

# Cartesian categories

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## 0 Categories

0. Recall the definition of *category* and provide some examples (e.g. **Set**, **Top**, **Vect**, **Grp**, **Rel**, **Cat**, etc.).

## 1 Cartesian categories

Suppose fixed a category  $\mathcal{C}$ . A *cartesian product* of two objects  $A$  and  $B$  is given by an object  $A \times B$  together with two morphisms

$$\pi_1 : A \times B \rightarrow A \quad \text{and} \quad \pi_2 : A \times B \rightarrow B$$

such that for every object  $C$  and morphisms  $f : C \rightarrow A$  and  $g : C \rightarrow B$ , there exists a unique morphism  $h : C \rightarrow A \times B$  making the diagram

$$\begin{array}{ccc} & C & \\ & \vdots & \\ & h & \\ & \downarrow & \\ & A \times B & \\ \swarrow & & \searrow \\ f & & g \\ \downarrow & & \downarrow \\ A & & B \\ \swarrow & \pi_1 & \searrow \\ & & \end{array}$$

commute. We also recall that a *terminal object* in a category is an object  $1$  such that for every object  $A$  there exists a unique morphism  $f_A : A \rightarrow 1$ . A category is *cartesian* when it has finite products, i.e. has a terminal object and every pair of objects admits a product.

1. Suppose that  $(E, \leq)$  is a poset. We associate to it category whose objects are elements of  $E$  and such that there exists a unique morphism between object  $a$  and  $b$  iff  $a \leq b$ . What is a terminal object and a product in this category?

*Solution.* A terminal object is an element  $b \in E$  such that, for every other element  $a \in E$ , there is a unique morphism  $a \rightarrow b$ , i.e.  $a \leq b$ . A terminal object is thus precisely a maximal element of the set.

Given two elements  $a$  and  $b$  of  $E$ , a product is an element  $a \times b$  equipped with two morphisms  $\pi_1 : a \times b \rightarrow a$  and  $\pi_2 : a \times b \rightarrow b$  such that for every element  $c$  equipped with two morphisms  $f : c \rightarrow a$  and  $f : c \rightarrow b$  there exists a unique morphism  $h : c \rightarrow a \times b$  making some diagrams commute. This is thus precisely an element  $a \times b$  such that  $a \times b \leq a$  and  $a \times b \leq b$  such that for every element  $c$  with  $c \leq a$  and  $c \leq b$ , we have  $c \leq a \times b$ . Otherwise said,  $a \times b$  is an infimum of  $a$  and  $b$ .

2. Show that the category **Set** of sets and functions is cartesian.

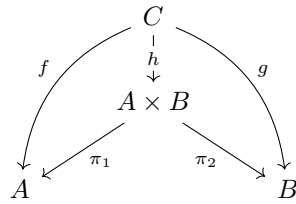
*Solution.* Given two sets  $A$  and  $B$ , we define

$$A \times B = \{(a, b) \mid a \in A \text{ and } b \in B\}$$

and

$$\begin{array}{ll} \pi_1 : A \times B \rightarrow A & \pi_2 : A \times B \rightarrow B \\ (a, b) \mapsto a & (a, b) \mapsto b \end{array}$$

Given a set  $C$  and functions  $f : C \rightarrow A$  and  $g : C \rightarrow B$ , suppose that there exists a function  $h : C \rightarrow A \times B$  making the following diagram commute:



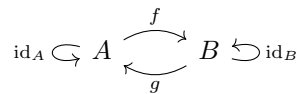
Given  $c \in C$ , we have  $h(c) \in A \times B$ , i.e.  $h(c)$  is of the form  $h(c) = (a, b)$ . The commutation of the left triangle imposes

$$a = \pi_1(a, b) = \pi_1 \circ h(c) = f(c)$$

and the one of the right that  $b = g(c)$ . Therefore, necessarily, we have  $h(c) = (f(c), g(c))$  (if  $h$  exists it is unique). Conversely, the function  $h$  thus defined makes the two triangle commutes ( $h$  actually exists).

3. Show that two terminal objects in a category are necessarily isomorphic.

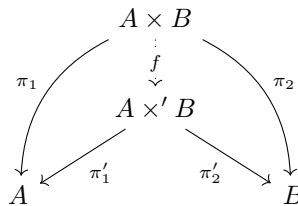
*Solution.* Suppose given two terminal objects  $A$  and  $B$ . Since  $B$  is terminal, we have a unique morphism  $f : A \rightarrow B$  and, since  $A$  is terminal, we have a unique morphism  $g : B \rightarrow A$ .



