THE PHYSICS OF PROGRAMS

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The Universe is executing a program.

Blame it on Feynman

I had some intuitions about analogies between physic and semantics of programs and then I read [Fey96]:



our aim is to exhibit some Hamiltonian for a system which could serve as a computer.

THERMODYNAMICS

Consider a set with *n* equiprobable elements.

The **information** brought by an element is

$\log_2(n)$

It can namely be identified by a string of $log_2(n)$ bits.

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Writing $p = \frac{1}{p}$ for the probability, the information is thus

$$l(p) = \log_2(n) = \log_2\left(\frac{1}{p}\right) = -\log_2(p)$$

If each element x_i has probability p_i , the mean information is called the **entropy**:

$$S = \sum_{i} P(x_i) I(x_i) = -\sum_{i} p_i \log_2(p_i)$$

A function I(p) for information should be

- 1. anti-monotonic: $p_1 \le p_2$ implies $I(p_1) \ge I(p_2)$
- 2. positive: $I(p) \ge 0$

3.
$$l(1) = 0$$

4.
$$l(p_1p_2) = l(p_1) + l(p_2)$$

The entropy is the only function satisfying this.

For instance, consider $\{a, b, c\}$ with

$$\rho(a) = \frac{1}{2}$$
 $\rho(b) = \frac{1}{4}$
 $\rho(c) = \frac{1}{4}$

These can be encoded as



and

$$\begin{split} S &= -\frac{1}{2}\log_2\left(\frac{1}{2}\right) - \frac{1}{4}\log_2\left(\frac{1}{4}\right) - \frac{1}{4}\log_2\left(\frac{1}{4}\right) \\ &= \frac{1}{2}\log_2(2) + \frac{1}{2}\log_2(4) \\ &= 1.5 \end{split}$$

Gas

Given a gas, the pressure P such that the resulting force is

$$F = PA$$

where A is the surface.

Given a piston



the work done on the gas is

$$dW = Fdx$$

or equivalently

$$dW = PA dx = P dV$$

We know that

$$PV = NkT$$

where

- ► N: number of particles
- ► $k = 1.38 \times 10^{-23} \text{ J K}^{-1}$
- ► T: temperature
- ► *kT*: (proportional to) the mean kinetic energy of particles

The work is thus

$$dW = \int_{V_1}^{V_2} NkT \frac{dV}{V}$$

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If the compression is *isothermal* (no change in *T*):

$$W = NkT\log \frac{V_2}{V_1}$$

(this work is energy given to the environment!)

Firs law of thermodynamics

The internal energy U satisfies

 $dU = \delta Q + dW$

where

- δQ is the heat brought to the system
- dW is the work done on the system

In our compression, dU = 0 so that

$$Q = -W$$

Second law of thermodynamics

There is a quantity called entropy with

$$dS \ge \frac{\delta Q}{T}$$

with equality for a reversible process:

- gas is at equilibrium,
- same temperature as surroundings,
- etc.

Suppose that we half the volume, with same temperature



We have

$$\Delta S = \frac{\Delta Q}{T}$$
$$= \frac{NkT\log\left(\frac{V}{2V}\right)}{T}$$
$$= -Nk\log(2)$$



In particular for N = 1 particle, halving the volume means

$$\Delta S = -k \log(2)$$

We have gained one bit of information! (up to a multiplicative constant)

Entropy

Boltzmann developed a statistical definition of entropy:

$$S = -k \sum_{i} p_{i} \ln(p_{i})$$

where

- ▶ *k* is the Boltzmann constant,
- *i* ranges over possible microstates giving rise to the macrostate.

In other words, S says how precisely we know what's in the box.

REVERSIBLE COMPUTING

Information in a tape

The information in a finite boolean tape

$$\cdots 0 1 1 ? ? 0 ? 1 0 ? \cdots$$

is the free energy required to reset the tape to 0 [Lan61]:

- ▶ if we know it is a 0, there is nothing to do
- ▶ if we know it is a 1 then it takes no energy to put it to 0:



if we don't know what it is it takes energy to put it to 0:

$$W = kT\log(2)$$

Reversible computations

A computation is **reversible** if no data if ever erased.

For those there is no lower limit on the consumed energy!

¹excepting in billiard ball, but here we are sensitive to fluctuations...

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This is wonderful but:

- bit erasure accounts for 1/10000-th of the consumed energy in a CPU today,
- they have to be run at infinitesimal speed, otherwise we loose energy: if we want the computer to have r times more chance to run forward than backwards then each step requires¹

$kT\log(r)$

¹excepting in billiard ball, but here we are sensitive to fluctuations...

Reversible boolean circuits

A boolean circuit is **reversible** if no information is lost.

Otherwise said, I can recover the inputs from the outputs.

Otherwise said, it computes an injective boolean function.

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Example

- NOT is reversible
- AND, OR, XOR are not reversible

The gate

$\mathsf{NOT} \quad : \quad \mathbb{B} \to \mathbb{B}$

is reversible:

The controlled not gate

 $\mathsf{CNOT} \quad : \quad \mathbb{B}^2 \to \mathbb{B}^2$

is reversible:

<i>i</i> ₀	<i>i</i> ₁	00	<i>O</i> ₁
0	0	0	0
0	1	0	1
1	0	1	1
1	1	1	0

The controlled controlled not (or Toffoli) gate

 $\mathsf{CCNOT} \quad : \quad \mathbb{B}^3 \to \mathbb{B}^3$

is reversible: last bit is inverted in both other are true.

<i>i</i> 0	i_1	i_2	00	o_1	O_2
0	0	0	0	0	0
0	0	1	0	0	1
0	1	0	0	1	0
0	1	1	0	1	1
1	0	0	1	0	0
1	0	1	1	0	1
1	1	0	1	1	1
1	1	1	1	1	0

The controlled swap (or Fredkin) gate

$\mathsf{CSWAP} \quad : \quad \mathbb{B}^3 \to \mathbb{B}^3$

is reversible: last two bits are exchanged if first is true.

<i>i</i> 0	i_1	i_2	00	o_1	O ₂
0	0	0	0	0	0
0	0	1	0	0	1
0	1	0	0	1	0
0	1	1	0	1	1
1	0	0	1	0	0
1	0	1	1	1	0
1	1	0	1	0	1
1	1	1	1	1	1

The swap gate $: \mathbb{B} \rightarrow \mathbb{B}^2$ SWAP is reversible: i_0 i_1 $o_0 | o_1$ 0 0 0 0 1 1

0

1 1 1

0

0

0 | 1

1 1

Perhaps more surprisingly, the duplicator

$$\mathsf{DUP} : \mathbb{B} \to \mathbb{B}^2$$

is also reversible:

$$\begin{array}{c|ccc} i_0 & o_0 & o_1 \\ \hline 0 & 0 & 0 \\ 1 & 1 & 1 \end{array}$$

Universal gates

A set of gates is **universal** if for every boolean function f, there is a reversible circuit r built from those such that f is obtained by discarding some outputs from r.

For instance,

- ▶ NOT + CNOT is not universal,
- CCNOT is universal,
- CSWAP is universal.

Reversible computing

In general, starting from a blank tape, we want our machines to give a blank tape (otherwise energy is required to reset it!) [Ben73]

Typically, starting from a blank tap and data, we

- 1. perform our computation, keeping intermediate results,
- 2. copy the final result to some memory,
- 3. reverse the computation to obtain again a blank tape.

We can implement reversible computations on physical devices! [Ben82]

The billiard-ball model

In the billiard-ball model [FT82], the switch:



and CSWAP:



In this model, only bijections which are conservative (preserve the number of 1) can be implemented: we code a bit as a pair 01 / 10.

A chemical machine



Fig. 6. Hypothetical enzymatic Turing machine. Macromolecular tape (a) consists of a structural backbone S-S-S bearing tape symbols 1,0 and head marker α . Macromolecule reacts (c,d) with enzyme (b) that catalyzes the transition $\alpha 0 \rightarrow 1R\beta$, via specific binding sites (tabs), thereby resulting in logical successor configuration (c). Enzyme is then prepared for reuse
A mechanical machine



Fig. 5. (a) Realization of $\Theta^{(3)}$, (b) Details of the AND mechanism.

