# Coproducts, pullbacks, monoids

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### 1 Coproducts

Notions in category theory can always be "dualized" in the following way.

- 1. Given a category C define the category  $C^{op}$  obtained by reversing the morphisms.
- A <u>cosomething</u> in a category  $\mathcal{C}$  is a something in  $\mathcal{C}^{\text{op}}$ .
  - 2. Show that **Set** is a cocartesian category, i.e. has coproducts and an initial object (an initial object is a coterminal object).
  - 3. Show that the usual categories are cocartesian : Set, Top, Rel, Vect, Cat.

#### 2 Pullbacks

Given two morphisms  $f : A \to C$  and  $g : B \to C$  with the same target, a *pullback* is given by an object D (sometimes abusively noted  $A \times_C B$ ) together with two morphisms  $p : D \to A$  and  $q : D \to B$  such that  $f \circ p = g \circ q$ , and for every pair of morphisms  $p' : D' \to A$  and  $q' : D' \to B$ (with the same source) such that  $f \circ p' = g \circ q'$ , there exists a unique morphism  $h : D' \to D$  such that  $p \circ h = p'$  and  $q \circ h = q'$ .



- 1. What is a pullback in the case where C is the terminal object?
- 2. What is a pullback in **Set**?

A pushout in a category  $\mathcal{C}$  is a pullback in  $\mathcal{C}^{\text{op}}$ .

- 3. What is a pushout in **Set**? In **Top**?
- 4. Show that the pushout of an isomorphism is an isomorphism.

#### 3 Monomorphisms

A monomorphism is a morphism  $f : A \to B$  such that for every morphisms  $g_1, g_2 : A' \to A$ , we have that  $f \circ g_1 = f \circ g_2$  implies  $g_1 = g_2$ :

$$A' \xrightarrow{g_1} A \xrightarrow{f} B$$

- 1. What is a monomorphism in **Set**?
- 2. Show that the pullback of a monomorphism along any morphism is a monomorphism.
- 3. Show that, in **Set**, the pushout of a monomorphism along any morphism is a monomorphism. Does this seem to be true in any category?
- 4. Define the dual notion of *epimorphism*. What is an epimorphism in **Set**?
- 5. In the category of posets, construct a morphism which is both a monomorphism and an epimorphism, but not an isomorphism.

## 4 (Co)monoids in cartesian categories

- 1. Given a cartesian category  $\mathcal{C}$ , show that the cartesian product induces a functor  $\mathcal{C} \times \mathcal{C} \to \mathcal{C}$ .
- 2. Generalize the definition of *monoid* to any cartesian category (a monoid in **Set** should be a monoid in the usual sense). When is a monoid commutative?
- 3. Generalize the notion of morphism of monoid.
- 4. A comonoid in C is a monoid in  $C^{\text{op}}$ . Make explicit the notion of comonoid.
- 5. What part of the cartesian structure on C did we really need in order to define the notion of monoid?
- 6. Show that in a cartesian category every object is a comonoid (with respect to product).
- 7. Given a category C, show that the category of commutative comonoids and morphisms of comonoids in C is cartesian.