The composite  $g \circ f : A \rightarrow A$  and  $\text{id}_A : A \rightarrow A$  are both morphisms with the same source and  $A$  as target: since  $A$  is terminal, we therefore have  $g \circ f = \text{id}_A$ . Similarly, we have  $f \circ g = \text{id}_B$  and we deduce that  $A$  and  $B$  are isomorphic.

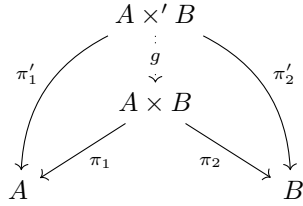
4. Similarly, show that the cartesian product of two objects is defined up to isomorphism.

*Solution.* Fix two objects  $A$  and  $B$  and suppose that they admit two products  $(A \times B, \pi_1, \pi_2)$  and  $(A \times' B, \pi'_1, \pi'_2)$ . From the following diagram,

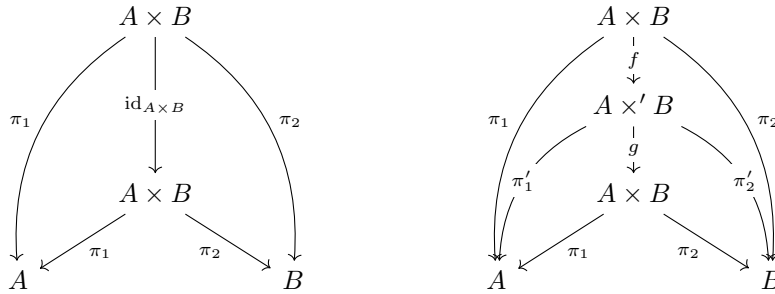


by definition of the product  $A \times' B$ , we deduce the existence of a unique morphism  $f : A \times B \rightarrow A \times' B$  such that  $\pi'_1 \circ f = \pi_1$  and  $\pi'_2 \circ f = \pi_2$ . Similarly, there exists a unique morphism

$g : A \times' B \rightarrow A \times B$  such that  $\pi_1 \circ g = \pi'_1$  and  $\pi_2 \circ g = \pi'_2$ :



Now, we have two morphisms from  $A \times B$  to  $A \times B$ , namely  $\text{id}_{A \times B}$  and  $g \circ f$ :



and those make the two triangles commute: we have

$$\pi_1 \circ \text{id}_{A \times B} = \pi_1$$

$$\pi_2 \circ \text{id}_{A \times B} = \pi_2$$

and

$$\pi_1 \circ g \circ f = \pi'_1 \circ f = \pi_1$$

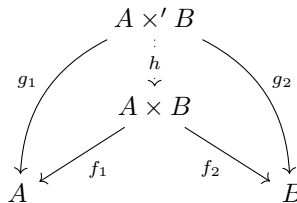
$$\pi_2 \circ g \circ f = \pi'_2 \circ f = \pi_2$$

Therefore by universal property of the product, we have  $g \circ f = \text{id}_{A \times B}$ . Similarly, we have  $f \circ g = \text{id}_{A \times' B}$  and thus  $A \times B$  and  $A \times' B$  are isomorphic.

5. How could you show previous question using question 3.?

*Solution.* Suppose fixed two objects  $A$  and  $B$  of our ambient category  $\mathcal{C}$ . The idea is to construct another category  $\mathcal{D}$  in which a terminal object is precisely a product of  $A$  and  $B$ . We therefore define the category

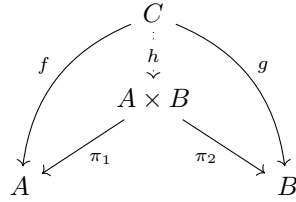
- whose objects are triple  $(C, f_1, f_2)$  where  $C$  is an object of  $\mathcal{C}$  and  $f_1 : C \rightarrow A$  and  $f_2 : C \rightarrow B$  are morphisms of  $\mathcal{C}$ ,
- a morphism in  $\mathcal{D}((C, f_1, f_2), (D, g_1, g_2))$  is a morphism  $h : C \rightarrow D$  of  $\mathcal{C}$  such that  $g_1 \circ h = f_1$  and  $g_2 \circ h = f_2$ :



- identities are identities of  $\mathcal{C}$ ,
- composition is the same as in  $\mathcal{C}$ .

