

# SUMMER SCHOOL ON ASPECTS OF OPTIMIZATION

Discrete Optimization  
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- 1 Tailored cutting-plane methods
  - Introduction
  - Strong inequalities
  - Polytope Dimension
  - Facet and Convex Hull Proofs
  - Lifting
  - Separation procedure
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## Decomposition :

$X = X^1 \cap X^2 \rightarrow$  focus on a simpler set, e.g.  $X^2$  because inequalities valid for  $X^2$  will be valid for  $X$ .



- Definition of strong inequalities

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# Strong valid inequalities

## Remark

Inequalities  $\pi^\top x \leq \pi^0$  and  $\lambda \pi^\top x \leq \lambda \pi^0$  are identical for any  $\lambda > 0$ .



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## Definition

If  $\pi^\top x \leq \pi^0$  and  $\mu^\top x \leq \mu^0$  are two valid inequalities for  $X \subseteq \mathbb{R}_+^n$ ,  $\pi^\top x \leq \pi^0$  dominates  $\mu^\top x \leq \mu^0$  if there exists  $u > 0$  such that  $\pi > u\mu$  and  $\pi^0 \leq u\mu^0$ .

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## Definition

An inequality  $\pi^\top x \leq \pi^0$  is redundant in the description of  $X \subseteq \mathbb{R}^n$ , if there exist  $k \geq 2$  valid inequalities  $(\pi^i)^\top x \leq \pi^{i0}$  for  $i = 1, \dots, k$  for  $X$  and weights  $u_i > 0$  for  $i = 1, \dots, k$  such that  $(\sum_{i=1}^k u_i \pi^i)^\top x \leq \sum_{i=1}^k u_i \pi^{i0}$  dominates  $\pi^\top x \leq \pi^0$ .

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- **Theoretically**: it is important to know which inequalities are nonredundant
- **Practically** : it is important to avoid using an inequality when one that dominates it is available

# Valid inequalities for the 01-KP problem

Classical formulation

$$K = \left\{ x \in \{0, 1\}^n \mid \sum_{j=1}^n w_j x_j \leq c \right\}$$

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The sets  $K$  and  $K^C$  coincide.



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Which formulation is stronger?

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$n = 3, w = (3, 3, 3), c = 5$ , thus

$$K = \{x \in \{0, 1\}^n \mid 3x_1 + 3x_2 + 3x_3 \leq 5\}$$

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Summing up the 3 inequalities of  $K^C$  we have  $2x_1 + 2x_2 + 2x_3 \leq 3$  which is implied by  $x_1 + x_2 + x_3 \leq 1$ , thus  $K$  is stronger than  $K^C$  (e.g.  $(\frac{1}{2}, \frac{1}{2}, \frac{1}{2})$ ).

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- Can we improve  $K^C$ ?  $\rightarrow$  strong valid inequalities
- $K^C$  has an exponential number of constraints  $\rightarrow$  no a-priori addition
- Use “some” constraints of  $K^C$  to improve  $K$   $\rightarrow$  separation procedure

# Strong Valid Inequalities

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## Theorem

*If  $P$  is a full-dimensional polyhedron, it has a unique minimal description*

$$P = \{x \in \mathbb{R}^n \mid (a^i)^\top x \leq b_i \text{ for } i = 1, \dots, m\}$$

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*where each inequality is unique within a positive multiple.*

Thus, any other valid inequality is redundant.

## Definition

The points  $x^1, x^2, \dots, x^k \in \mathbb{R}^n$  are affinely independent if the only solution to the linear system

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## Definition

The dimension of  $P$ ,  $\dim(P)$ , is one less than the maximum number of affinely independent points in  $P$ .

## Definition

- $F$  defines a face of the polyhedron  $P$  if  $F = \{x \in P \mid \pi^\top x = \pi^0\}$  for some valid inequality  $\pi^\top x \leq \pi^0$  of  $P$ . The valid inequality  $\pi^\top x \leq \pi^0$  is said to represent or define a face.

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- Alternatively,  $F$  is a face of the polyhedron  $P$  if there is a vector  $c$  for which  $F$  is the set of vectors attaining  $\max\{c^\top x \mid x \in P\}$  provided that this maximum is finite.

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- $F$  is a facet of  $P$  if  $F$  is a face of  $P$  of  $\dim(F) = \dim(P) - 1$ .

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Thus, for full-dimensional polyhedra,  $\pi^\top x \leq \pi^0$  defines a facet of  $P$  if and only if there are  $n$  affinely independent points of  $P$  satisfying it at equality.

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# Polytope dimension

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## Method 2:

- Find a system  $Ax = b$  of equations such that  $S \subseteq \{x \in \mathbb{R}^n \mid Ax = b\}$  and  $\text{rank}(A) = n - k$
- Prove that every equation  $\alpha^\top x = \beta$  satisfied by all  $x \in S$  is a linear combination of  $Ax = b$  (i.e., there exists  $u$  s.t.  $\alpha = uA$  and  $\beta = ub$ ). Thus,  $Ax = b$  is the affine hull of  $S$ , thus  $S$  has dimension  $k$ .



Example: 01-KP

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$$S = \text{conv}(\{x \in \{0,1\}^n \mid w^\top x \leq c\})$$

## Proposition

The dimension of  $S$  is  $n - |J|$ , with  $J = \{j \mid w_j > c \text{ for } j = 1, \dots, n\}$ .

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Hp.  $\text{conv}(X)$  is bounded and full-dimensional.

- Given  $X \in \mathbb{Z}_+^n$  and a valid inequality  $\pi^\top x \leq \pi^0$  for  $X$ , how to prove that  $\pi^\top x \leq \pi^0$  is a facet of  $\text{conv}(X)$ ?

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## Approach 1:

- Find  $n$  points in  $X$  that satisfy  $\pi^\top x = \pi^0$  and prove they are affinely independent.

# Facet and Convex Hull Proofs

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## Approach 2:

- Select  $t \geq n$  points  $x^1, \dots, x^t \in X$  satisfying  $\pi^\top x = \pi^0$ . Suppose that all these points lie on a generic hyperplane  $\mu^\top x = \mu^0$ .
- Solve the linear equation system

$$\sum_{j=1}^n \mu_j x_j^k = \mu^0 \quad \text{for } k = 1, \dots, t$$

in the  $n + 1$  unknowns  $(\mu, \mu^0)$

- If the only solution is  $(\mu, \mu^0) = \alpha(\pi, \pi^0)$  for  $\alpha \neq 0$ , then the inequality  $\pi^\top x = \pi^0$  is facet-defining.



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- Given  $X \in \mathbb{Z}_+^n$  and a polyhedron  $P = \{x \in \mathbb{R}^n \mid Ax \leq b\}$ , show  $P$  describes the  $\text{conv}(X)$ .

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## Approach 2:

- Show that points  $x^* \in P$  with a fractional element are not extreme points of  $P$ .

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## Approach 4:

- Show that if  $\pi^\top x \leq \pi^0$  defines a facet of  $\text{conv}(X)$  then it must be identical to one of the inequalities  $a_i^\top x \leq b_i$  defining  $P$ .

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## Approach 5:

- Show that for any  $c \in \mathbb{R}^n, c \neq 0$  the set of optimal solutions to the problem  $\min\{c^\top x \mid x \in X\}$  lies in  $\{x \mid a_i^\top x = b_i\}$  for some  $i = 1, \dots, m$ .



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Etc.

## Some remarks

Consider the decomposition  $X = X^1 \cap X^2 \rightarrow$  focus on a simpler set, e.g.  $X^2$  because inequalities valid for  $X^2$  will be valid for  $X$ .

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Computational results to show method effectiveness.

# Strengthening Cover Inequalities: Extended Covers

Yesterday we have seen the Extended Cover Inequalities for the 01-KP problem.

## Proposition

If  $C$  is a cover for  $X$ , the feasible set of the 01-KP problem, the extended cover inequality

$$\sum_{j \in E(C)} x_j \leq |C| - 1$$

is valid for  $X$ , where  $E(C) = C \cup \{j \mid w_j \geq w_i \text{ for all } i \in C\}$ .

# Strengthening Cover Inequalities: Extended Covers

Yesterday we have seen the Extended Cover Inequalities for the 01-KP problem.

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Let  $C$  be a cover for  $K$ . The cover inequality associated with  $C$  is facet-defining for  $P_C = \text{conv}(K) \cap \{x \in \mathbb{R}^n \mid x_j = 0 \ \forall j \in \{1, \dots, n\} \setminus C\}$  if and only if  $C$  is a minimal cover.

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Given a cover  $C$  for the 01-KP problem, find the best possible values for  $\alpha_j$  for  $j \in \{1, \dots, n\} \setminus C$  such that

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Iteratively solve the problem

$$\xi_t = \max \sum_{i=1}^{t-1} \alpha_{j_i} x_{j_i} + \sum_{j \in C} x_j$$
$$\sum_{i=1}^{t-1} \alpha_{j_i} x_{j_i} + \sum_{j \in C} w_j x_j \leq b - w_{j_t}$$
$$x \in \{0, 1\}^{|C|+t-1}$$

for  $t$  the current iteration and  $j_1, \dots, j_r$  be an ordering of  $N \setminus C$ .

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The resulting Lift Cover Inequalities are **facet-defining**.

# Strengthening Cover Inequalities: Lifting

Generalization



# Strengthening Cover Inequalities: Lifting

## Generalization

### Proposition

Consider a set  $S \subseteq \{0, 1\}^n$  such that  $S \cap \{x \mid x_n = 1\} \neq \emptyset$  and let  $\sum_{j=1}^{n-1} \alpha_j x_j \leq \beta$  be a valid inequality for  $S \cap \{x \mid x_n = 0\}$ . Then

$$\alpha_n = \beta - \max \left\{ \sum_{j=1}^{n-1} \alpha_j x_j \mid x \in S, x_n = 1 \right\}$$

is the largest coefficient such that  $\sum_{j=1}^{n-1} \alpha_j x_j + \alpha_n x_n \leq \beta$  is valid for  $S$ . Furthermore, if  $\sum_{j=1}^{n-1} \alpha_j x_j \leq \beta$  defines a  $d$ -dimensional face of  $\text{conv}(S) \cap \{x_n = 0\}$ , then  $\sum_{j=1}^n \alpha_j x_j \leq \beta$  defines a face of  $\text{conv}(S)$  of dimension at least  $d + 1$ .

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Potentially as difficult as solving the 01-KP problem  $\rightarrow$  heuristics!



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