Name:
SSN:

CSE 428

Fall 1998

Final Exam 16 December 1998

The exam consists of 10 problems on 9 pages, totaling 200 points. Read each question carefully and use your time judiciously.

Write your name/number on every page.

1. For each of the following grammars,

(20 pts)

- state whether or not it is ambiguous; if it is ambiguous, give a string with more than one parse tree;
- state any operator precedences which are enforced;
- state any operator associativities which are enforced;

Note that even if a grammar is ambiguous, it can still enforce operator precedences and associativities.

(a)

$$\begin{array}{lll} E & :: = & E \text{``*'}F \mid F \\ F & :: = & E \text{``+''}F \mid F \text{``-''}G \mid G \\ G & :: = & N \mid Id \mid \text{``("E")"} \end{array}$$

(b)

$$E ::= E ** F \mid F$$
 $F ::= F *+ G \mid F *- G \mid G$ 
 $G ::= N \mid Id \mid "("E")"$ 

2. Consider the following program fragment:

(20 pts)

```
program main
   x : integer;

procedure p(a : integer)
begin
    ??? := ???
end p;

begin main
   x := 1;
   p(x);
   write(x);
end main;
```

(a) Replace ??? := ??? above with one assignment statement which yields a program that has different behaviors (values for x printed out) for call-by-value and call-by-reference parameter-passing modes, but identical behaviors for call-by-value-result and call-by-reference parameter-passing modes.

(b) Give a simple, but general, description for code which can be inserted as the body of p such that the resulting program has *identical* behavior under *call-by-value*, *call-by-value*, *call-by-value*, and *call-by-reference* parameter passing. Give your description in terms of the sets of variables that can be read from and written to. Your description should be general enough to cover the most cases of code which could be inserted in p and yield identical programs.

3. Give the most general types for each of the following function declarations.

(20 pts)

(a) fun com (f,g) x = f(g(x));

(b) fun mix(x,y) (f,g) = (f(x),g(y));

(c) fun Curry3 f x y z = f(x,y,z);

(d) fun unCurry3 f (x,y,z) = f x y z;

4. Recall the definition of iterators from Assignment 6:

```
(20 pts)
```

```
fun iter(f,v) 0 = v
    |iter(f,v) n = f(iter(f,v) (n-1), n);

fun gen_iter(f,g,h) (0,a) = g(a)
    |gen_iter(f,g,h) (n,a) = f(gen_iter(f,g,h) (n-1,h(n,a)), n,a);
```

For each of the following functions, give a brief explanation of what it does. (Each performs a function which can be described in a few words.)

5	Consider	the	following	informal	description	of IS-trees
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(30 pts)

A tree is either (i) a leaf containing a *string*; or (ii) a node containing an *integer* and two IS-trees.

(a) Give an inductive definition of IS-trees.

(b) Give an ML datatype declaration for IS-trees.

(c) Give the principle of induction for proving properties about IS-trees. (I.e., to prove a property *P* holds for all IS-trees, what must be shown.)

6.	An implementation of a statically-scoped, block-structured language might use activation records containing both static links and dynamic links. Briefly explain why a dynamically-scoped, block-structured language would not need static links. Describe what is needed in an activation record in this case and how these things are used.	(10 pts)
7.	Call-by-value and call-by-reference are two common modes of parameter passing. For each mode, give an advantage it has over the other mode:	(10 pts)
	(a) Advantage of CBV over CBR:	
	(b) Advantage of CBR over CBV:	

8. Consider the following typing rule for select expressions:

(30 pts)

$$\frac{\Gamma \vdash e_0 : \tau_0 \qquad \Gamma[x : \tau_0] \vdash e_1 : \tau \qquad \Gamma[y : \tau_0] \vdash e_2 : \tau}{\Gamma \vdash \mathtt{select}(e_0, \, x.e_1, \, y.e_2) : \tau}$$

(a) Below is a fragment the SML typechecker for expressions (from Analyze1.sml) which has been extended with an appropriate constructor for select expressions. Fill in the case for typechecking select expressions. (Recall that the type of Cxt is ((string \* tp) list).)

(b) Below is a fragment of the Prolog typechecker for expressions (from hw9-Starter.pl). Fill in the case for typechecking select expressions.

```
infer(_,N,int) :- integer(N).
infer(G,plus(E1,E2),int) :- infer(G,E1,int), infer(G,E2,int).
...
infer(G,select(E0,X,E1,Y,E2)) :-
```

- 9. Each of the following Prolog programs defines a relation p between two lists. In each (20 pts) case, briefly and concisely describe the relationship between lists xs and ys if p(xs,ys) is provable. For example, an answer might be "if p(xs,ys) is provable then ys is the reverse of xs. (Note: each of these programs is independent of the others.)
  - (a) p([],Ys). p([X|Xs],[X|Ys]) := p(Xs,Ys).

 $\begin{array}{ll} \text{(b)} & \quad \text{p([],Ys).} \\ & \quad \text{p([X|Xs],[Y|Ys]) :- p(Xs,Ys).} \end{array}$ 

(c) p([],[]).p([X|Xs],[X|Ys]) := p(Xs,Ys).

 $\begin{array}{lll} (\mathrm{d}) & & \mathtt{p}(\mathtt{X}\mathtt{s},\mathtt{Y}\mathtt{s})\,. \\ & & \mathtt{p}([\mathtt{X}|\mathtt{X}\mathtt{s}],[\mathtt{X}|\mathtt{Y}\mathtt{s}]) \,:-\, \mathtt{p}(\mathtt{X}\mathtt{s},\mathtt{Y}\mathtt{s})\,. \end{array}$ 

p :- q, s.	
p :- q, r.	
q :- t, u.	
q :- t.	
r :- r, v.	
r :- w.	
s :- t.	
t.	
t :- v.	
u :- w, u.	
v.	
w :- fail.	
x :- y,p.	
y :- x,q.	
(The proposition fail is never provable.)	
Prolog) will answer yes, no, or do something else (not terminate or report an error when runtime space is exceeded). Write yes, no, or other to indicate the behavior.  (a) p.	
(b) q.	
(c) r.	
(d) u.	
(e) x.	
11. If your life was a game of <i>Jeopardy!</i> what would the categories be?	(0 pts)

(20 pts)

10. Consider the following Prolog program.