[t-1]}) \end{cases} \quad \forall t$$

(Note that x_t is decided before D_t is revealed.)

Parameters:

- T : horizon
- S : maximal inventory
- $C^{\max}$: maximal inventory variation
- D_t : random demand
- P_t : random price
- ψ : final cost

Decision variables:

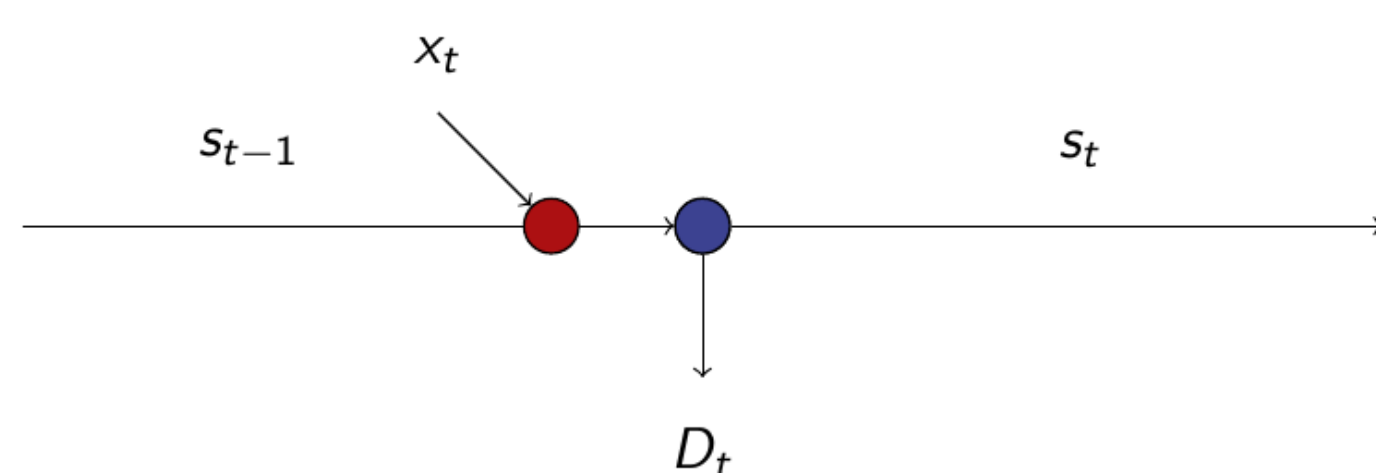
- s_t : inventory
- x_t : order size

MOTIVATION AND DYNAMICS

The motivation comes from a collaboration with the French electricity provider EDF : **how to manage optimally a battery on the electricity “intraday” market?**

Interpretation in the electricity context:

- Horizon $\iff T$ time steps of one hour
- Inventory $\iff$ energy level
- Order size $\iff$ electricity charge or discharge
- Demand $\iff$ reserve activation: the battery can be activated by the transmission system operator (TSO) to stabilize the network.
- $\psi \iff$ to anticipate the next day



PRELIMINARY REMARKS

- The optimal value is convex with respect to the initial inventory (easy to prove).
- Feasible initial inventories form an interval (easy to prove as well).

NUMERICAL RESULTS

Realistic instances: P_t auto-regressive, D_t uniform

Lower bound given by Theorem 2 (better than removing non-anticipativity)

s_0	$C^{\max}$	Anticipate	Thm 2
30	20	-12521	-2501
30	80	-20055	-2501
50	20	-13939	-2696
50	80	-22526	-2696
70	20	-15270	-2892
70	80	-22274	-2892

Table 1: Experimental results for S

CONTRIBUTIONS

For the following results, we assume that the essential supremum and infimum of D_t are known. The policy $\pi_t(s_{t-1}, P_{[t-1]})$ decides x_t .

Proposition 1 (Meunier and S. E. 2024). *Deciding whether s_t is a feasible inventory at time t can be done in linear time.*

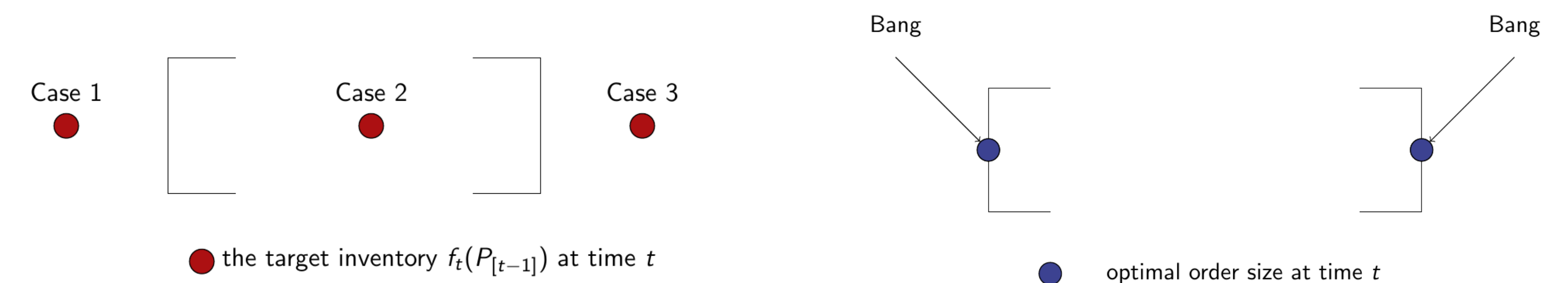
Theorem 1 (Meunier and S. E. 2024). *There exist measurable functions $\ell_t, u_t: [0, S] \rightarrow \mathbb{R}$ and $f_t: \mathbb{R}^{t-1} \rightarrow \mathbb{R}$ such that an optimal policy is given when s_{t-1} is feasible by*

$$\pi_t^*(s_{t-1}, P_{[t-1]}) = \begin{cases} \ell_t(s_{t-1}) - s_{t-1} & \text{if } \ell_t(s_{t-1}) > f_t(P_{[t-1]}), \\ u_t(s_{t-1}) - s_{t-1} & \text{if } u_t(s_{t-1}) < f_t(P_{[t-1]}), \\ f_t(P_{[t-1]}) - s_{t-1} & \text{otherwise.} \end{cases}$$

Up to an extra assumption, the previous theorem can be made more explicit.

Theorem 2 (Meunier and S. E. 2024). *Suppose $C^{\max} = +\infty$. Then the value function depends affinely on the current inventory and an optimal policy when $t < T$ is given by*

$$\pi_t^*(s_{t-1}, P_{[t-1]}) = \begin{cases} \text{ess sup}(D_t) - s_{t-1} & \text{if } \mathbb{E}[P_{t+1} - P_t | P_{[t-1]}] \geq 0, \\ S + \text{ess inf}(D_t) - s_{t-1} & \text{otherwise.} \end{cases}$$



Theorem 1

Theorem 2

Proposition 2 (Meunier and S. E. 2024). *Suppose $\psi = 0$. Assume furthermore that the prices form a martingale and all have the same sign with probability one. Up to a linear-time precomputation, an optimal policy at time t can be computed in constant time.*

OPEN QUESTION

The functions ℓ_t and u_t of Theorem 1 actually admit closed-form expressions. This is not the case for the functions f_t .

Open question. *Under which conditions the quantity $f_t(P_{[t-1]})$ in Theorem 1 can be efficiently computed?*

Theorem 2 and Proposition 2 provide examples of such conditions.