A mechanical machine



Fig. 7. Brownian Turing machine made of rigid, frictionless, loosely fitting clockwork. This figure shows the Turing machine tape (a, b, c) and the read-write-shift equipment. The machine is scanning square b. Each tape square has a disk (d) which interlocks with several E-shaped bit-storage blocks (e), holding them in the up (1) or down (0) position. A framework (f) fits over the scanned tape square, engaging the disks of the two adjacent squares (via their grooves g), to keep them from rotating when they are not supposed to. After the bits are read (cf. next figure) they must in general be changed. In order to change a bit, its knob (k) is first arranged by the manipulator (m) then the potch (a) is rotated into alignment by the according to the scanned tape square (m) then the potch (b) is first severed.

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A mechanical machine



Fig. 9. Above: master camshaft of the clockwork Turing machine's control unit. Below: top view of camshaft, and the transition rules to which its tunnels are isomorphic.

Note that the brownian movement will make it move randomly, so we only need a bit of energy to push it forward. If the function was not reversible, we would get lost in possible backward paths.

Relations

We know the gates to generate boolean circuits (do we?).

But what are the relations?

Lafont [Laf03] has investigated presentation for the corresponding PROP.

PROP = symmetric monoidal category with \mathbb{N} as objects

The PROP of functions

The PROP with morphisms functions $\mathbb{B}^m \to \mathbb{B}^n$ is generated by



The PROP of functions

Subject to relations:



The PROP of functions



For the PROP **B** of boolean bijections, we could hope that it is generated by the CCNOT gate:



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Theorem

The PROP B is not finitely generated.

Proof.

This follows from the remark that if

$$f : \mathbb{B}^n \to \mathbb{B}^n$$

is a bijection, then

$$f \times \mathrm{id}_{\mathbb{B}}$$
 : $\mathbb{B}^{n+1} \to \mathbb{B}^{n+1}$

is an even permutation.

However,

Theorem

The PROP **A** of even permutations is contained in the sub-PROP of **B** generated by τ , T_1 , T_2 and T_3 . It is thus finitely generated.

The relations are not known though.

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As a corollary, we have

Theorem

The PROP **B** is generated by τ and the T_n for $n \ge 1$.

The relations are not known though.

Question

Since we have duplication, we are more interested in the PROP for boolean injections.

It could also be interesting to consider conservative bijections.

I don't know of any presentation.

Unfoldings

► RCCS

reversing = universal cover (we take a cofibrant replacement)

There is another fundamental physical computer based on quantum mechanics.

We need some more material to appreciate it.

CLASSICAL MECHANICS

Newton's law

Recall Newton's law of motion

$$m\ddot{q} = F$$

(we write q for the position of the particle).

It can be derived from first principles!

Newton's law

Recall Newton's law of motion

$$\dot{\rho} = F$$

(we write $p = m\dot{q}$ for the **momentum** of the particle).

It can be derived from first principles!

Lagrangian mechanics

Axiom

Any mechanical system (e.g. a particle) is characterized by a function, called **Lagrangian**

 $L(q,\dot{q})$

depending on

- the position q,
- ▶ the velocity ġ,
- ▶ and not anything else (not \ddot{q} , a hidden variable, etc.)!

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Remark

Here \dot{q} is a formal parameter (not a derivative), i.e.,

$$L$$
 : $TQ \rightarrow \mathbb{R}$

for some manifold Q.

Action

The action is

$$S \quad = \quad \int_{t_1}^{t_2} L(q, \dot{q}) \, \mathrm{d}t$$

Axiom

The **principle of least action**: a mechanical system has minimal (or, more precisely, stationary) action.

The principle of least action

If we change q to $q + s \delta q$, the first-order variation of the action should be zero:



The principle of least action

If we change *q* to $q + s \delta q$, the first-order variation of the action should be zero:



A **variation** of *q* is a function δq such that

$$\delta q(t_1) = \delta q(t_2) = 0$$

and we formally consider

$$\delta S = \frac{\mathsf{d}}{\mathsf{ds}} S(q_{\mathsf{s}}) \bigg|_{\mathsf{s}=0}$$

with

 $q_{s} = q + s \delta q$

Euler-Lagrange equation

We have

$$\delta S = \delta \int_{t_1}^{t_2} L \, dt$$

$$= \int_{t_1}^{t_2} \left(\frac{\partial L}{\partial q} \delta q + \frac{\partial L}{\partial \dot{q}} \delta \dot{q} \right) dt$$

$$= \left[\frac{\partial L}{\partial \dot{q}} \delta q \right]_{t_1}^{t_2} + \int_{t_1}^{t_2} \left(\frac{\partial L}{\partial q} - \frac{d}{dt} \frac{\partial L}{\partial \dot{q}} \right) \delta q \, dt$$

(by an i.p.p. since $\delta \dot{q} = d\delta q / dt$).

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(by an i.p.p. since $\delta \dot{q} = d\delta q / dt$).

Therefore $\delta S = 0$ for every δq if and only if

$$\frac{\mathrm{d}}{\mathrm{d}t}\frac{\partial L}{\partial \dot{q}} = \frac{\partial L}{\partial q}$$

called the Euler-Lagrange equation.

The Euler-Lagrange equation

<u>1</u>

 $\frac{\mathrm{d}}{\mathrm{d}t}\frac{\partial L}{\partial \dot{q}} = \frac{\partial L}{\partial q}$

We define

We have:

▶ the momentum:

the force:

$$p = \frac{\partial L}{\partial \dot{q}}$$
$$F = \frac{\partial L}{\partial q}$$

The equation becomes

$$\frac{\mathrm{d}\rho}{\mathrm{d}t} = F$$

A force *F* is **conservative** if there is a potential $V : Q \to \mathbb{R}$ such that

$$F = -\frac{\partial V}{\partial q}$$

(or $F = -\nabla V$ in general).

This is equivalent to the work

$$V(q) = -\int_{q_0}^{q} F \cdot dq = -\int_{t_1}^{t_2} F(q) \cdot \dot{q} dt$$
$$q(t_1) = q_0$$
$$q(t_2) = q$$

being well-defined, i.e., not depending on the chosen path.

We will see that energy is conserved for those.

Axiom All forces are conservative.

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Example For the harmonic oscillator,



we have

$$F = -kx \qquad \qquad V = \frac{1}{2}kx^2$$

-1

Axiom All forces are conservative.

Example A friction force

$$F = -\alpha \dot{x}$$

is not conservative, but there are missing parts in our system!

The Lagrangian

In general, the Lagrangian looks like

$$L = E - V = \frac{1}{2}m\dot{q}^2 - V$$

so that

$$\frac{\mathrm{d}}{\mathrm{d}t}\frac{\partial L}{\partial \dot{q}} = m\ddot{q} = F = \frac{\partial L}{\partial q}$$

We recover Newton's equation!

Conservation of energy

The laws of physics are invariant by translation in time:

$$\frac{\partial L}{\partial t} = 0$$

Therefore,

$$\begin{aligned} \frac{\mathrm{d}L}{\mathrm{d}t} &= \frac{\partial L}{\partial q} \dot{q} + \frac{\partial L}{\partial \dot{q}} \ddot{q} + \frac{\partial L}{\partial t} \\ &= \frac{\mathrm{d}}{\mathrm{d}t} \frac{\partial L}{\partial \dot{q}} \dot{q} + \frac{\partial L}{\partial \dot{q}} \ddot{q} \\ &= \frac{\mathrm{d}}{\mathrm{d}t} \left(\frac{\partial L}{\partial \dot{q}} \dot{q} \right) \end{aligned}$$

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We define the energy (or hamiltonian) as

$$H = \frac{\partial L}{\partial \dot{q}} \dot{q} - L$$

and have

$$\frac{\mathrm{d}H}{\mathrm{d}t} = 0$$

Theorem

If there is a continuous symmetry of the system then there is a conserved quantity.

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We suppose given a map

$$egin{array}{cccc} \mathbb{Q}^{\mathbb{R}} imes \mathbb{R} & o & Q^{\mathbb{R}}\ (q,s) & \mapsto & q_{ ext{s}} \end{array}$$

(

Typically,

$$q_{\mathrm{s}}(t) = q(t+\mathrm{s})$$
 $q_{\mathrm{s}}(t) = q(t) + \mathrm{sv}$ etc.

Consider and infinitesimal transformation and its effect on S:

$$\delta q = \left. \frac{\mathrm{d}}{\mathrm{ds}} q_{\mathrm{s}} \right|_{\mathrm{s}=0} \qquad \delta L = \left. \frac{\mathrm{d}}{\mathrm{ds}} S(q_{\mathrm{s}}, \dot{q}_{\mathrm{s}}) \right|_{\mathrm{s}=0}$$

We have (we already went through this)

$$\delta S = \dots$$

$$= \left[\frac{\partial L}{\partial \dot{q}} \delta q \right]_{t_1}^{t_2} + \int_{t_1}^{t_2} \underbrace{\left(\frac{\partial L}{\partial q} - \frac{d}{dt} \frac{\partial L}{\partial \dot{q}} \right)}_{=0} \delta q \, dt$$

Therefore $\delta S = 0$ implies that the quantity

$$\frac{\partial L}{\partial \dot{q}} \delta q = p \delta q$$

is invariant with respect to t.

More generally, if

$$\delta L = \frac{\mathrm{d}\ell}{\mathrm{d}t}$$

then the quantity

$$p \, \delta q - \ell$$

is conserved over time.
Noether's theorem

A general galilean transformation is given by

symmetry	invariant
translation in time (s)	energy (hamiltonian)
translation in space (a)	momentum
uniform move (v)	center of mass
rotation (R)	angular momentum
gauge invariance of the e.m. field	conservation of charge
:	:

Reversibility

There is another kind of symmetry which is not covered:

$$egin{array}{cccc} q & \mapsto & q \ \dot{q} & \mapsto & -\dot{q} \ t & \mapsto & -t \end{array}$$

Time is reversible!

In fact, since the evolution is described by a differential equation it is **reversible**: two states cannot lead to the same state.

Parity inversion

If we negate all the coordinates, i.e.,

$$q \quad \rightsquigarrow \quad -q$$

This generally gives rise to a discrete symmetry, but is violated in weak interactions (Wu experiment)!

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We can also use the **phase space** $(q, p) \in T^*Q$.

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$$H(q,p) = p\dot{q} - L(q,\dot{q})$$

which is valid if

$$\det\left(\frac{\partial p_i}{\partial \dot{q}_j}\right) = \det\left(\frac{\partial^2 L}{\partial \dot{q}_i \partial \dot{q}_j}\right) \quad \neq \quad 0$$

Hamilton's equations are satisfied:

$$\frac{\mathrm{d}q}{\mathrm{d}t} = \frac{\partial H}{\partial p} \qquad \qquad \frac{\mathrm{d}p}{\mathrm{d}t} = -\frac{\partial H}{\partial q}$$

(the second by Euler-Lagrange)

The Poisson bracket

Given $A, B : T^*Q \to \mathbb{R}$, their **Poisson bracket** is

$$\{A,B\} = \sum_{i} \frac{\partial A}{\partial q_{i}} \frac{\partial B}{\partial p_{i}} - \frac{\partial A}{\partial p_{i}} \frac{\partial B}{\partial q_{i}}$$

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This is a Lie bracket:

- linear
- antisymmetric: $\{B,A\} = -\{A,B\}$
- ▶ satifying Jacobi: $\{\{A,B\},C\} + \{\{B,C\},A\} + \{\{C,A\},B\} = 0$

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In particular

$$\{q_i, q_j\} = \{p_i, p_j\} = 0$$
 $\{q_i, p_j\} = \delta_{ij}$

hamiltonian = evolution

The hamiltonian encodes the evolution of a physical quantity $A: T^*Q \to \mathbb{R}$:

$$\frac{dA}{dt} = \frac{dA}{dq}\frac{dq}{dt} + \frac{dA}{dp}\frac{dp}{dt}$$
$$= \frac{dA}{dq}\frac{\partial H}{\partial p} - \frac{dA}{dp}\frac{\partial H}{\partial q}$$
$$= \{A, H\}$$

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$$= \{A, H\}$$

In particular, Hamilton equations are

$$\frac{\mathrm{d}q}{\mathrm{d}t} = \{q, H\} \qquad \qquad \frac{\mathrm{d}p}{\mathrm{d}t} = \{p, H\}$$

TODO: Noether's theorem.....

QUANTUM MECHANICS

A Hilbert space ${\mathcal H}$ is a ${\mathbb C}\text{-vector}$ space equipped with an inner product such that

- $\blacktriangleright \langle \phi | \rangle$ is linear
- \blacktriangleright $\langle -|\psi \rangle$ is antilinear
- $\blacktriangleright \langle \phi | \psi \rangle = \langle \psi | \phi \rangle^*$
- $\blacktriangleright \ \langle \psi | \psi \rangle \geq 0$

We write

- $\blacktriangleright ~ \|\psi\| = \sqrt{\langle \psi |\psi \rangle}$ for the norm
- $\blacktriangleright |\psi
 angle \in \mathcal{H}$ for a vector

$$\blacktriangleright \ \langle \psi | = \langle \psi | - \rangle \in \mathcal{H}^*$$

The bra-ket notation is nice:

$$\langle \phi | | \psi \rangle = \langle \phi | \psi \rangle$$

and more generally, for $A:\mathcal{H}\to\mathcal{H}$

$$\langle A^* \phi | \psi \rangle = \langle \phi | A | \psi \rangle = \langle \phi | A \psi \rangle$$

where A^* is the conjugate transpose of A:

$$\langle \phi | A | \psi \rangle = \langle \psi | A^* | \phi \rangle$$

An Hilbert space \mathcal{H} admits an orthonormal basis (e_i):

$$\langle \mathbf{e}_i | \mathbf{e}_j \rangle = \delta_{ij}$$

a vector: $|\psi\rangle = \sum_{i} \psi_{i} |e_{i}\rangle$ **a** operator $A : \mathcal{H} \to \mathcal{H}$ $A_{ij} = \langle e_{j} | A | e_{i} \rangle$ and $A = \sum_{i} |e_{i}\rangle A_{ii} \langle e_{i} \rangle$

$$egin{array}{rcl} \mathsf{A} &=& \sum_{i,j} \ket{e_j} \mathcal{A}_{ij} ig \langle e_i
vert \end{array}$$

We do as if we were in finite-dimensional spaces (and trust smart people for details).

This is not a good idea since a basis for our space typically consists in all possible states $q \in Q$, or all possible values for momentum.

Self-adjoint operators

An operator A is self-adjoint when

$$A = A^*$$

This implies that

- its eigenvalues are real
- eigenvectors with distinct eigenvalues are orthogonal
- we can find an orthonormal basis of eigenvectors²

²in good cases (e.g. finite dimension)

In quantum mechanics, the state of a system is given by a vector

 $|\psi\rangle$

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 $|\psi\rangle$

and an **observable** (= physical quantity) is represented by a self-adjoint operator

Α

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(supposing $\langle \psi | \psi \rangle = 1$, otherwise divide by it) where

 $|\psi\rangle\langle\psi|$

is the **density operator** associated to $|\psi\rangle$.

Density operators

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More generally, we can consider as a **density operators** any ρ which is self-adjoint, nonnegative ($\langle \psi | \rho | \psi \rangle \ge 0$) and of unit trace, those as above are *pure states*.