(We leave the reader check the composite is well-defined, i.e. that the composite of two morphisms is still a morphism, and that identities are actually morphisms). A terminal object in  $\mathcal{D}$  is an object  $(A \times B, \pi_1, \pi_2)$  such that for every other object  $(C, f, g)$  there is a

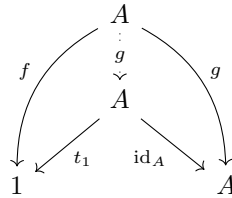
unique morphism  $h : (C, f, g) \rightarrow (A \times B, \pi_1, \pi_2)$  in  $\mathcal{D}$ , i.e. there exists a unique morphism  $h : C \rightarrow A \times B$  in  $\mathcal{C}$  making the following diagram commute:



i.e.  $A \times B$  is a product of  $A$  and  $B$ . By question 3, two such objects are isomorphic in  $\mathcal{D}$ , and they will thus be isomorphic in  $\mathcal{C}$  since composition and identities in  $\mathcal{D}$  are induced by those of  $\mathcal{C}$ .

6. Show that for every object  $A$  of a cartesian category, the objects  $1 \times A$ ,  $A$  and  $A \times 1$  are isomorphic.

*Solution.* Of course the same reasoning “by hand” as above can be performed here. Another way to proceed in order to show that  $1 \times A$  and  $A$  are isomorphic is to show that  $(A, t_A, \text{id}_A)$  is a product of  $1$  and  $A$  (where  $t_A : 1 \rightarrow A$  is the terminal map) and conclude by question 4. Namely, given two morphisms  $f : C \rightarrow 1$  and  $g : C \rightarrow A$ , we have the morphism  $g$  which make the following diagram commute:



The left triangle commutes because  $1$  is terminal and therefore  $\pi_1 \circ g = f$ , and the right triangle commutes by definition of identities:  $\text{id}_A \circ g = g$ . Conversely,  $g$  is the only such morphism by commutation of the right triangle.

7. Show that for every objects  $A$  and  $B$ ,  $A \times B$  and  $B \times A$  are isomorphic.  
 8. Show that for every objects  $A$ ,  $B$  and  $C$ ,  $(A \times B) \times C$  and  $A \times (B \times C)$  are isomorphic.

## 2 Examples of cartesian categories

1. Show that the category **Rel** of sets and relations is cartesian.

*Solution.* Let us first properly define the category **Rel**. An object is a set, a morphism in  $\mathbf{Rel}(A, B)$  is a relation between  $A$  and  $B$ , i.e. a subset of  $A \times B$ :

$$\mathbf{Rel}(A, B) = \mathcal{P}(A, B)$$

Composition of  $R : A \rightarrow B$  and  $S : B \rightarrow C$ , i.e.  $R \subseteq A \times B$  and  $S \subseteq B \times C$ , is the following subset of  $A \times C$ :

$$S \circ R = \{(a, c) \mid \exists b \in B. (a, b) \in R \wedge (b, c) \in S\}$$

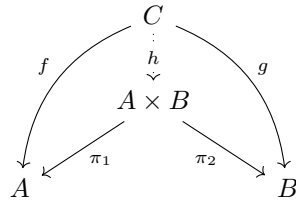
the identity on  $A$  is the diagonal subset of  $A \times A$ :

$$\text{id}_A = \{(a, a) \mid a \in A\}$$

It turns out that that product of  $A$  and  $B$  is the disjoint union  $A \sqcup B$  of the sets  $A$  and  $B$  (we write  $\iota_1(a)$ , resp.  $\iota_2(b)$ , for the canonical injections of an element  $a \in A$ , resp.  $b \in B$ , in  $A \sqcup B$ ). The projections are

$$\pi_1 = \{(\iota_1(a), a) \mid a \in A\} \subseteq (A \sqcup B) \times A \quad \pi_2 = \{(\iota_2(b), b) \mid b \in B\} \subseteq (A \sqcup B) \times B$$

Given a pair of morphisms  $f : C \rightarrow A$  and  $g : C \rightarrow B$ ,



one can check that the unique mediating morphism  $h : C \rightarrow A \times B$  is

$$h = \{(c, \iota_1(a)) \mid (c, a) \in f\} \cup \{(c, \iota_1(b)) \mid (c, b) \in g\} \subseteq C \times (A \sqcup B)$$

2. We write **Vect** for the category of  $\mathbb{k}$ -vector spaces (where  $\mathbb{k}$  is a fixed field) and linear functions. Show that this category is cartesian. Given a basis for  $A$  and  $B$ , describe a basis for  $A \times B$ .
3. Show that the category **Cat** is cartesian.

### 3 Cartesian product as a functor

1. Recall the definition of a *functor* and provide some examples.
2. Define the category **Cat** of categories and functors.
3. Given a cartesian category  $\mathcal{C}$ , show that the cartesian product induces a functor  $\mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ .