

# Mots et racines dans les groupes de Coxeter

– 2, 9 et 16 octobre 2013 au LIX –

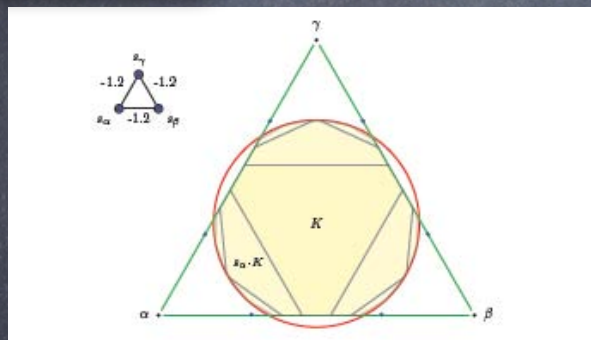
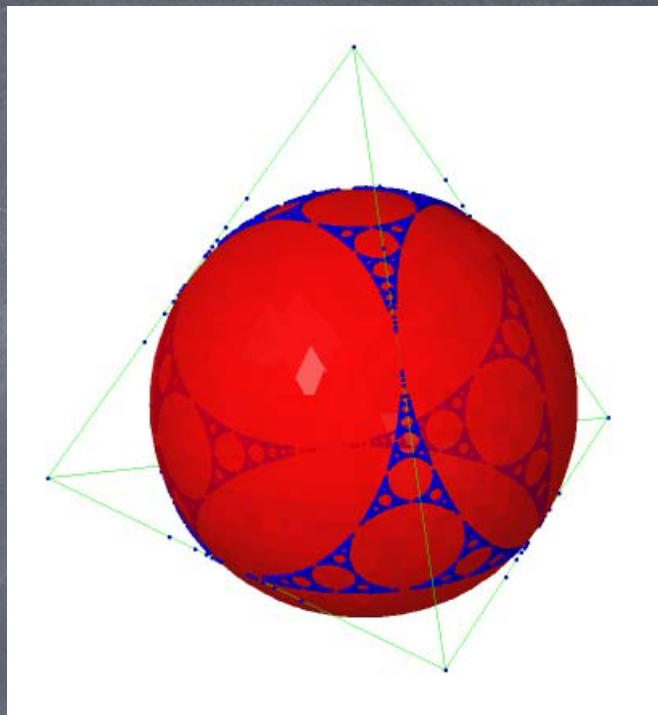
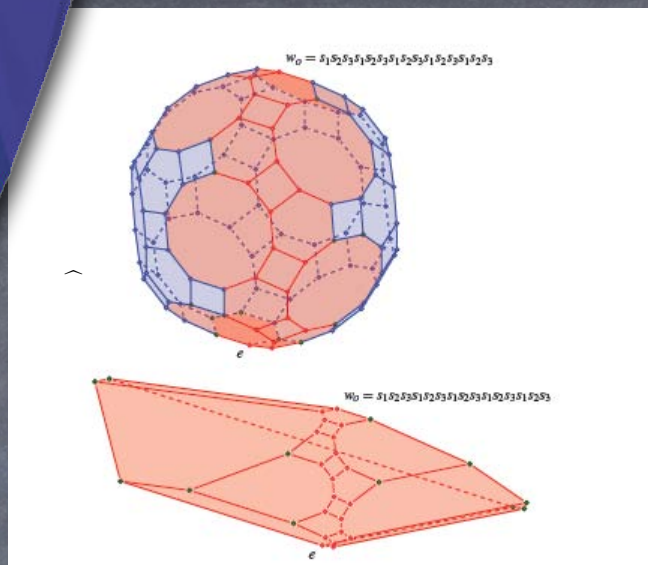
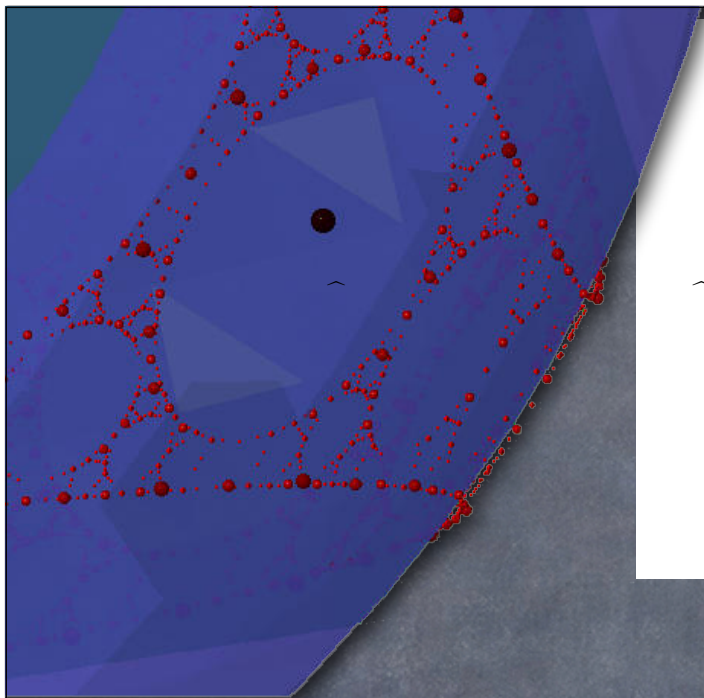
Christophe Hohlweg, LaCIM, UQAM  
(et LIX pour un mois)



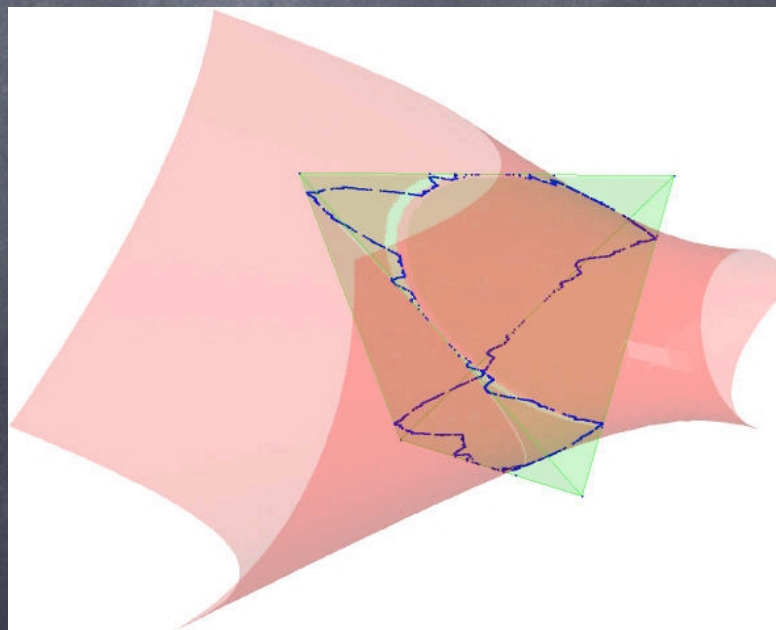
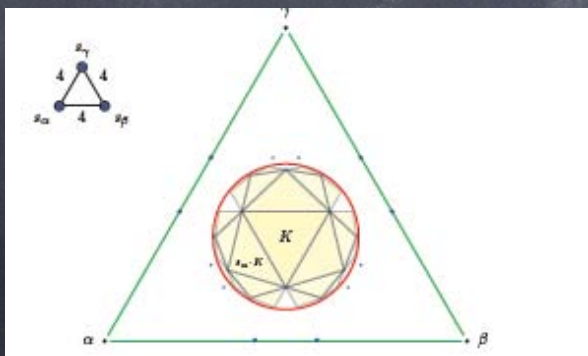
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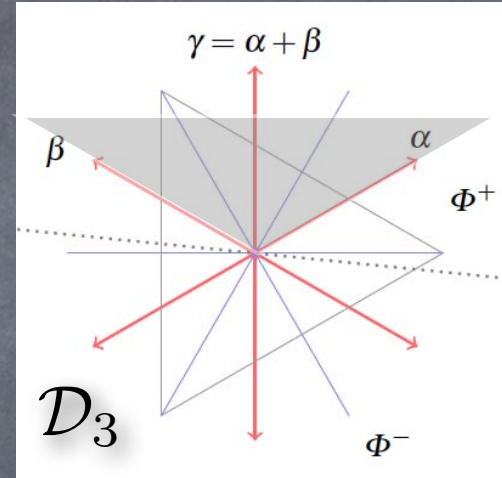


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# In the last episode



•  $W \leq O(V)$  FRG  $\longleftrightarrow \Phi$  root system in  $V$

• Separating  $\Phi$  by a (linear) hyperplane we have:

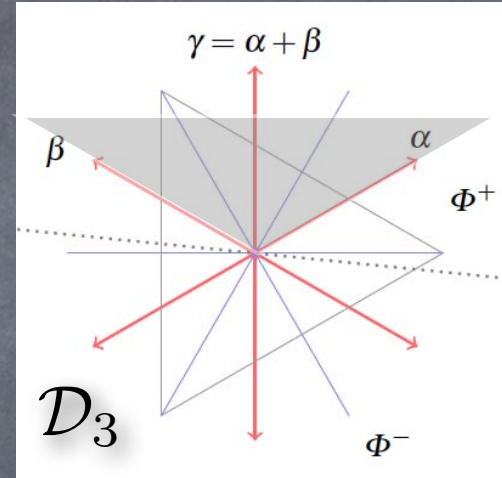
reflections	$T$	$\longleftrightarrow^{1:1}$	$\Phi^+$ positive roots
	$s_\beta$	$\longleftarrow$	$\beta$

simple reflections	$S \subseteq T$	$\longleftrightarrow^{1:1}$	$\Delta$ basis of $\text{cone}(\Phi^+)$
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**Theorem.**  $W$  is generated by  $S = \{s_\alpha \mid \alpha \in \Delta\}$



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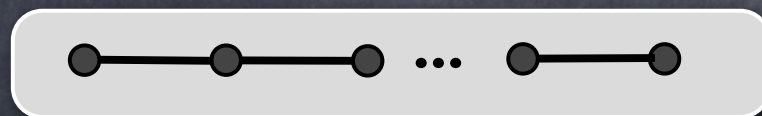
**Theorem.**  $W$  is generated by  $S = \{s_\alpha \mid \alpha \in \Delta\}$

**Problem:** find the relations for  $W = \langle S \rangle$ !

**Examples:**  $\mathcal{D}_m = \langle s, t \mid s^2 = t^2 = (st)^m = e \rangle$



$\mathcal{S}_n = \langle \tau_i \mid \tau_i^2 = (\tau_i \tau_j)^2 = (\tau_i \tau_{i+1})^3 = e, 1 \leq i < n, |j - i| > 1 \rangle$





# Length, inversions and relations in FRG

**Theorem.**  $W$  is generated by  $S = \{s_\alpha \mid \alpha \in \Delta\}$

- any  $w \in W$  is a word in the alphabet  $S$  ;
- Length function  $\ell : W \rightarrow \mathbb{N}$  with  $\ell(e) = 0$  and
$$\ell(w) = \min\{k \mid w = s_1 s_2 \dots s_k, s_i \in S\}$$

How to relate words on  $S$  representing  $w$ ? Is a word  $s_1 s_2 \dots s_k$  a **reduced expression** for  $w$  (i.e.  $k = \ell(w)$ ) ?

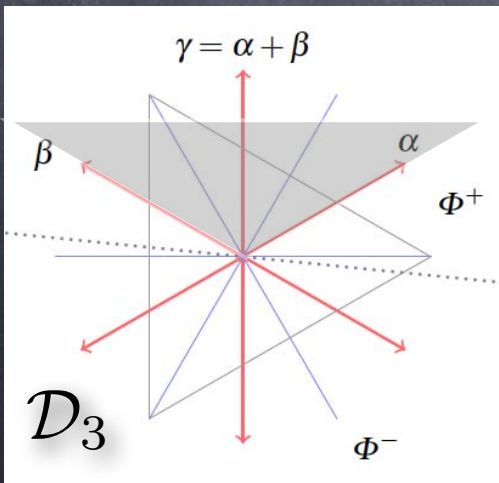


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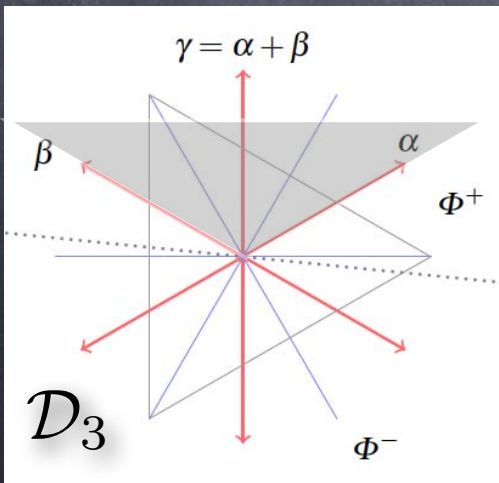


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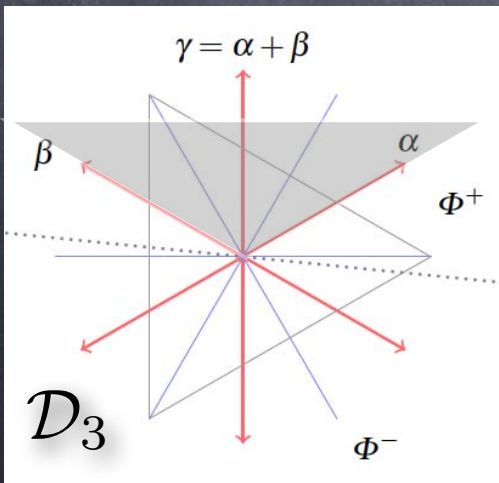


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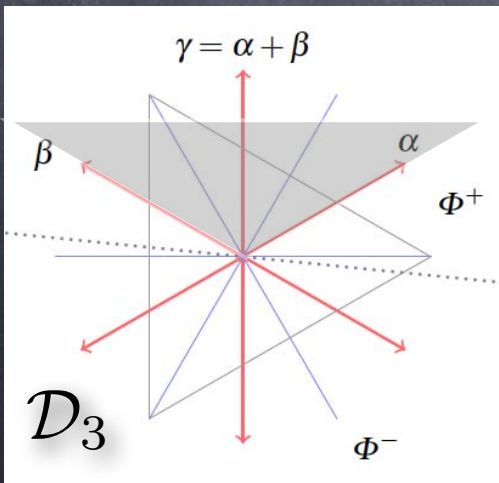


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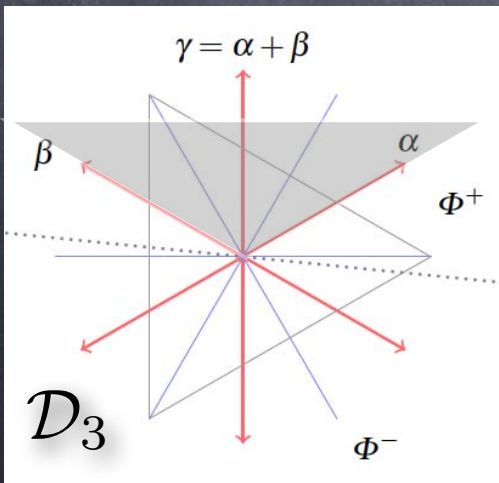


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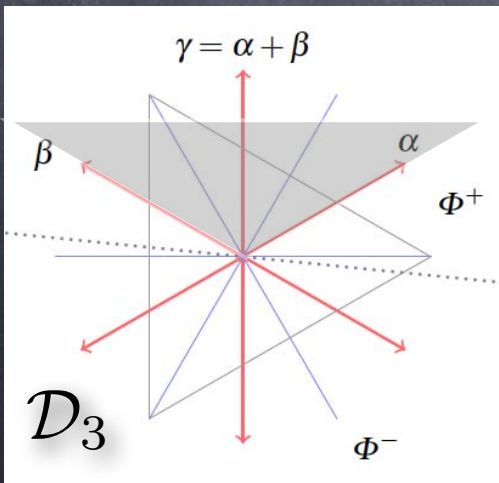


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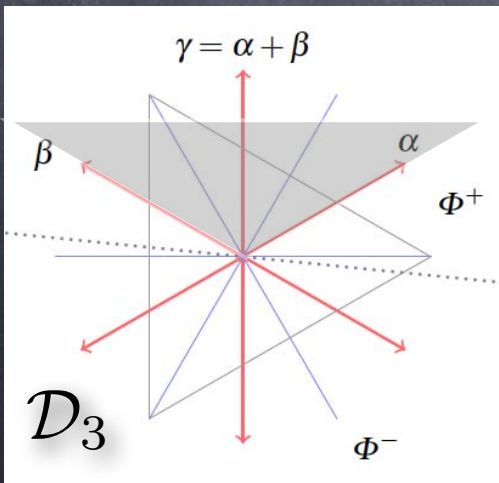


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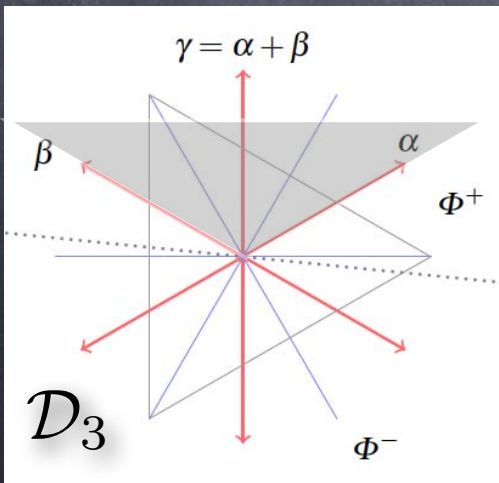


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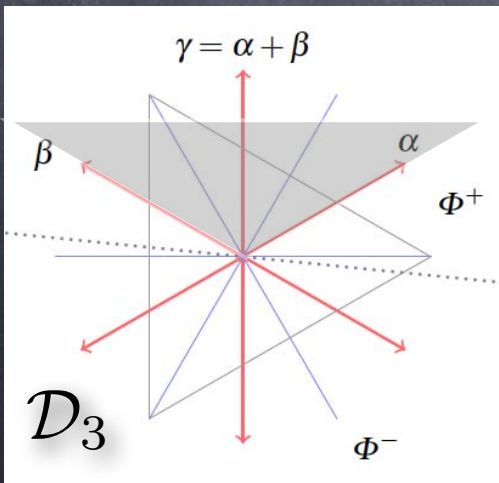


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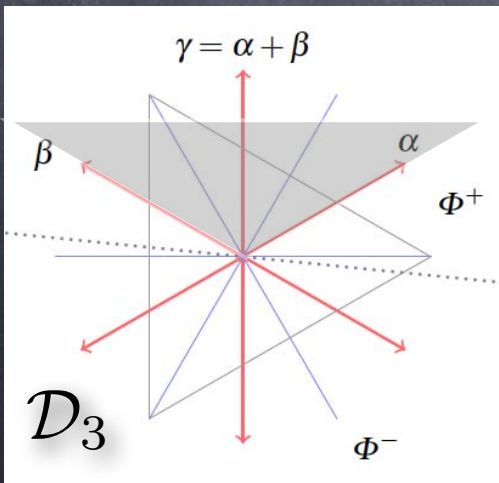


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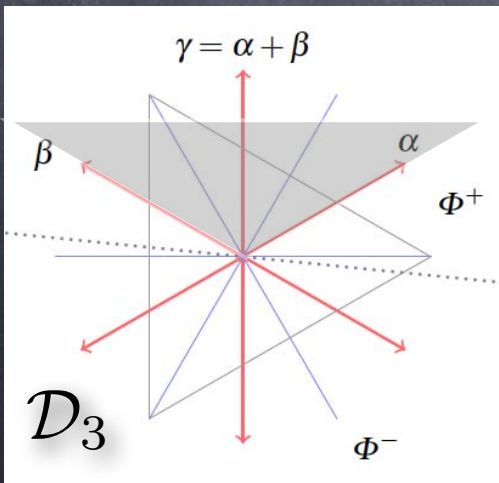


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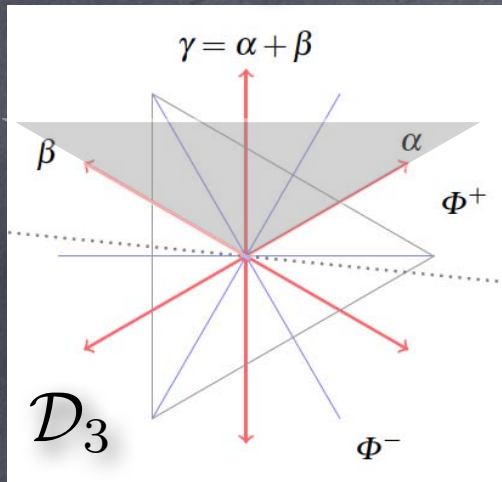


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$$\ell(w) = |\{\nu \in \Phi^+ \mid w(\nu) \in \Phi^-\}|$$



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**Definition.** The inversion set of  $w \in W$  is

$$\text{inv}(w) = \Phi^+ \cap w^{-1}(\Phi^-) = \{\nu \in \Phi^+ \mid w(\nu) \in \Phi^-\}$$

and  $\text{des}(w) = \text{inv}(w) \cap \Delta$  is its descent set.

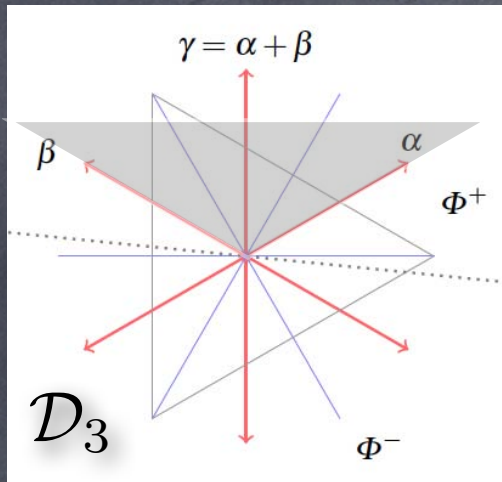
• If  $W = S_n$  then those “are” the natural inversion and descent statistics:

$$\text{inv}(\sigma) = \{e_j - e_i \mid 1 \leq i < j \leq n, e_{\sigma(j)} - e_{\sigma(i)} \in \Phi^-\}$$

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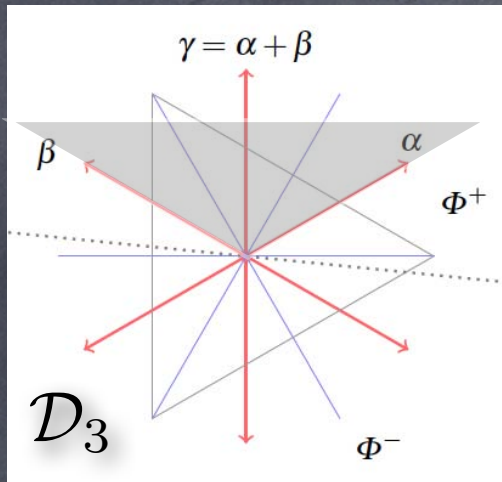
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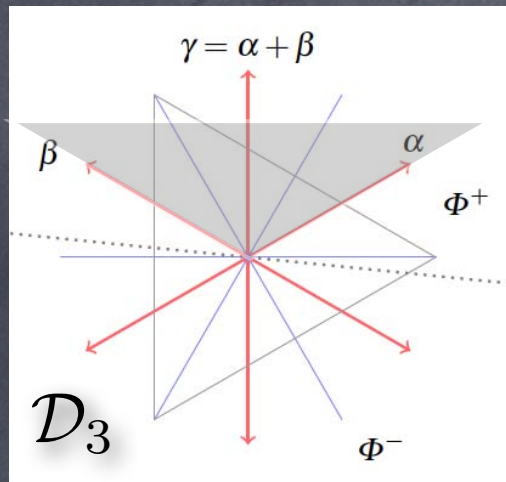
$$\text{inv}(\sigma) \simeq \{(i, j) \mid 1 \leq i < j \leq n, \sigma(i) > \sigma(j)\}$$

$$\text{des}(\sigma) \simeq \{i \mid 1 \leq i < n, \sigma(i) > \sigma(i+1)\}$$





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**Proposition.** Let  $w \in W$  and  $\alpha \in \Phi^+$ , then:

- (i)  $\ell(ws_\alpha) < \ell(w)$  if and only if  $\alpha \in \text{inv}(w)$  (i.e.  $w(\alpha) \in \Phi^-$ ).  
Otherwise,  $\ell(ws_\alpha) > \ell(w)$  if and only if  $w(\alpha) \in \Phi^+$ .
- (ii) If  $\alpha \in \Delta$ , then  $s_\alpha$  is a bijection on  $\Phi^+ \setminus \{\alpha\}$  and

$$\ell(ws_\alpha) = \begin{cases} \ell(w) - 1 & \text{if } \alpha \in \text{des}(w) \text{ i.e. } w(\alpha) \in \Phi^- \\ \ell(w) + 1 & \text{if } \alpha \notin \text{des}(w) \text{ i.e. } w(\alpha) \in \Phi^+ \end{cases}$$

- (iii)  $\ell(w) = |\text{inv}(w)|$ .

**N.B.:**  $(-w) : \Phi^+ \cap w^{-1}(\Phi^-) \rightarrow \Phi^+ \cap w(\Phi^-)$  is a bijection, so  $\ell(w) = \ell(w^{-1})$ .



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(iii)  $\ell(w) = |\text{inv}(w)|$ .

Critical tools for the proof are the following equivalent statements for a word  $w = s_1 \dots s_k \in W$

• **Exchange condition.** if  $\alpha \in \text{des}(w)$  then  $w$  may be rewritten with  $s_\alpha$  as the last letter:  $w = s_1 \dots \hat{s}_i \dots s_k s_\alpha$

• **Deletion condition.** the word  $w = s_1 \dots s_k$  has a subword that is a reduced expression for  $w$  obtained by deleting pairs of letters



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$$\ell(ws_\alpha) = \begin{cases} \ell(w) - 1 & \text{if } \alpha \in \text{des}(w) \text{ i.e. } w(\alpha) \in \Phi^- \\ \ell(w) + 1 & \text{if } \alpha \notin \text{des}(w) \text{ i.e. } w(\alpha) \in \Phi^+ \end{cases}$$

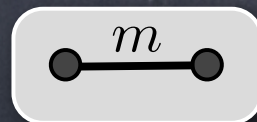
(iii)  $\ell(w) = |\text{inv}(w)|$ .

And finally

**Theorem.**  $W = \langle S \mid (st)^{m_{st}} = e \rangle$  where  $m_{st} = m_{ts}$  is the order of the rotation  $st$  (and  $m_{ss} = 1$ )



$\mathcal{D}_m$  is



and  $S_n$  is



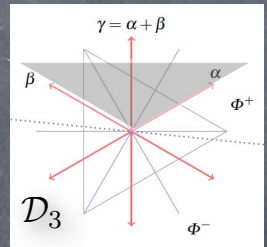


# Conclusion: words and roots in FRG

- $W \leq O(V)$  FRG  $\longleftrightarrow \Phi$  root system in  $V$
- Separating  $\Phi$  by a (linear) hyperplane we have:



reflections  $T$   $\xleftrightarrow{1:1}$   $\Phi^+$  positive roots  
 $s_\beta$   $\xleftarrow{\quad}$   $\beta$



simple reflections  $S \subseteq T$   $\xleftrightarrow{1:1}$   $\Delta$  basis of  $\text{cone}(\Phi^+)$

**Theorem.**  $W$  is generated by  $S = \{s_\alpha \mid \alpha \in \Delta\}$

- Length function  $\ell(w)$   $\longleftrightarrow \text{inv}(w)$  inversion set

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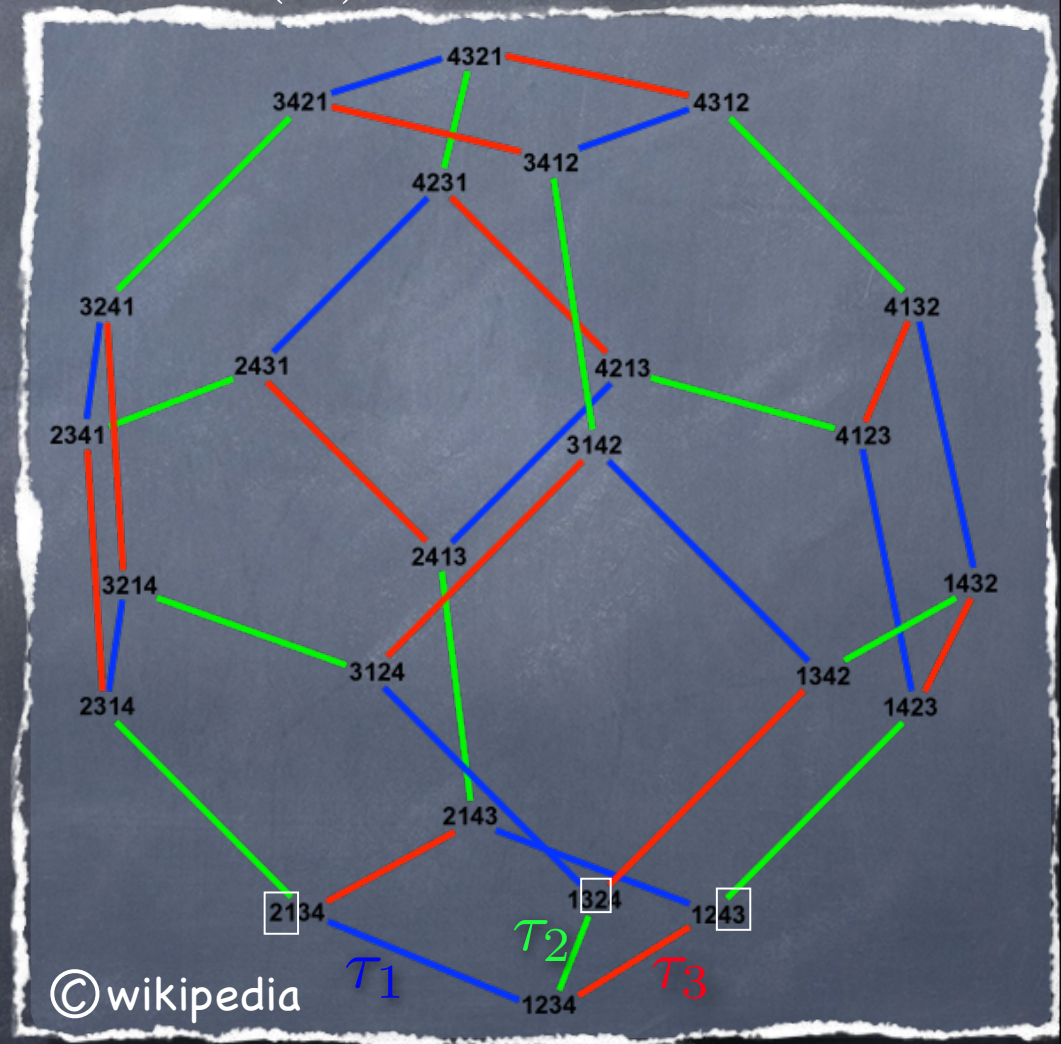
# On the Cayley graph of FRG

Theorem.  $W$  is generated by  $S = \{s_\alpha \mid \alpha \in \Delta\}$

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Cayley graph of  $W = \langle S \rangle$ :

- vertices  $W$
- edges  $w \xrightarrow{s} ws$   
( $s \in S$ )





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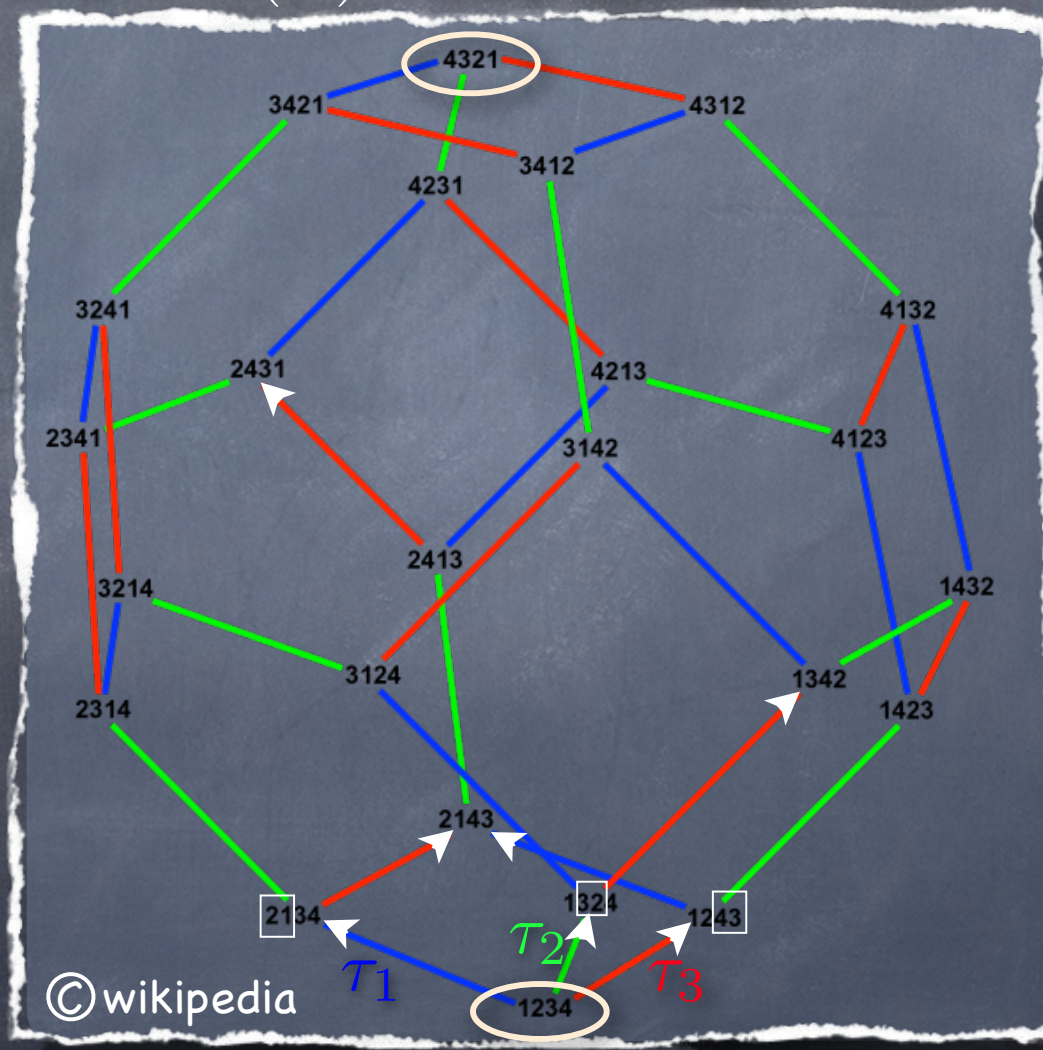
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The Cayley graph is naturally oriented by the (right) weak order:

$w < ws$  if  $\ell(w) < \ell(ws)$

write:  $w \xrightarrow{s} ws$





# On the Cayley graph of FRG

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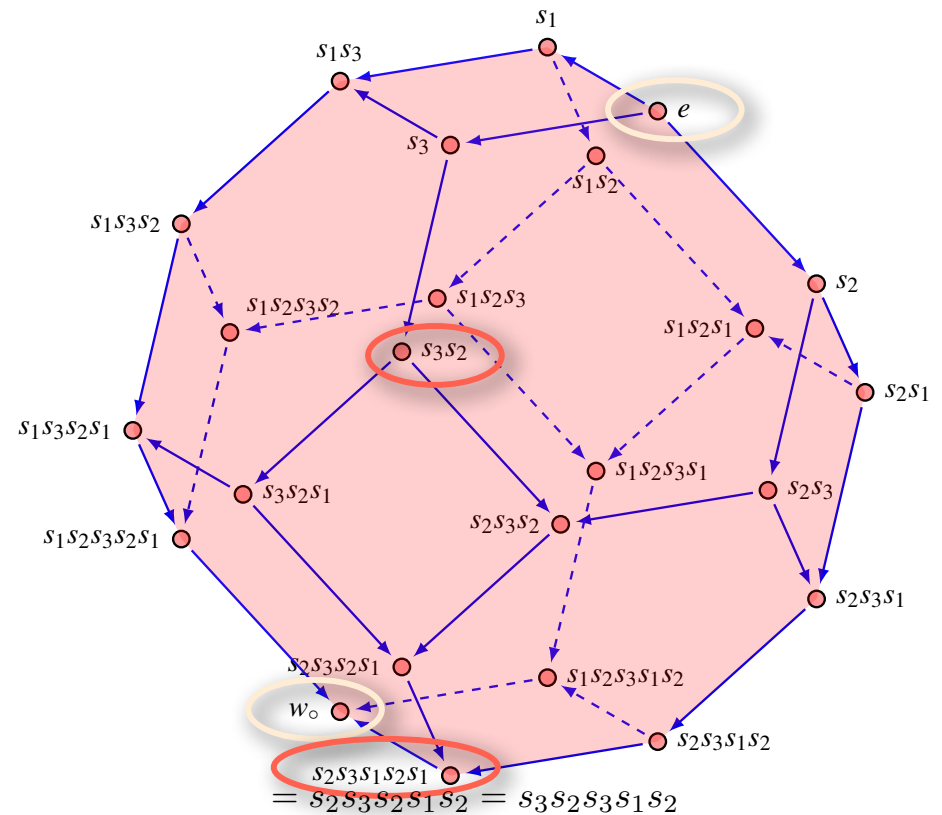
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**write:**  $w \xrightarrow{s} ws$





# On the Cayley graph of FRG

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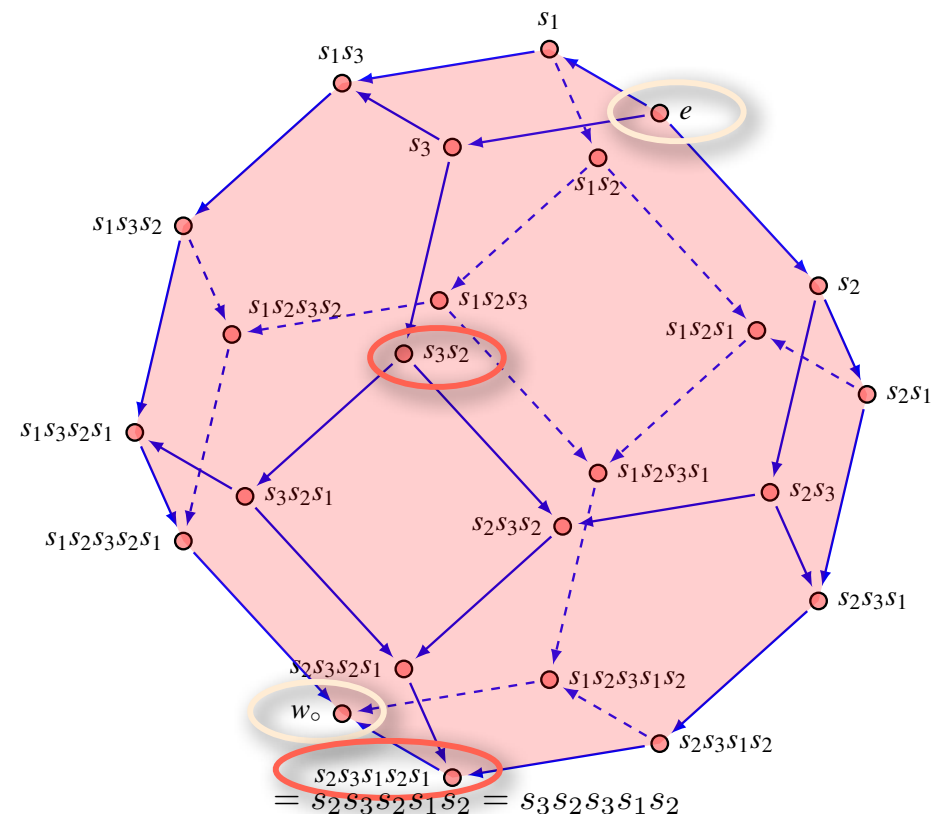
**Theorem.** The weak order is a lattice. Moreover:

□ reduced expression of  $w$  are in bijection with maximal chains in the interval  $[e, w]$ .

□  $u \leq w \iff \ell(u^{-1}w) = \ell(w) - \ell(u)$

□  $u \leq w$  iff a reduced expression of  $u$  is a prefix of a reduced expression of  $w$ .

□  $w_\circ$  is the unique element of maximal length





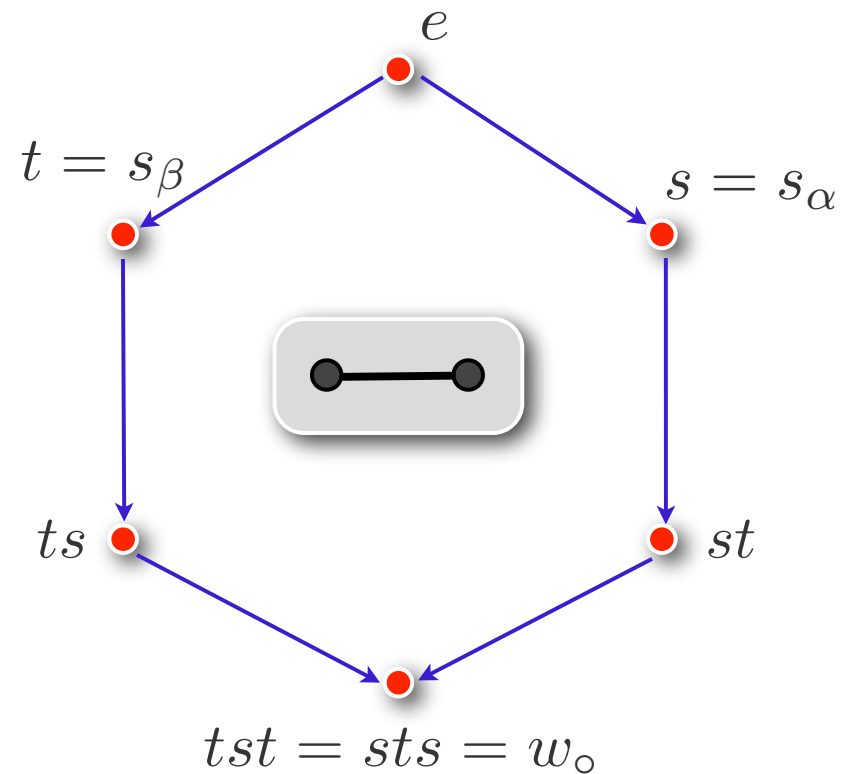
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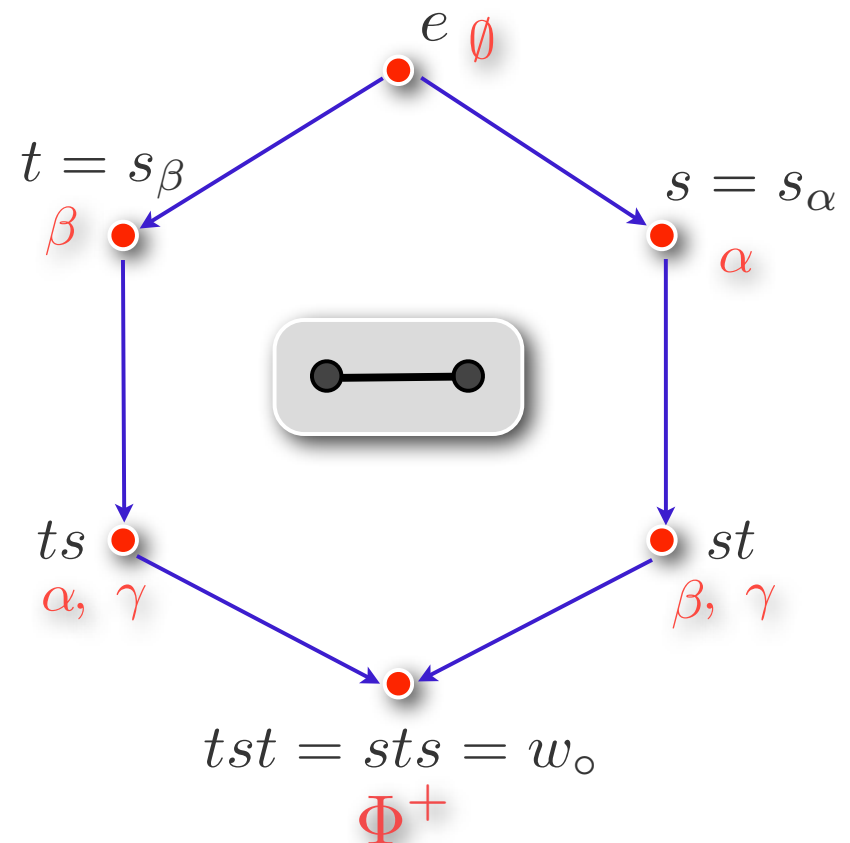




	$e$	$s = s_\alpha$	$t = s_\beta$	$st$	$ts$	$sts = tst$
$\ell$	0	1	1	2	2	3
$\alpha$	$\alpha$	$-\alpha$	$\gamma$	$\beta$	$-\gamma$	$-\beta$
$\beta$	$\beta$	$\gamma$	$-\beta$	$-\gamma$	$\alpha$	$-\alpha$
$\alpha + \beta = \gamma$	$\gamma$	$\beta$	$\alpha$	$-\alpha$	$-\beta$	$-\gamma$

**Theorem.** The weak order is a lattice. Moreover:

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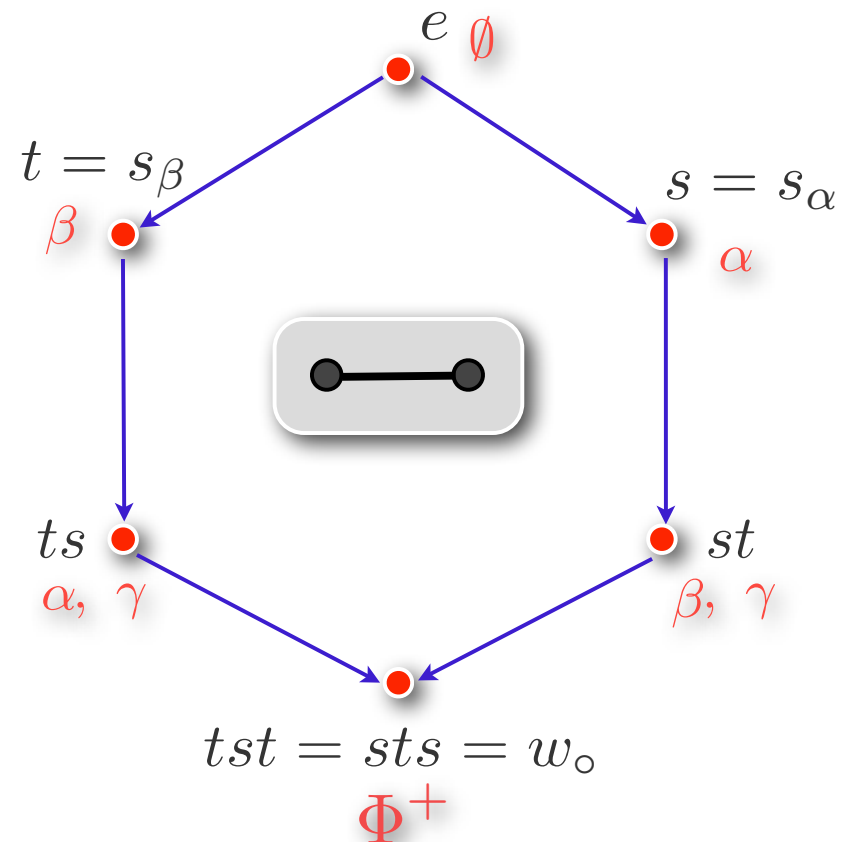
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# On the Cayley graph of FRG

**Theorem.**  $W$  is generated by  $S = \{s_\alpha \mid \alpha \in \Delta\}$

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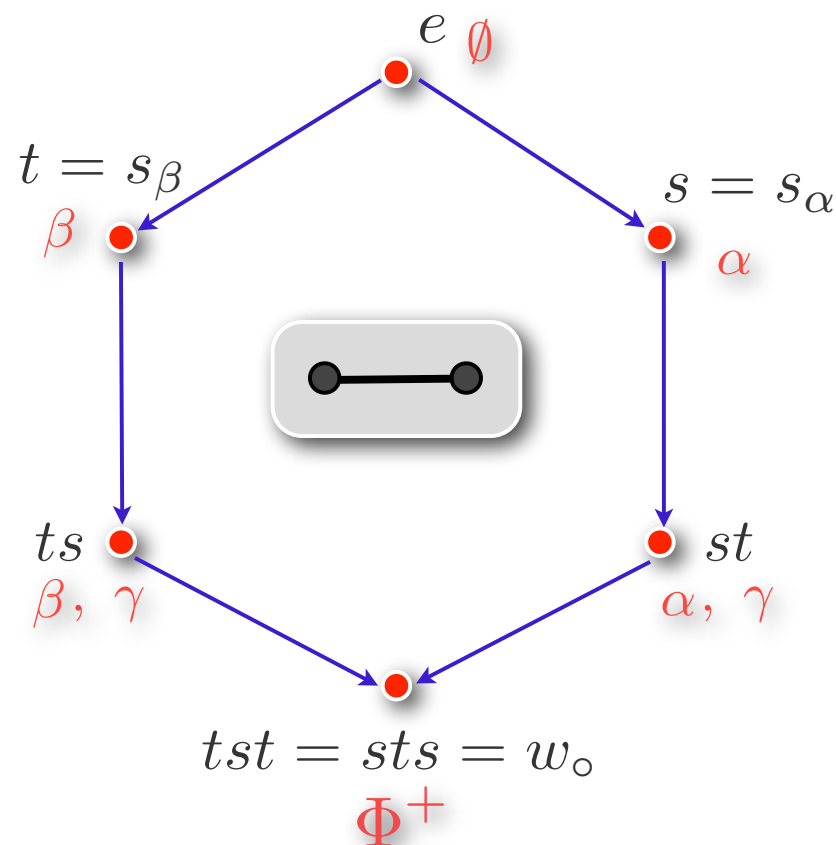
**Theorem.** The weak order is a lattice. Moreover:

$$\square u \leq w \iff \text{inv}(u^{-1}) \subseteq \text{inv}(w^{-1})$$

$\square w_o$  is the unique element of maximal length:  $\ell(w_o) = |\Phi^+|$

$\square \vee \neq \cup; \wedge \neq \cap$  so ... ?

$\square$  Problem: to understand reduced expressions as maximal chains in intervals of inversion sets? Count them!




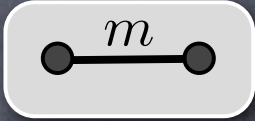
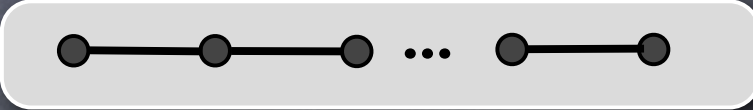


# Coxeter groups and Reflection groups

**Theorem.**  $W = \langle S \mid (st)^{m_{st}} = e \rangle$  where  $m_{st} = m_{ts}$  is the order of the rotation  $st$  (and  $m_{ss} = 1$ )

- So the presentation  $W = \langle S \mid (st)^{m_{st}} = e \rangle$  is illustrated with a **Coxeter graph**  $\Gamma_S$  :

- vertices  $S$  (i.e.  $\Delta$ )
- edges  for  $m_{st} \geq 3$


- $D_m$  is  and  $S_n$  is 

- A **Coxeter system**  $(W, S)$  is a group  $W = \langle S \rangle$  obtained from a Coxeter graph (allowing  $m_{st} = \infty$  if there is no relation between  $s$  and  $t$ , and  $m_{ss} = 1$ )

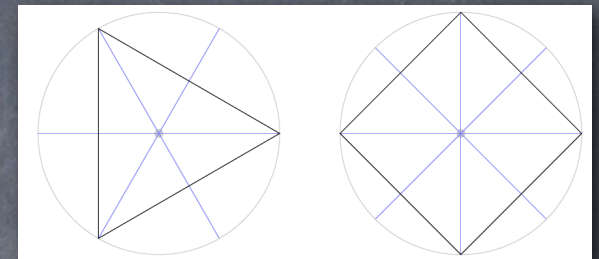


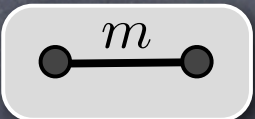

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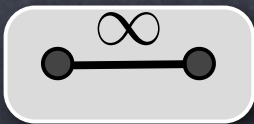

- $S_n$  is  of rank  $|S| = n - 1$

- Rank 2 Coxeter groups



- $D_m$  is  or  ( $m = 2$ ) - FRG

- A infinite example: the infinite dihedral group

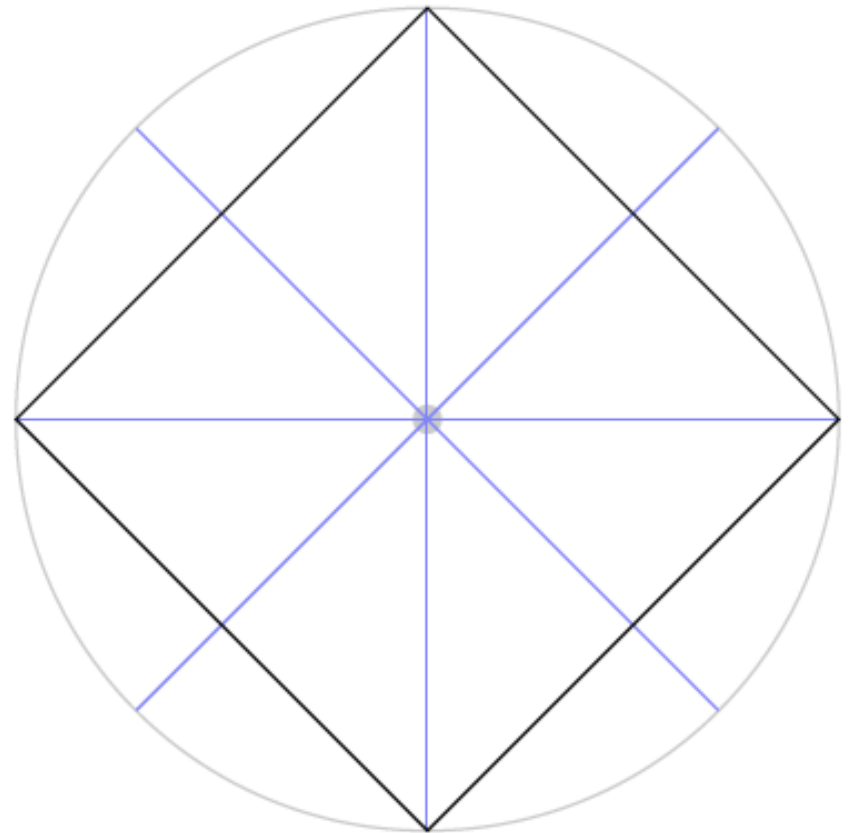
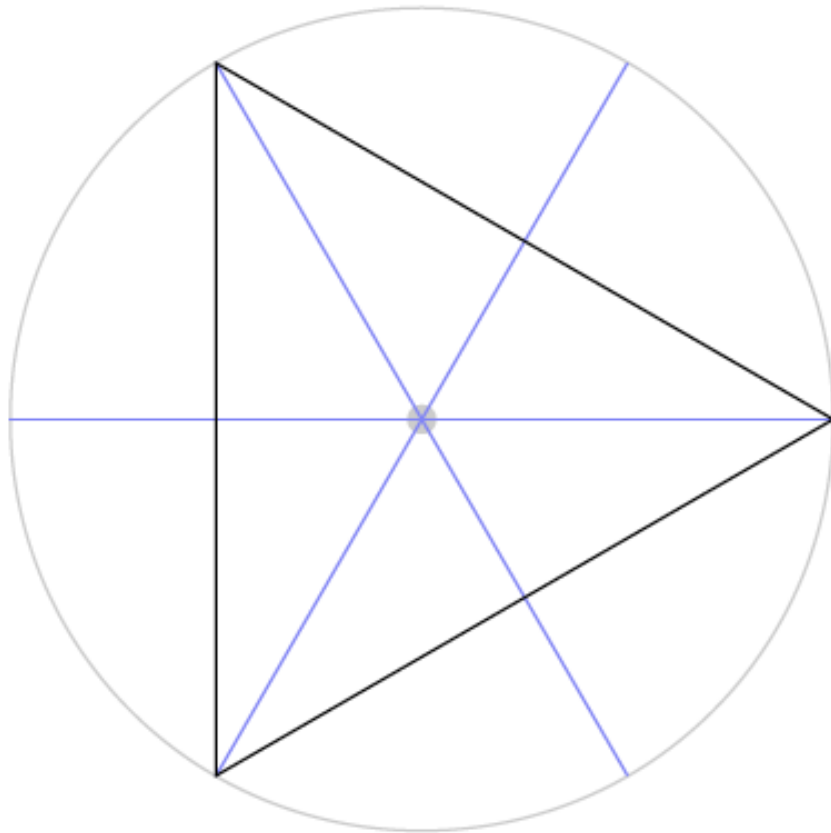
- $D_\infty$  is  (  ; Infinite number of words: )  
 $(ts)^p, (st)^p, s(ts)^p, t(st)^p, p \in \mathbb{N}$



# Coxeter groups and Reflection groups


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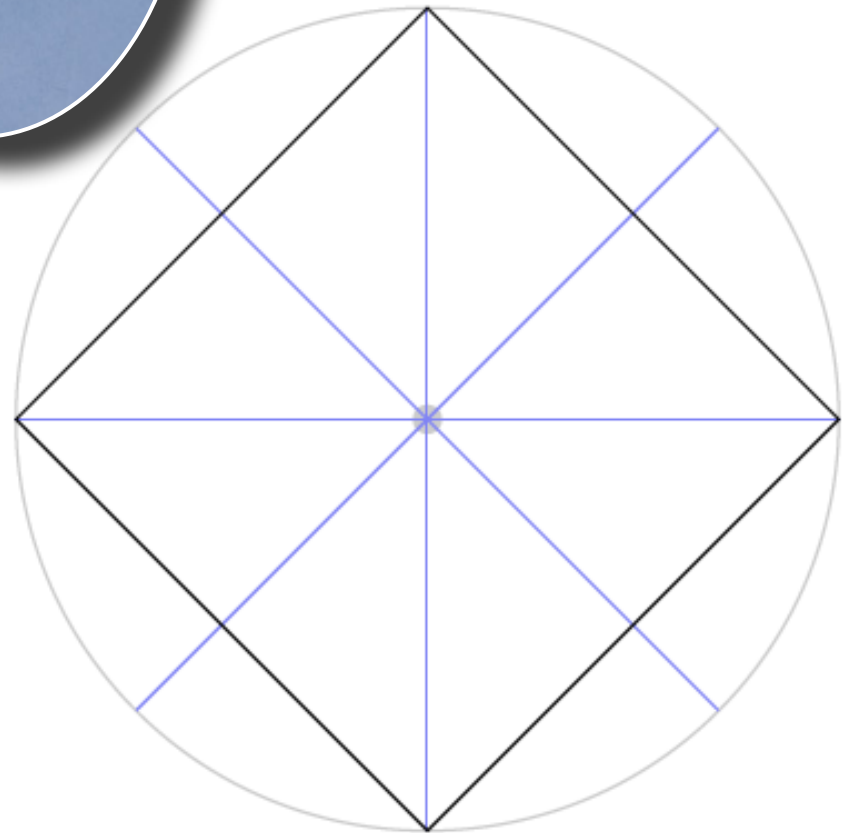
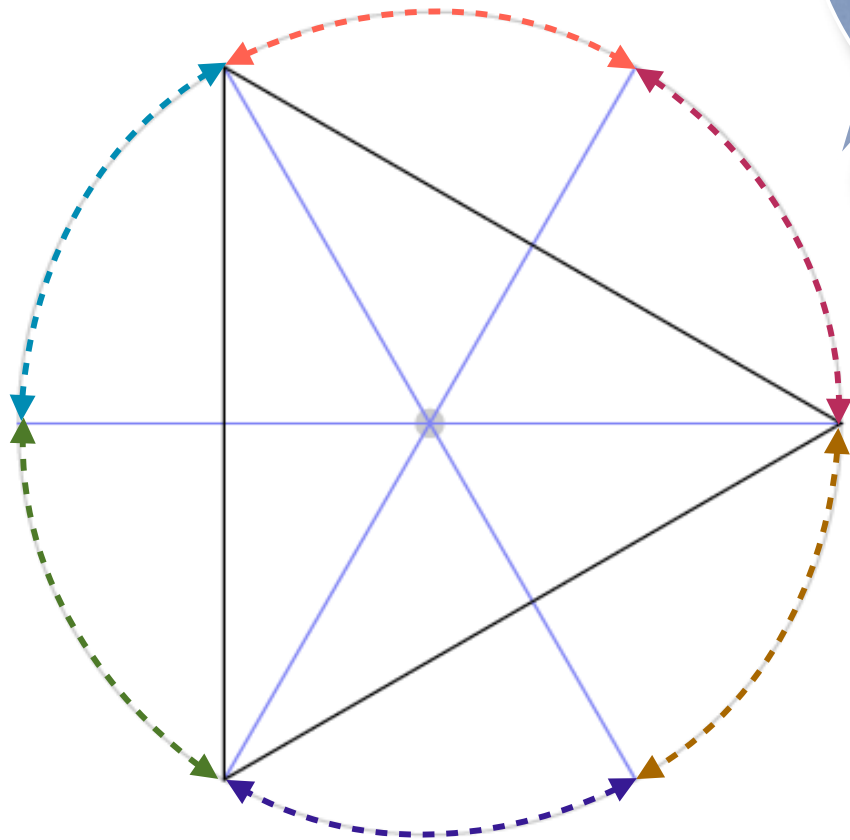




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Tilings in spherical geometry

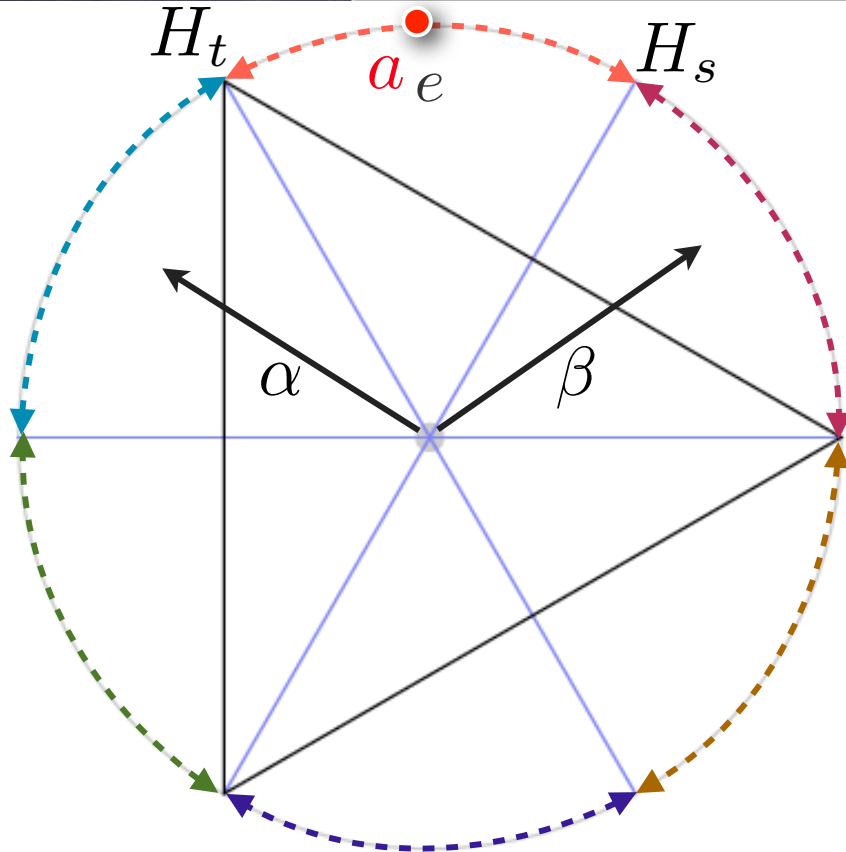




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# Tilings in spherical geometry, roots and words

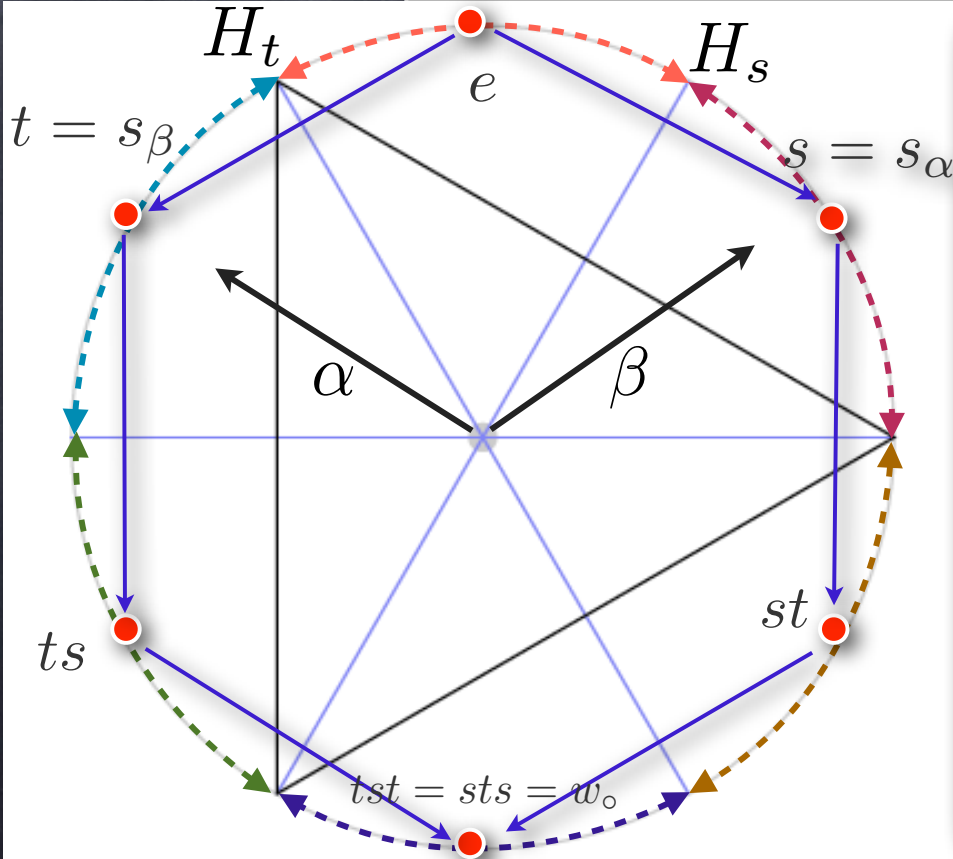
- $\Delta = \{\alpha, \beta\}$  simple system;
- $S = \{s = s_\alpha, t = s_\beta\}$ ;
- Choose a **generic** on a tile s.t.  
 $\langle a, \alpha \rangle > 0, \quad \langle a, \beta \rangle > 0$
- Label the corresp. tile by  $e$ .
- Then label by acting ...



# Coxeter groups and Reflection groups

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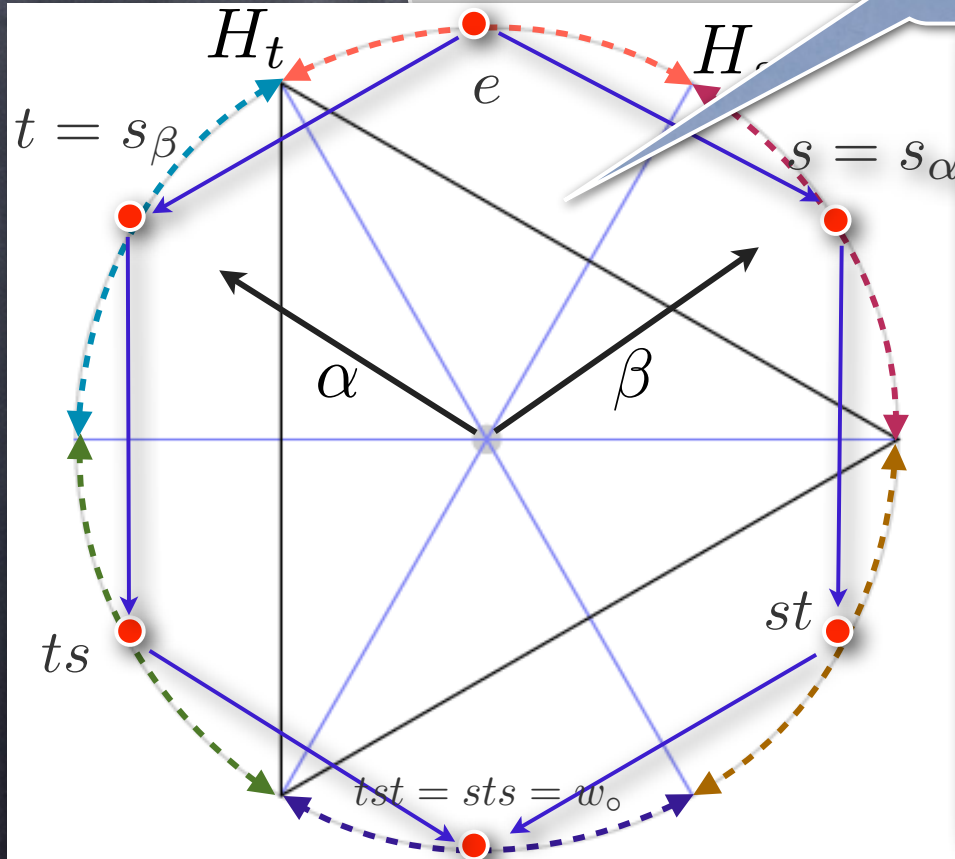
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# Coxeter groups and Reflection groups

- A **Coxeter system**  $(W, S)$  obtained from a Coxeter group if there is no relation

- $S_n$  is



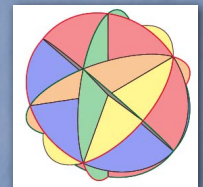
The convex hull of

$$w(a), w \in W$$

is a convex polytope called a **permutahedron**

$$\text{Perm}^a(W) = \text{conv} \{w(a) \mid w \in W\}$$

Tilings in spherical geometry,  
roots and words (general case)



©Pilaud-Stump

- $\Delta$  simple system;
- $S = \{s_\alpha \mid \alpha \in \Delta\}$ ;
- Choose a **generic** on a tile s.t.  

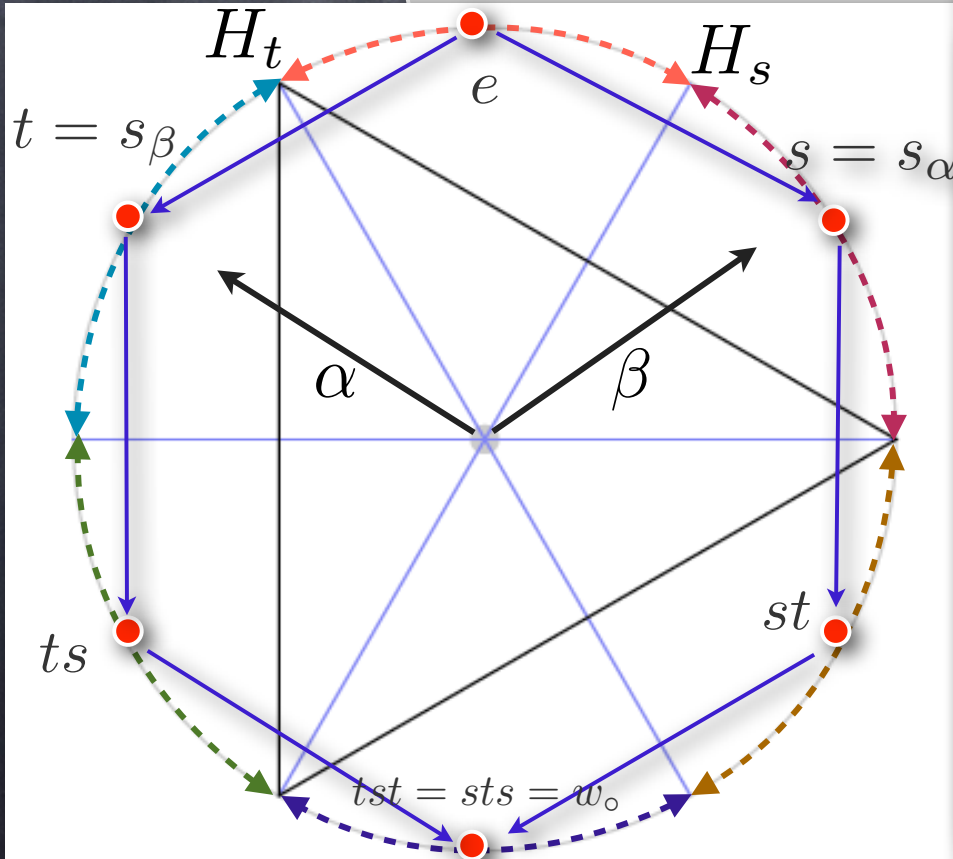
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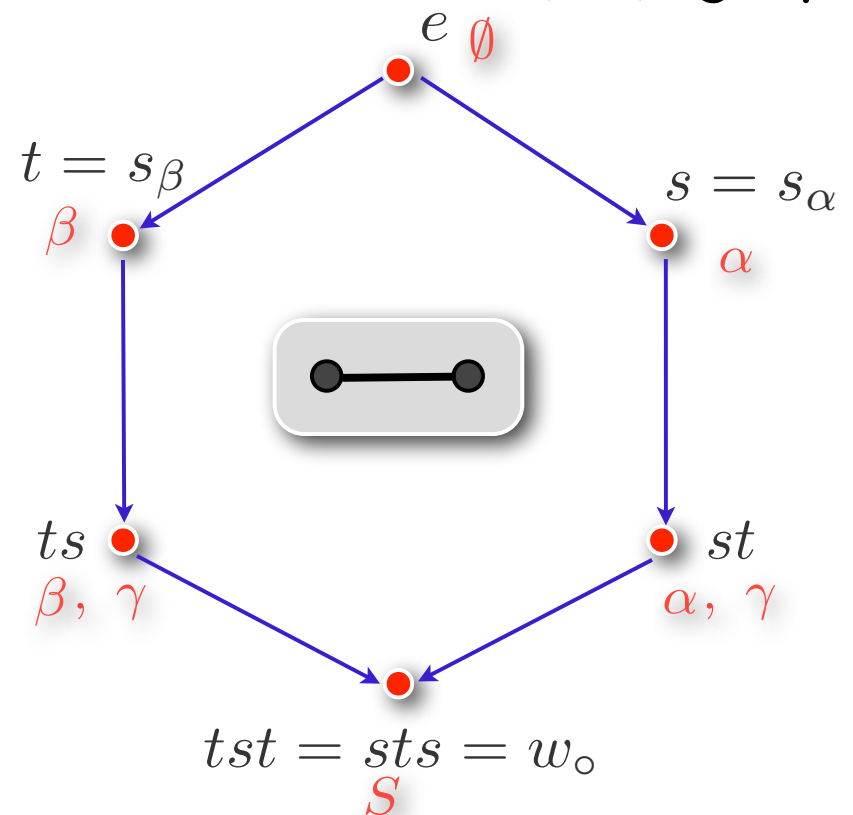
# Coxeter groups and Reflection groups

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$S_n$  is




Weak order & Cayley graph!



