M2 BIM – STRUCT – Lecture 1 Folding RNA *in silico*

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Outline

Introduction

Dynamic programming 101 Why RNA? RNA folding RNA Structure(s) Some representations of RNA structure

Some flavours of folding prediction

Thermodynamics vs Kinetics Dynamic programming: Reminder

Free-energy minimization Nussinov-style RNA folding

... or how to make a million bucks by giving change parsimoniously!!

Problem: You have access to unlimited amount of **1**, **20** and **50** cents coins. A client prefers to travel light, i.e. to minimize the #coins. How to give N cents back in change without losing a customer?

Strategy #1: Start with heaviest coins, and then complete/fill-up with coins of decreasing value.

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Strategy #1: Start with heaviest coins, and then complete/fill-up with coins of decreasing value.



Problem *a priori* (?!) non-solvable using such a *greedy* approach, as a (simpler) problem is already NP-complete (thus Efficient solution \Rightarrow 1M\$).

Strategy #2: Brute force enumeration \rightarrow #Coins^N (Ouch!)

Strategy #3: The following recurrence gives the minimal number of coins:

$$Min \# Coins(N) = Min \begin{cases} \bigcirc & \rightarrow & 1 + Min \# Coins(N-1) \\ \bigcirc & \rightarrow & 1 + Min \# Coins(N-20) \\ \bigcirc & \rightarrow & 1 + Min \# Coins(N-50) \end{cases}$$

With some memory (N intermediate computations), the minimum number of coins can be obtained after $N \times \#$ Coins operations. An actual set of coins can be reconstructing by tracing back the choices performed at each stage, leading to the minimum.

Remark: We still haven't won the million, as *N* has exponential value compared to the length of its encoding, so the algorithm does not qualify as *efficient* (i.e. polynomial).

Still, this approach is much more efficient than a brute-force enumeration: \Rightarrow Dynamic programming.

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Targeting system for DNA Editing CRISPR therapies Sickle-cell anemia, β-thalassamia, Leber congenital amaurosis (LCA), cancers...



Hendel et al, 2015; Agrotis & Ketteler, 2015







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Sensor of metabolites Riboswitches







Hendel et al. 2015: Agrotis & Ketteler, 2015



Riboswitches

Quantitative expression Transcriptomic signatures Cancer diagnosis/prognosis/relapse...



[NGuyen et al, 2021]







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Sensor of metabolites Riboswitches

RiboNucleic Acids (RNAs)

 Quantitative expression

 Transcriptomic signatures

 Cancer diagnosis/prognosis/relapse...

 A TOON Deservey dataset:

 IND (N=50)

 TOS (N=50)

[NGuyen et al, 2021]



(a)

Non-coding mutations IncRNAs, mIRNAs, structure-associated (RiboSnitches) β-thalassemia, duchenne muscular dystrophy, Cystlc fibrosis, Rett syndrome...

Solem et al. 2015



Targeting system for DNA Editing

Hendel et al, 2015; Agrotis & Ketteler, 2015





Quantitative expression Transcriptomic signatures Cancer diagnosis/prognosis/relapse...

TCGA Discovery dataset:
 NO (N +50)
 TES (N +50)
 Dog fold change of mean normalized expression

[NGuyen et al, 2021]



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Encodes proteins mRNA Vaccines COVID-19, Malaria (Zika, CMV, Cancers?)

Genomic material for Human pathogens HIV-1, SARS-CoV 2, HCoVs, MERS

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Sensor of metabolites Riboswitches

RiboNucleic Acids (RNAs)

Quantitative expression Transcriptomic signatures Cancer diagnosis/prognosis/relapse...

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Non-coding mutations IncRNAs, mIRNAs, structure-associated (RiboSnitches) β-thalassemia, duchenne muscular dystrophy, Cystlc fibrosis, Rett syndrome...

Regulation of gene expression

RNAi therapies (FDA approved) Primary hyperoxaluria type 1 (PH1), Hereditary transthyretin amyloidosis (ATTRv), Acute hepatic porphyria (AHP)



Hendel et al. 2015: Agrotis & Ketteler, 2015

Encyclopaedia Brittannica, Inc 2013



Encodes proteins mRNA Vaccines COVID-19, Malaria (Zika, CMV, Cancers?)

