## Algebras

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## I Algebras for an endofunctor

An algebra for an endofunctor  $F : \mathcal{C} \to \mathcal{C}$  is a pair (A, f) where A is an object of  $\mathcal{C}$  and  $f : FA \to A$  a morphism of  $\mathcal{C}$ . A morphism  $h : (A, f) \to (B, g)$  between two such algebras consists of a morphism  $h : A \to B$  such that

$$\begin{array}{ccc} FA & \xrightarrow{Fh} & FB \\ f \downarrow & & \downarrow^g \\ A & \xrightarrow{h} & B \end{array}$$

In the following, we mostly consider algebras in **Set**.

- 1. Define inductively the functions
  - length : 'a list -> int giving the length of a list,
  - map : ('a -> 'b) -> 'a list -> 'b list applying a function to all elements of a list,
  - double : 'a list -> 'a list which duplicates every successive element, for instance double [1;2;3] = [1;1;2;2;3;3].
- 2. Suppose given a type 'a ilist of infinite lists with elements of type 'a. Define coinductively
  - even : 'a ilist -> 'a ilist keeping elements of a list at even positions,
  - merge : 'a ilist -> 'a ilist taking alternatively elements from one
    of two lists.
- 3. We write  $S : \mathbb{N} \to \mathbb{N}$  for the successor function. Show that  $[0, S] : 1 + \mathbb{N} \to \mathbb{N}$  is an initial algebra for the endofunctor TX = 1 + X of **Set**.
- 4. Use this fact to define the function  $h : \mathbb{N} \to \mathbb{Q}$  such that  $h(n) = 2^{-n}$ .
- 5. Show that two initial algebras of an endofunctor are isomorphic (via morphisms of algebras).
- 6. Show that an initial algebra  $f: FA \to A$  of an endofunctor F is an isomorphism.
- 7. Solve the equation x = 1 + ax and develop the solution in power series.
- 8. Show that the set  $A^* = \biguplus_{n \in \mathbb{N}} A^n$ , which can be seen as the set of lists of elements of A, is an initial algebra for  $TX = 1 + A \times X$ .
- 9. Use this fact to define the length function  $\ell : A^* \to \mathbb{N}$  and the double function  $d : A^* \to A^*$ . Show that  $\ell \circ d(l) = 2\ell(l)$  for every  $l \in A^*$ .
- 10. Explain briefly how we could interpret simple inductive types of OCaml by using initial algebras.
- 11. What is the initial algebra for  $TX = 1 + X \times X$ ? For  $TX = X^*$ ? For  $TX = A \times X$ ?

## II Coalgebras for an endofunctor

A coalgebra for  $F : \mathcal{C} \to \mathcal{C}$  is a pair (A, f) with  $f : A \to FA$ . Morphisms are defined similarly as previously.

- 1. Show that the set  $A^{\mathbb{N}}$  of *streams* is a final coalgebra for the endofunctor  $TX = A \times X$ .
- 2. Use this to define,
  - given  $a \in A$ , the constant stream equal to a,
  - the function  $\mathbb{N} \to \mathbb{N}^{\mathbb{N}}$  which to *n* associates the stream  $(n, n+1, n+2, \ldots)$ ,
  - the function  $A^{\mathbb{N}} \times A^{\mathbb{N}} \to A^{\mathbb{N}}$  which merges two streams,
  - the function  $A^{\mathbb{N}} \to A^{\mathbb{N}}$  keeping even elements.
- 3. Show that final coalgebras are unique up to isomorphism and are isomorphisms.
- 4. Show that merge(even(u), odd(u)) = u for every  $u \in A^{\mathbb{N}}$ , where odd(l) = even(tail(l)).

A bisimulation on  $A^{\mathbb{N}}$  is a relation  $R \subseteq A^{\mathbb{N}} \times A^{\mathbb{N}}$  such that R(x :: u, x' :: u') implies x = x' and R(u, u'). The coinductive proof principle says that if R(u, u') for some bisimulation R then u = u'.

- 5. Assuming this principle, show again the result of previous question.
- 6. Show the coinductive proof principle (hint: show that R has a coalgebra structure).
- 7. Generalize the coinductive proof principle to an arbitrary endofunctor.
- 8. What is the final coalgebra of  $TX = 1 + A \times X$ ? of TX = 1 + X?
- 9. Show that automatas can be seen as coalgebras.

## References

[1] Bart Jacobs and Jan Rutten. An introduction to (co)algebra and (co)induction, page 38–99. Cambridge Tracts in Theoretical Computer Science. Cambridge University Press, 2011.