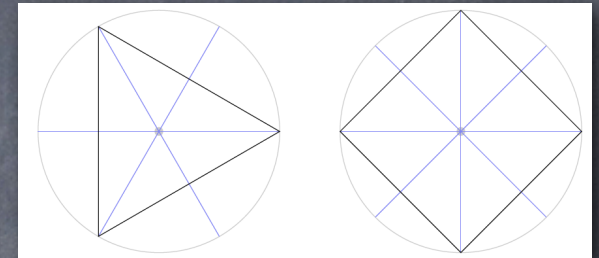


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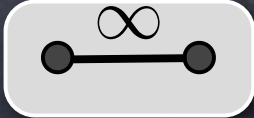

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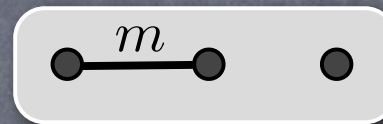
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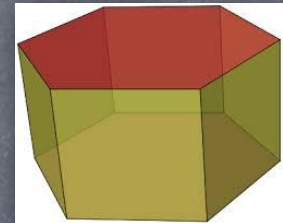
# Coxeter groups and Reflection groups

Rank 3 finite reflection groups: isometry groups of

□ m-gonal regular prisms



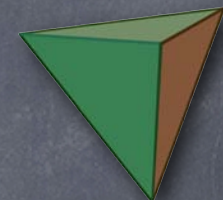
$$\mathcal{D}_m \times S_2$$



□ Regular polyhedra

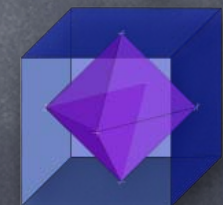
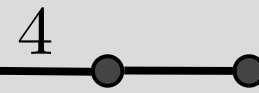
✓ tetrahedron

$$S_4(A_3)$$



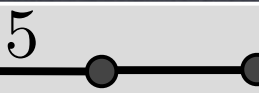
✓ cube/  
octahedron

$$B_3$$



✓ dodecahedron/  
icosahedron

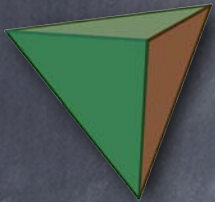
$$H_3$$



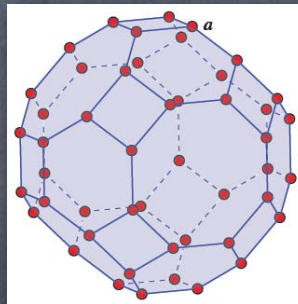
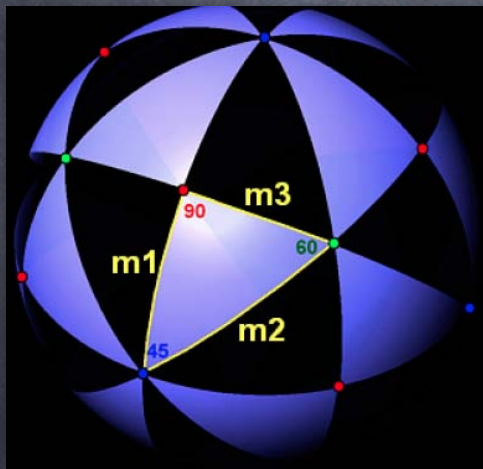
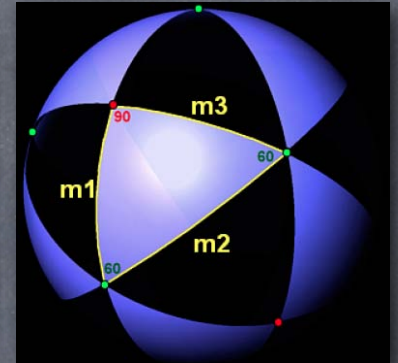
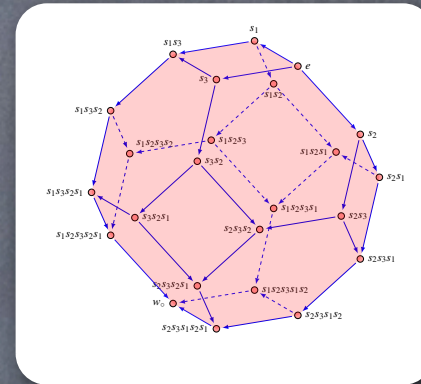


# Coxeter groups and Reflection groups

- Regular polyhedra, Permutohedra and tilings in rk 3

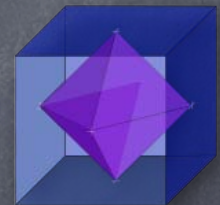


$$S_4(A_3)$$



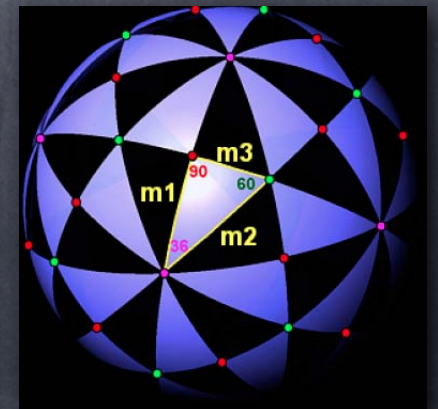
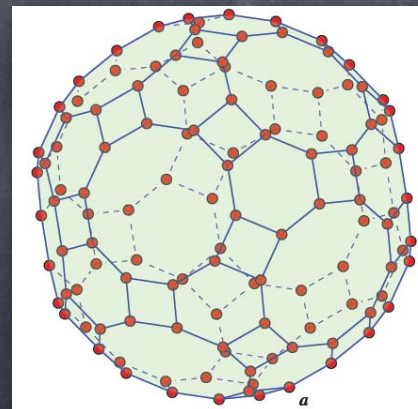
$$B_3$$

4



$$H_3$$

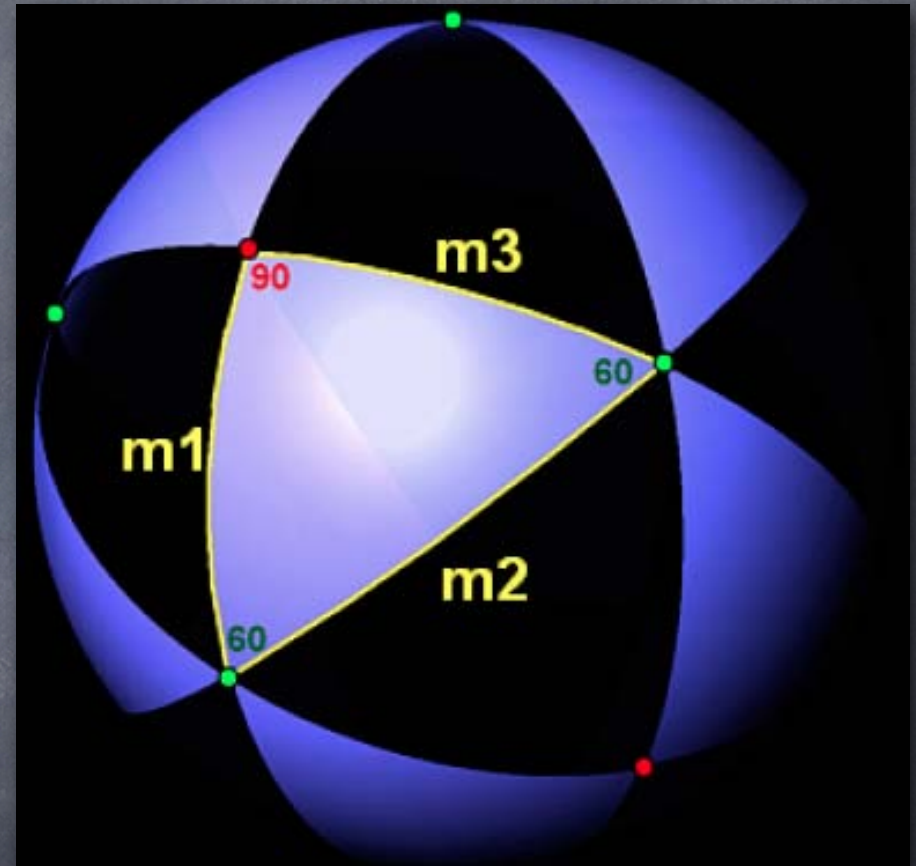
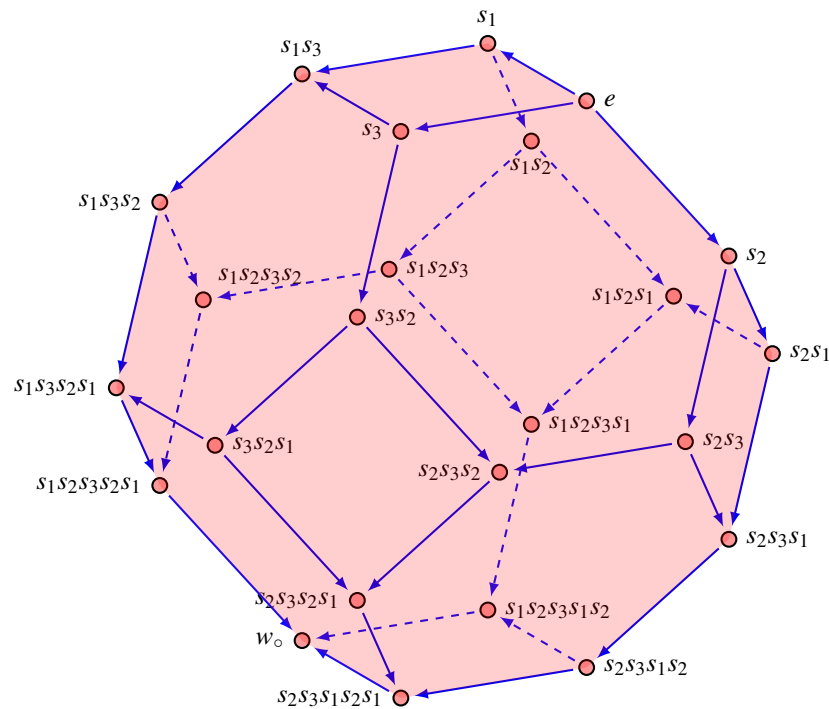
5



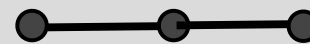


# Coxeter groups and Reflection groups

- Regular polyhedra, Permutohedra and tilings in rk 3



$$S_4(A_3)$$





# Coxeter groups and Reflection groups

How to find all Coxeter graphs that correspond to Finite Reflection groups (FRG)? to Finite Coxeter groups?

world of roots



world of words

Coxeter groups/graphs

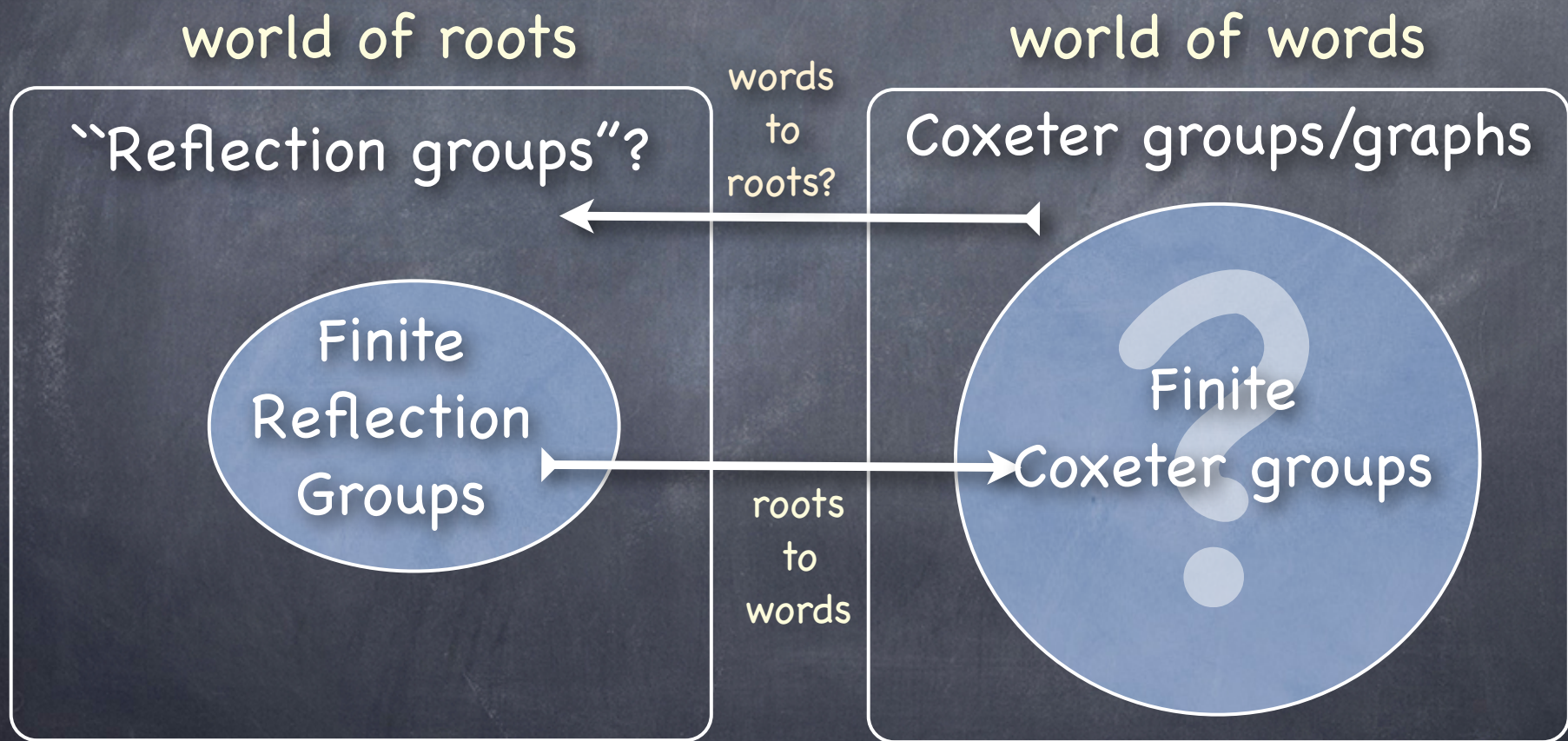


roots  
to  
words



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How to find all Coxeter graphs that correspond to Finite Reflection groups (FRG)? to Finite Coxeter groups?





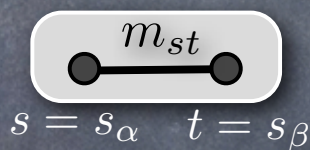
# Root systems for Coxeter groups ?

An observation

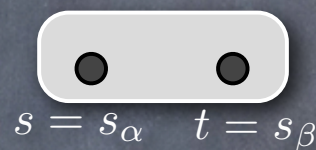
If  $(W, S)$  is a **Finite Reflection Group** with  $\Delta \subseteq \Phi^+ \subseteq \Phi$ .

• **Dihedral (standard) parabolic subgroups:**  $I = \{s, t\} \subseteq S$

□  $W_I = \langle I \rangle \leq W$  corresponds to the subgraphs:



or



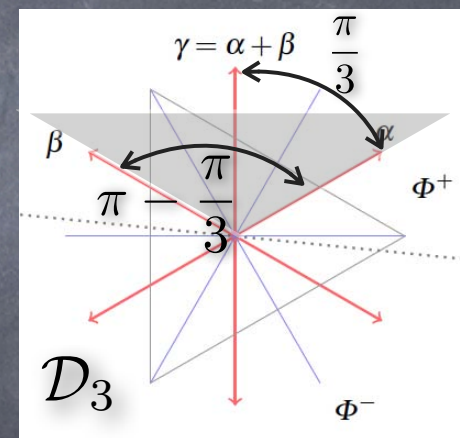
□  $W_I = \mathcal{D}_{m_{st}}$  acts on  $V_I = \text{span}(\alpha, \beta)$ :

$$s_\alpha(\beta) = \beta - 2\langle \alpha, \beta \rangle \alpha$$

□ We have:  $\langle \alpha, \beta \rangle = -\cos\left(\frac{\pi}{m_{st}}\right)$

• the scalar product is given on the basis  $\Delta$  by

$$(\langle \alpha, \beta \rangle)_{\alpha, \beta \in \Delta} = \left( -\cos\left(\frac{\pi}{m_{st}}\right) \right)_{s, t \in S}$$





# Root systems for Coxeter groups !

## • Geometric representations of a Coxeter system $(W, S)$

- $V$  real vector space with basis  $\Delta = \{\alpha_s \mid s \in S\}$
- $B$  symmetric bilinear form defined by:

$$B(\alpha_s, \alpha_t) = \begin{cases} -\cos\left(\frac{\pi}{m_{st}}\right) & \text{if } m_{st} < \infty \\ a \leq -1 & \text{if } m_{st} = \infty \end{cases}$$

- $W$  acts on  $V$ :  $s(v) = v - 2B(v, \alpha)\alpha$ ,  $s \in S$
- **Root system**:  $\Phi = W(\Delta)$ ,  $\Phi^+ = \text{cone}(\Delta) \cap \Phi$

**Theorem.** (i)  $W \leq O_B(V)$  “ $B$ -isometries”

(ii) All the properties words/roots for FRG hold in this case (inversion sets/length, weak order etc).



# Root systems for Coxeter groups !

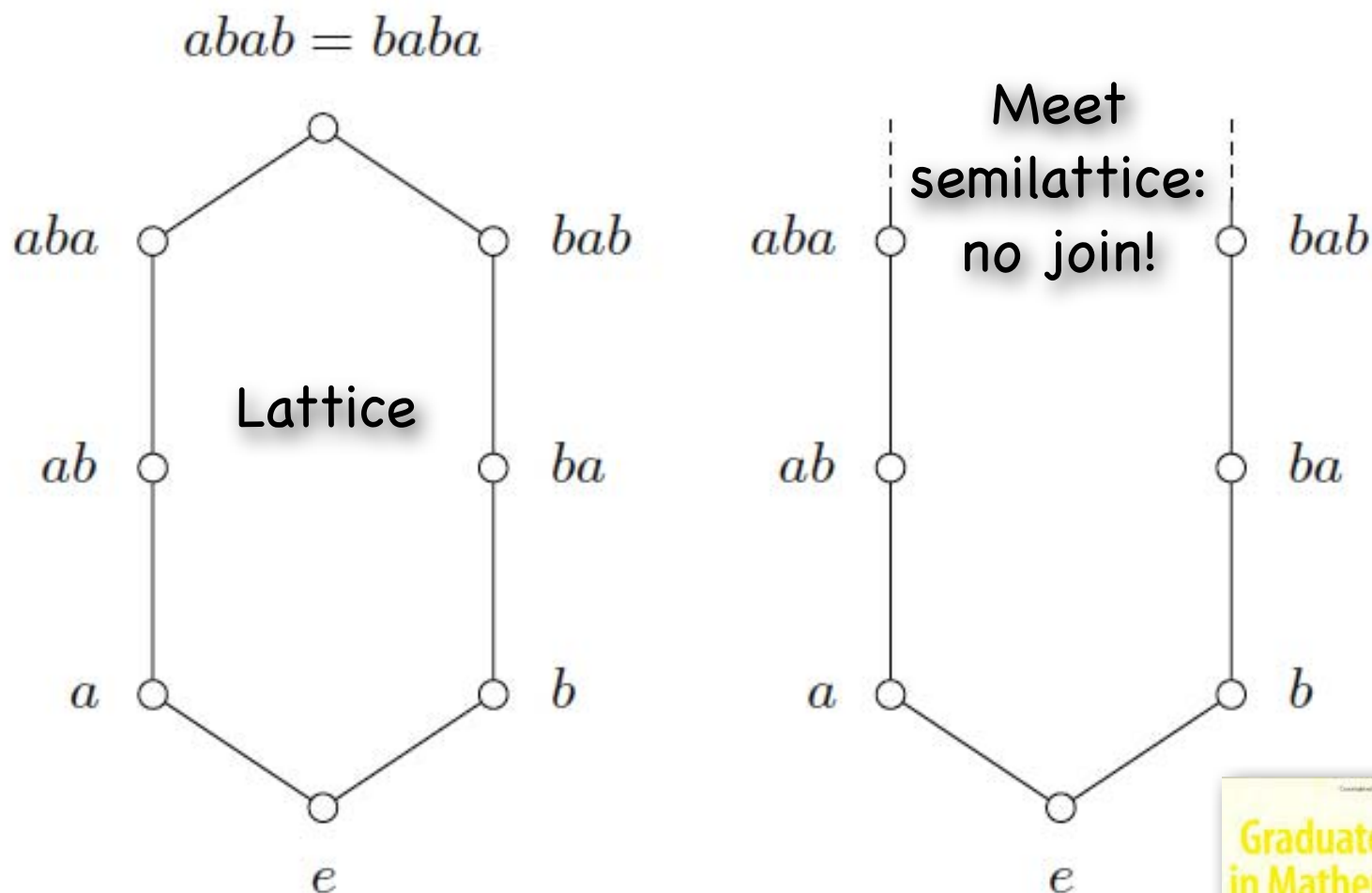
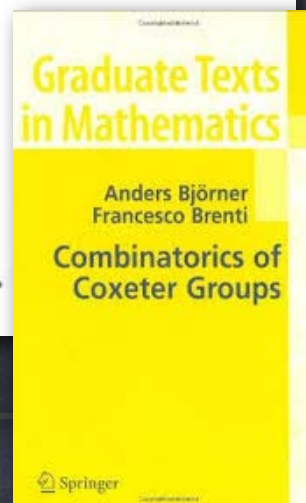


Figure 3.1. Weak order of dihedral groups.

this case (inversion sets/length, weak



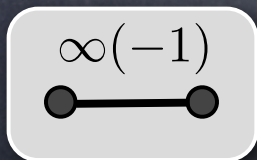
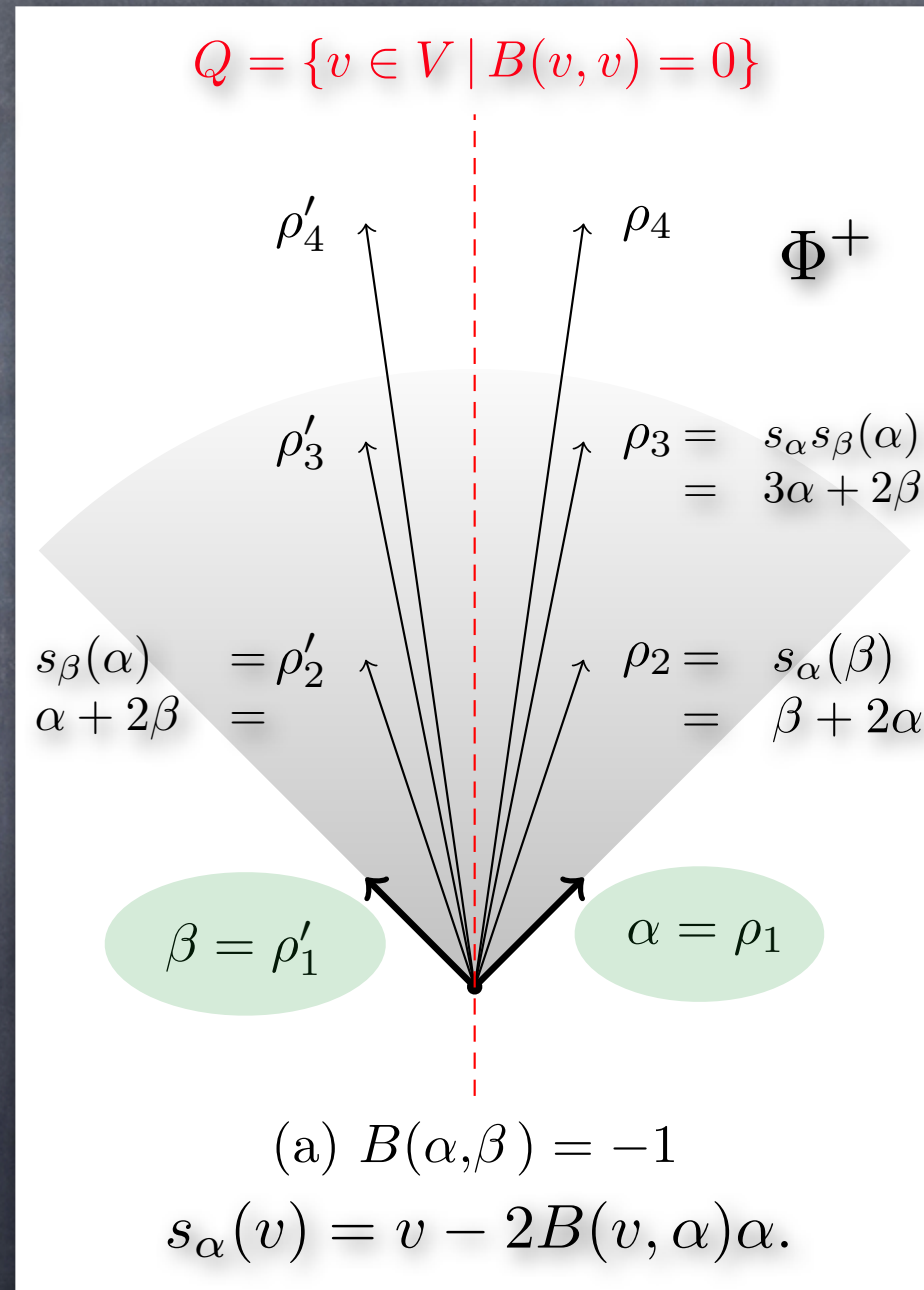


# Root systems for Coxeter groups

$$\rho'_n = n\alpha + (n+1)\beta$$

Infinite  
dihedral  
group I

$$\rho_n = (n+1)\alpha + n\beta$$





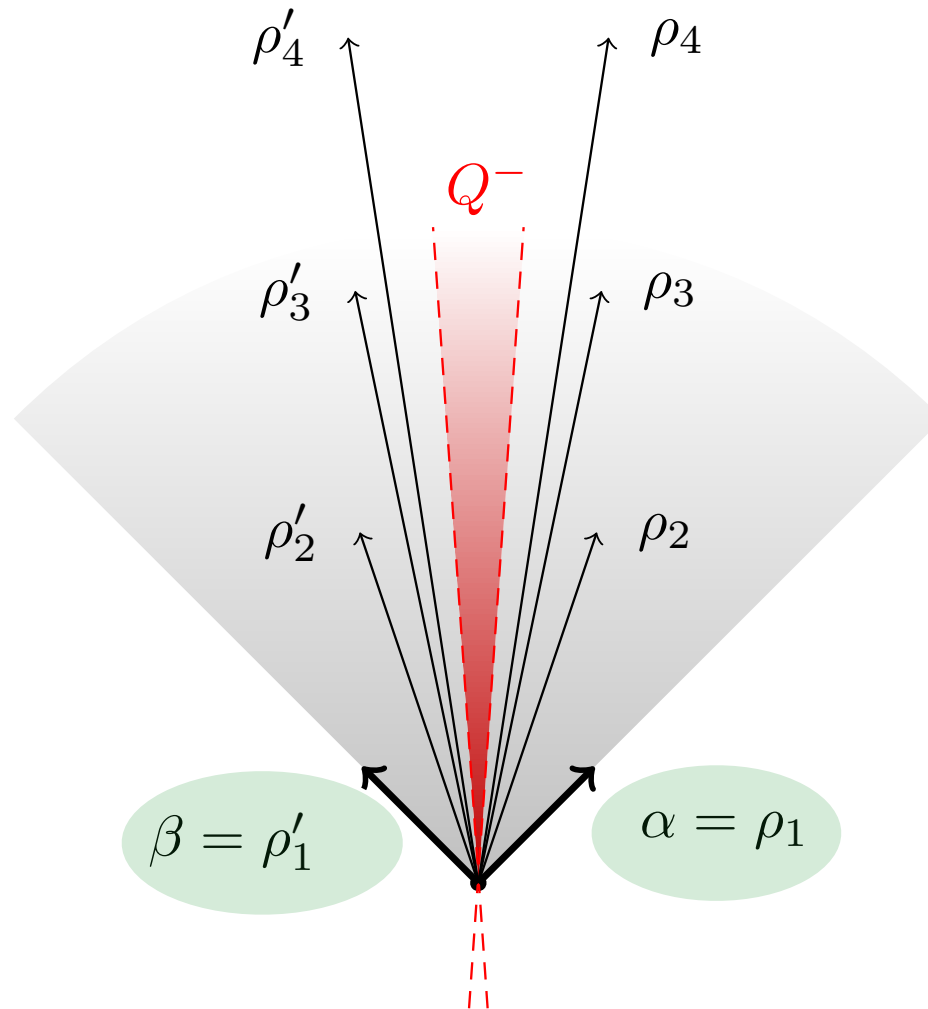
# Root systems for Coxeter groups

Infinite  
dihedral  
group II

$\infty(-1, 01)$



$$Q^- = \{v \in V \mid B(v, v) \leq 0\}$$

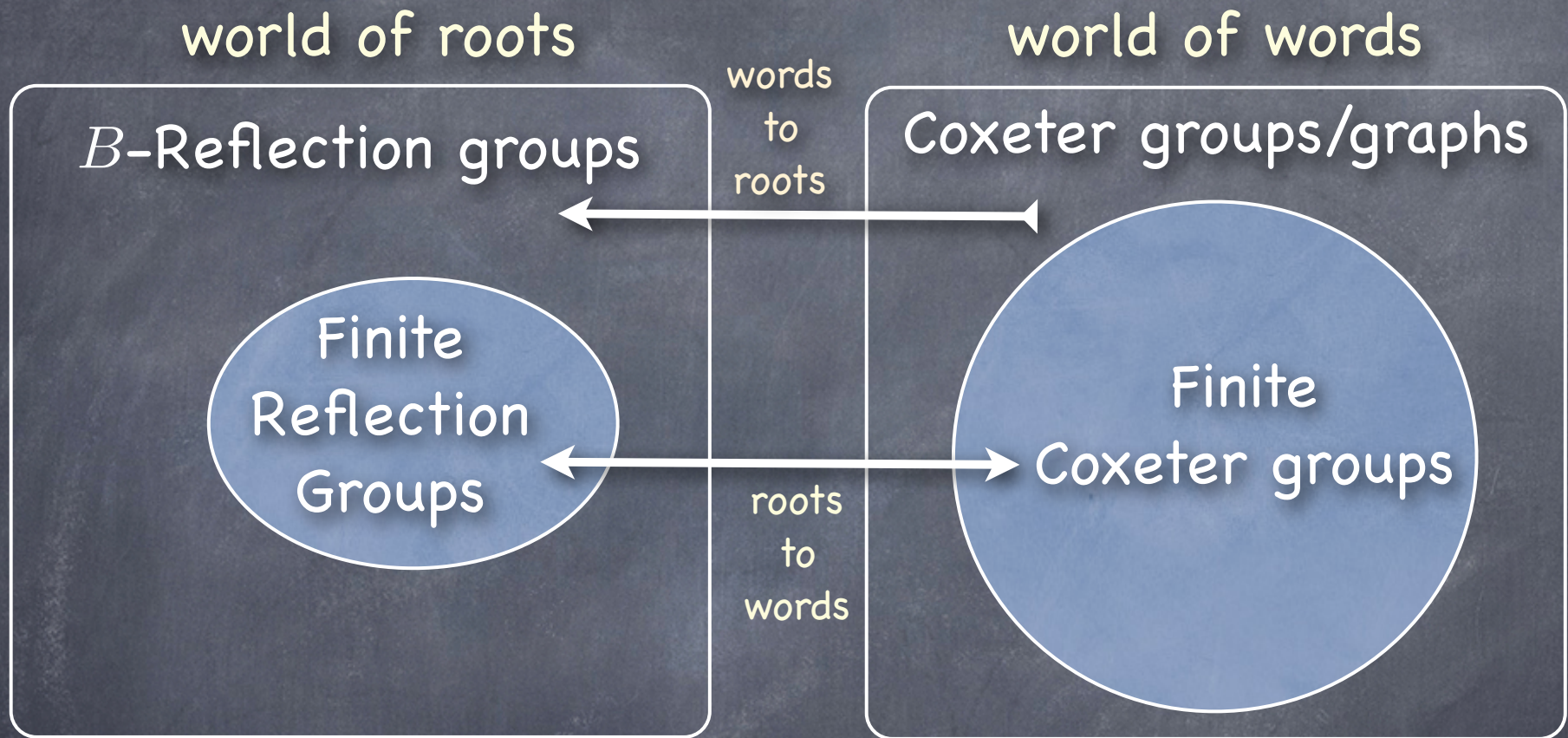


$$(b) \quad B(\alpha, \beta) = -1.01 < -1$$

$$s_\alpha(v) = v - 2B(v, \alpha)\alpha.$$



# Classification of Finite Reflection Groups



**Theorem.** The following assertions are equivalent:

- (i)  $(W, S)$  is a finite Coxeter system;
- (ii)  $B$  is a scalar product and  $W \leq O_B(V)$ ;
- (iii)  $W$  is a finite reflection group.

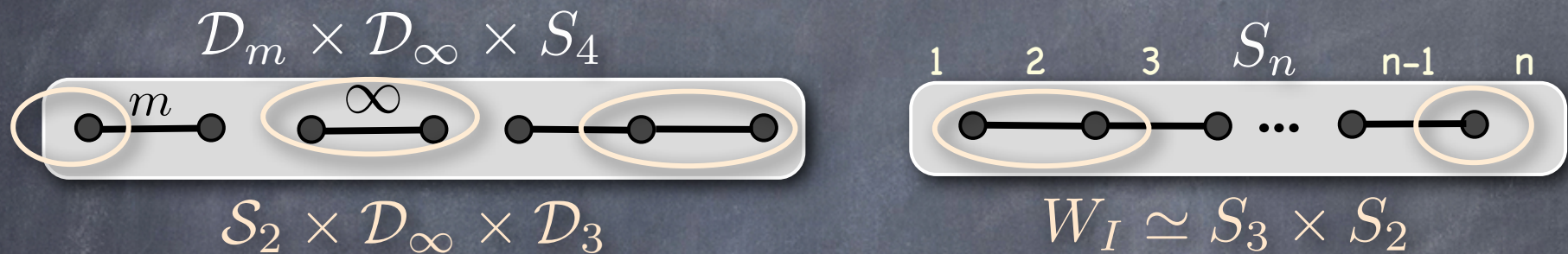


# Classification of Finite Reflection Groups

- Subgraphs and standard parabolic subgroups

$$I \subseteq S \iff \Gamma|_I ; \quad (W_I, I) \text{ is a Coxeter system}$$

- $W$  is irreducible iff  $\Gamma_S$  is connected



**Proposition.** If  $I_1, \dots, I_k$  corresponds to the connected components of  $\Gamma_I$  ( $I$  may be  $S$ ), then

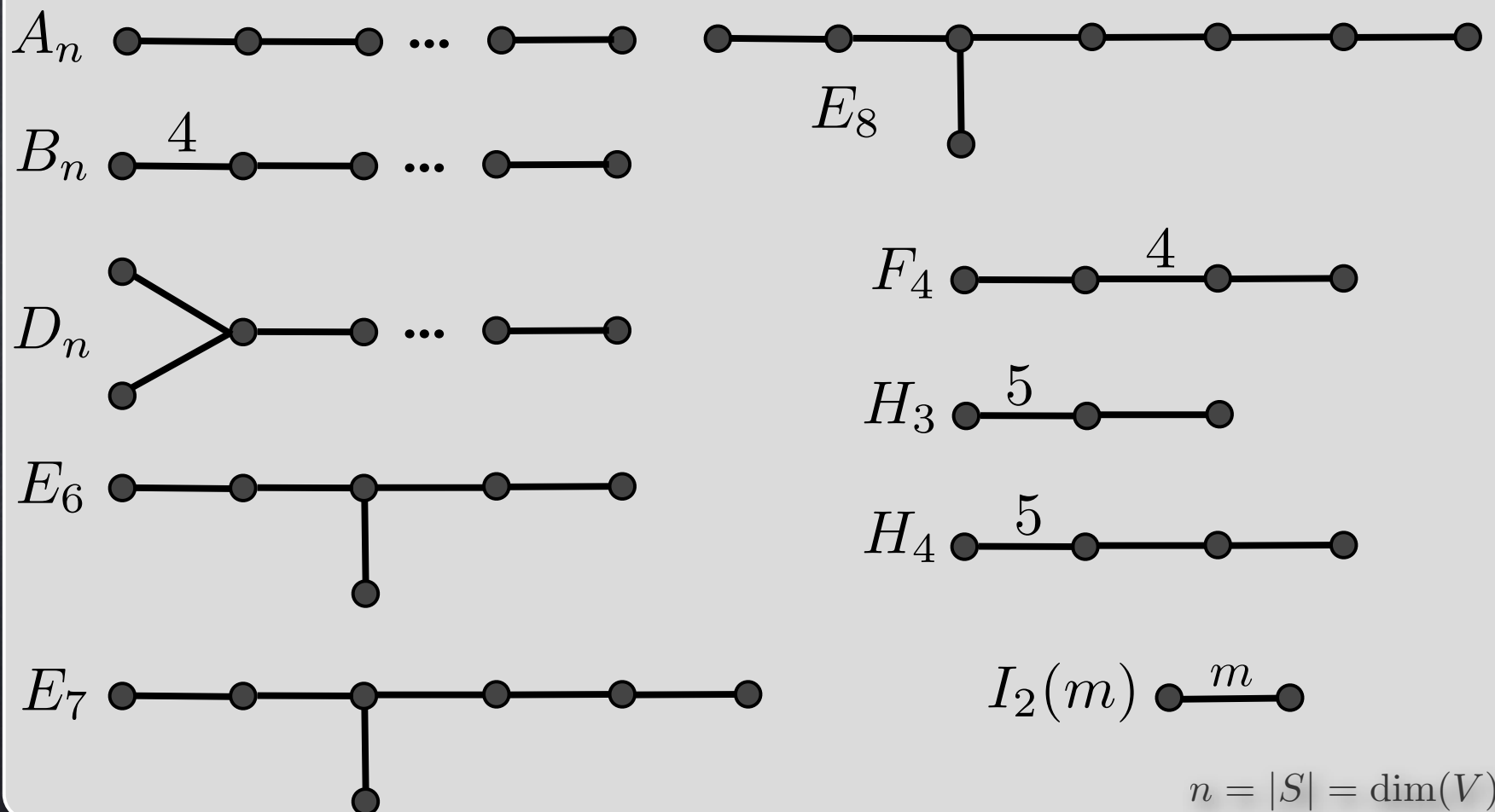
$$W_I \simeq W_{I_1} \times \dots \times W_{I_k}$$

To classify finite reflection groups, i.e., finite Coxeter groups, we just have to find all connected Coxeter graphs that correspond to scalar product!



# Classification of Finite Reflection Groups

**Theorem.** The irreducible FRG are precisely the finite irreducible Coxeter groups. Their graphs are:





# Conclusion

world of roots

world of words

$B$ -Reflection groups

signature  $(p, q, r)$  of  $B$

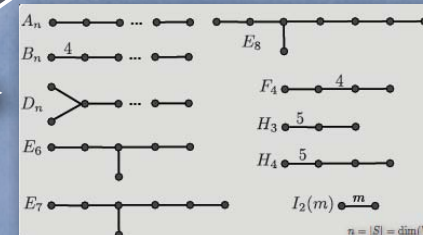
words  
to  
roots

Coxeter groups/graphs

$\Gamma_W$  allowing  $\infty(a \leq -1)$

Finite  
Reflection  
Groups

roots  
to  
words



**Problem:** Let  $p, q, r \in \mathbb{N}$ , classify all the Coxeter graphs with signature  $(p, q, r)$ . Count them?

N.B.: Known for  $(n, 0, 0)$  - FRG -;  $(n - 1, 0, 1)$  - affine type - and partially for  $(n - 1, 1, 0)$  - "weakly hyperbolic" type

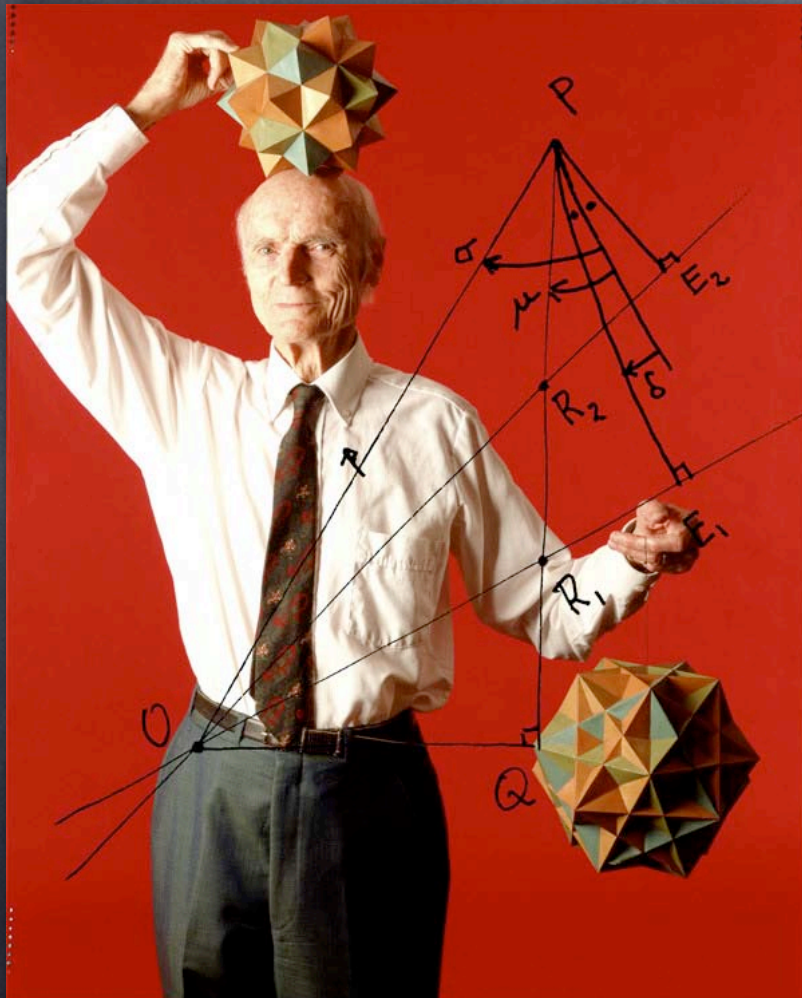


# Donald Coxeter

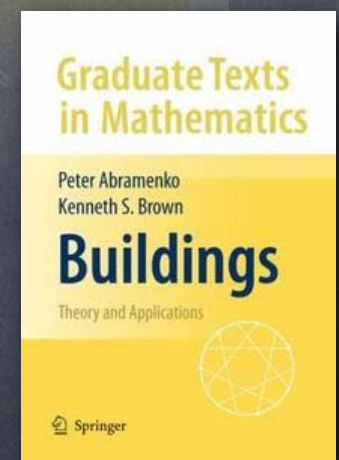
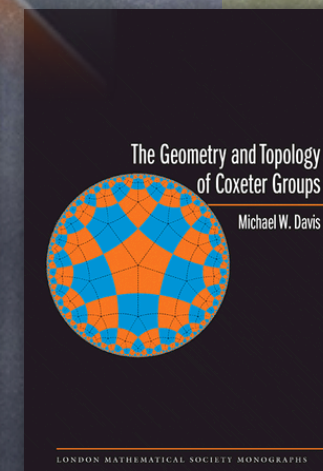
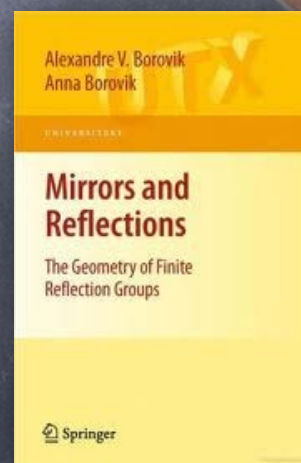
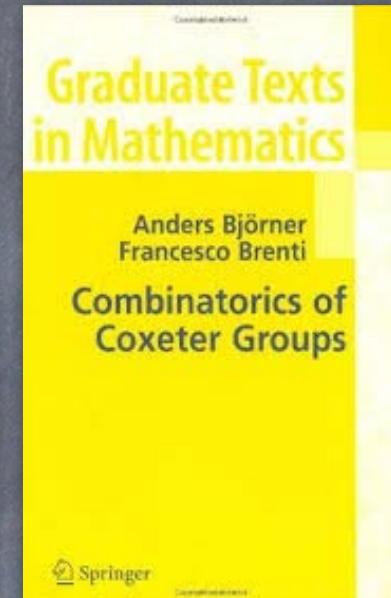
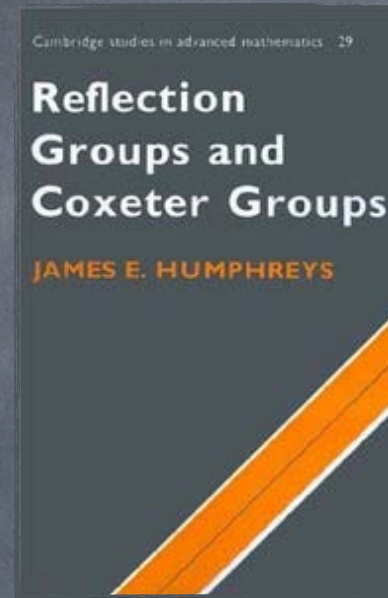
(London 1907, Toronto 2003)

Professor at University of Toronto  
(1936–2003)

## Selected biblio of Part 1 ...

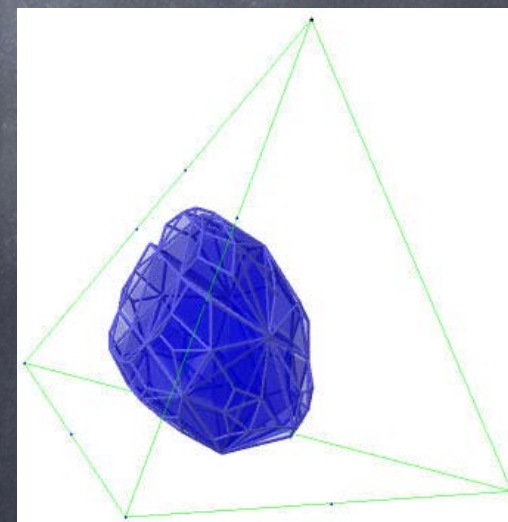
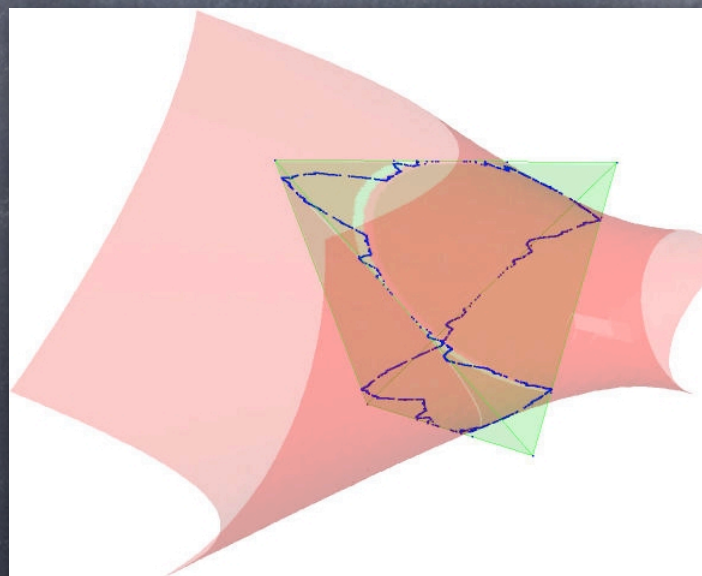
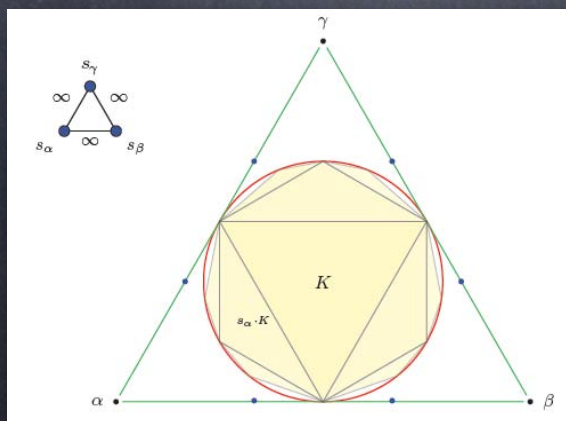
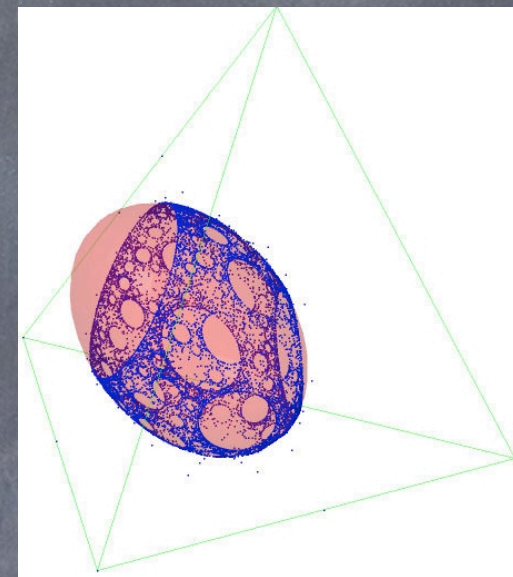
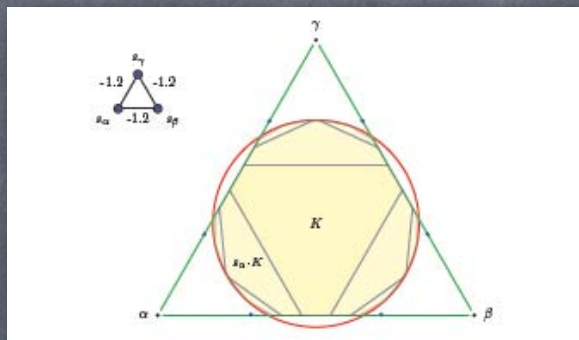
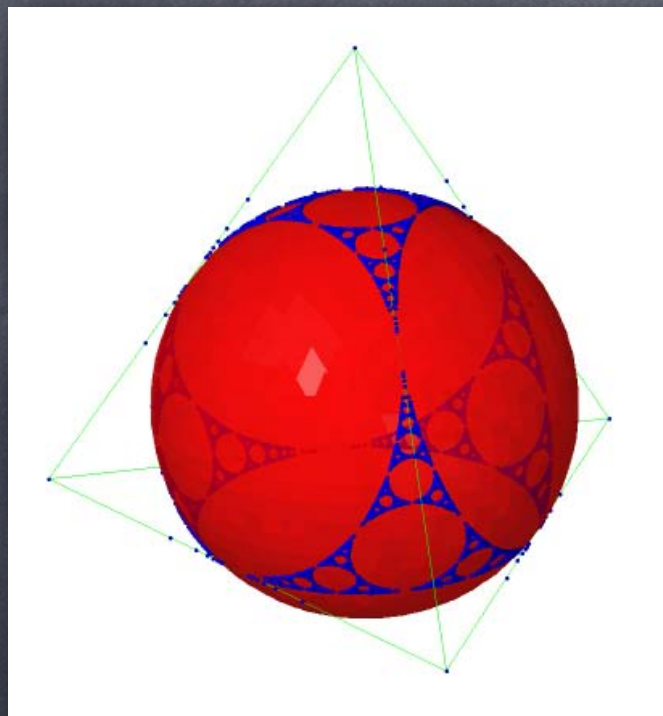


Mandatory photo credit. Mathematics genius Donald Coxeter is the subject of a public talk by journalist Siobhan Roberts, *The Man Who Saved Geometry*, on Sunday, July 31. Photo courtesy of The Banff Centre.





# Part 3 – Roots and Words in infinite Coxeter groups





# In the last Episode

world of roots

world of words

$B$ -Reflection groups

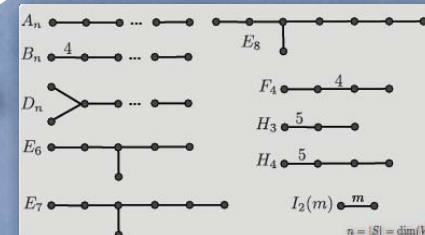
words  
to  
roots

Coxeter groups/graphs

signature  $(p, q, r)$  of  $B$   $\longleftrightarrow$   $\Gamma_W$  allowing  $\infty (a \leq -1)$

Finite  
Reflection  
Groups

roots  
to  
words



The Cayley graph of  $(W, S)$  is naturally oriented by the (right) weak order:  $w < ws$  if  $\ell(w) < \ell(ws)$ .

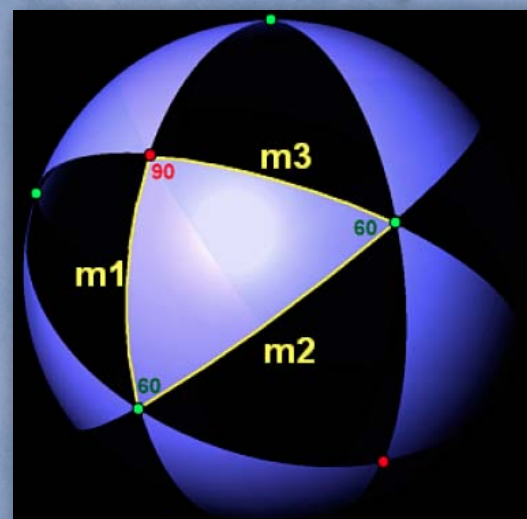
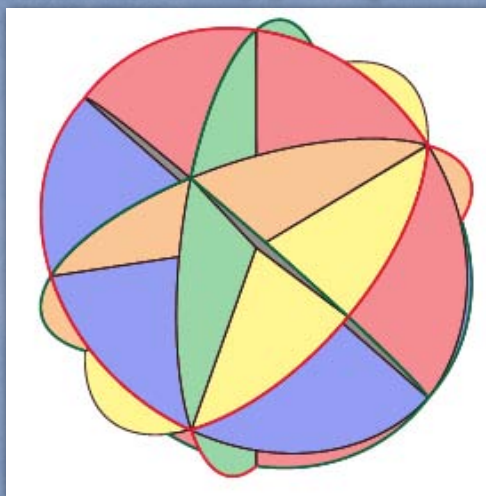
The weak order is a meet-semilattice and

$$u \leq w \iff \text{inv}(u^{-1}) \subseteq \text{inv}(w^{-1}) \quad (\text{inv}(w^{-1}) = \Phi^+ \cap w(\Phi^-))$$



# In the last Episode

In the spherical, euclidean and hyperbolic case, all finitely generated discrete B-reflection groups are Coxeter groups (models for these geometry exist in  $V$  or its dual; 'cut' these models by the hyperplanes of reflections)



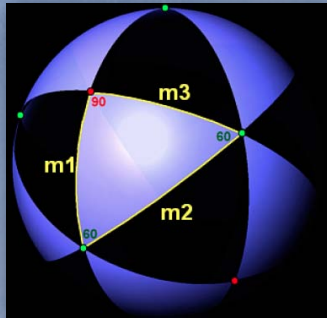
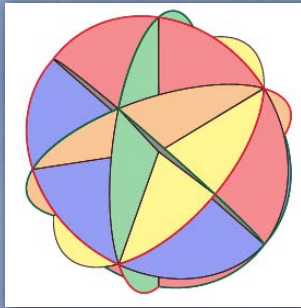
Finite case i.e.  $B$  is a scalar product ( $V = V^*$ ):  
the model is the unit sphere

$$||v||^2 = B(v, v) = 1$$

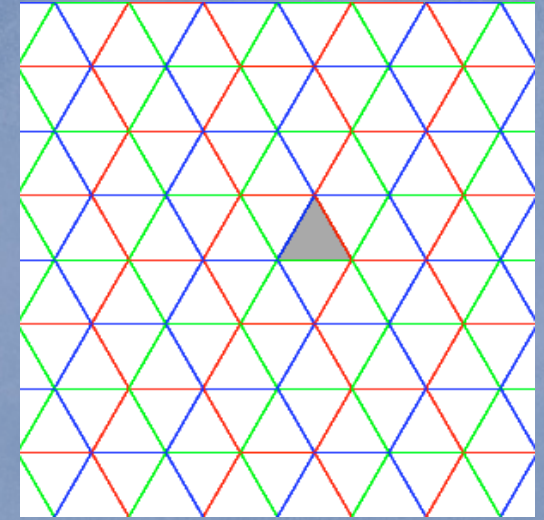
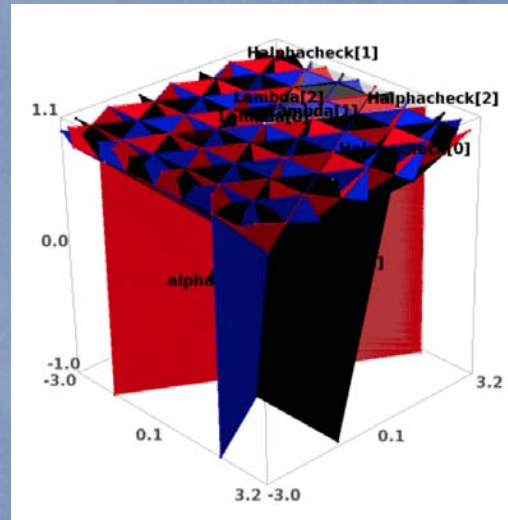


# In the last Episode

In the spherical, euclidean and hyperbolic case, they are all Coxeter groups (models for these geometry exist in  $V$  or its dual; 'cut' these models by the hyperplanes of reflections)



Finite case i.e.  $B$  is a scalar product  
 $\text{sgn}(B) = (n, 0, 0)$



**Affine case** i.e.  $B$  is positive degenerate. Its radical is a line:  $\text{Rad}(B) = \{v \in V \mid B(v, \alpha) = 0, \forall \alpha \in \Delta\} = \mathbb{R}x$   
 The model is an affine hyperplane in the dual  $V^*$ :

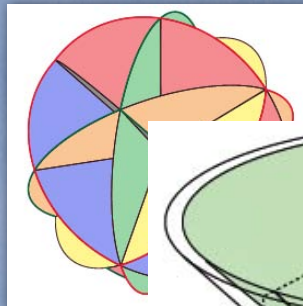
$$H = \{\varphi \in V^* \mid \varphi(x) = 1\}$$

N.B: reflection hyperplanes leave in the dual here.

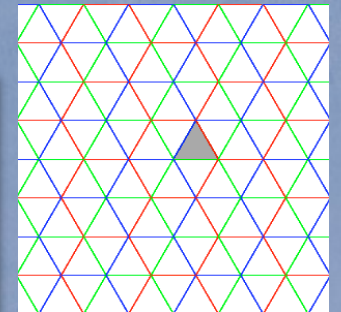
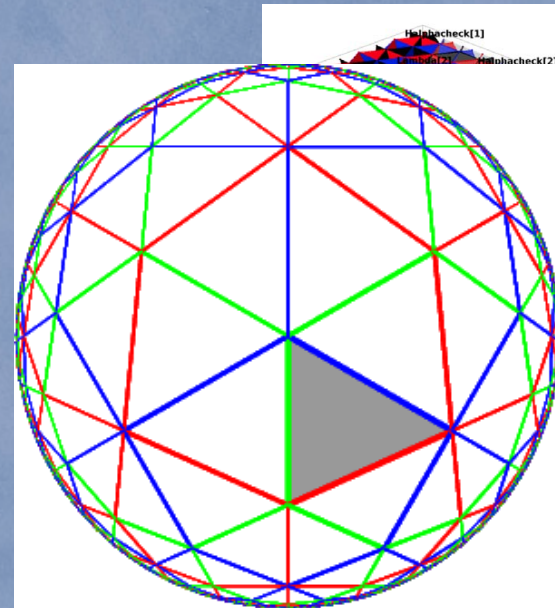
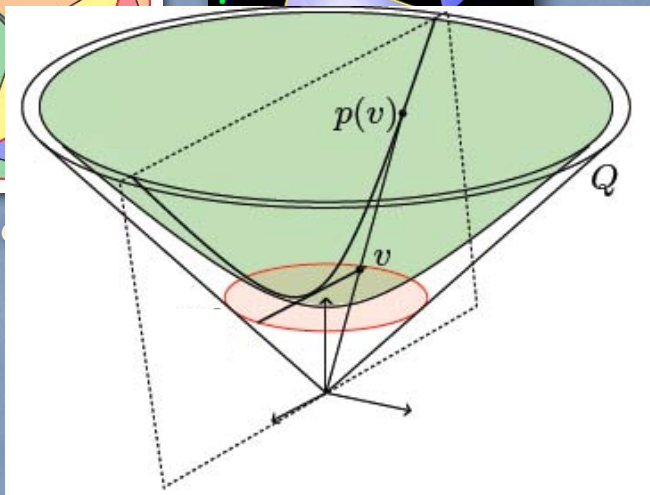
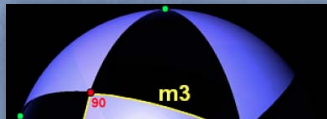


# In the last Episode

In the spherical, euclidean and hyperbolic case, they are all Coxeter groups (models for these geometry exist in  $V$  or its dual; 'cut' these models by the hyperplanes of reflections)



Finite



positive degenerate.  
( $n - 1, 0, 1$ )

**Hyperbolic case** i.e.  $\text{sgn}(B) = (n - 1, 1, 0)$  ( $V = V^*$ ). Many models exists: projective (non conformal), hyperboloid or the ball model

$$H^{n-1} = \{x \in V \mid B(x, x) = -1\}$$



# In the last Episode

world of roots

world of words

$B$ -Reflection groups

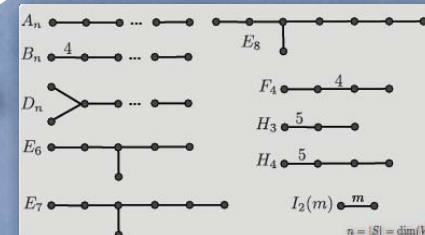
words  
to  
roots

Coxeter groups/graphs

signature  $(p, q, r)$  of  $B$   $\longleftrightarrow$   $\Gamma_W$  allowing  $\infty (a \leq -1)$

Finite  
Reflection  
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roots  
to  
words



The Cayley graph of  $(W, S)$  is naturally oriented by the (right) weak order:  $w < ws$  if  $\ell(w) < \ell(ws)$ .

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$$u \leq w \iff \text{inv}(u^{-1}) \subseteq \text{inv}(w^{-1}) \quad (\text{inv}(w^{-1}) = \Phi^+ \cap w(\Phi^-))$$



# An Illustration: Words, Roots and Generalized Associahedra



©someone on the internet

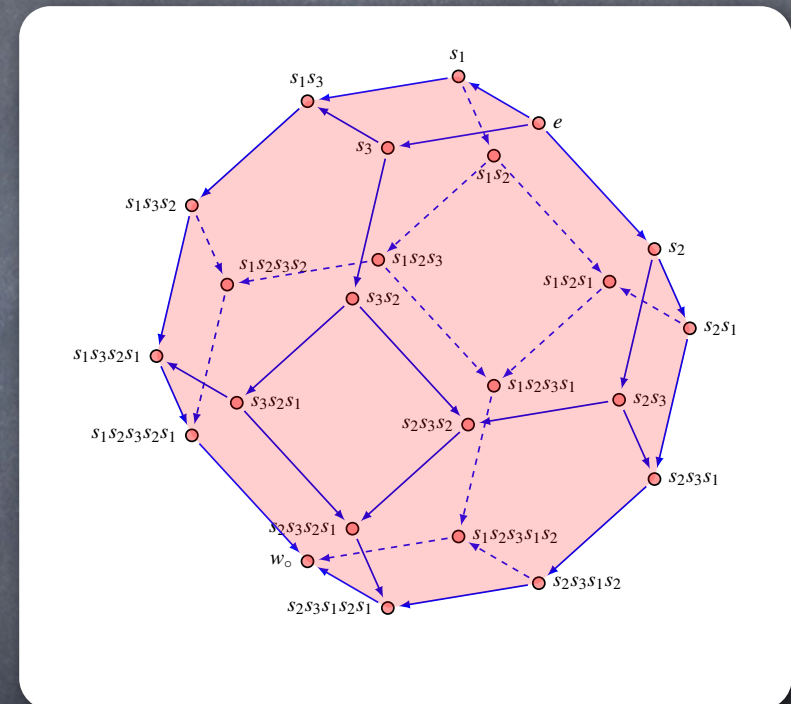
$(W, S)$  finite Coxeter system, so  $W \leq O(V)$

# Permutahedra

- $\Delta$  simple system;
- $S = \{s_\alpha \mid \alpha \in \Delta\}$ ;
- Choose a generic i.e.

$$\langle \mathbf{a}, \alpha \rangle > 0, \quad \forall \alpha \in \Delta$$

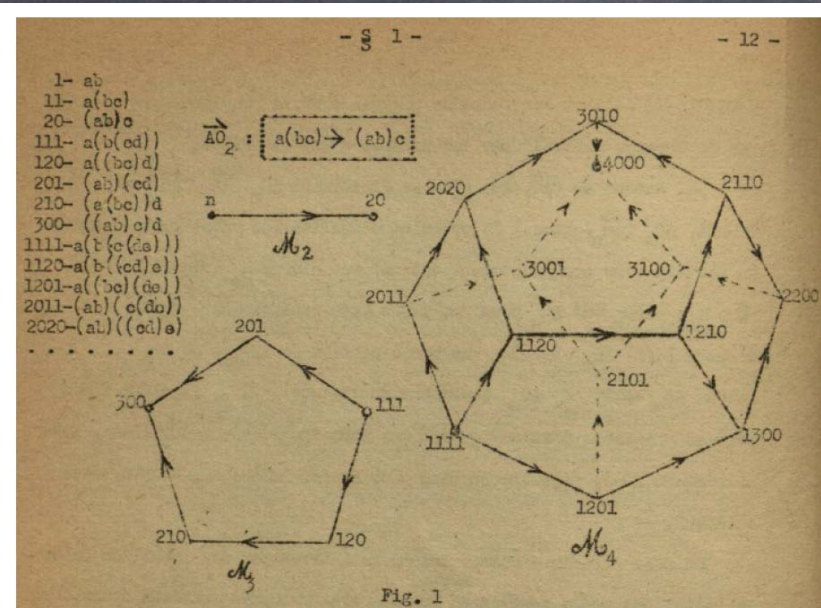
$$\text{Perm}^{\mathbf{a}}(W) = \text{conv} \{w(\mathbf{a}) \mid w \in W\}$$



**Proposition.**  $\text{Perm}^a(W)$  is a simple polytope whose oriented 1-skeleton is the graph of the (right) weak order.



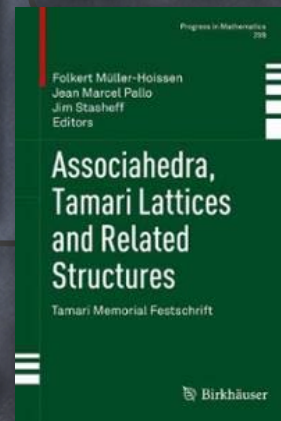
# Building Generalized Associahedra



Tamari's associahedron

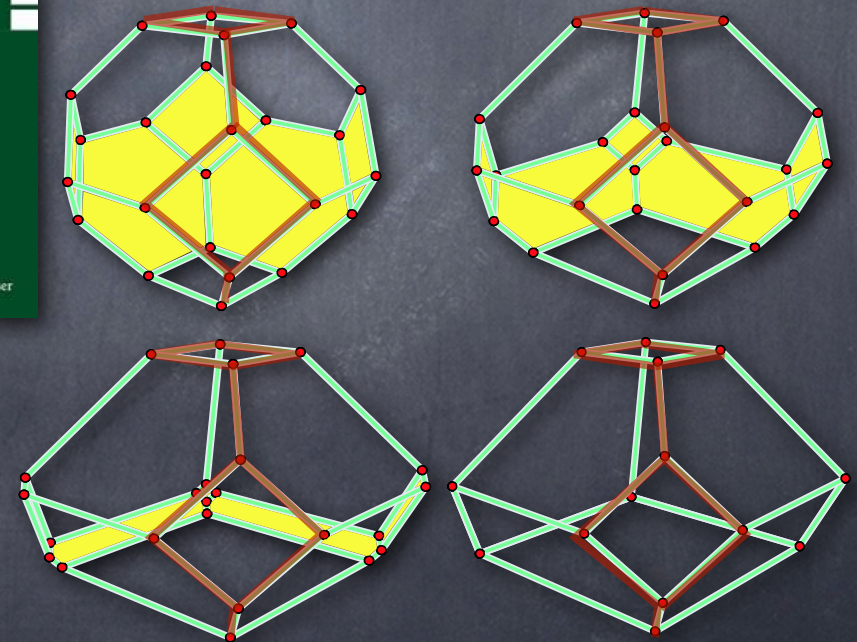
## Associahedra (lattices/complexes):

- Lattice (Tamari, 1951)
- Cell complex (Stasheff, 1963)
- Cluster complex (Fomin-Zelevinsky, 2003)
- Cambrian lattices (Reading 2007, 2007 )and more ...



## Associahedra (Convex polytopes):

- Type A (Haiman 1984, Lee, Loday, ... )
- Type B - cyclohedra (Bott-Taubes 1994, ...)
- Weyl groups (Chapoton-Fomin-Zelevinsky, 2003)
- from permutahedra of finite Coxeter groups (CH-Lange-Thomas 2011, ...)



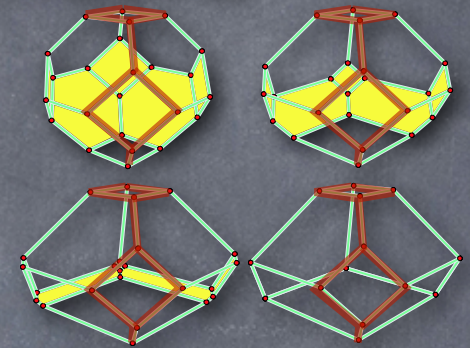
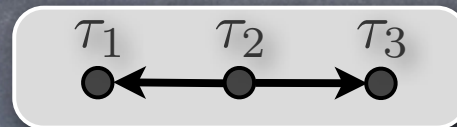


# Building Generalized Associahedra

Hohlweg, C. Lange, H. Thomas (2009)

- **Data:**  $\text{Perm}^a(W)$  and an orientation of  $\Gamma_W$

$$W = S_4$$



- $c$  Coxeter element associated to this orientation i.e product without repetition of all the simple reflections;

$$c = \tau_2 \tau_3 \tau_1$$

- $c_{(I)}$  subword with letters  $I \subseteq S$

$$I = \{\tau_1, \tau_2\} \subseteq S \Rightarrow c_{(I)} = \tau_2 \tau_1$$

- $c$  - word of  $w_o$ :  $w_o(c) = c_{(K_1)} c_{(K_2)} \dots c_{(K_p)}$  reduced expression s.t.  $S \supseteq K_1 \supseteq K_2 \supseteq \dots \supseteq K_p \neq \emptyset$

$$w_o(\tau_1 \tau_2 \tau_3) = \tau_1 \tau_2 \tau_3 \cdot \tau_1 \tau_2 \cdot \tau_1 = c_{(S)} c_{(\{\tau_1, \tau_2\})} c_{(\{\tau_1\})}$$

$$w_o(\tau_2 \tau_3 \tau_1) = \tau_2 \tau_3 \tau_1 \cdot \tau_2 \tau_3 \tau_1 = c_{(S)} c_{(S)}.$$



# Building Generalized Associahedra

Hohlweg, C. Lange, H. Thomas (2009)

□  $C$  - **word** of  $w_o$ :  $w_o(c) = c_{(K_1)} c_{(K_2)} \cdots c_{(K_p)}$  **reduced** expression s.t.  $S \supseteq K_1 \supseteq K_2 \supseteq \cdots \supseteq K_p \neq \emptyset$

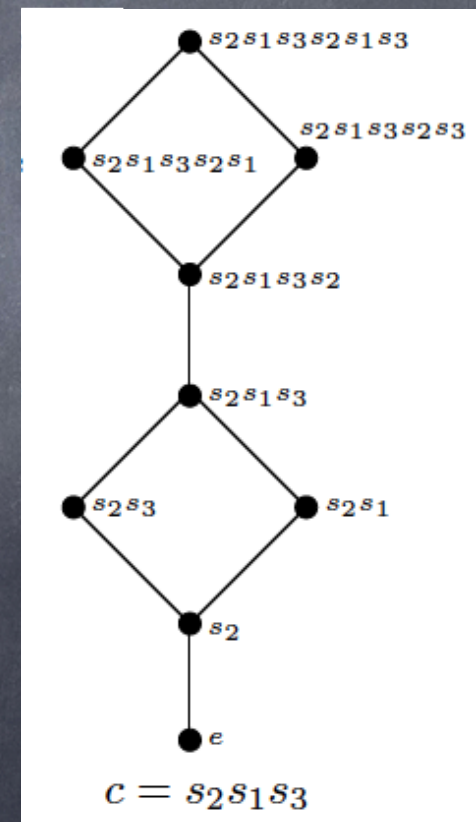
$$w_o(\tau_1 \tau_2 \tau_3) = \tau_1 \tau_2 \tau_3 \cdot \tau_1 \tau_2 \cdot \tau_1 = c_{(S)} c_{(\{\tau_1, \tau_2\})} c_{(\{\tau_1\})}$$

$$w_o(\tau_2 \tau_3 \tau_1) = \tau_2 \tau_3 \tau_1 \cdot \tau_2 \tau_3 \tau_1 = c_{(S)} c_{(S)}.$$

□  $C$  - **singletons** are the prefixes of  $w_o(c)$  up to commutations

$e,$	$\tau_2 \tau_3,$	$\tau_2 \tau_3 \tau_1 \tau_2 \tau_3,$
$\tau_2,$	$\tau_2 \tau_3 \tau_1,$	$\tau_2 \tau_3 \tau_1 \tau_2 \tau_1,$ and
$\tau_2 \tau_1,$	$\tau_2 \tau_3 \tau_1 \tau_2,$	$w_o = \tau_2 \tau_1 \tau_3 \tau_2 \tau_1 \tau_3.$

**Proposition.**  $C$  - singletons form a distributive sublattice of the weak order.





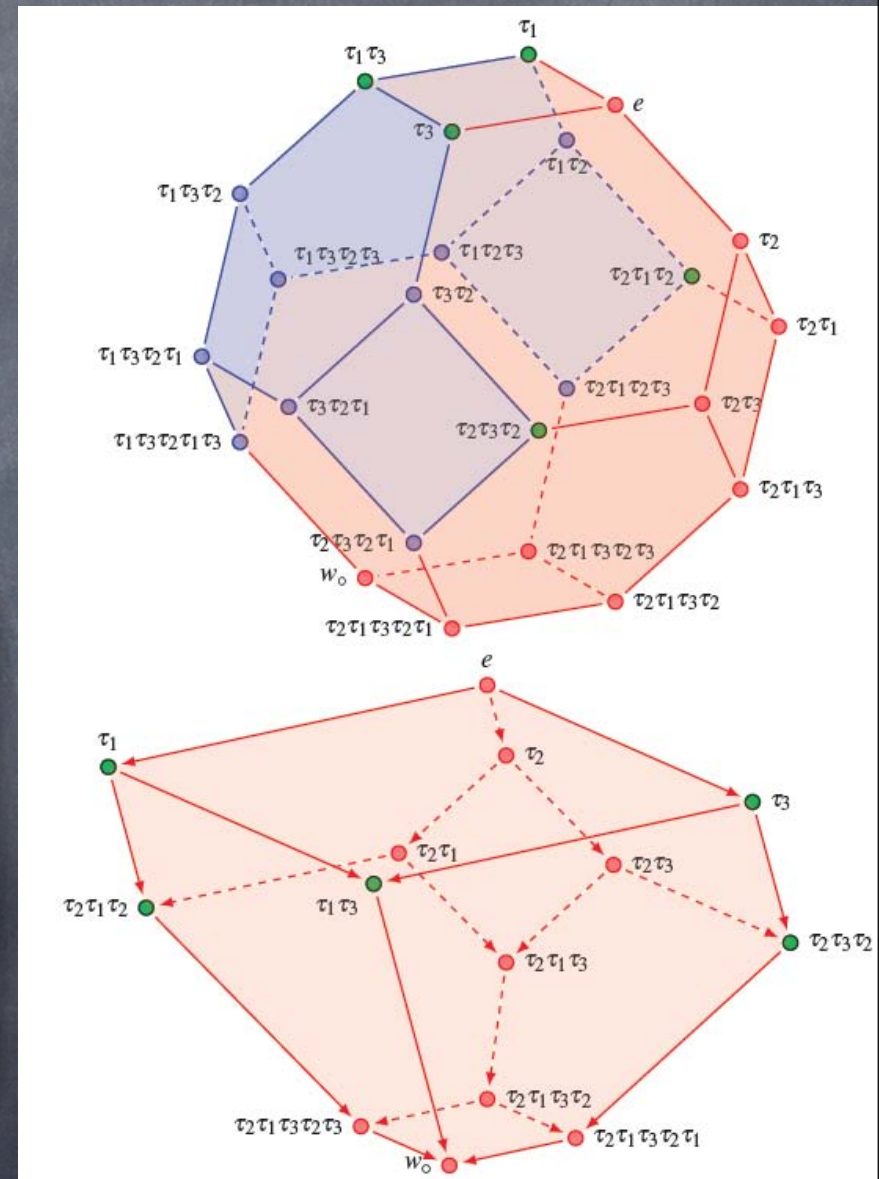
# Building Generalized Associahedra

Hohlweg, C. Lange, H. Thomas (2009)

□  $C$  - generalized associahedron  
is the polytope  $\text{Asso}_C^a(W)$   
obtained from  $\text{Perm}^a(W)$  by  
keeping only the facets  
containing a  $C$  - singleton

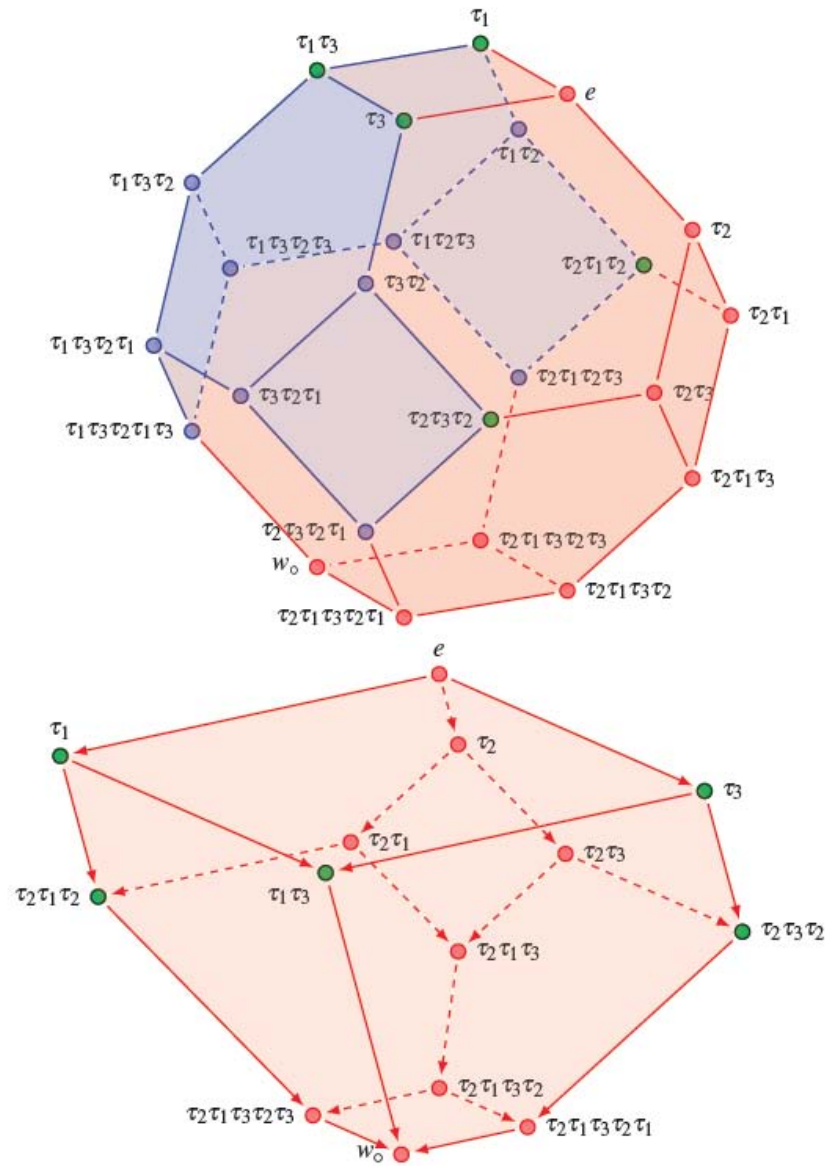
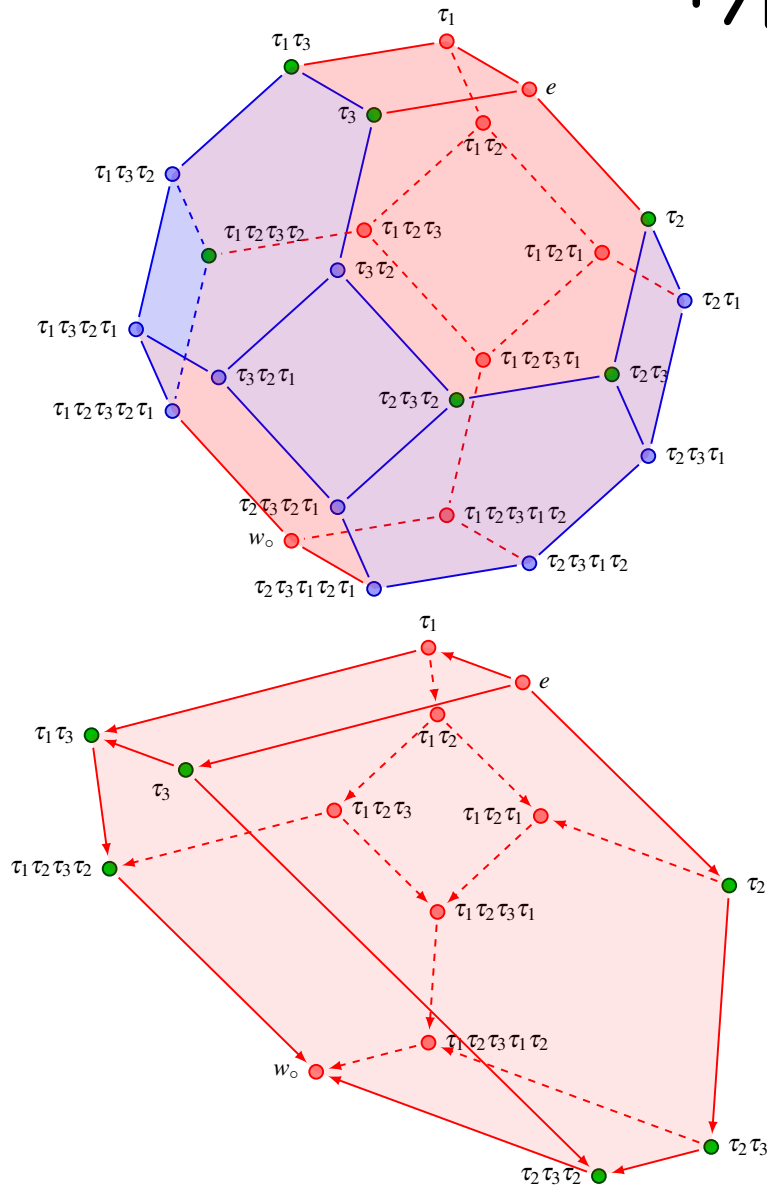
**Theorem.** The 1-skeleton of  
 $\text{Asso}_C^a(W)$

is N. Reading's  $C$ -Cambrian lattice;  
its normal fan is the corresponding  
Cambrian fan studied in detailed by  
N. Reading & D. Speyer.