Genomic material for Human pathogens HIV-1, SARS-CoV 2, HCoVs, MERS






RNA world: Resolving the chicken vs egg paradox at the origin of life...



A gene big enough to specify an enzyme would be too big to replicate accurately without the aid of an enzyme of the very kind that it is trying to specify. So the system *apparently cannot get started*.

[...] This is the RNA World. To see how plausible it is, we need to look at why proteins are good at being enzymes but bad at being replicators; at why DNA is good at replicating but bad at being an enzyme; and finally why RNA might just be good enough at both roles to break out of the Catch-22.

R. Dawkins. The Ancestor's Tale: A Pilgrimage to the Dawn of Evolution

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RNA folding

RNA is single-stranded and folds on itself, establishing complex 3D structures that are essential to its function(s).

RNA structures are stabilized by base-pairs, each mediated by hydrogen bonds.



RNA Structure(s)

Three¹ levels of representation:

UUAGGCGGCCACAGC GGUGGGUUGCCUCC CGUACCCAUCCCGAA CACGAAGAUAAGCC CACGAAGAUACUGGAGUCCGGG GAGUACUGGAGUCCGGG CGAGCCUCUGGGAAA CCCGGUUCGCCGCCA CC

Primary structure

Secondary structure

Seales and

Tertiary structure

Source: 5s rRNA (PDB 1K73:B)

¹Well, mostly...

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UUAGGCGGCCACAGC GGUGGGUUGCCUCC CGUACCCAUCCCGAA CACGAAGAUAAGCC CACGAAGAUAAGCC CACCAGCGUUCCGGG GAGUACUGGAGUGG CGAGCCUCUGGAAA CCCGGUUCGCCGCCA CC



Secondary⁺ structure

"Const

Tertiary structure

Source: 5s rRNA (PDB 1K73:B)

¹Well, mostly...

Ignored by secondary structure

Non-canonical base-pairs

Any base-pair other than {(A-U), (C-G), (G-U)} Or interacting on non-standard edge (\neq WC/WC-Cis) [LW01].



Canonique CG pair(WC/WC-Cis)

Non-canonique CG pair (Sugar/WC-Trans)

Pseudoknots (PKs)



Pseudoknoted structure of group I ribozyme (PDBID: 1Y0Q:A)

Considering PKs may lead to better predictions, but:

- Some PK conformations are simply unfeasible;
- Folding in silico with general pseudoknots is NP-complete [LP00];

Still, folding on restricted classes of conformations seems promising [CDR+04].



Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*

Supporting intuitions

Different representations

Common combinatorial structure

* Additional steric constraints



Outer-planar graphs Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*



Dot plots Adjacency matrices*

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Dot plots Non-crossing arc diagrams* Adjacency matrices*

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Dot plots Non-crossing arc diagrams* Adjacency matrices*

Motzkin words*

Supporting intuitions

Different representations Common combinatorial structure

*Additional steric constraints

Adjacency matrices*



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Various representations for a versatile biomolecule



Adjacency matrices*

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Thermodynamics aparté

At the nanoscopic scale, RNA structure *fluctuates* (\approx Markov process).



Convergence towards a stationary distribution at the Boltzmann equilibrium, where the probability of a conformation only depends on its free-energy. Corollary: Initial conformation does not matter.

Questions: For a given conformation space and free-energy model:

- A. Determine most stable (Minimum Free-Energy) structure at equilibrium;
- B. Compute average properties of Boltzmann ensemble;

Transcription: RNA synthesized, supposedly without structure²



²Except for co-transcriptional folding...

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²Except for co-transcriptional folding...

Transcription: RNA synthesized, supposedly without structure²



But most mRNAs are degrade before 7h (Org.: Souris [SSN+09]).

- A. Determine most stable (Minimum Free-Energy) structure at equilibrium;
- B. Compute average properties of Boltzmann ensemble;
- C. Determine most likely structure at finite time T.

(c.f. H. Isambert through simulation, NP-complete deterministically [MTSC09])

²Except for co-transcriptional folding...

Dynamic programming: General principle

Dynamic programming = General optimization technique.

Prerequisite: Optimal solution for problem P can be derived from solutions to strict sub-problems.

Bioinformatics :

Discrete solution space (alignments, structures...)

- + Additively-inherited objective function (cost, log-odd score, energy...)
- \Rightarrow Efficient dynamic programming scheme

Example: Local Alignment(Smith/Waterman)



Algorithmic details

Dynamic programming scheme defines a space of (sub)problems and a recurrence that relates the score of a problem to that of smaller problems.

Given a scheme, two steps :

- Matrix filling: Computation and tabulation of best scores (Computed from smaller problems to larger ones).
- Traceback: Reconstruct best solution from contributing subproblems.

Complexity of algorithm depends on:

- Cardinality of sub-problem space
- Number of alternatives considers at each step (#Terms in recurrence)

Smith&Waterman example:

- ► $j: 1 \rightarrow m + 1 \Rightarrow \Theta(m)$
- 3 operations at each step
- $\Rightarrow \Theta(m.n)$ time/memory

$$W(i,0) = 0$$

$$W(0,j) = 0$$

$$W(i,j) = \max \begin{cases} W(i-1,j-1) + m_{i,j} \\ W(i-1,j) + p_i \\ W(i,j-1) + p_d \end{cases}$$

Example: Local alignment of AGCACACA and ACACACTA

			А	С	А	С	А	С	Т	Α	
W(i, 0) = 0 W(0, j) = 0		0	0	0	0	0	0	0	0	0	
$W(i,j) = \max \begin{cases} W(i-1,j-1) + m_{i,j} \\ W(i-1,j) + p_i \\ W(i,j-1) + p_d \end{cases}$	A	0									
	G	0									
	с	0									
	А	0									
	с	0									
	A	0									
	с	0									
	A	0									

Example: Local alignment of AGCACACA and ACACACTA

M((+ 0) 0		А	С	А	С	А	С	Т	Α	
W(0,j) = 0 W(0,j) = 0		0 0	0	0	0	0	0	0	0	
$W(i,j) = \max \begin{cases} W(i-1,j-1) + m_{i,j} \\ W(i-1,j) + p_i \\ W(i,j-1) + p_d \end{cases}$	A	$0 \rightarrow 2$								
	G	0								
	с	0								
	А	0								
	с	0								
	А	0								
	с	0								
	A	0								

Example: Local alignment of AGCACACA and ACACACTA

W(i, 0) = 0 W(0, i) = 0		0	A 0	C 0	A 0	C 0	A 0	C 0	Т 0	A 0
$W(i, i) = \max \begin{cases} W(i-1, j-1) + m_{i,j} \\ W(i-1, j) + p_{i} \end{cases}$	А	0	2 -	∖ ↓ → 1						
$W(i,j) = \max \begin{cases} W(i-1,j) + p_i \\ W(i,j-1) + p_d \end{cases}$	G	0								
	с	0								
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	с	0								
	A	0								
	с	0								
	A	0								

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$$\begin{split} & \mathcal{W}(i,0) = 0 \\ & \mathcal{W}(0,j) = 0 \\ & \mathcal{W}(i,j) = \max \left\{ \begin{array}{cccccc} \mathcal{W}(i-1,j-1) + m_{l,j} \\ \mathcal{W}(i-1,j) + p_i \\ \mathcal{W}(i,j-1) + p_d \end{array} \right| & \mathcal{K}(i-1,j) + p_i \\ & \mathcal{K}(i,j-1) + p_d \\ & \mathcal{K}(i,j-1) + p_d \end{array} \right. \\ & \mathcal{K}(i,j-1) = \mathcal{K}(i,j-1) \\ & \mathcal{K}(i,j-1) + \mathcal{K}(i,j-1) + \mathcal{K}(i,j-1) \\ & \mathcal{K}(i$$

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Complete example

Example: Local alignment of AGCACACA and ACACACTA

$$\begin{array}{rcl} W(i,0) &=& 0 \\ W(0,j) &=& 0 \\ W(i,j) &=& \max \left\{ \begin{array}{rrrr} W(i-1,j-1) + m_{i,j} \\ W(i-1,j) + p_i \\ W(i,j-1) + p_d \end{array} \right. & \left. \begin{array}{cccccc} A & C & T & A \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ A & 0 & 2 \rightarrow 1 & 2 \rightarrow 1 \rightarrow 2 \rightarrow 1 \rightarrow 0 \rightarrow 2 \\ G & 0 \\ C & 0 \\ A & 0 \end{array} \right. \\ \end{array}$$

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Complete example

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$$A = C = A = C = A = C = T = A$$

$$0 = 0 = 0 = 0 = 0$$

$$A = 0 = 2 \Rightarrow 1 = 2 \Rightarrow 1$$

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Properties of DP schemes

Necessary properties:

► Correctness: ∀ sub-problem, the computed value must indeed maximize the objective function . Proofs usually inductive, and quite technical, but very systematic.

Desirable properties of DP schemes:

- Completeness of space of solutions generated by decomposition. Algorithmic tricks, by *cutting branches*, may violate this property.
- Unambiguity: Each solution is generated at most once.
- \Rightarrow Under these properties, one can enumerate solution space.

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Folding by minimizing free-energy

Problem A: Determine Minimum Free-Energy structure (MFE).

Ab initio folding prediction =

Predict RNA structure from its sequence ω only.



- Conformations: Set S_{ω} of secondary structures compatible (w.r.t. base-pairing constraints) with primary structure ω .
- Free-Energy: Function $E_{\omega,S}$ (KCal.mol⁻¹), additive on motifs occurring in any sequence/conformation couple (ω, S).
- Native structure: Functional conformation of the biomolecule.

Remarks:

- Not necessarily unique (Kinetics, or bi-stable structures);
- In presence of PKs → Ambiguous: Which is the native conformation?

Nussinov/Jacobson model

Nussinov/Jacobson energy model (NJ)

Base-pair maximization (with a twist):

- Additive model on independently contributing base-pairs;
- Canonical base-pairs only: Watson/Crick (A/U,C/G) and Wobble (G/U)

$$\Rightarrow \textit{E}_{\omega,\textit{S}} = -\#\textit{Paires}(\textit{S})$$

Folding in NJ model \Leftrightarrow Base-pair (weight) maximization

Example:



Variant: Weight each pair with -#Hydrogen bonds $\Delta G(G \equiv C) = -3$ $\Delta G(A = U) = -2$ $\Delta G(G - U) = -1$

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Nussinov/Jacobson DP scheme



$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \begin{cases} j & i \text{ unpaired} \\ \min_{k=i+\theta+1} \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{cases}$$



$$N_{i,t} = 0, \quad \forall t \in [i, i+\theta]$$

$$N_{i,j} = \min \begin{cases} j & N_{i+1,j} & i \text{ unpaired} \\ \min_{k=i+\theta+1} \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{cases}$$

Correctness. Goal = Show that MFE over interval [i, j] is indeed found in $N_{i,j}$ after completing the computation. Proceed by induction:

- Assume that property holds for any [i', j'] such that j' i' < n.
- Consider [i, j], j i = n. Let MFE_{*i*,*j*} := Base-pairs of best struct. on [i, j]. Then first position *i* in MFE_{*i*,*j*} = is either:
 - ► Unpaired: MFE_{*i*,*j*} = MFE_{*i*+1,*j*} → free-energy = $N_{i+1,j}$ → free-energy = $N_{i+1,j}$

(Indeed, any BP between [i + 1, k - 1] and [k + 1, j] would cross (i, k))



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► Unpaired:
$$MFE_{i,j} = MFE_{i+1,j}$$
 \rightarrow free-energy = $N_{i+1,j}$
► Paired to k: $MFE_{i,j} = \{(i,k)\} \cup MFE_{i+1,k-1} \cup MFE_{k+1,j}$.
(Indeed, any BP between $[i+1, k-1]$ and $[k+1, j]$ would cross (i, k))
 \rightarrow free-energy = $\Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j}$.



$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \begin{cases} j & i \text{ unpaired} \\ \min_{k=i+\theta+1} \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{cases}$$

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	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•		•
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
С			_		_					4	$\langle \rangle$	e^{θ}				0	0	0
G	i			j	_	i i+1			j	Ťi			k	j			0	0
A	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
А				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
А						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
А													0	0	0	0	0	0
G														0	0	0	0	0
A											_	_			0	0	0	0
С	-				=		~~~~			4	$\langle \rangle$	$\theta $				0	0	0
G	i			j		i i+1			j	' i			k	j			0	0
А	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
А				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
А						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
А													0	0	0	0	0	0
G														0	0	0	0	0
A											-	_			0	0	0	0
С			~		=						$\langle \rangle$	θ				0	0	0
G	i			j		i i+1			j	' i			k	j			0	0
А	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•		
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A											_	_			0	0	0	0
С					=		~~~~			+ 4	<u> </u>	$\theta $				0	0	0
G	i			j		i i+1			j	i i			k	j			0	0
A	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	(•	•	•	•	•	•	•	•	•	•	•	•	•	•	•)	
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A											_	_			0	0	0	0
С			_		_					4	< ≥	e^{θ}				0	0	0
G	i			j	_	i i+1			j	ī			k	j			0	0
A	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	(•		•	•	•	•	•	•	•	•	•	•	•	•	•)	•
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A											_	_			0	0	0	0
С					=					+ 4	\leq	e^{θ}				0	0	0
G	i			j	-	i i+1			j	Ťi			k	j			0	0
A	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	(•	•	•	•	•	•	•	•	•	•	•	•	•	•	•)	
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
А				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
А						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A											_	_			0	0	0	0
С					_		~~~~			_ <i>4</i>	2	θ				0	0	0
G	i			j		i i+1			j	' ī			k	j			0	0
А	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	(•		•	•	•	•	•	•	•	•	•	•	•	•)	
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
А				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
А						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A											-	_			0	0	0	0
С					_					4	< ≥	θ				0	0	0
G	i			j	_	i i+1			j	ī			k	j			0	0
А	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	(•	•	•	•	•	•	•	•	•	•	•	•	•	•	•)	
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A											_	_			0	0	0	0
С					_					_ 4	\leq	θ				0	0	0
G	i			j	_	i i+1			j	ī			k	j			0	0
A	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	(•	•	•	•	•	•	•	•	•	•	•	•	•	•	•)	
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
А				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
А						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
А													0	0	0	0	0	0
G														0	0	0	0	0
A											-	_			0	0	0	0
С			_		_					4	< ≥	θ				0	0	0
G	i			j	_	i i+1			j	ī			k	j			0	0
А	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	(•	•	•	•	•	•	•	•	•	•	•	•	•	•	•)	
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
А				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
А						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
А													0	0	0	0	0	0
G														0	0	0	0	0
A											-	_			0	0	0	0
С			_		_					4	< ≥	e^{θ}				0	0	0
G	i			j	-	i i+1			j	ī			k	j			0	0
А	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	(•	•	•	•	•	•	•	•	•	•		•	•	•	•)	
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
А				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
А						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
А													0	0	0	0	0	0
G														0	0	0	0	0
A											-	_			0	0	0	0
С			_		_					4	\geq	e^{θ}				0	0	0
G	i			j	_	i i+1			j	ī			k	j			0	0
А	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	(•	•	•	•	•	•	•	•	•	•	•	•	•	•	•)	•
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
А				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
А						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A											_				0	0	0	0
С			_		_					4	<u> </u>	e^{θ}				0	0	0
G	i			j	-	i i+1			j	Ťi			k	j			0	0
А	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	А	С	G	A
	((•	•	•	•	•	•	•	•	•	•	•	•))	
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
Α				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
А						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A											_	_			0	0	0	0
С					=		~~~~			+ 4	2	θ				0	0	0
G	i			j		i i+1			j	' ī			k	j			0	0
A	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	А	С	G	A
	((•	•	•	•	•	•	•	•	•	•	•	•	•))	
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
А				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
А						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
А													0	0	0	0	0	0
G														0	0	0	0	0
A											_	_			0	0	0	0
С					_					+ 4	\leq	e^{θ}				0	0	0
G	i			j	_	i i+1			j	Τ i			k	j			0	0
А	L																	0

	С	G	G	A	U	A	С	U	U	С	U	U	A	G	A	С	G	A
	((•	•	•	•	•	•	•	•	•	•	•	•))	
С	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
С							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
С										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A											-	_			0	0	0	0
С					=					+ 4	2	θ				0	0	0
G	i			j		i i+1			j	i i			k	j			0	0
A	L																	0
























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