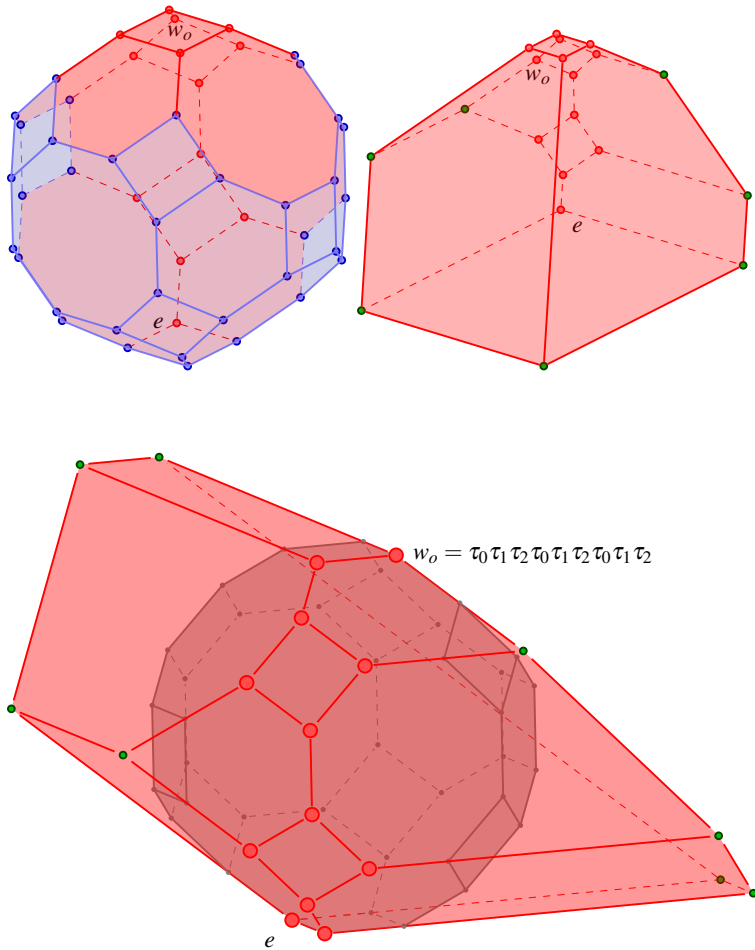


# Type A

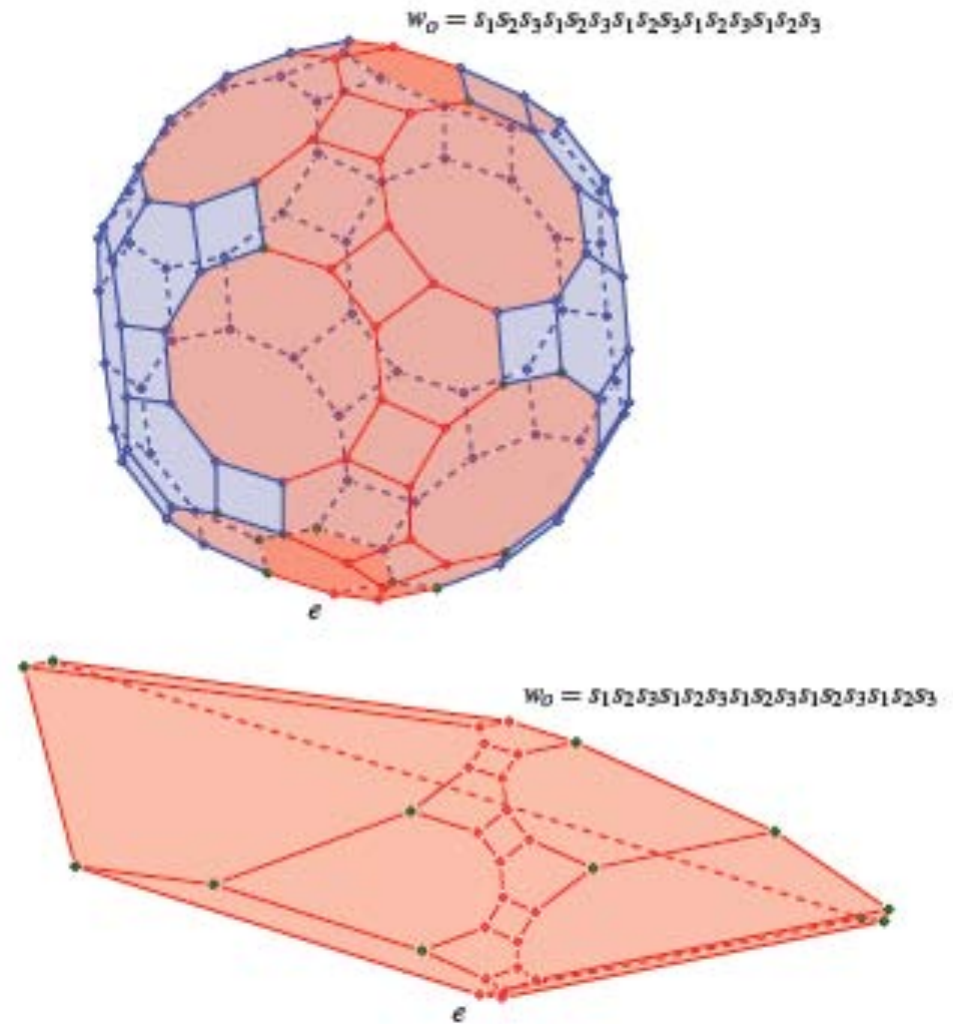




## Type B



## Type H





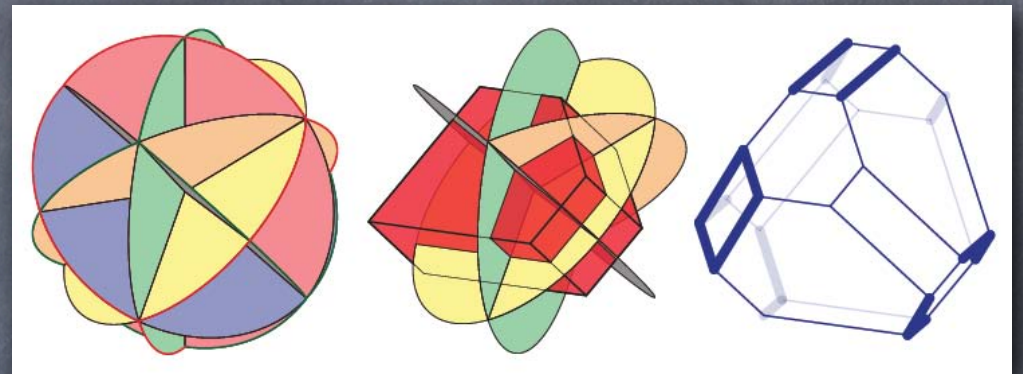
# Selected developments on the subject

- Convex hull of the vertices: brick polytopes. Barycenter identical to the permutahedron:

V. Pilaud and C. Stump:

1. Brick polytopes of spherical subword complexes: A new approach to generalized associahedra (2012)
2. Vertex barycenter of generalized associahedra (2012)

© Pilaud–Stump

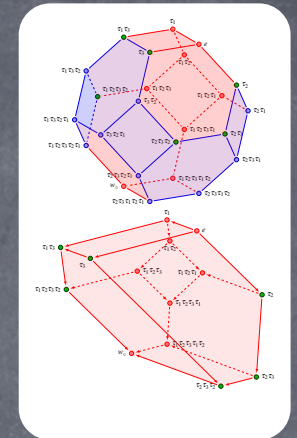


- Classification of isometry classes in term of the lattices of  $C$ —singletons (N. Bergeron, Hohlweg, C. Lange, H. Thomas, 2009)
- Recovering the corresponding cluster algebra:  
S. Stella, Polyhedral models for generalized associahedra via Coxeter elements (2013)

## ASSOCIAHEDRA IN INFINITE CASE ?



# ASSOCIAHEDRA IN INFINITE CASE ?



My original motivation (2010): to generalize this approach in the infinite case ...

- Infinite case: Cambrian meet-semilattices (Sortable Elements in Infinite Coxeter Groups, N. Reading and D. Speyer, 2011) are not big enough ...

**Problem:** is it possible to «enlarge» Coxeter groups to have reasonable candidates with a weak order that is a complete lattice ? An answer may lie on the side of inversion sets!

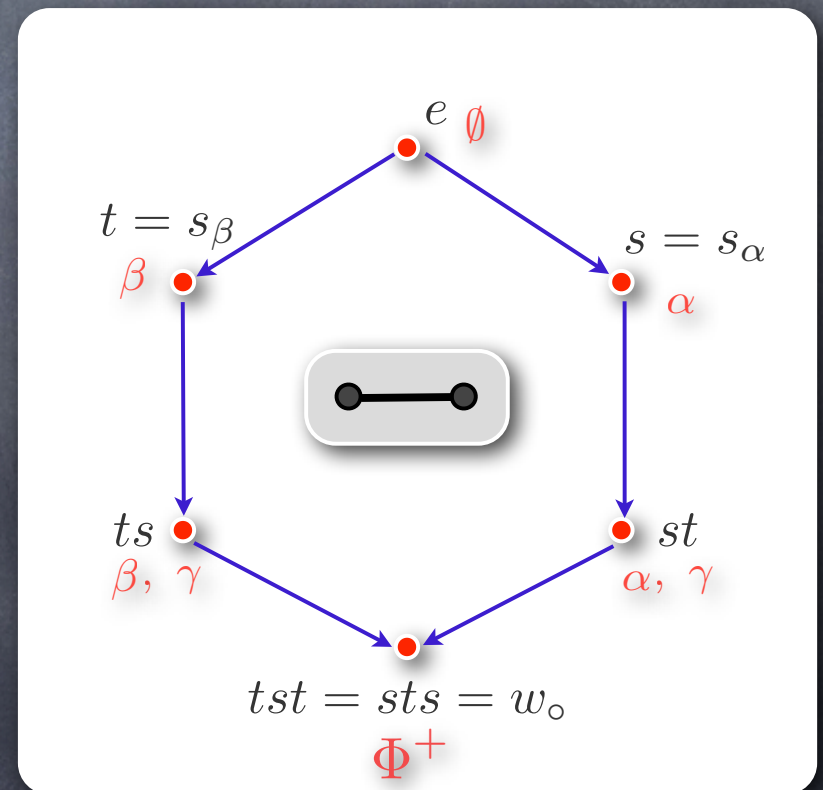
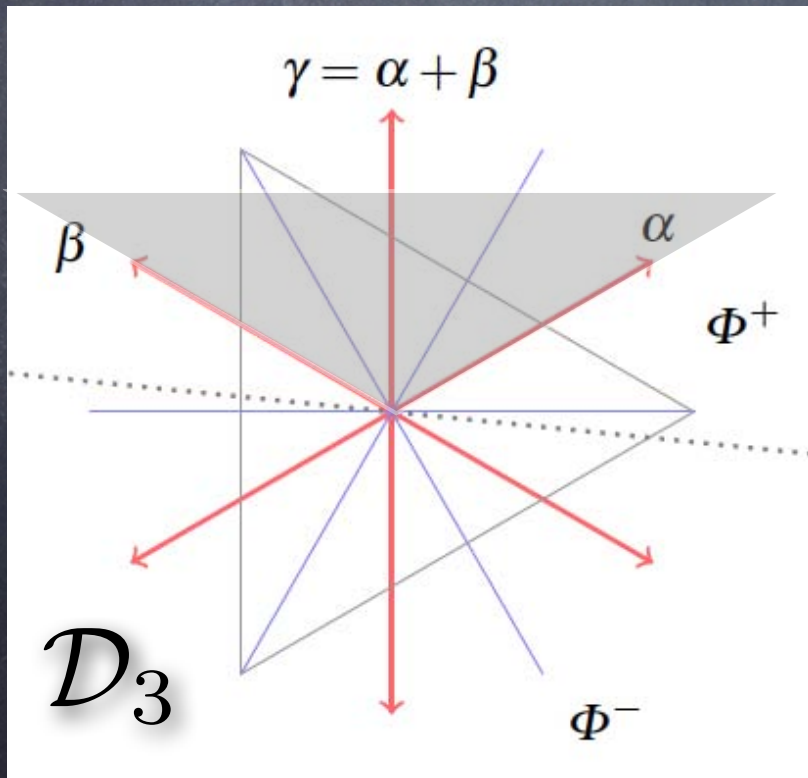


# More on the weak order

Weak order: write  $N(w) = \text{inv}(w^{-1})$  then  $u \leq v \iff N(u) \subseteq N(v)$

**Proposition.** The map  $N : W \rightarrow \mathcal{P}(\Phi^+)$  is an injective morphism of meet-semilattice. Reduced expressions 'are' chains in intervals.

What is  $\text{Im}(N)$  ?



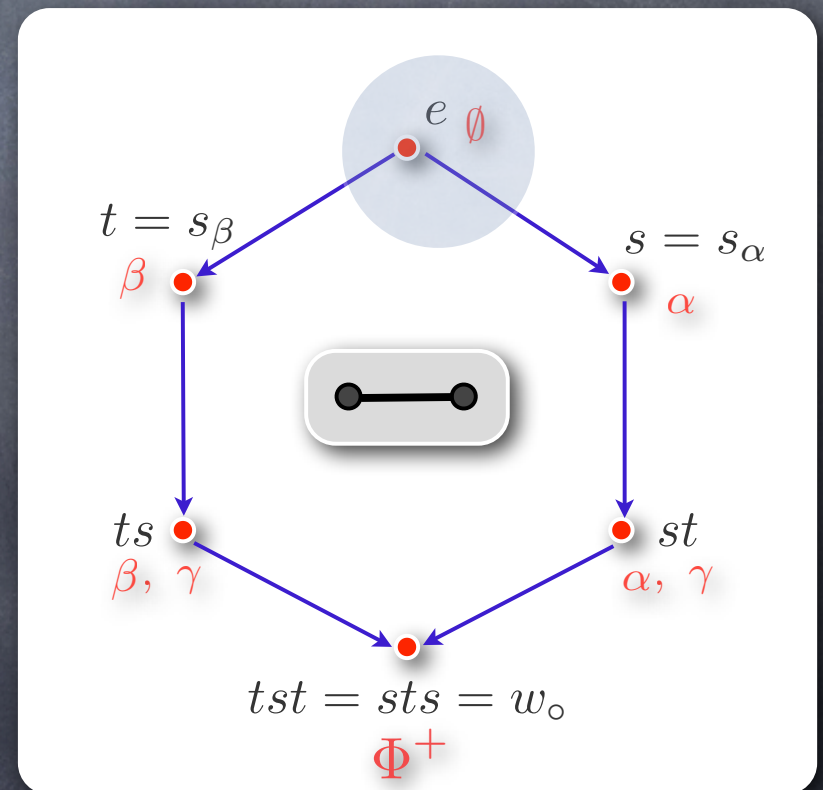
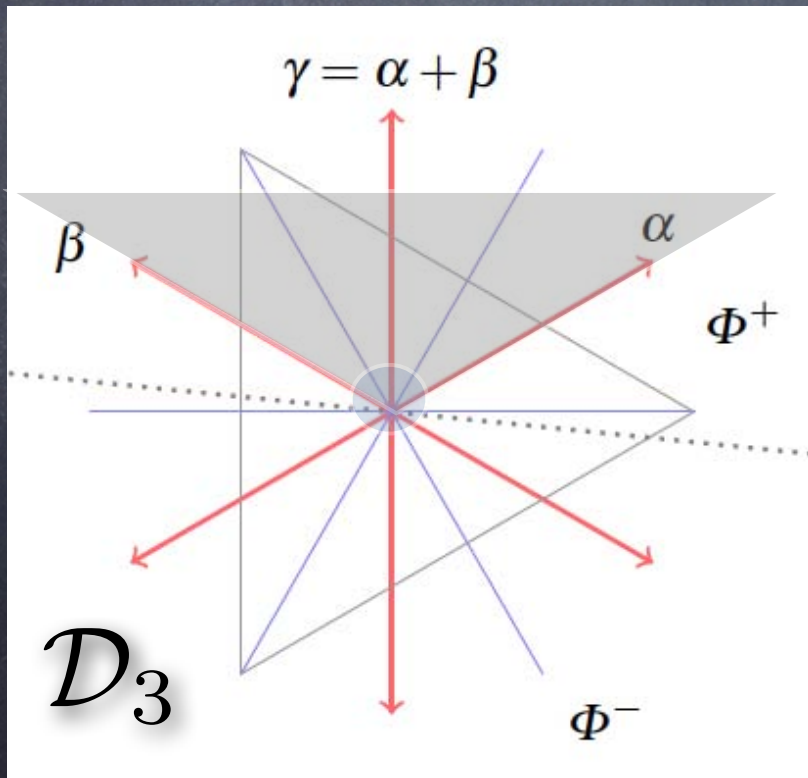


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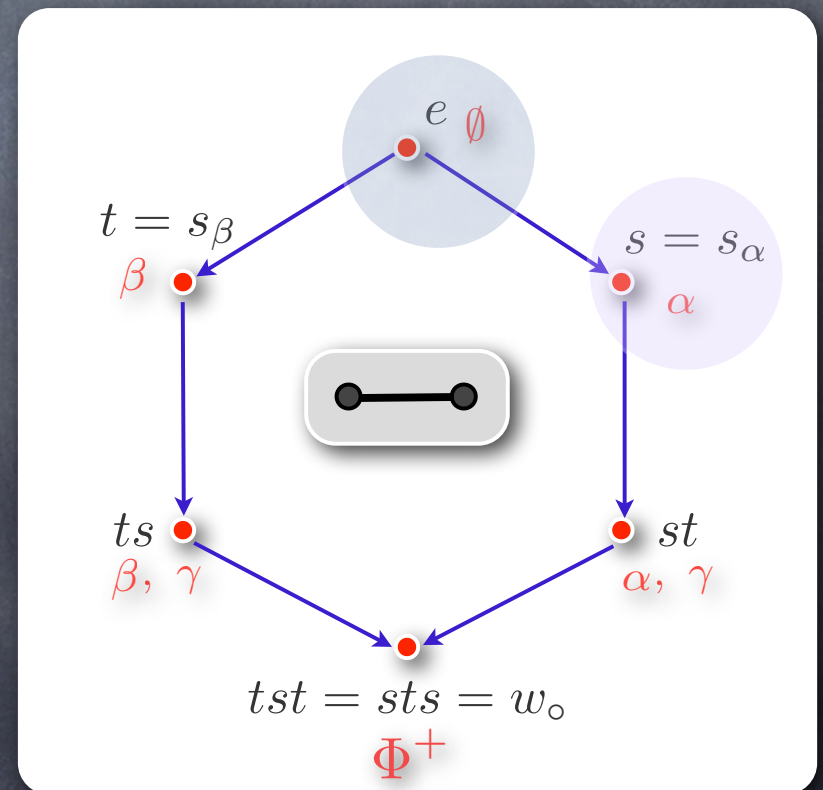
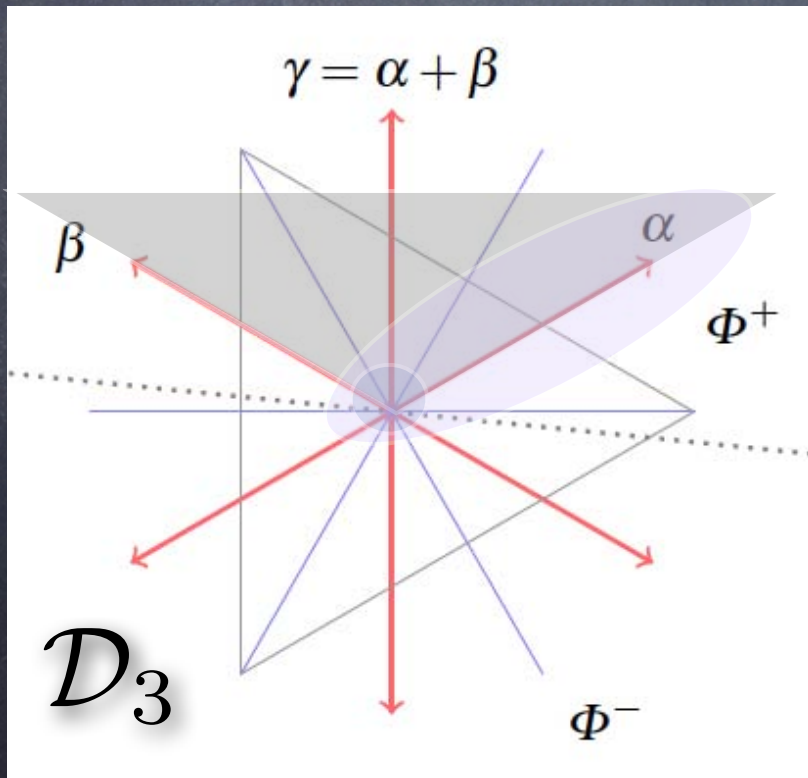


# More on the weak order

Weak order: write  $N(w) = \text{inv}(w^{-1})$  then  $u \leq v \iff N(u) \subseteq N(v)$

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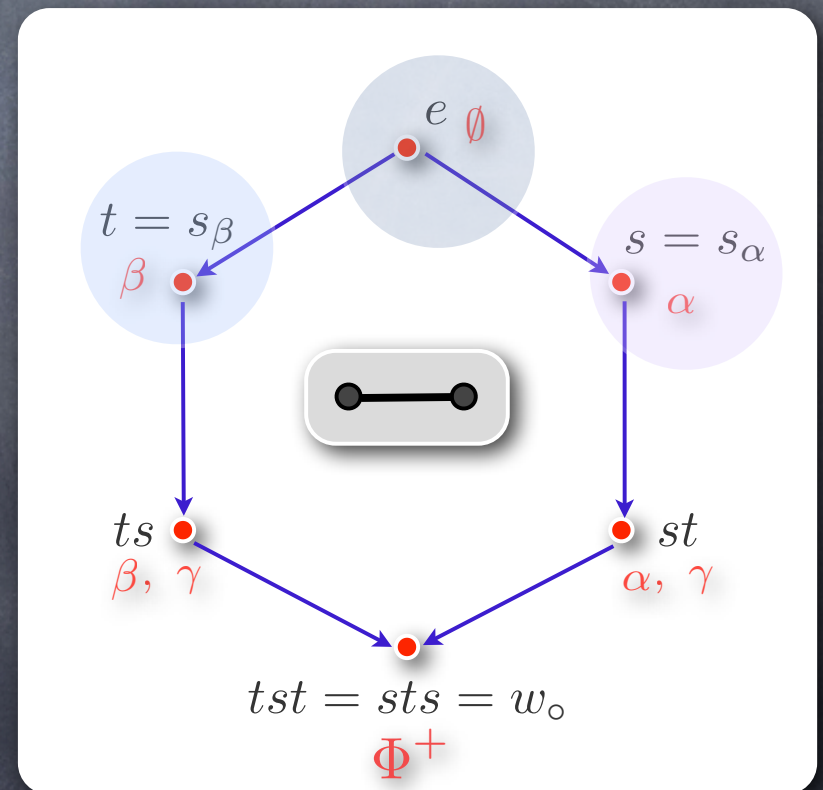
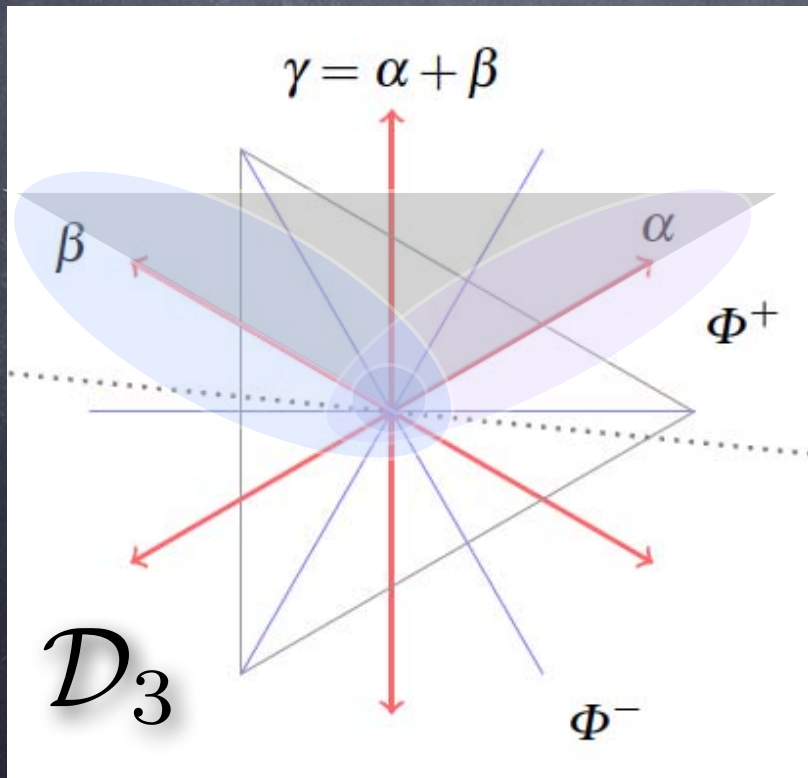


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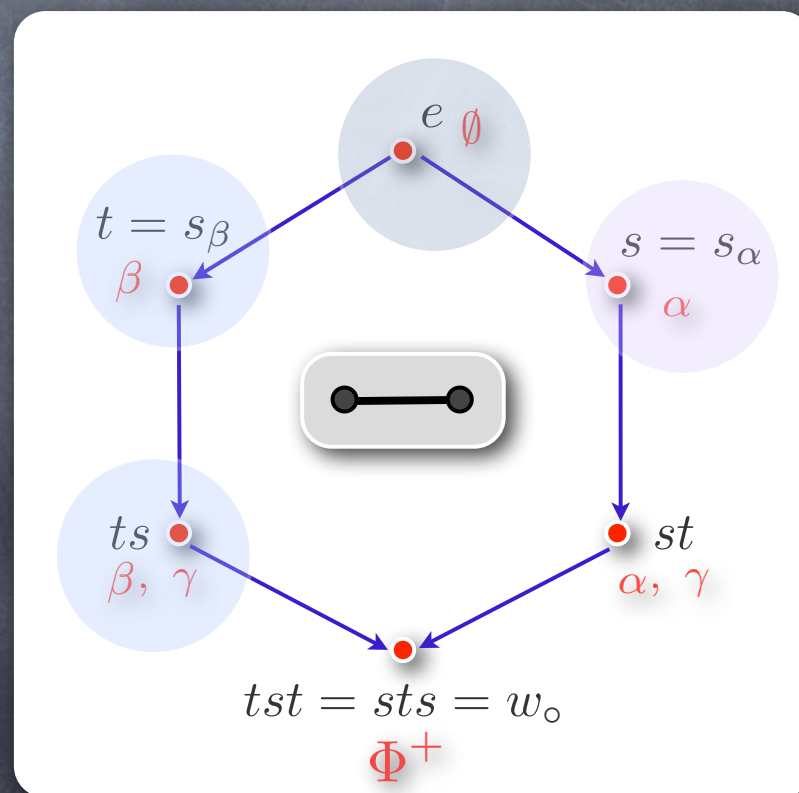
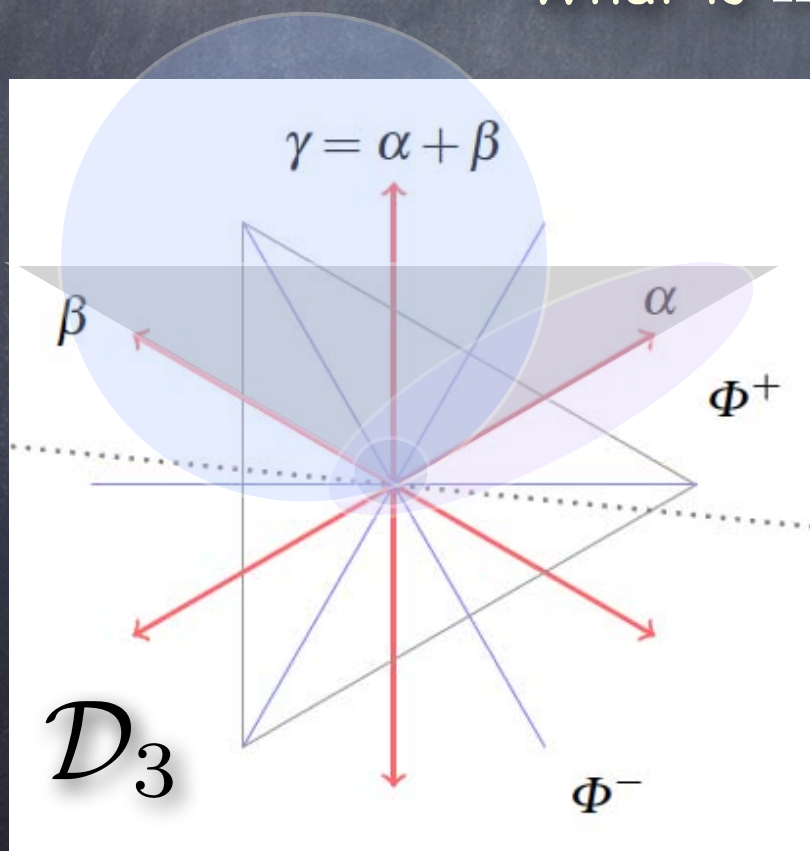


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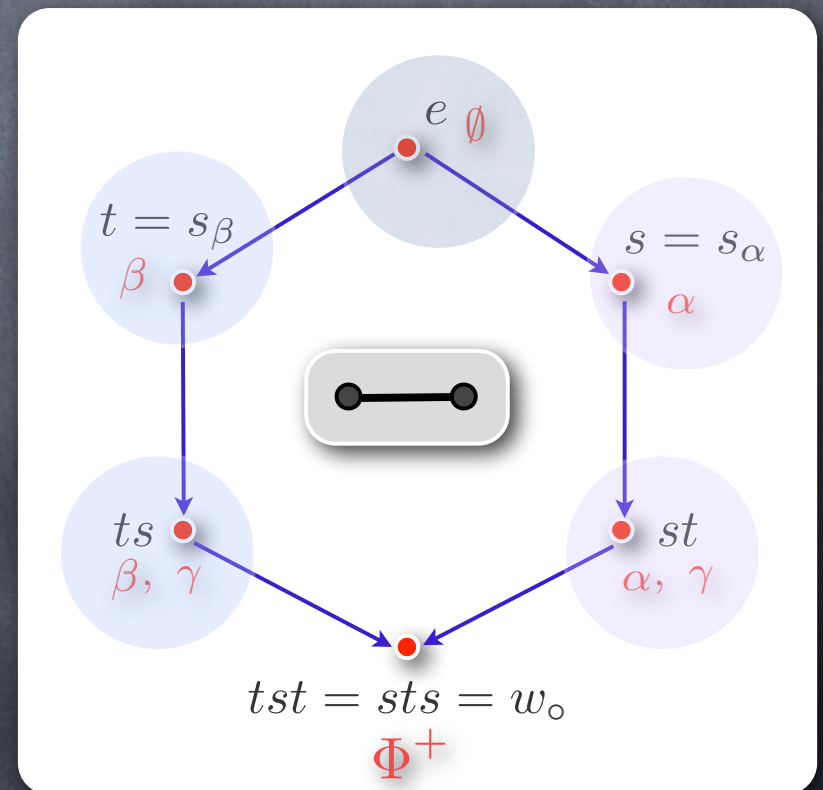
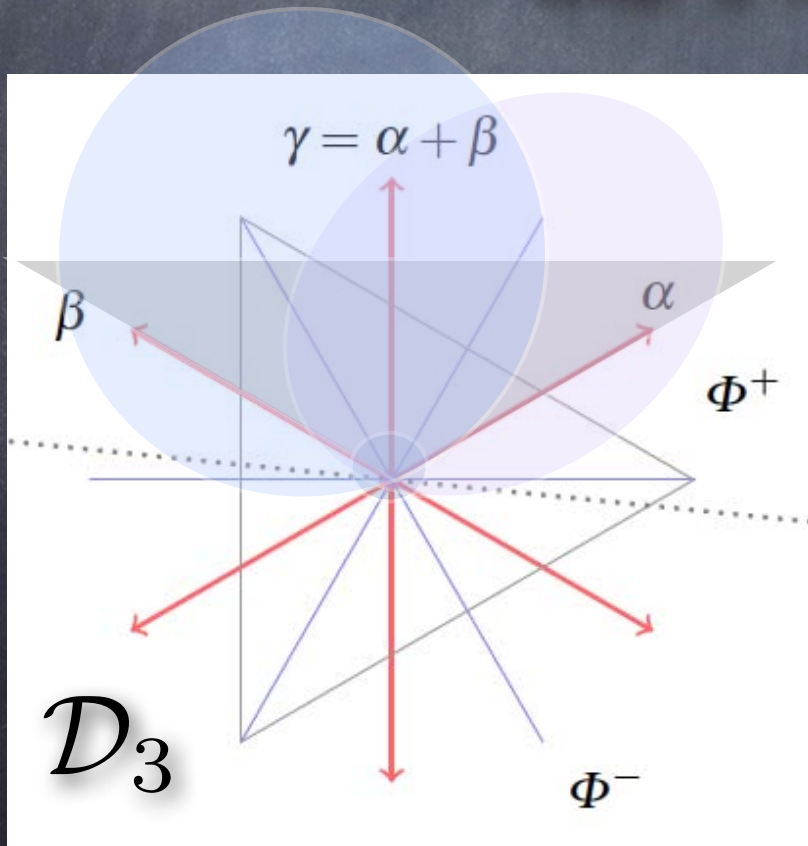


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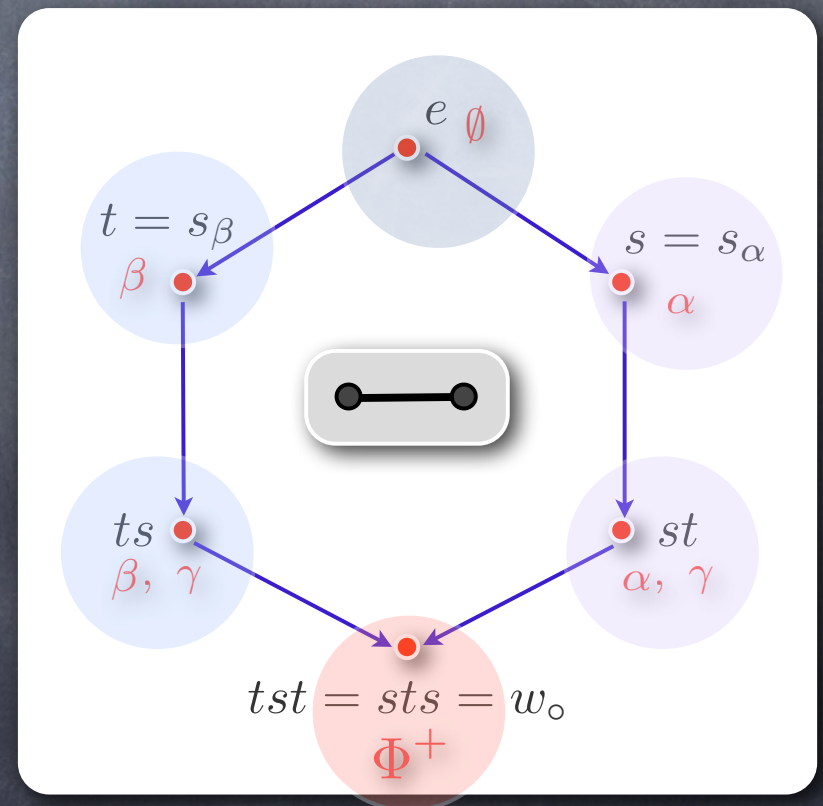
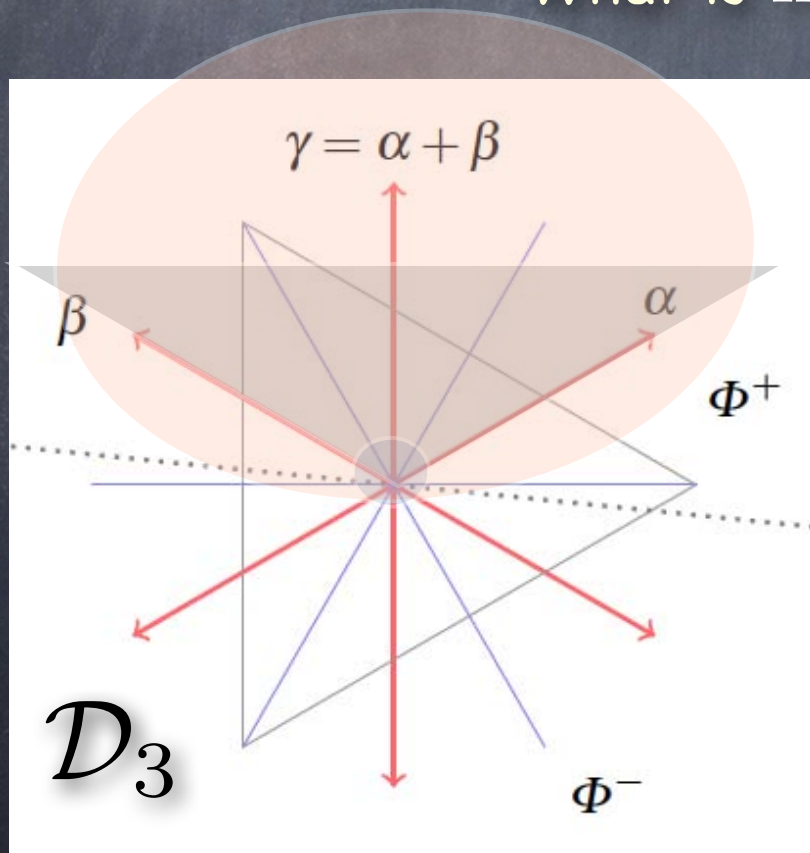


# A recap from the other way around

Weak order: write  $N(w) = \text{inv}(w^{-1})$  then  $u \leq v \iff N(u) \subseteq N(v)$

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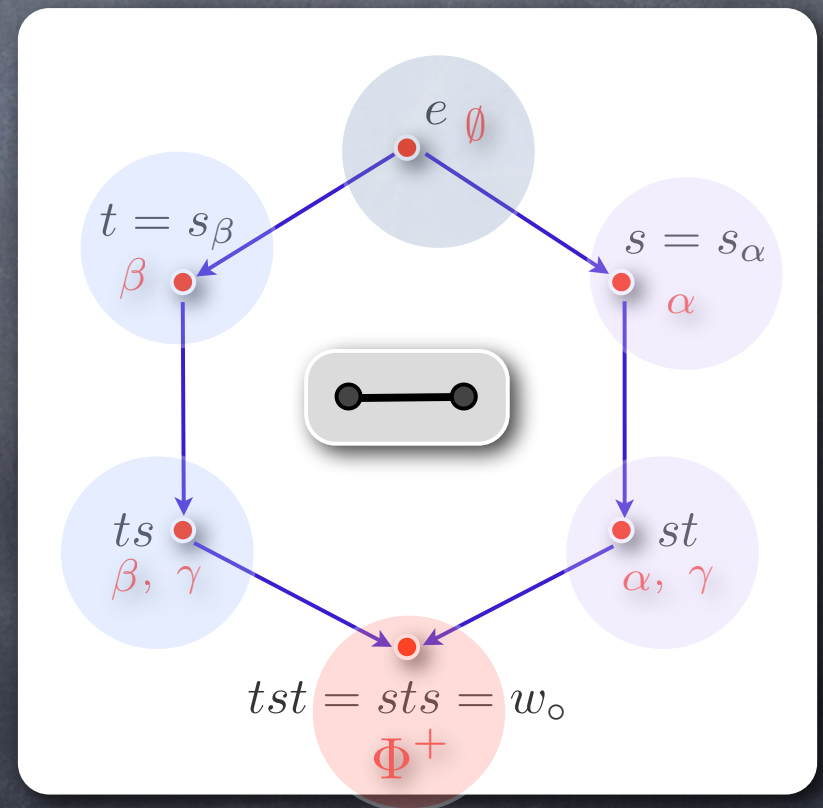
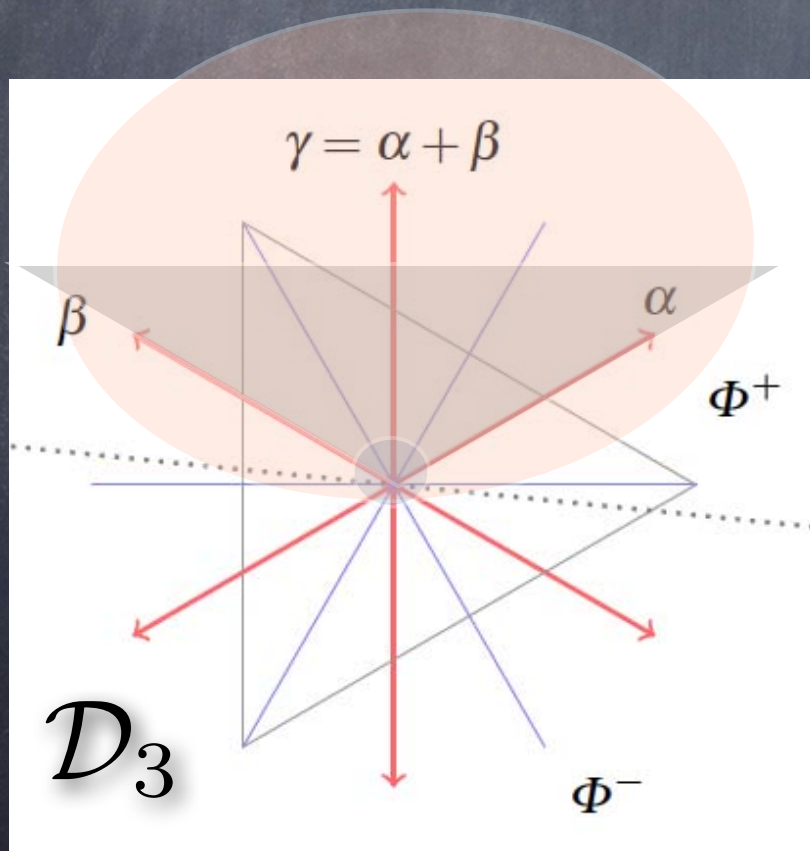




# A 'weak order lattice' in general?

**Proposition.**  $\text{Im}(N) = \{\text{finite Biclosed sets in } \Phi^+\}$

- $A \subseteq \Phi^+$  is **closed** if for all  $\alpha, \beta \in A$ ,  $\text{cone}(\alpha, \beta) \cap \Phi \subseteq A$  ;
- $A \subseteq \Phi^+$  is **biclosed** if  $A, \Phi^+ \setminus A$  are closed.
- $\mathcal{B}(W) = \{\text{biclosed sets}\}$





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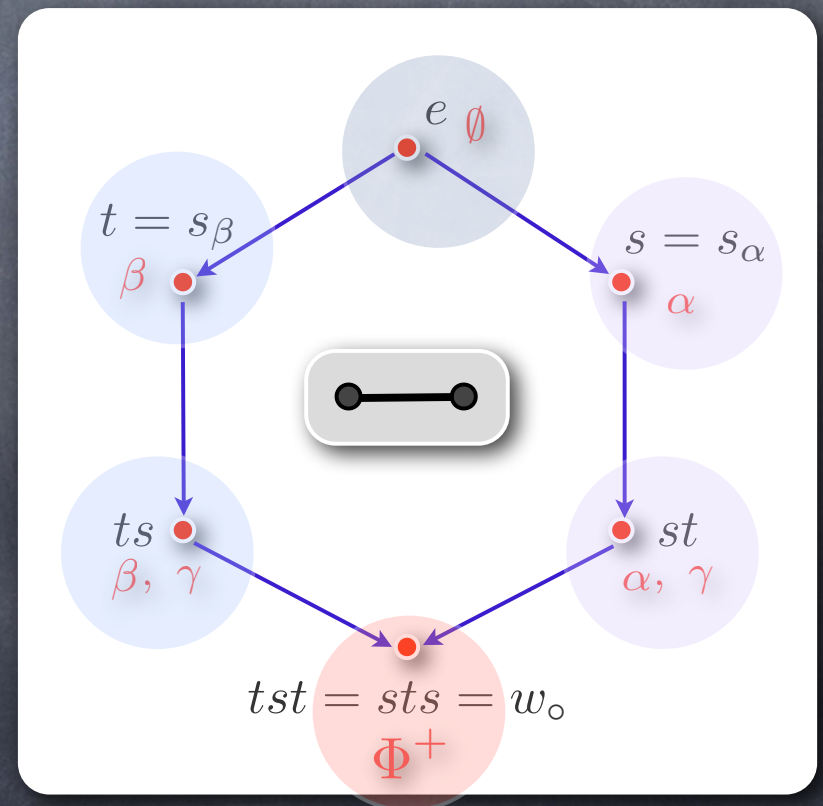
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- $\mathcal{B}(W) = \{\text{biclosed sets}\}$

**Conjecture (M. Dyer, 2011).**

$(\mathcal{B}(W), \subseteq)$  is a lattice (with minimal element  $\emptyset$  and maximal element  $\Phi^+$ )

- $\vee \neq \cup; \wedge \neq \cap$  so how to understand them geometrically?
- Biclosed sets are the candidate for «generalized words»



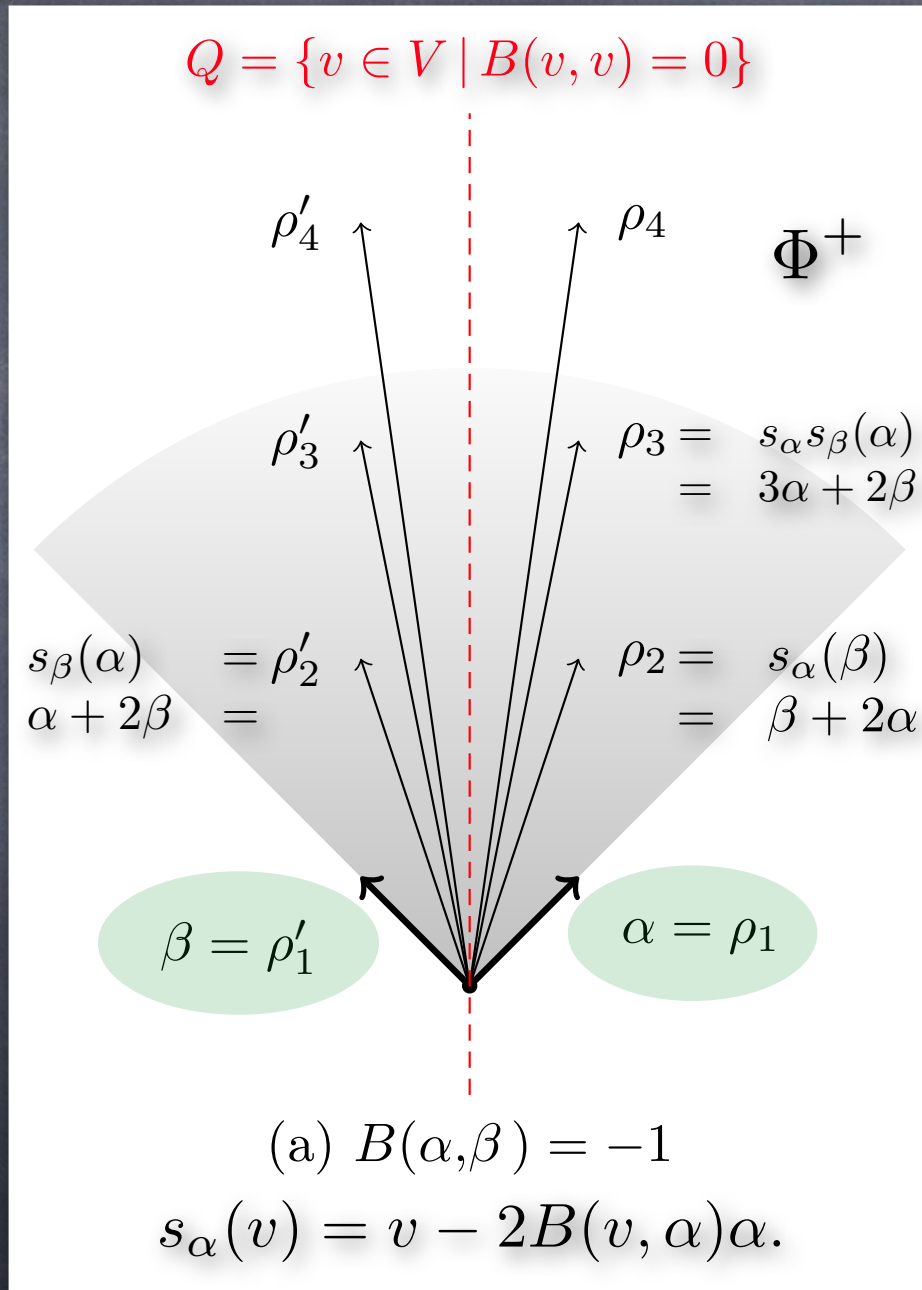
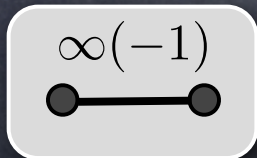


# The example of the infinite dihedral group

$$\rho'_n = n\alpha + (n+1)\beta$$

$$\rho_n = (n+1)\alpha + n\beta$$

Infinite  
dihedral  
group I



The biclosed are:

- the finite ones;
- their complements;
- and two infinite ones: the left and right side of  $Q$ !

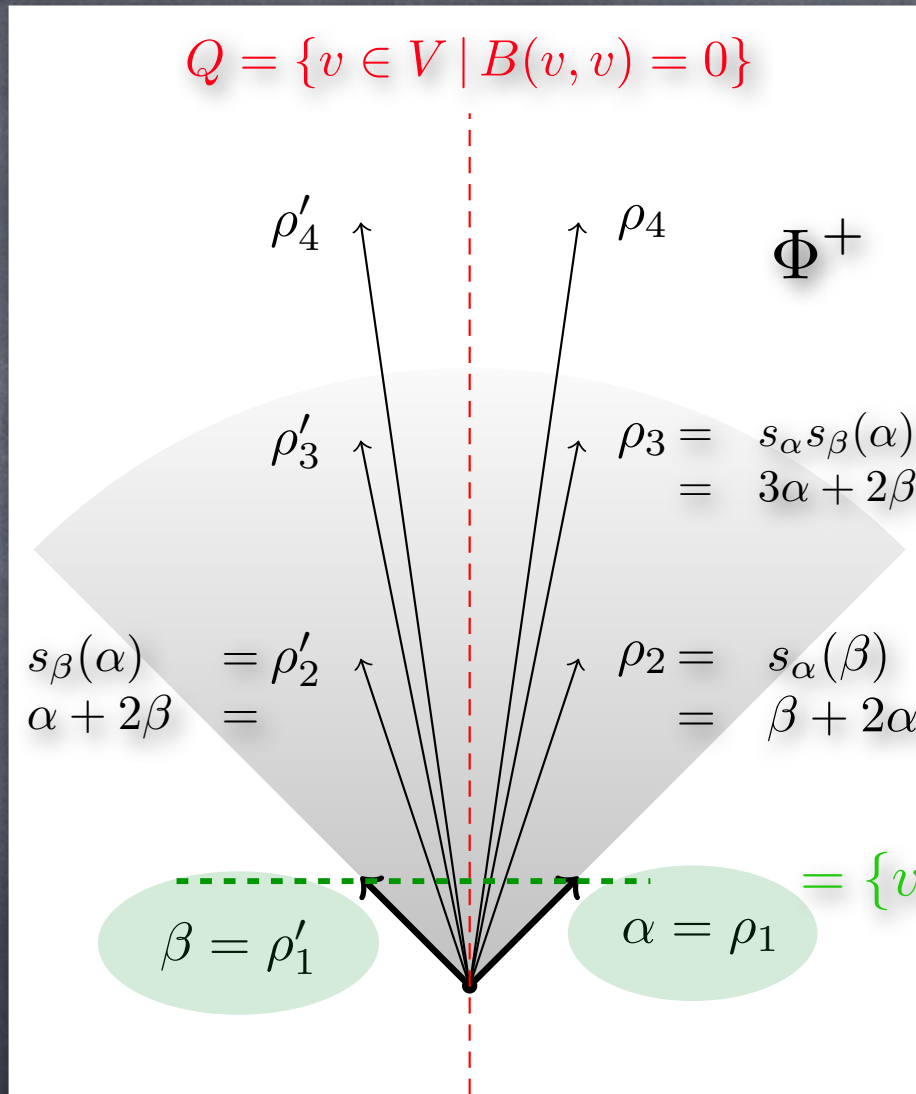
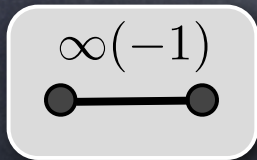


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$$(a) \quad B(\alpha, \beta) = -1$$

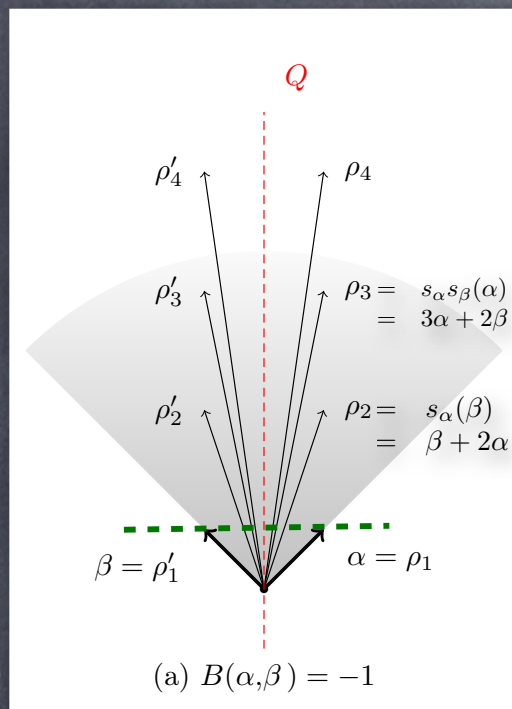
$$s_\alpha(v) = v - 2B(v, \alpha)\alpha.$$

More  
examples:  
'Cut'  $\Phi^+$  by  
an affine  
hyperplane

$$= \{v \in V \mid \sum_{\alpha \in \Delta} v_\alpha = 1\}$$



# Other examples of infinite root systems?



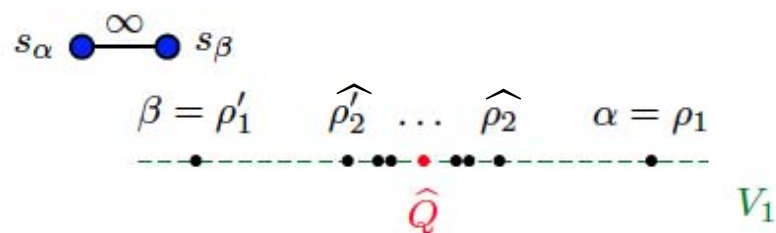
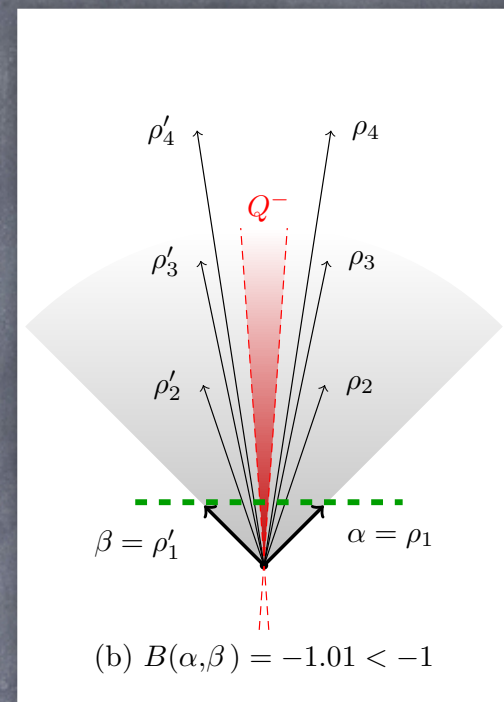
Affine hyperplane

$$V_1 = \{v \in V \mid \sum_{\alpha \in \Delta} v_\alpha = 1\}$$

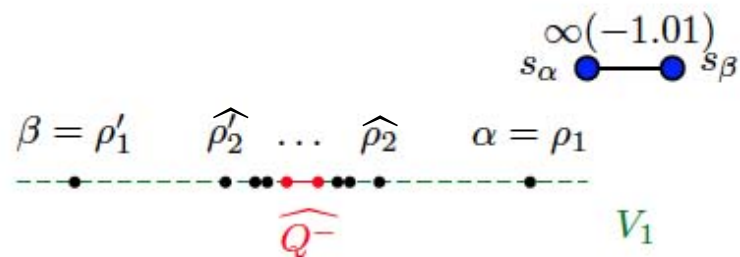
Normalized isotropic cone:  $\hat{Q} := Q \cap V_1$

Normalized roots

$$\hat{\rho} := \rho / \sum_{\alpha \in \Delta} \rho_\alpha$$



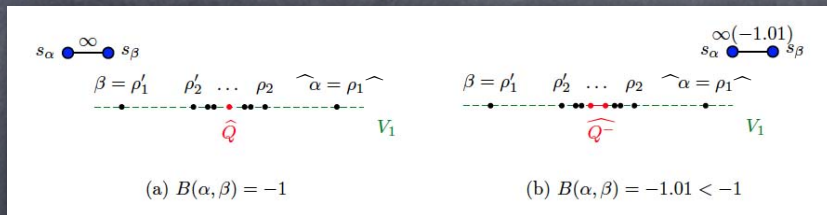
(a)  $B(\alpha, \beta) = -1$



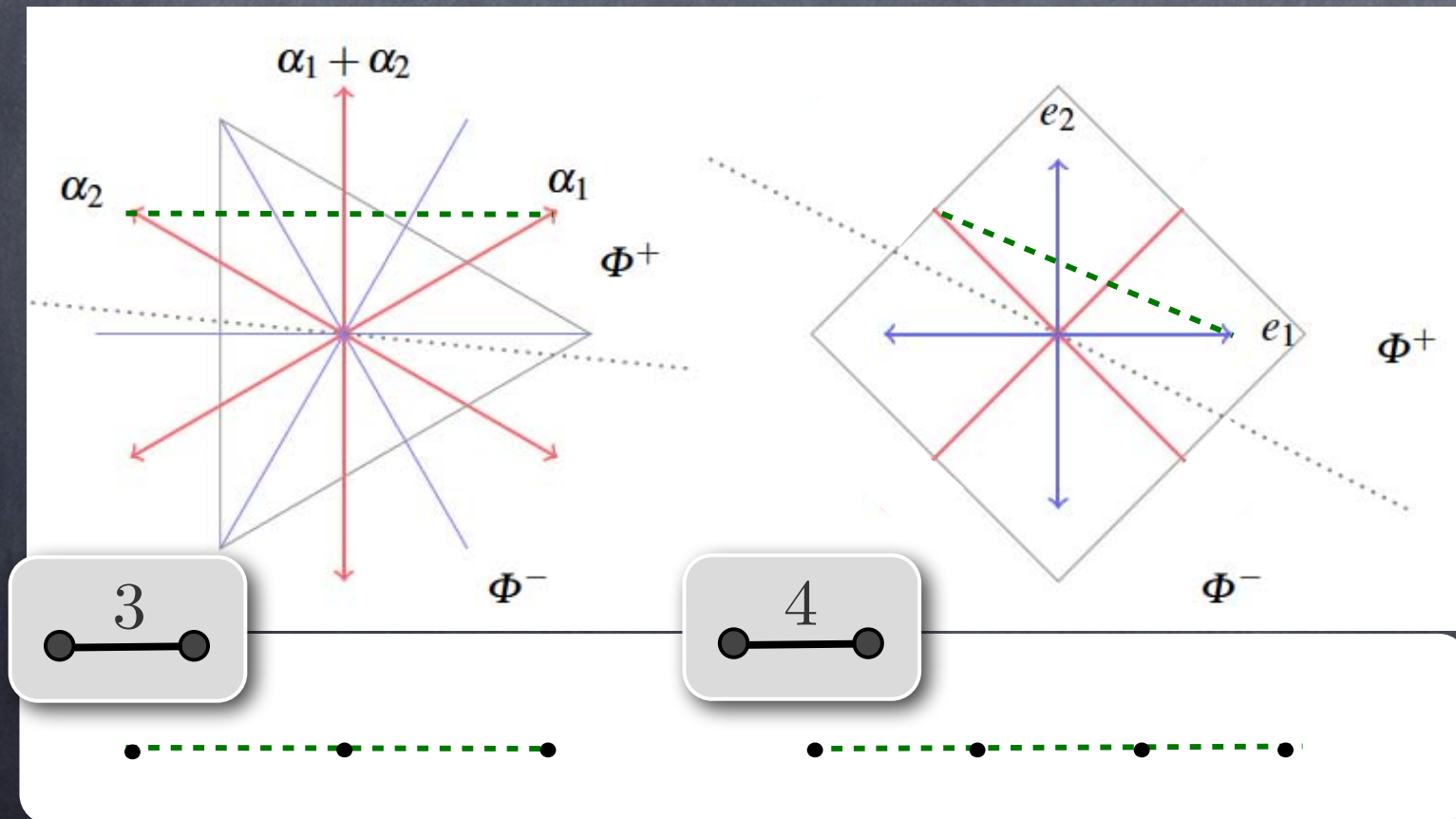
(b)  $B(\alpha, \beta) = -1.01 < -1$



# Other examples of infinite root systems ...

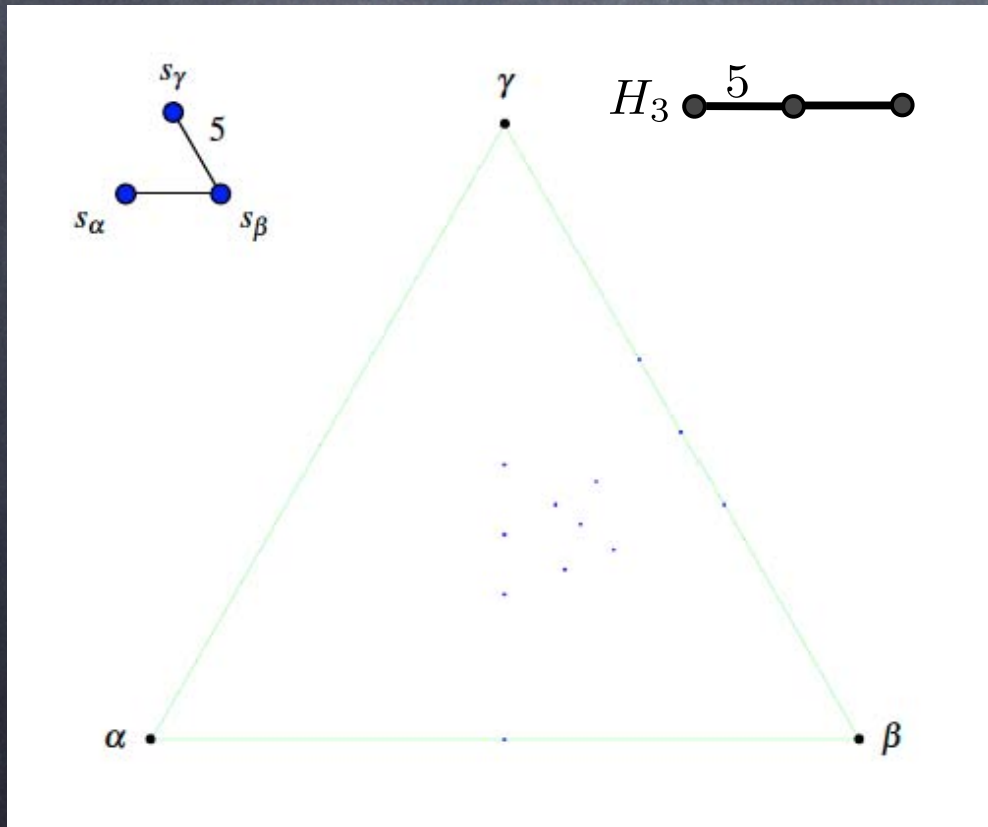
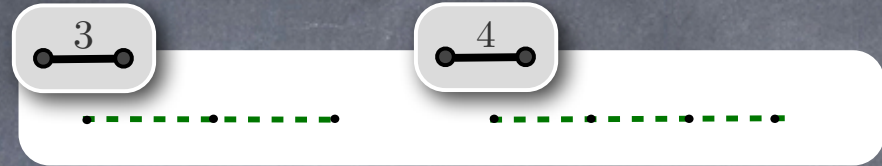
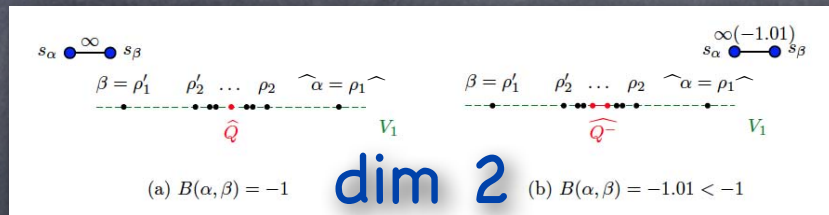


finite type of rank 2

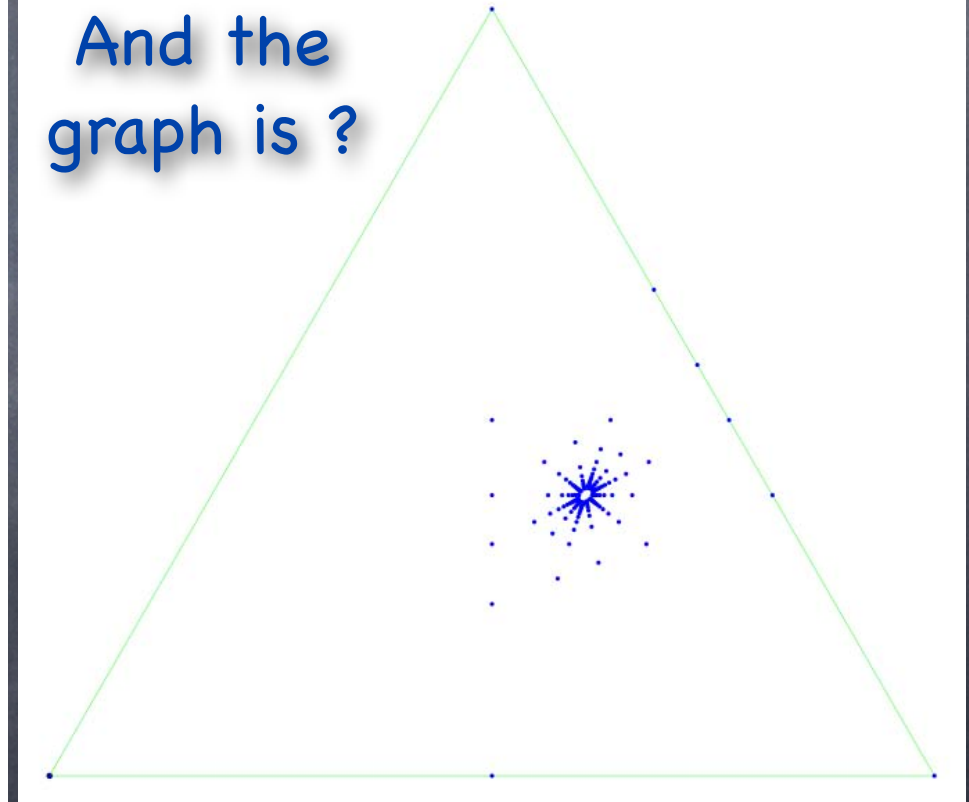




# Other examples of infinite root systems (with SAGE)

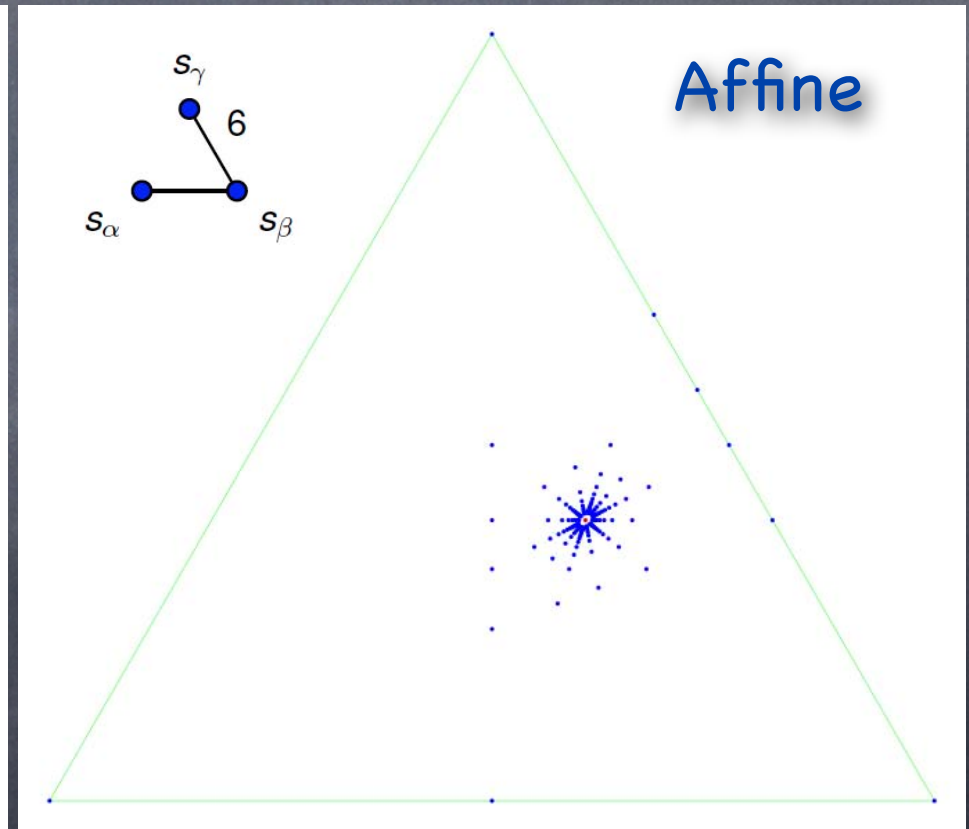
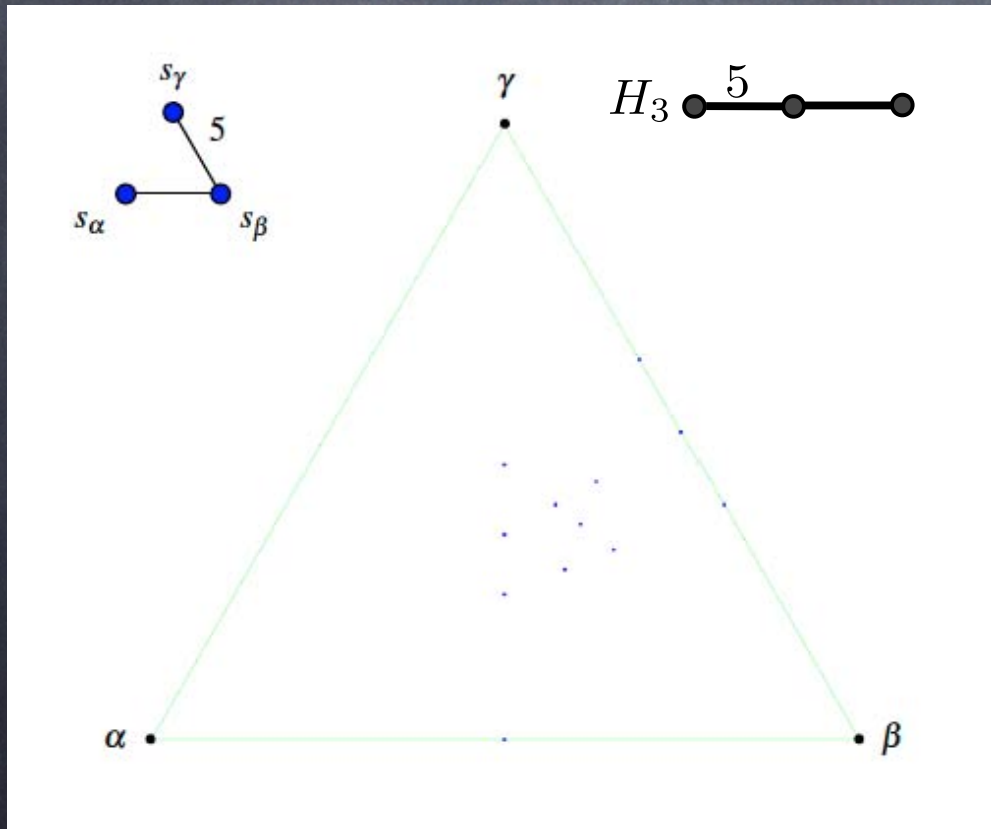
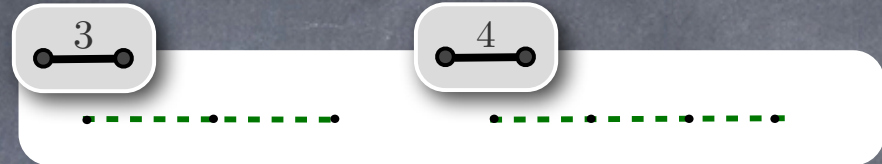
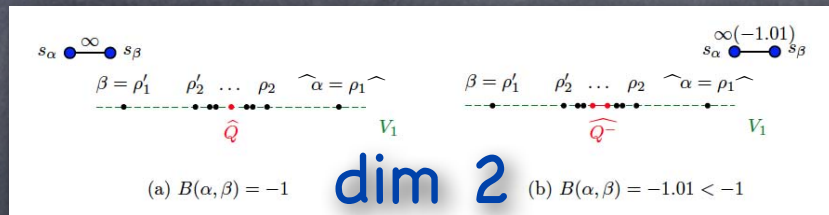


And the graph is ?



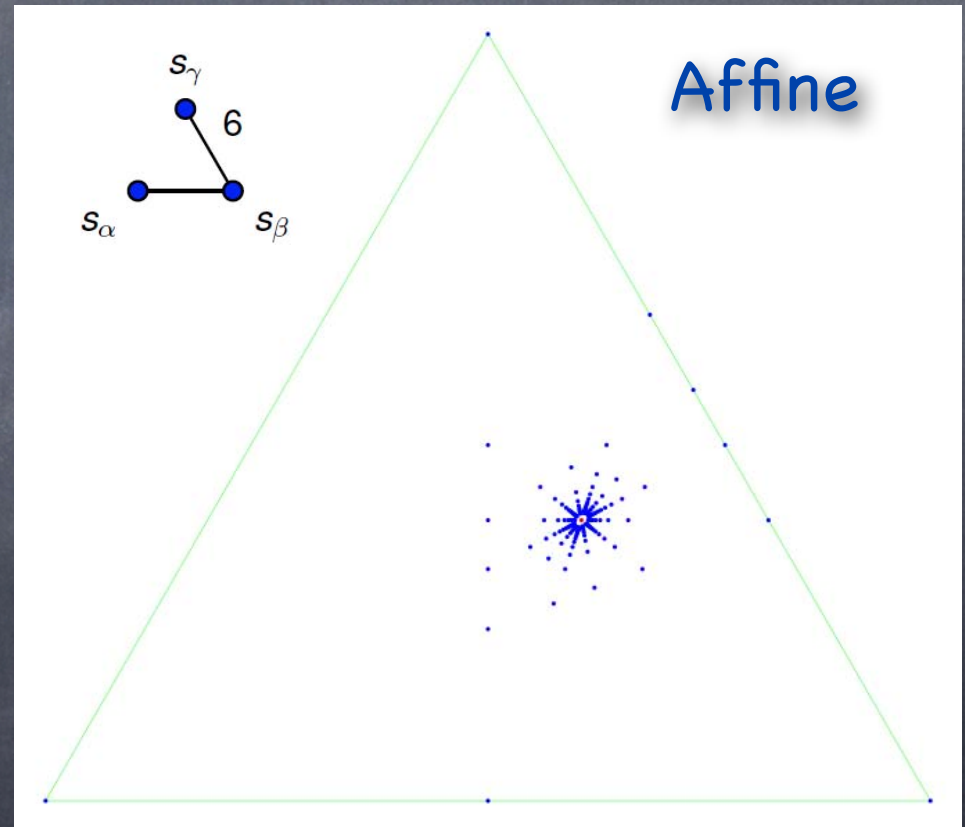
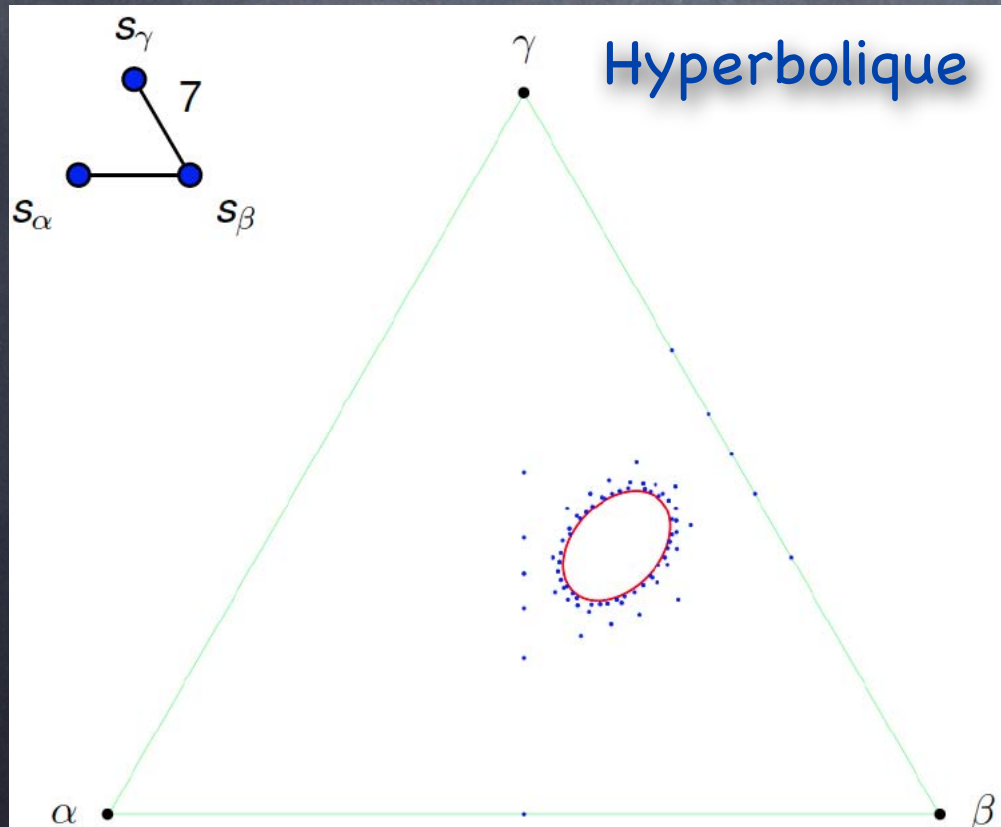
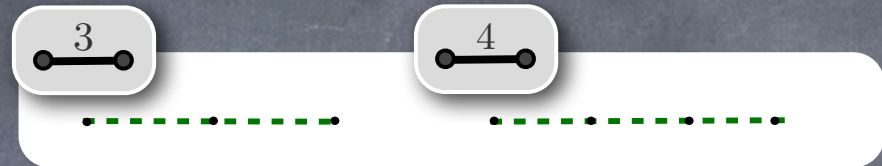
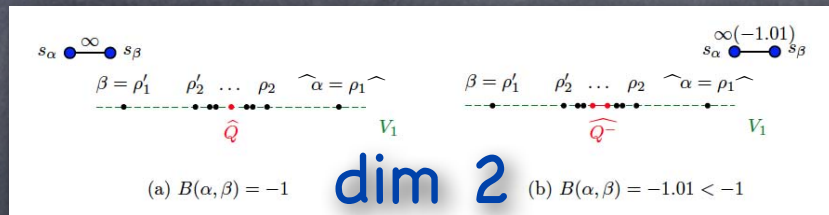


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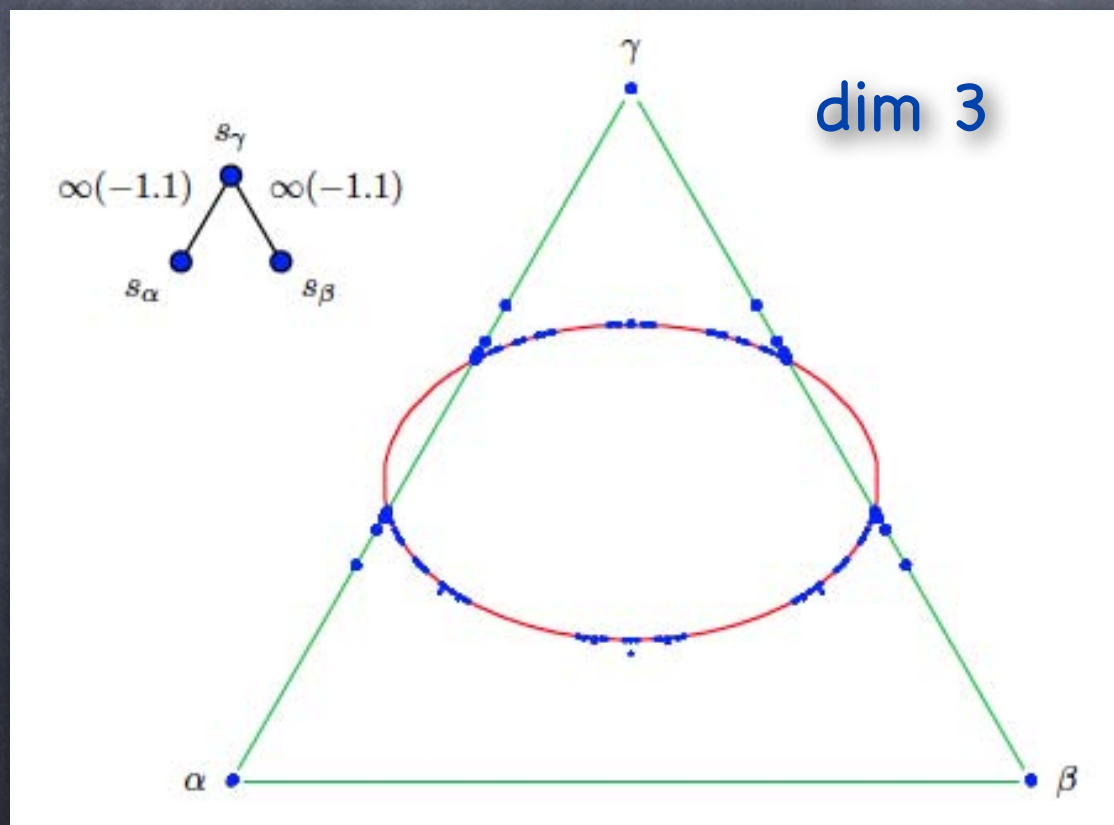
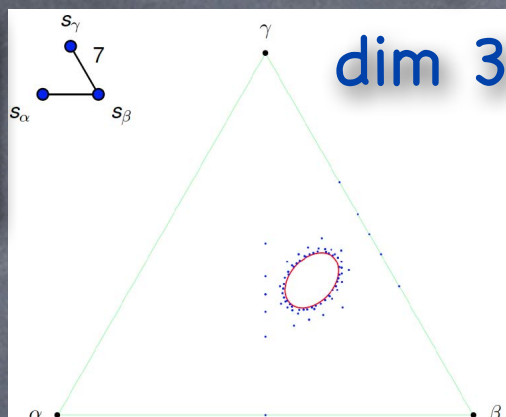
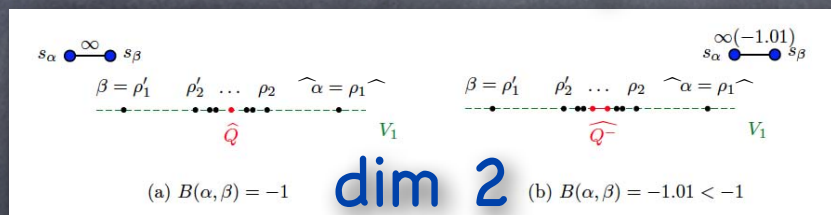


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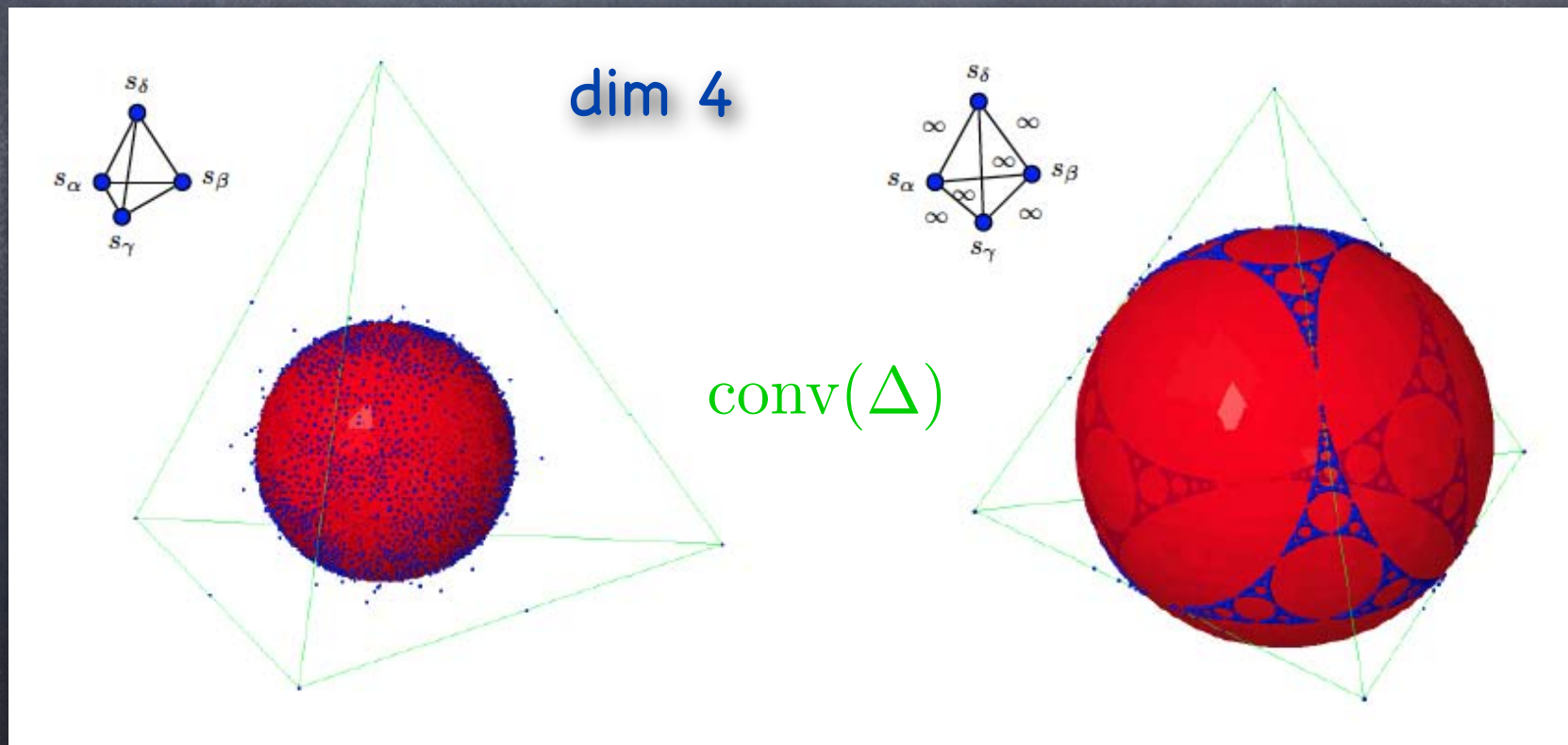
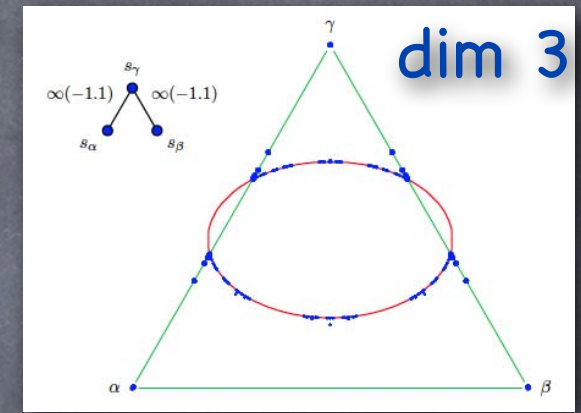
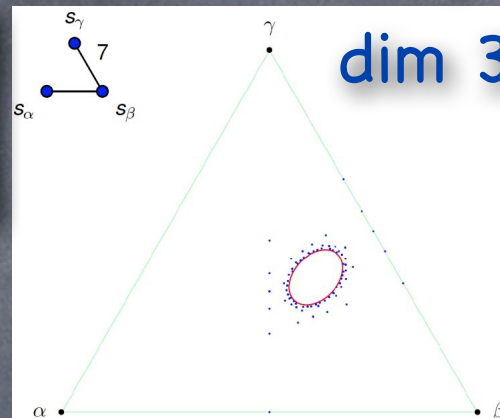
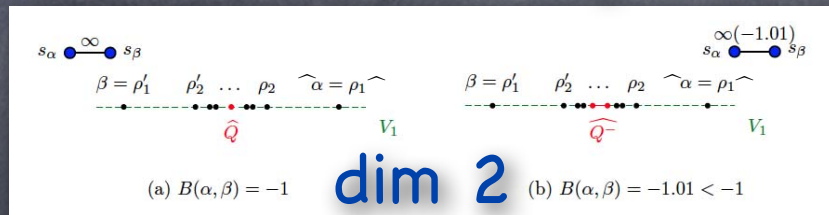
# Other examples of infinite root systems (with SAGE)



**Observation:** a dihedral subgroup group is infinite iff the associated line cuts  $Q$

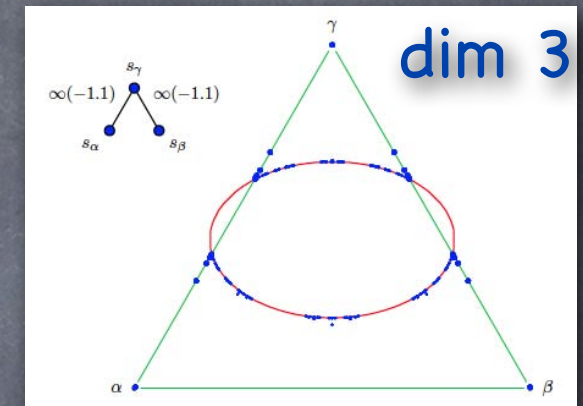
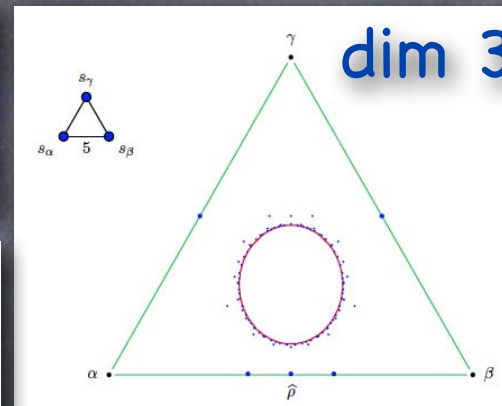
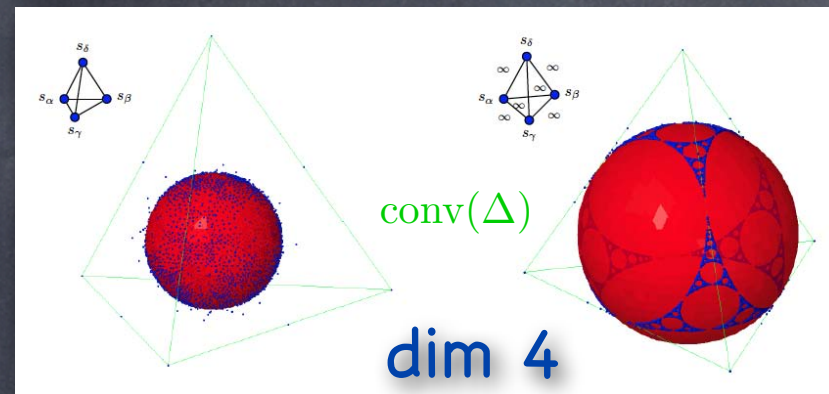
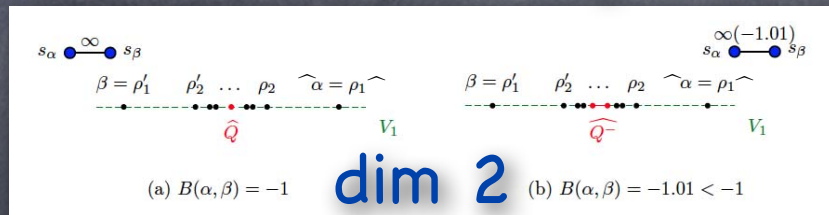


# Other examples of infinite root systems (with SAGE)





# Other examples of infinite root systems (with SAGE)



Sgn is (2, 2)

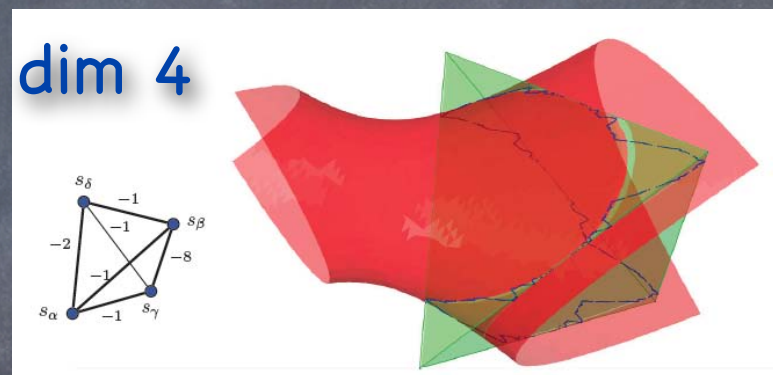
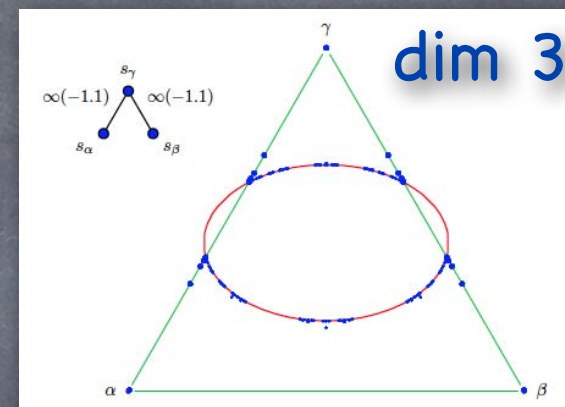
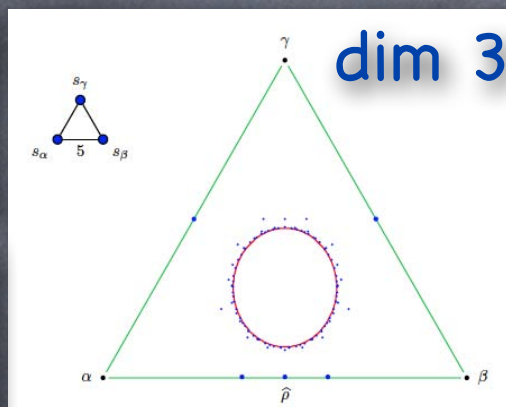
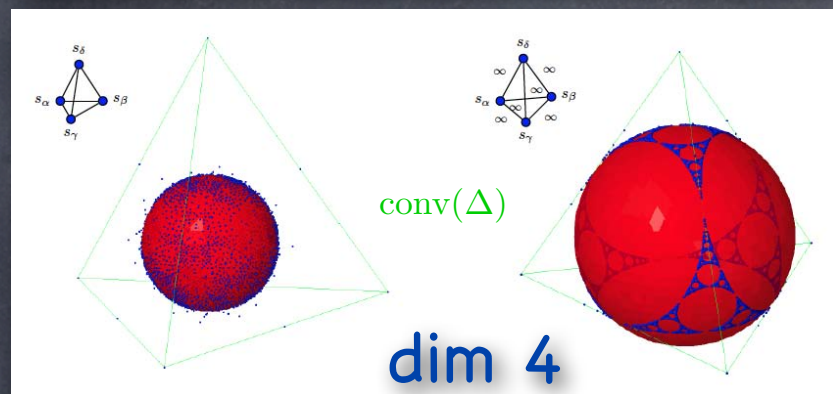
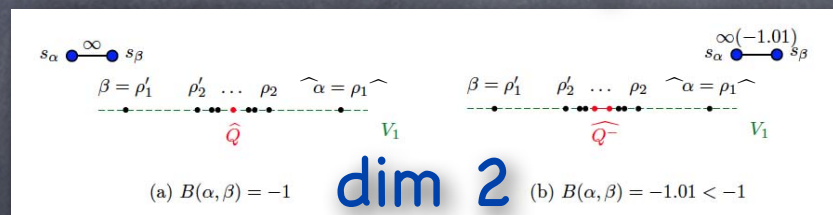


From joint works  
with:

- J.P Labbé and V. Ripoll (2012)
- M. Dyer and V. Ripoll (2013)



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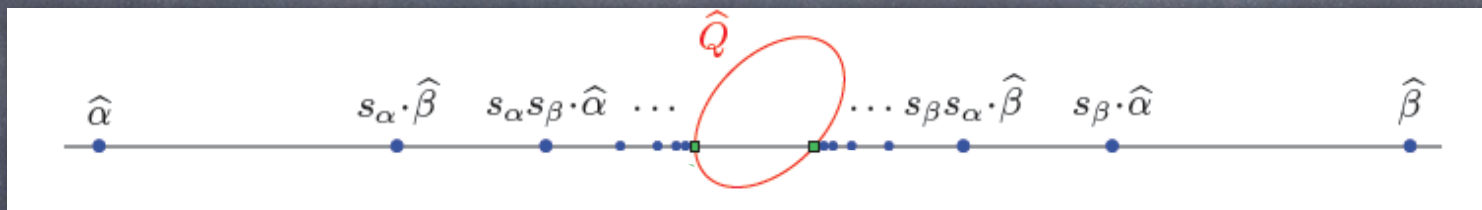
Problem still there: what can we say about these pictures that help understand biclosed sets?

Actually, at this point, not that much about biclosed but...



# Illustration of some known combinatorics on infinite root systems

How to see the action of  $W$  on  $\widehat{\Phi}$  :  $s_\alpha \cdot \beta = \widehat{s_\alpha(\beta)} \in L(\widehat{\alpha}, \widehat{\beta})$   
is a barycenter of  $\widehat{\alpha}$  and  $\widehat{\beta}$ .



Depth of a root is  $\text{dp}(\rho) = 1 + \min\{k \mid \rho = s_{\alpha_1} s_{\alpha_2} \cdots s_{\alpha_k}(\alpha_{k+1}),$   
 $\alpha_1, \dots, \alpha_k, \alpha_{k+1} \in \Delta\}.$

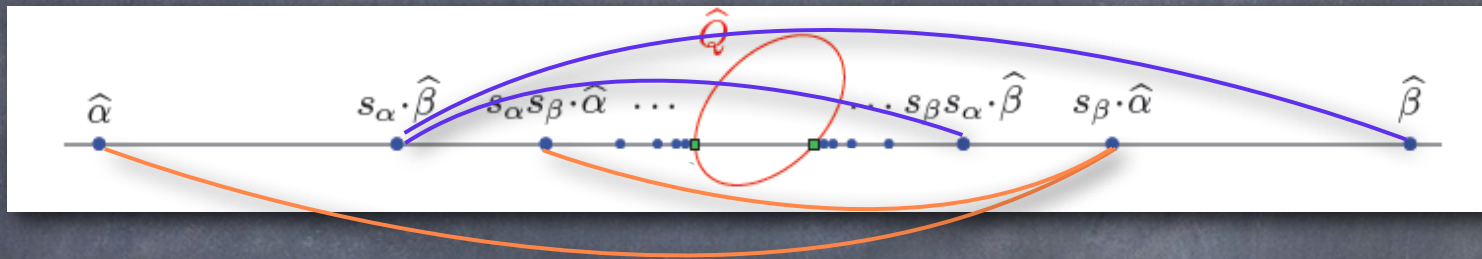
Root poset on  $\Phi^+$  : transitive closure of the relation

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**Small root:** roots obtained from  $\Delta$  along a path in the root poset corresp. to finite dihedral reflection subgroups (i.e. the lines does not cut  $Q$ ).

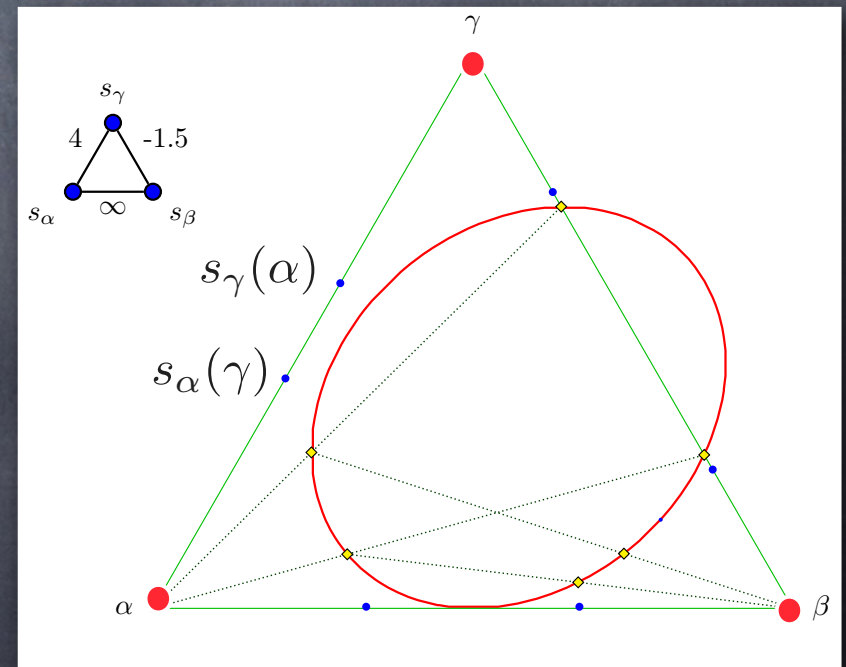
**Theorem** (Brink-Howlett, 1993)  
The set  $\Sigma$  of small roots is finite.

**Problem:**

$|\Sigma|$  from  $\Gamma_W$  ?

A finite state automaton that recognize reduced expressions:

- everything depends of the combinatorics of the small descent set  $D_\Sigma(w) = \text{inv}(w) \cap \Sigma$
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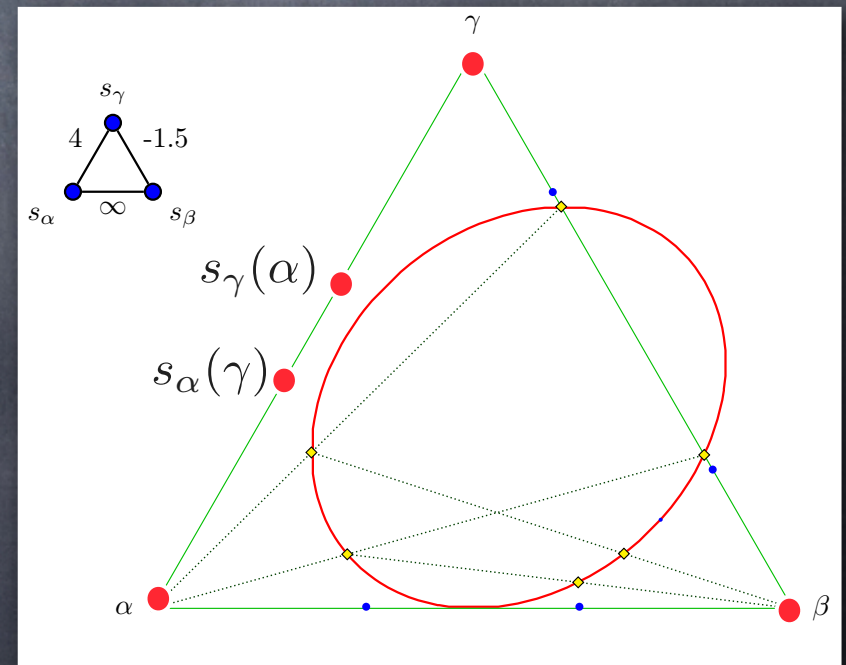
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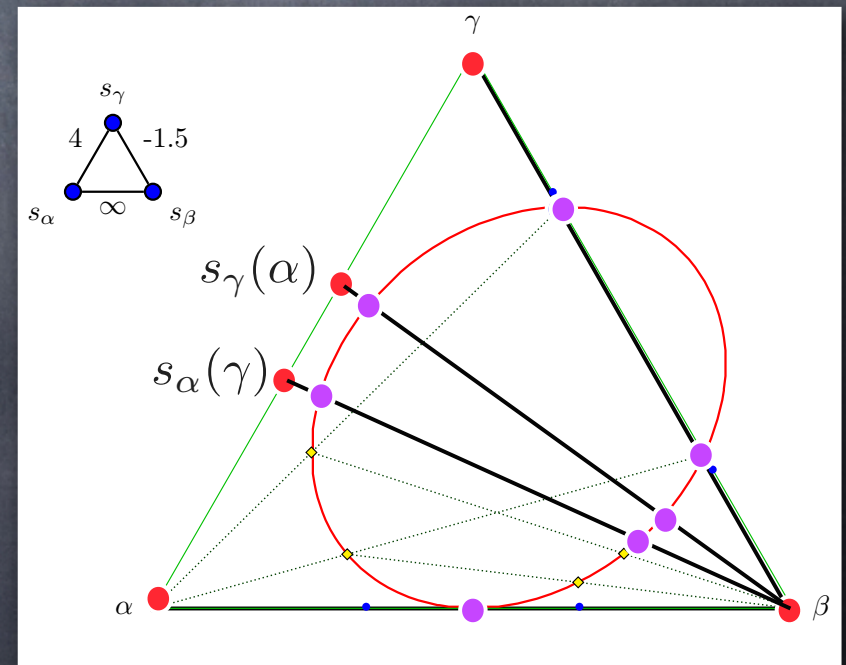
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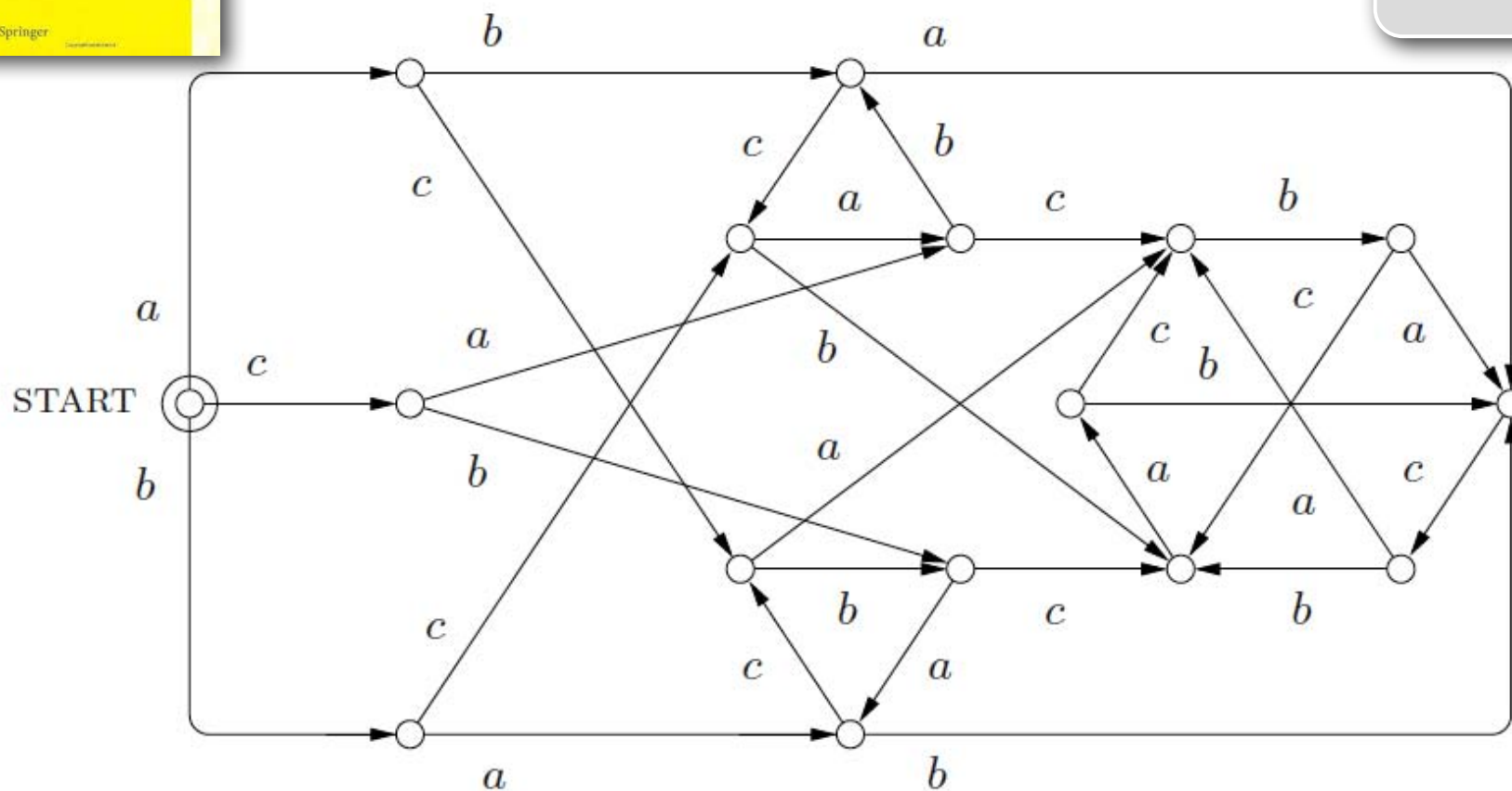
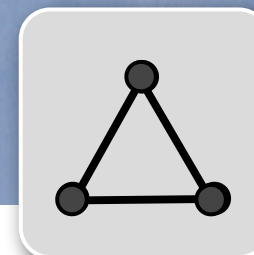
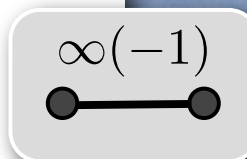
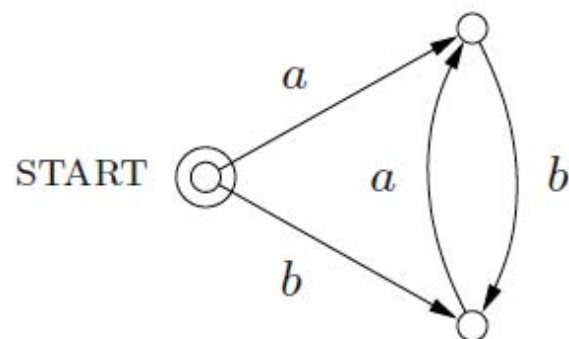
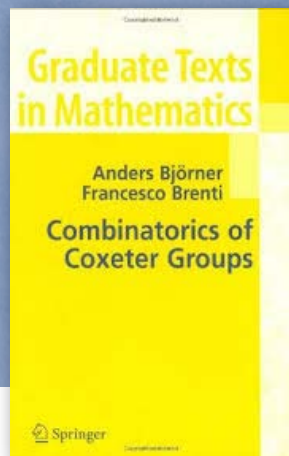
**Problem:**  
 $|\Sigma|$  from  $\Gamma_W$  ?

A finite state automaton that recognize reduced expressions:

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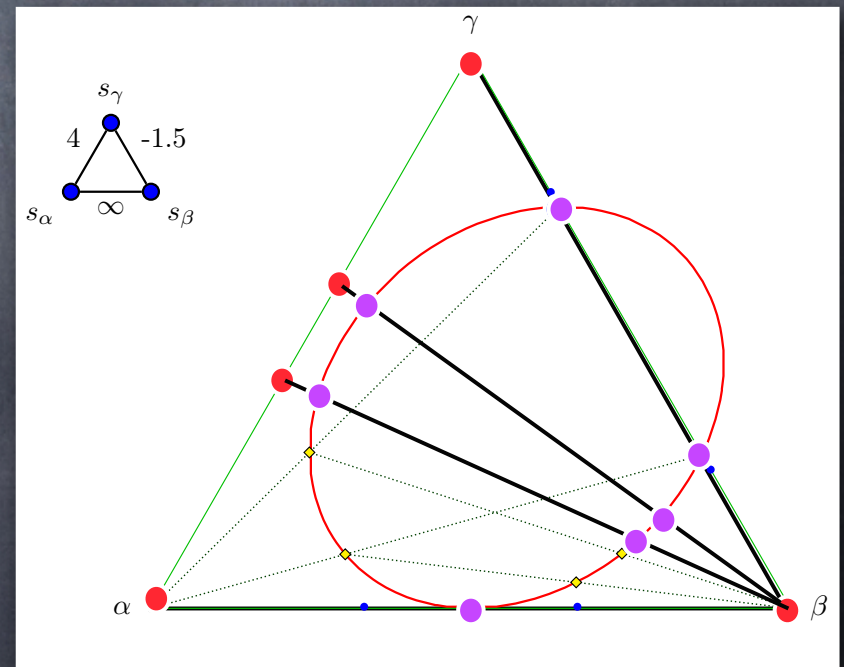
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For building a finite state automaton that recognize reduced expressions:

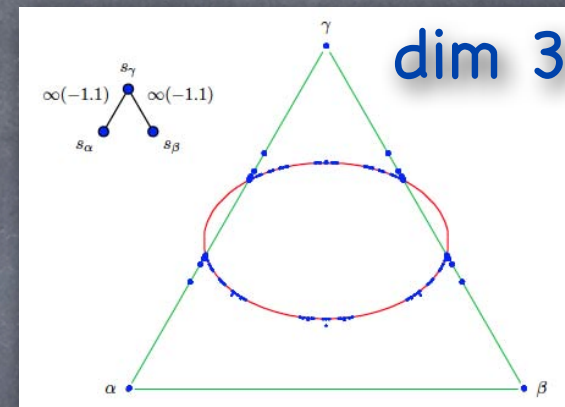
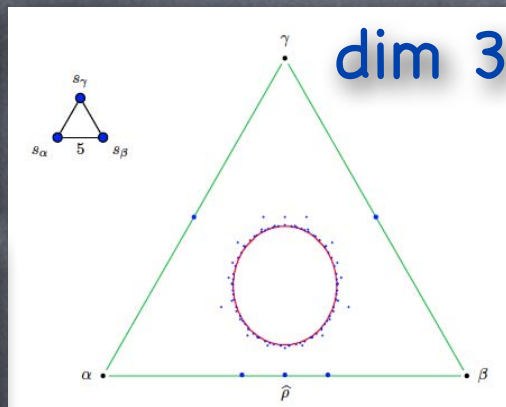
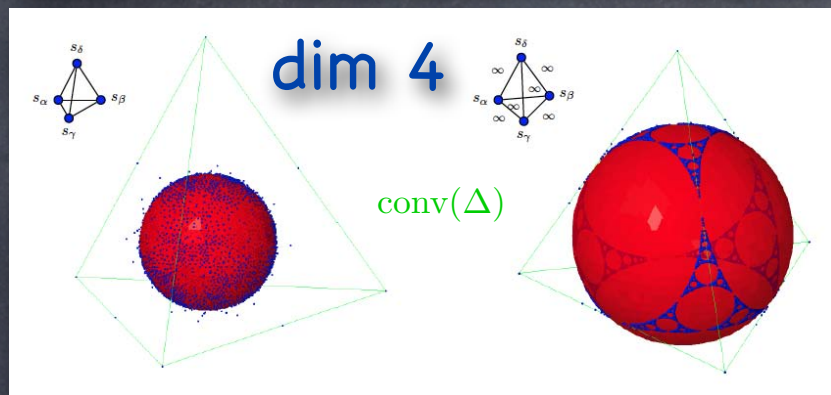
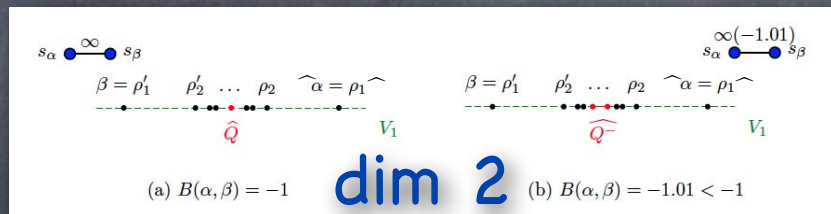
**Question:**  
Is it possible to recognize biclosed sets???

If  $A$  is biclosed, properties  
of  $\sum_{\beta \in A} q^{\text{dp}(\beta)}$ ?



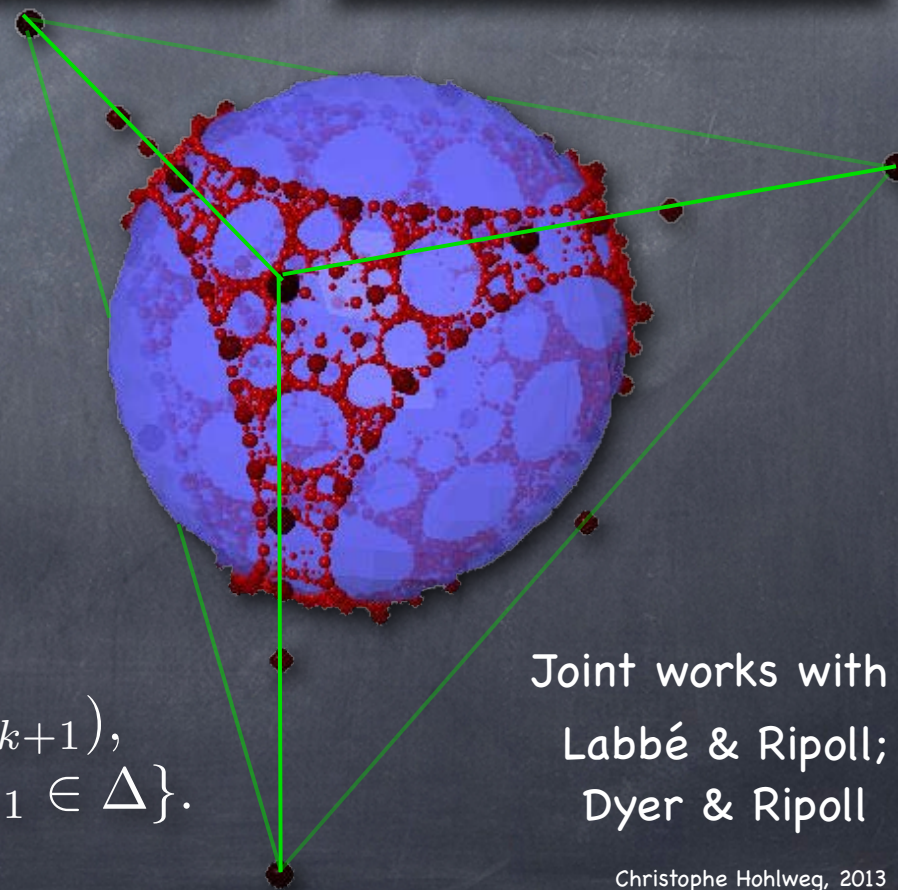


# What can we say about these pictures?



Observation: The 'size' of a generalized root (in red in this last picture) is decreasing as the depth of the root is increasing.

$$\text{dp}(\rho) = 1 + \min\{k \mid \rho = s_{\alpha_1} s_{\alpha_2} \dots s_{\alpha_k}(\alpha_{k+1}), \alpha_1, \dots, \alpha_k, \alpha_{k+1} \in \Delta\}.$$



Joint works with  
Labbé & Ripoll;  
Dyer & Ripoll



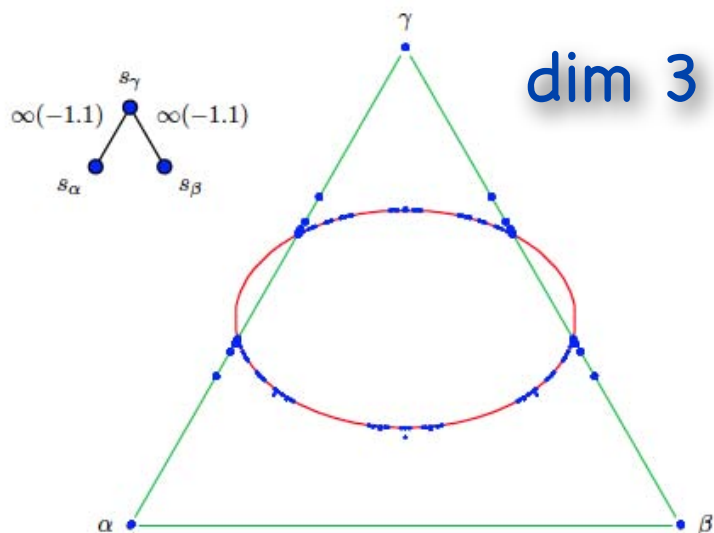
# A look at limits of roots

Joint works with :  
Labbé & Ripoll;  
Dyer & Ripoll

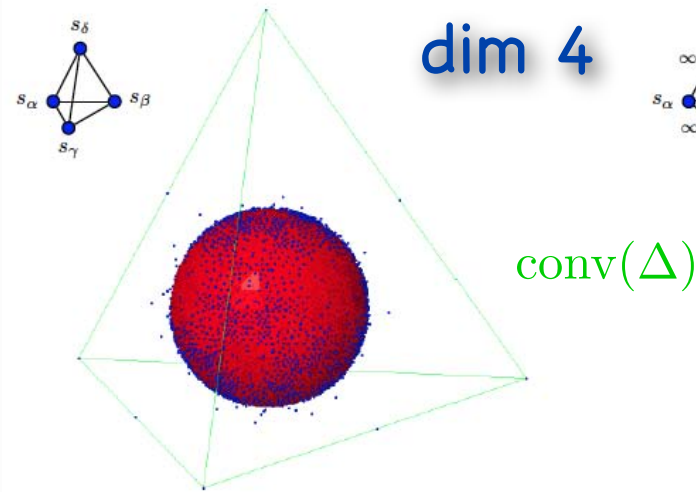
**Definition/Proposition:** the set of limit roots is:

$$E(\Phi) = \text{Acc}(\widehat{\Phi}) \subseteq Q \cap \text{conv}(\Delta)$$

- A 'fractal phenomenon'? How  $W$  acts on  $E(\Phi)$ ?
- Link with hyperbolic geometry (hyperbolic reflection groups) and with Apollonian gasket (Kleinian groups) – story in CH, JP-Préaux and V. Ripoll (2013)

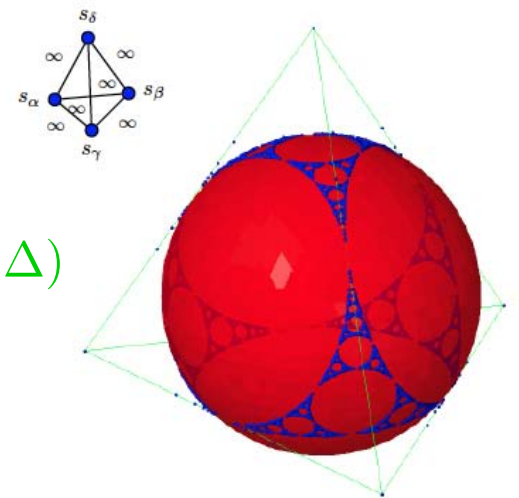


dim 3



dim 4

$\text{conv}(\Delta)$





# A geometric action on $E(\Phi) = \text{Acc}(\widehat{\Phi})$

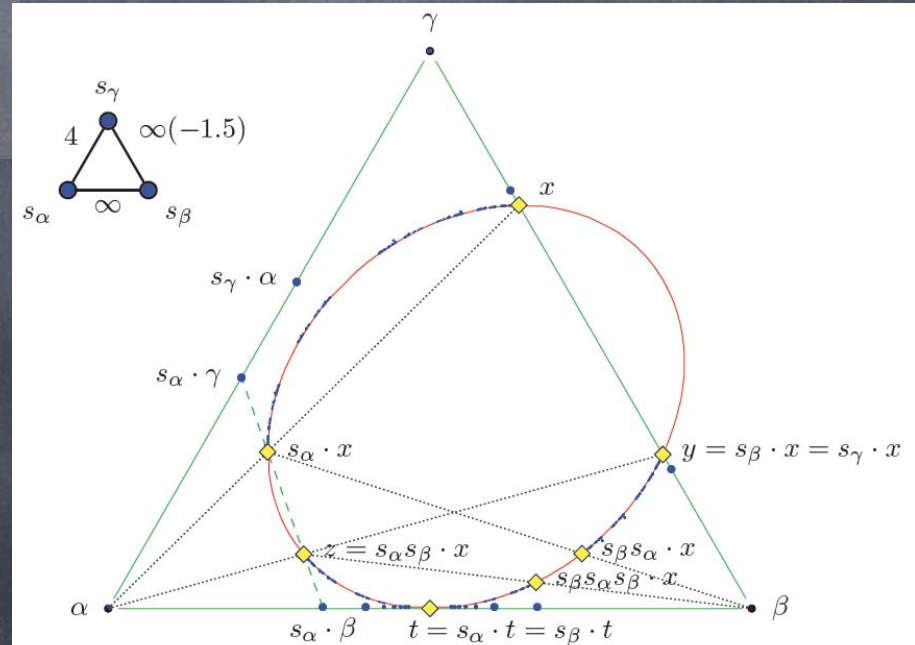
Extending the 'barycentric' action  $W \cdot \widehat{\Phi}$

👁 **New action:**  $w \cdot v = \widehat{w(v)}$  on the set  $\widehat{\Phi} \sqcup E$  given on  $E$

by:  $\widehat{Q} \cap L(\alpha, x) = \{x, s_\alpha \cdot x\}$

**Theorem** (Dyer, CH, Ripoll 2013)

Action on  $E$  faithful if irreducible  
not affine nor finite of rank  $> 2$ .



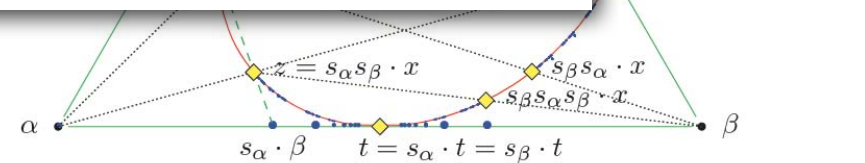
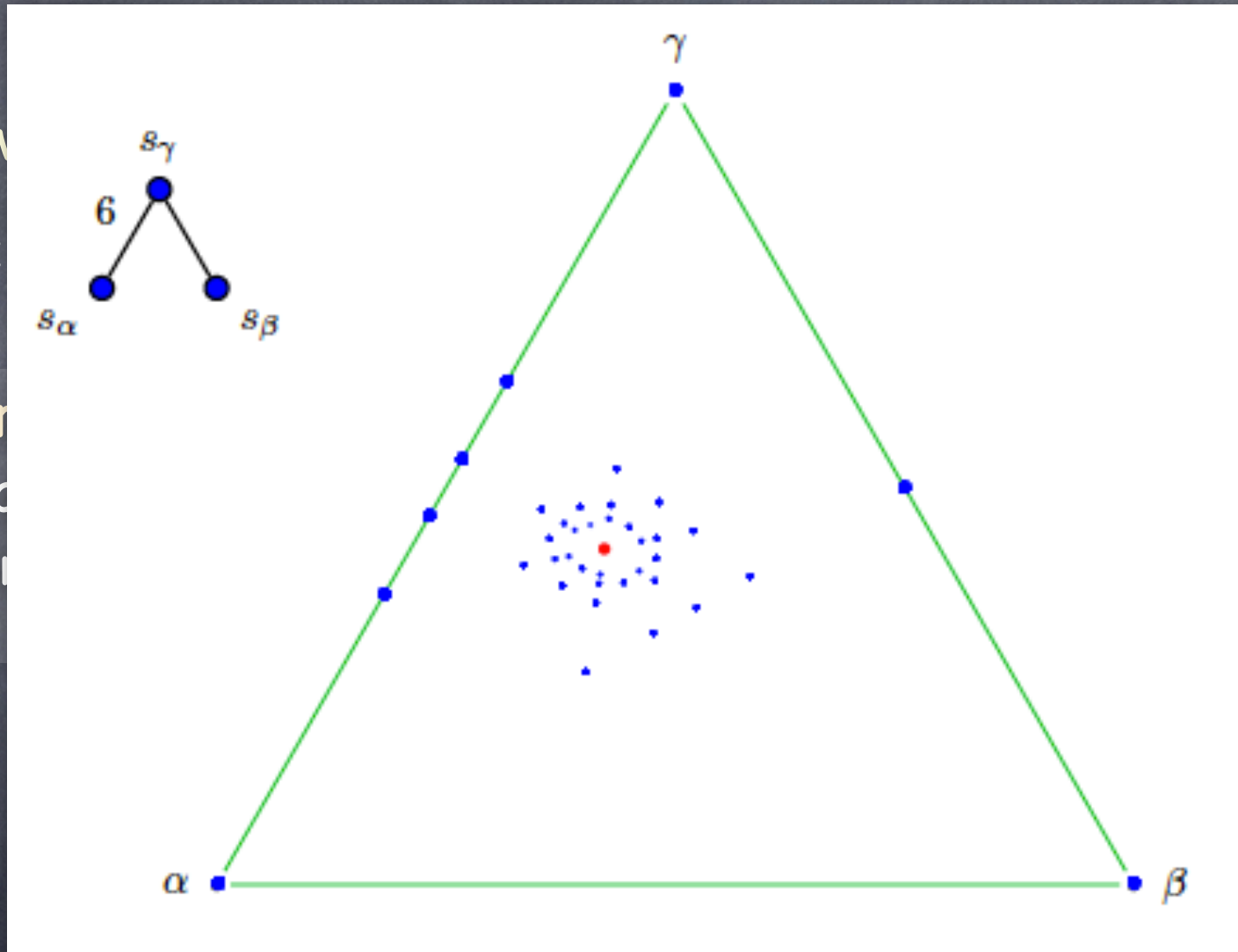


# A geometric action on $E(\Phi) = \text{Acc}(\hat{\Phi})$

New  
by:

Theorem  
Action of  
not affi

on  $E$





# A geometric action on $E(\Phi) = \text{Acc}(\widehat{\Phi})$

Remark:  $V_1$  is not stable under  $W$ .

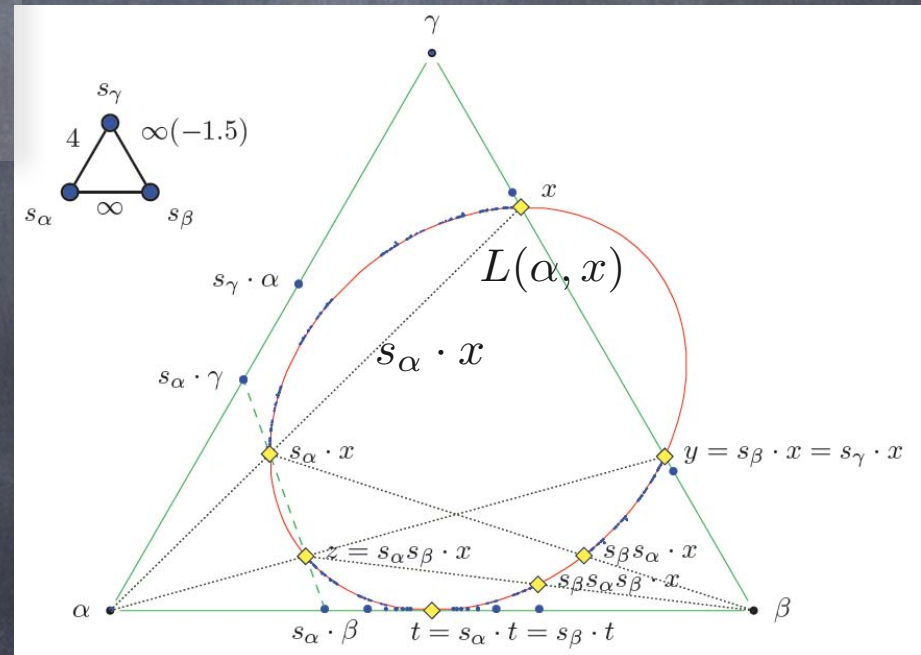
👁 **New action:**  $w \cdot v = \widehat{w(v)}$  on the set  $\widehat{\Phi} \sqcup E$  given on  $E$

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**Theorem** (Dyer, CH, Ripoll 2013)

Action on  $E$  faithful if irreducible  
not affine nor finite of rank  $> 2$ .

**Corollary:** to build  $E$  ...



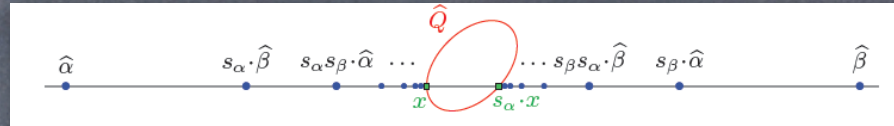


# Remarkable dense subsets of $E(\Phi) = \text{Acc}(\widehat{\Phi})$

**Dihedral reflection subgroups:**  $W' = \langle s_\rho, s_\gamma \rangle, \rho, \gamma \in \Phi^+$

**Associated root system:**  $\Phi' = W'(\{\rho, \gamma\})$

**Observation:**  $E(\Phi') = \widehat{Q} \cap L(\widehat{\rho}, \widehat{\gamma})$



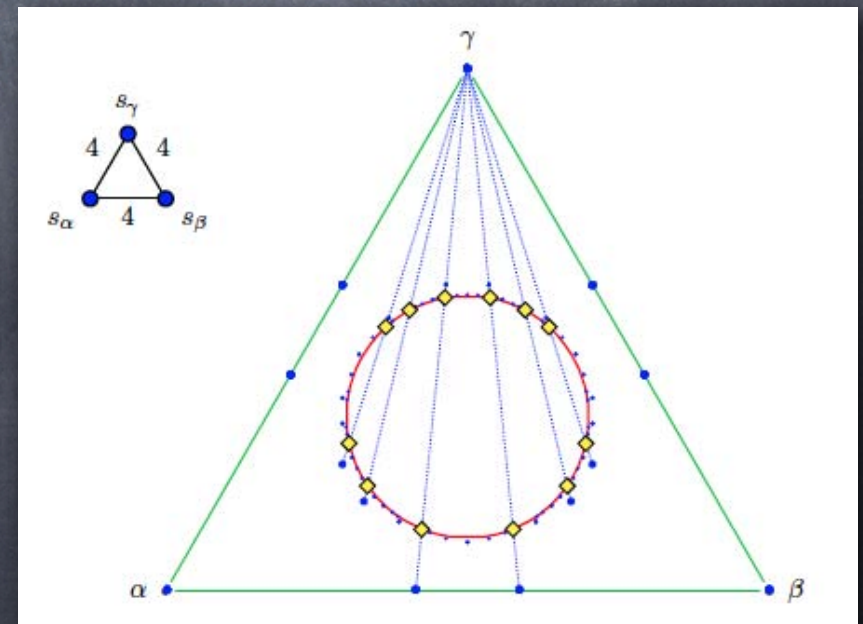
Limits of roots of dihedral reflection subgroups:

•  $E_2 = W \cdot E_2^\circ$  where

$$E_2^\circ := \bigcup_{\substack{\alpha \in \Delta \\ \rho \in \Phi^+}} L(\alpha, \widehat{\rho}) \cap \widehat{Q}$$

**Theorem** (CH, Labbé, Ripoll 2012)

The sets  $E_2$  and  $E_2^\circ$  are dense  
in  $E(\Phi)$ .





# The action on $E$ is minimal

## Theorem (Dyer, CH, Ripoll 2013)

The closure of  $W \cdot x$  is dense in  $E(\Phi)$  for  $x \in E(\Phi)$

## Theorem (Dyer, CH, Ripoll, 2013)

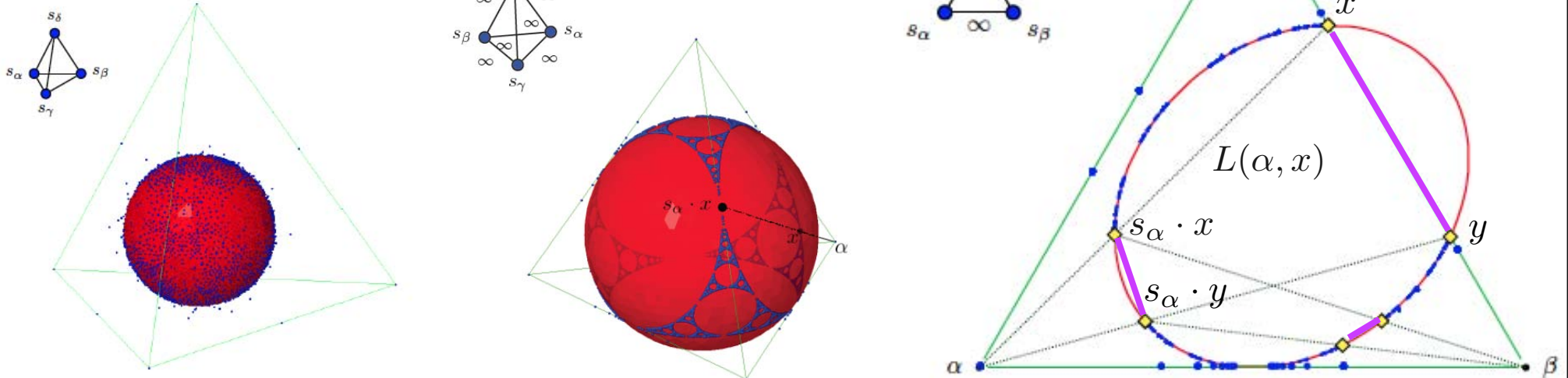
$$E = \hat{Q} \iff \hat{Q} \subseteq \text{conv}(\Delta)$$

Moreover, in this case,

$$\text{sgn}(B) = (n, 1, 0)$$

## Corollary (Dyer, CH, Ripoll, 2013)

## A first fractal Phenomenon.

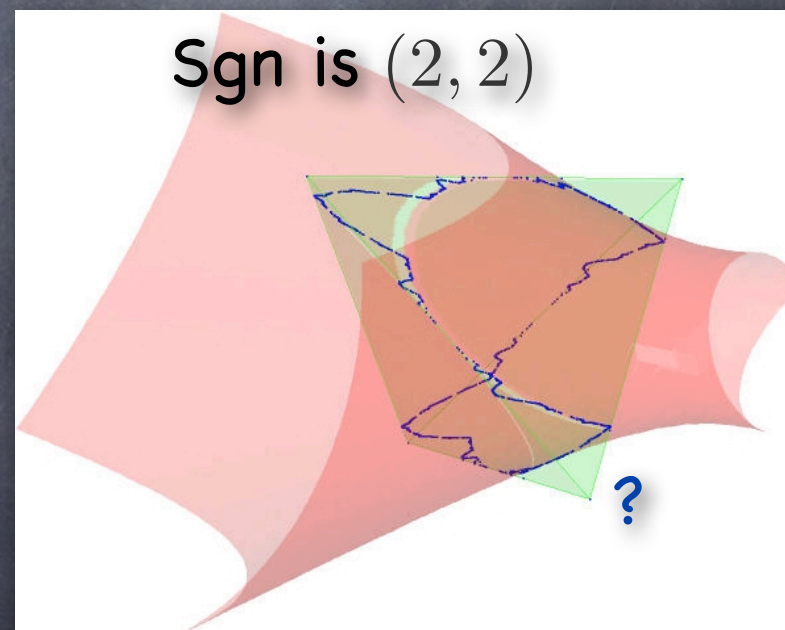
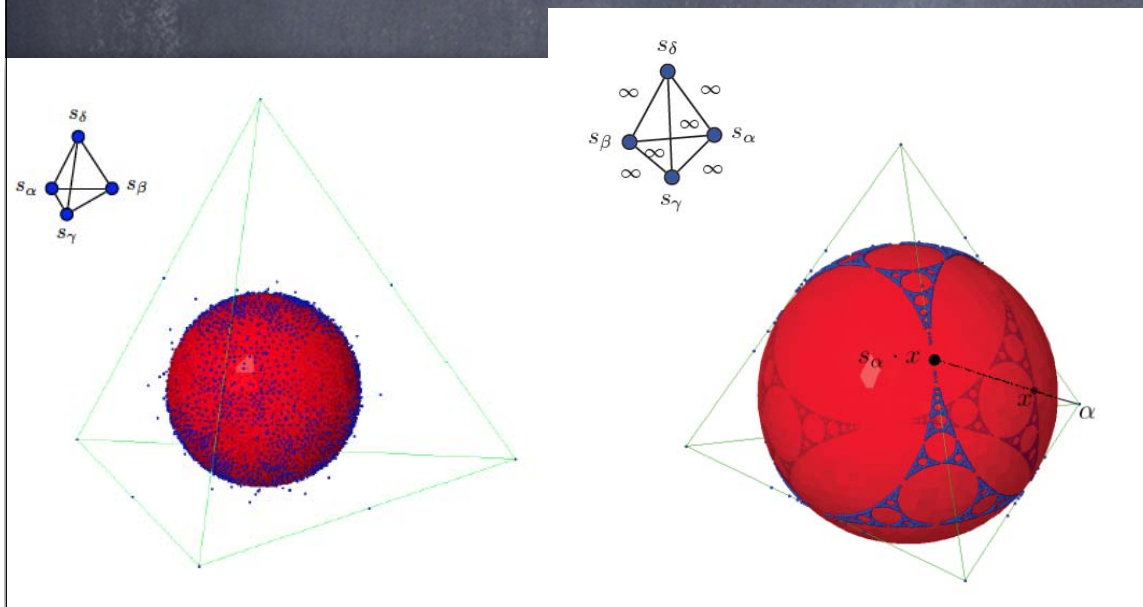




# A second fractal phenomenon

**Theorem** (Dyer, CH, Ripoll 2013) For irreducible root of signature  $(n, 1, 0)$  we have:  $E = \text{conv}(E) \cap Q$

**Problem (second fractal phenomenon):** is it true for other indefinite types?





# Imaginary cone and tiling of $\text{conv}(E)$

**Proposition** (Dyer, CH, Ripoll 2013). The action of  $W$  on  $E$  extends to an action of  $W$  on  $\text{conv}(E)$ . So  $W$  acts on  $\widehat{\Phi} \sqcup \text{conv}(E)$

🌀 **Imaginary convex body**  $\mathcal{I}$  is the  $W$ -orbit of the polytope

$$K = \{v \in \text{conv}(\Delta) \mid B(v, \alpha) \leq 0, \forall \alpha \in \Delta\}$$

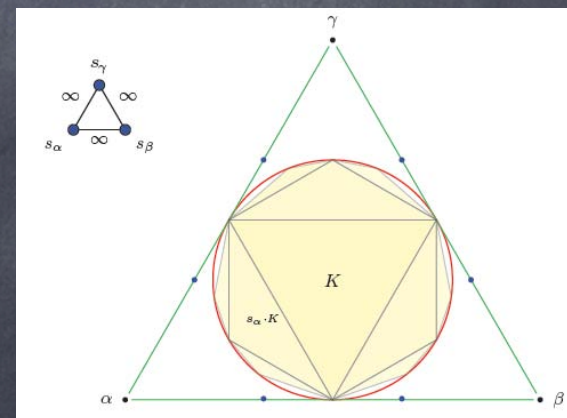
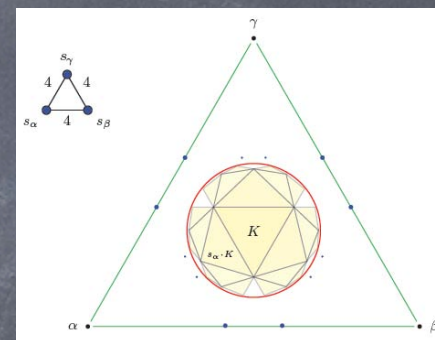
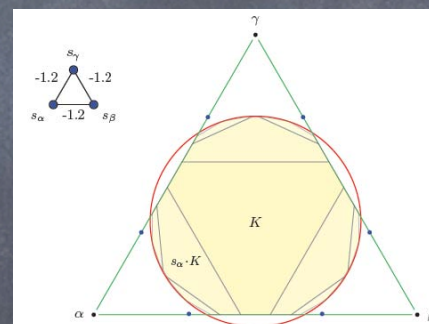
**Theorem** (Dyer, 2012).  $\overline{\mathcal{I}} = \text{conv}(E)$

**Theorem** (Dyer, CH, Ripoll 2013).

$$E \subseteq \overline{W \cdot z}, \quad \forall z \in \overline{\mathcal{I}}$$

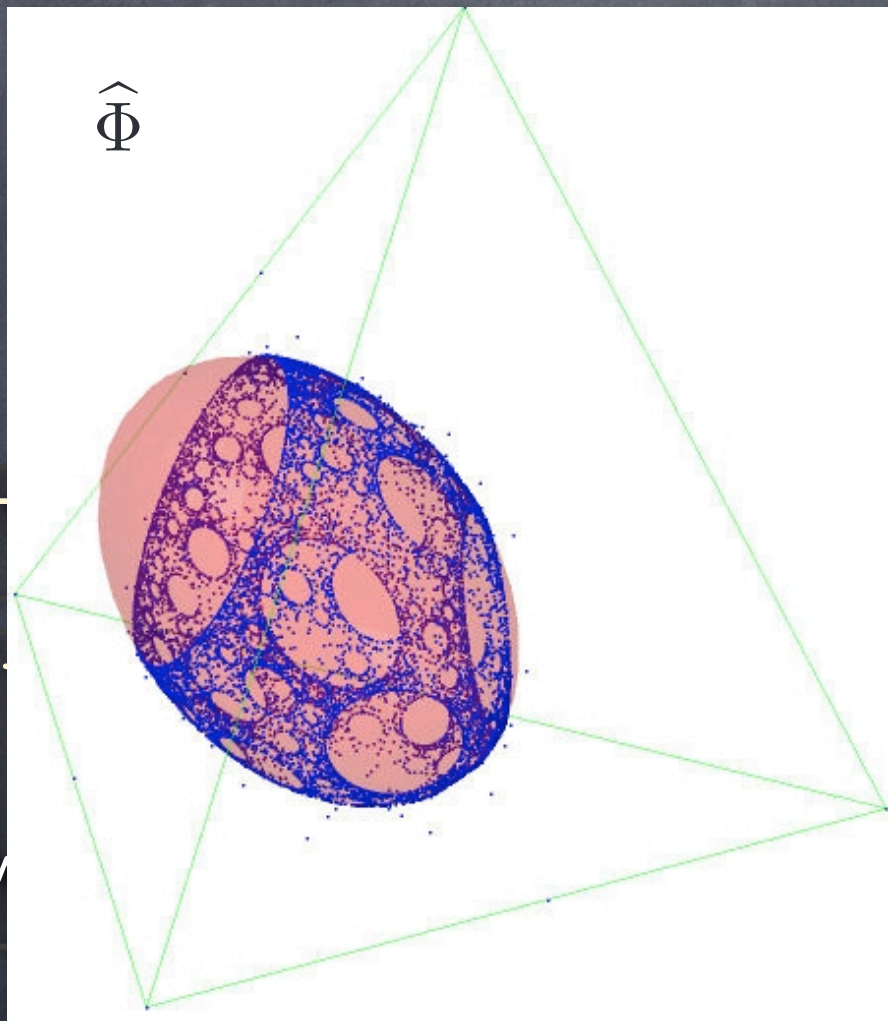
with equality for  $\text{sgn}(n, 1, 0)$

**Problem:** prove equality in general!

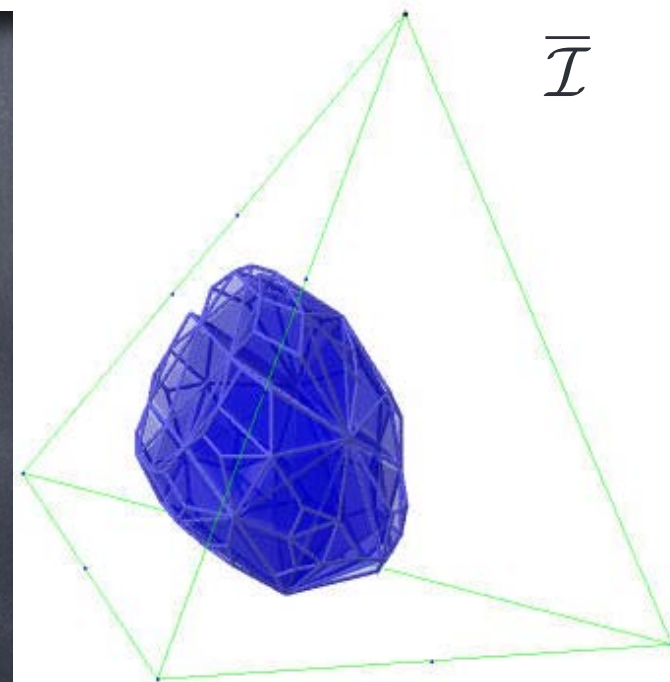
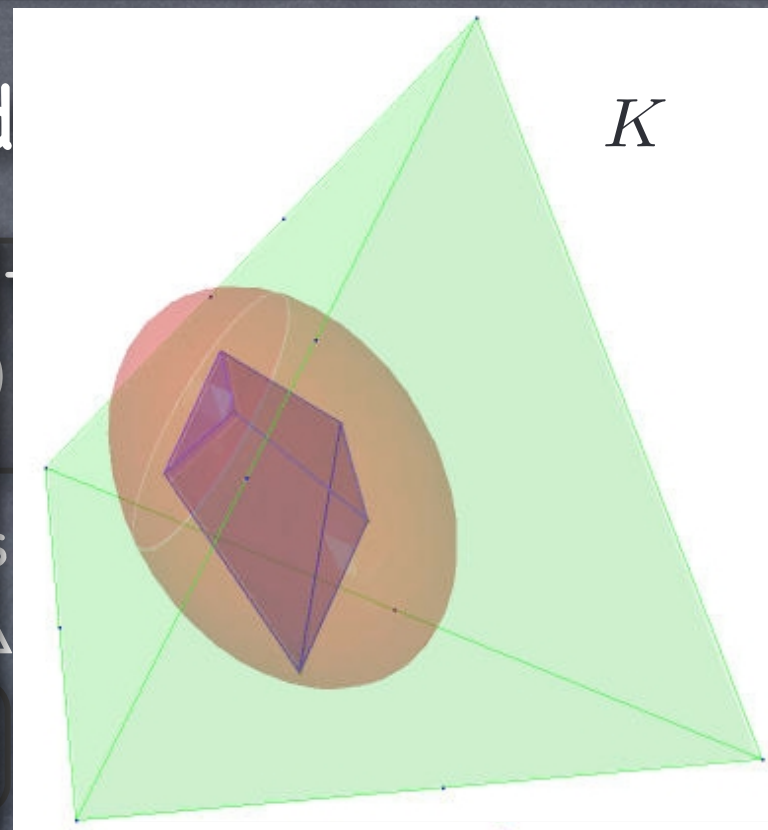




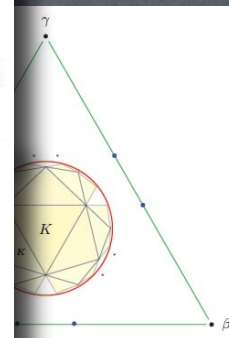
# Imaginary cone and



Problem: prove equality in general!



nds  
( $E$ )  
ope





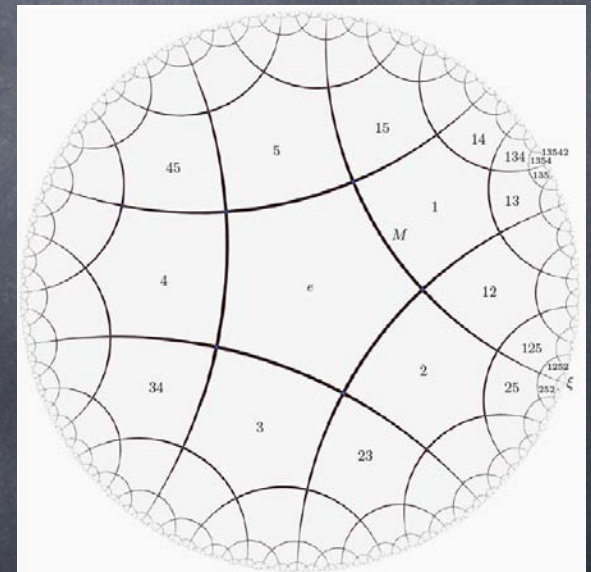
# A step toward biclosed sets: Infinite words, their inversion sets and limit weak order

- Thomas Lam & Anne Thomas: in “Infinite Reduced Words and the Tits Boundary of a Coxeter Group”
  - Limit weak order on infinite words (modulo braid relations) as the topology of the visual boundary of the Davis complex

**Proposition.** The inversion sets of infinite words are biclosed.

$$N(s_\alpha s_\beta s_\alpha \dots) = \{\alpha, s_\alpha(\beta), s_\alpha s_\beta(\alpha), \dots\}$$

**Work in progress** (CH 2013). The imaginary convex body is a geometric realization of the Davis complex and  $E$  is the visual boundary. Biclosed and their boundary!



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# A step toward biclosed sets: Infinite words, their inversion sets and limit weak order

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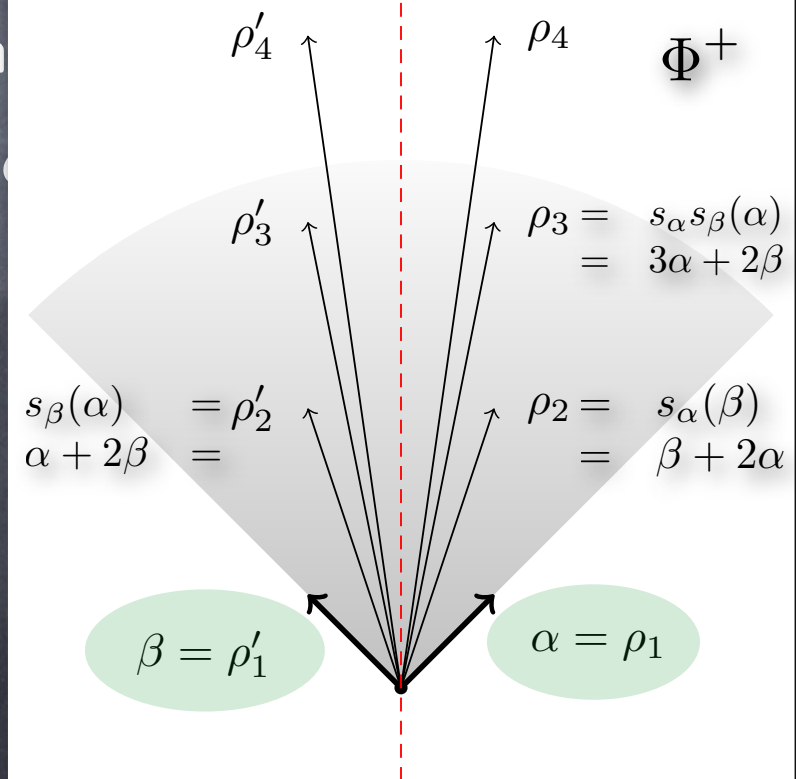
- Limit weak order on infinite words (motivated by the topology of the visual boundary)

**Proposition.** The inversion sets of infinite words are biclosed.

$$N(s_\alpha s_\beta s_\alpha \dots) = \{\alpha, s_\alpha(\beta), s_\alpha s_\beta(\alpha), \dots\}$$

Work in progress (CH 2013). The imaginary convex body is a geometric realization of the Davis complex and  $E$  is the visual boundary. Biclosed and their boundary!

$$Q = \{v \in V \mid B(v, v) = 0\}$$



$$(a) B(\alpha, \beta) = -1$$

$$s_\alpha(v) = v - 2B(v, \alpha)\alpha.$$



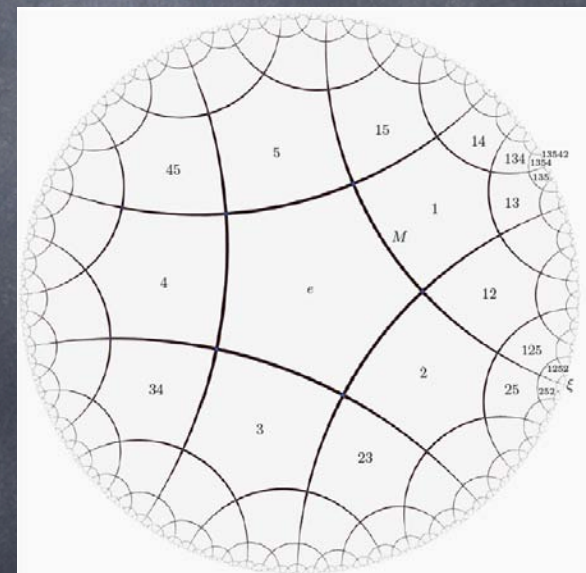
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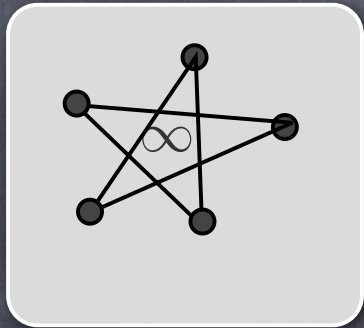
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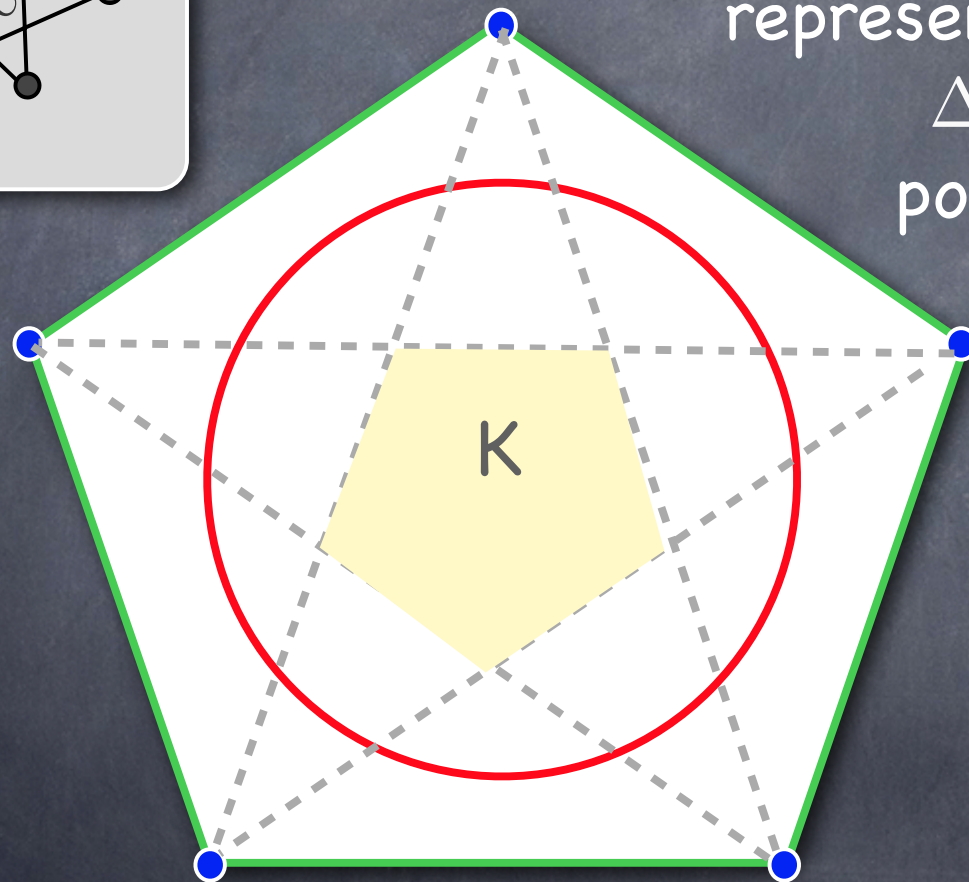


# A step toward biclosed sets: Infinite words, their inversion sets and limit weak order



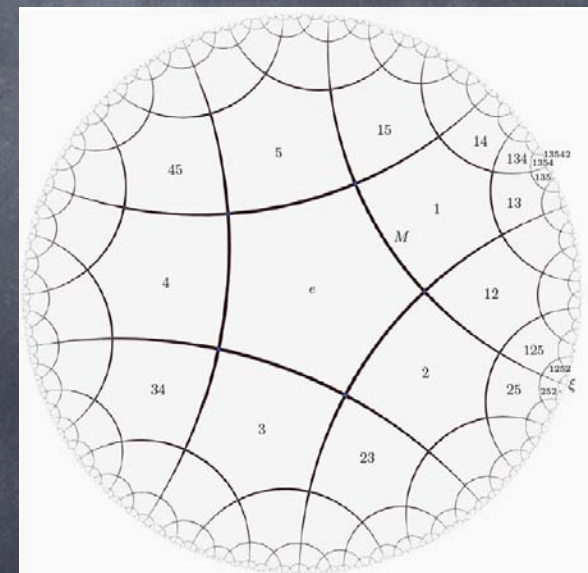
Here a rank 5 Coxeter group is represented in dim 3:

$\Delta$  is not a basis but is positively independent.



Roots and imaginary convex body model

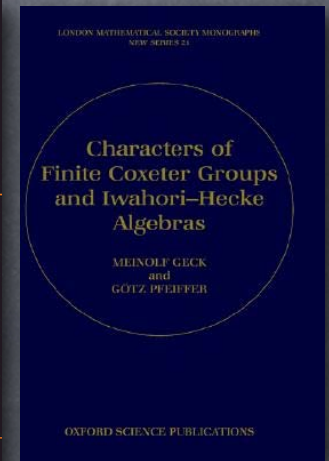
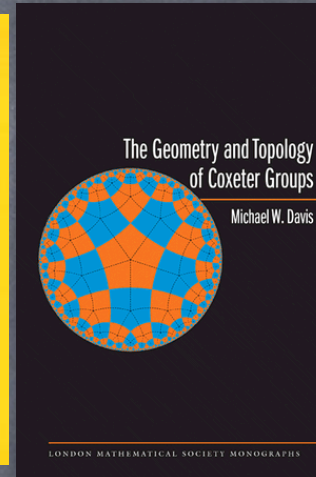
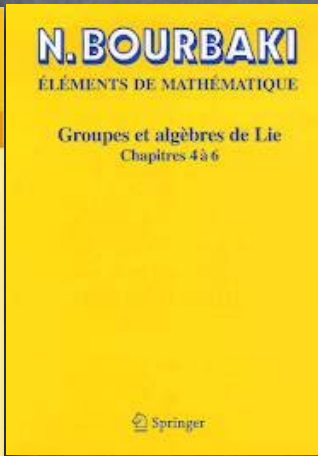
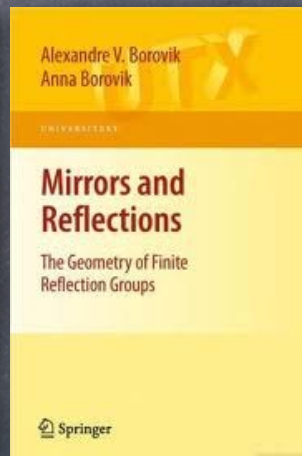
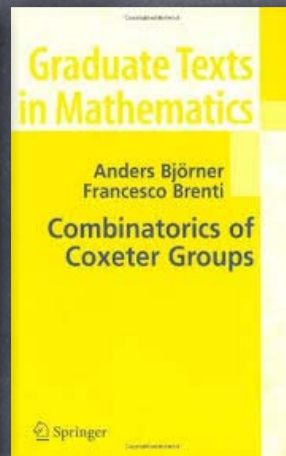
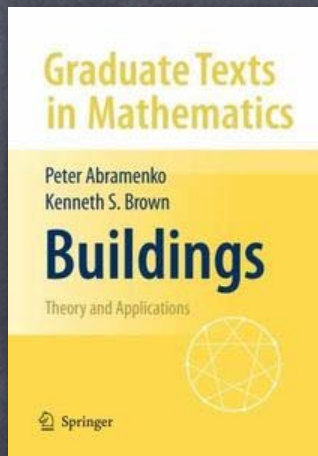
Ball model



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# Selected bibliography and other readings



And articles already cited + from

- Brigitte Brink, Bill Casselman, Fokko du Cloux, Bob Howlett, Xiang Fu (regarding automaton and comb.)
- Matthew Dyer (imaginary cones, weak order(s))
- CH & coauthors (Matthew Dyer, Jean-Philippe Labbé, Jean-Philippe Préaux, Vivien Ripoll). A good start for limit of roots and imaginary convex bodies is the survey of the case of Lorentzian spaces (CH, Ripoll, Préaux)
- Paolo Pappi and Ken Ito (limit weak order)
- ...