Expectation generalizes as expected

$$\langle A \rangle = \frac{\operatorname{tr}(\rho A)}{\operatorname{tr} \rho}$$

A *symmetry* for a system has to preserve inner product and thus be of the form

$$\begin{array}{cccc} |\psi\rangle & \mapsto & U \, |\psi\rangle \\ A & \mapsto & UAU^* \end{array}$$

for some **unitary** transformation U

$$UU^* = U^*U = I$$

Remark

Most operations are (sub)unitary because we want to preserve normalization of states: $\langle\psi|\psi\rangle=1.$

Consider a continuous family of transformations U(s) such that

$$U(0) = I$$
 $U(s_1 + s_2) = U(s_1)U(s_2)$

The associated infinitesimal transformation is

$$\delta U = \frac{\mathrm{d}U}{\mathrm{ds}}\Big|_{\mathrm{s}=0}$$

i.e.,

$$U = I + s \,\delta U + O(s^2)$$

By unitarity, we have

$$I = UU^* = I + s(\delta U + \delta U^*) + O(s^2)$$

thus

$$\delta U + \delta U^* = 0$$

and therefore

$$\delta U = iK$$

for some self-adjoint (real) K.

It determines U since

$$\left. \frac{\mathrm{d}U(s_0 + s)}{\mathrm{d}s} \right|_{s=0} = \delta(U(s_0 + s)) = \delta(U(s_0)U(s)) = U(s_0)\,\delta U(s)$$

thus

$$\frac{\mathrm{d}U(s)}{\mathrm{d}s}\Big|_{s=s_0} = U(s_0)\,\mathrm{i}K$$

 $\frac{dU}{ds} = UiK$ $U = e^{iKs}$

thus

i.e.,

Operators from symmetries

A general galilean transformation is

symmetry	operator
rotation	$e^{-i\theta_k J_k}$
translation	$e^{-ia_kP_k}$
velocity	$e^{iv_kG_k}$
time	e ^{isH}

Note: we should divide all operators by \hbar but nevermind.

A translation in time

$$t \quad \leadsto \quad t' = t + s$$

induces a transformation of a vector

$$|\psi(t)\rangle \quad \rightsquigarrow \quad |\psi'(t)\rangle = |\psi(t-s)\rangle = \mathrm{e}^{\mathrm{i} \mathrm{s} \mathrm{H}} |\psi(t)\rangle$$

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$$\frac{\mathrm{d}}{\mathrm{d}t} \left| \psi(t) \right\rangle \quad = \quad -\mathrm{i} H \left| \psi(t) \right\rangle$$

or even

$$H = i \frac{d}{dt}$$

The position operator Q is

$$Q|x\rangle = x|x\rangle$$

The position operator Q is

$$Q\left|x\right\rangle = x\left|x\right\rangle$$

The velocity operator V has to satisfy

$$\langle V \rangle = \frac{d}{dt} \langle Q \rangle$$

and thus

$$\langle \psi | V | \psi \rangle = \frac{\mathrm{d}}{\mathrm{d}t} \left(\langle \psi | Q | \psi \rangle \right) = \left(\frac{\mathrm{d}}{\mathrm{d}t} \langle \psi | \right) Q | \psi \rangle + \langle \psi | Q \left(\frac{\mathrm{d}}{\mathrm{d}t} | \psi \rangle \right)$$
$$= \mathrm{i} \left\langle \psi | HQ | \psi \right\rangle - \mathrm{i} \left\langle \psi | QH | \psi \right\rangle$$

i.e.,

$$V = i(HQ - QH) = i[H, Q]$$
Pictures of motion

The Heisenberg equation of motion is

$$i\frac{dA}{dt} = [A, H]$$

whose solution is

$$A(t) = U^{-1}(t)A(0)U(t)$$

with $U(t) = e^{-iHt}$.

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In the Heisenberg picture

$$\langle A \rangle = \langle \psi | (U^{-1}A_0U) | \psi \rangle$$

In the Schrödinger picture

$$\langle A \rangle = (\langle \psi | U^{-1}) A_0 (U | \psi \rangle)$$

Noether's theorem

Given a bounded self-adjoint operator O,

$$[O,H] = 0$$

if and only if for every state $|\psi\rangle$ such that

$$\frac{\mathrm{d}}{\mathrm{d}t} \left| \psi(t) \right\rangle = -\mathrm{i} H \left| \psi(t) \right\rangle$$

the expected value

 $\langle \psi(t) | O | \psi(t) \rangle$

is constant.

Propagators

The evolution of the state is given by

$$|\psi(t)\rangle = U(t,t_0) |\psi(t_0)\rangle$$

and we have

$$|\psi(x,t)\rangle = \int \langle x| U(t,t_0) |x_0\rangle \psi(x_0,t_0) dx_0$$

where the propagator

$$G(x, t, x_0, t_0) = \langle x | U(t, t_0) | x_0 \rangle$$

is (proportional to) the probability of finding the particle in x at time t if it was in x_0 at t_0 .

Feynman's formula

The propagator can be expressed as a path integral

$$G(x, t, x_0, t_0) = \int_{\substack{q \text{ path} \\ \text{from } (x_0, t_0) \\ \text{to } (x, t)}} e^{iS(q)/\hbar}$$

where S is the **action**

$$S(q) = \int L(q,\dot{q}) dt$$

The least action principle

How does light knows where the shortest path is?



(and also even more wired paths)

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How does light knows where the shortest path is?



(and also even more wired paths)

It does not, it tries everything:

$$G(x,t,x_0,t_0) \quad = \quad \int \mathrm{e}^{\mathrm{i} S(q)/\hbar} \, \mathrm{D} q$$

However, the phases are pretty much random, excepting when *S* is stable: we have recovered the least action principle!

Quantizing physics

When quantizing the laws of physics

- we replace functions by operators
- the Poisson brackets become commutators

$$[Q_i, Q_j] = [P_i, P_j] = 0 \qquad [Q_i, P_j] = i\hbar\delta_{ij}$$

etc.

The momentum coordinates

Let's work out a bit the classical harmonic oscillator where

$$H = \frac{P^2}{2m} + \frac{1}{2}m\omega^2 Q^2$$

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We introduce the creation and anihilation operators

$$a = \sqrt{\frac{m\omega}{2\hbar}} \left(\mathbf{Q} + \frac{1}{m\omega} \mathbf{i} \mathbf{P} \right) \qquad a^* = \sqrt{\frac{m\omega}{2\hbar}} \left(\mathbf{Q} - \frac{1}{m\omega} \mathbf{i} \mathbf{P} \right)$$

and the number operator

$$N = aa^*$$

The Hamiltonian is

$$H = \frac{1}{2}\hbar\omega\left(aa^* + a^*a\right) = \hbar\omega\left(aa^* - \frac{1}{2}\right) = \hbar\omega\left(N - \frac{1}{2}\right)$$

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From [Q, P] = i follows

$$[a, a^*] = 1$$
 $[N, a] = -a$ $[N, a^*] = a$

Given an eigenvector $|n\rangle$ of *N* of eigenvalue $n \in \mathbb{R}$, i.e.,

$$N|n\rangle = n|n\rangle$$

we have

$$Na |n\rangle = a(N-I) |n\rangle = (n-1)a |n\rangle$$

We have constructed an eigenvector for the value n - 1.

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We have constructed an eigenvector for the value n - 1.

However,

$$n = \langle n | N | n \rangle = \langle a | a^* | n \rangle \ge 0$$

Therefore the smallest eigenvalue can only be $|0\rangle$.

Therefore, the eigenvectors for the hamiltonian are

$$H|n\rangle = E_n|n\rangle$$

with

$$E_n = \hbar\omega\left(n+\frac{1}{2}\right)$$

with $n \in \mathbb{N}$: energy levels are quantified.

Note that the relation

$$[a,a^*] = aa^* - a^*a = 1$$

means

$$\sum_{n} ((n+1) - 1) - \sum_{n} ((n-1) + 1) = 1$$

HAMILTONIAN FOR PROGRAMS

Truth values

Consider one of the following possible rings of truth values:

- classical: B
- ▶ probabilistic: [0,1]
- stochastic: N
- ► chemical: ℝ
- ▶ quantum: ℂ

States

Consider a boolean circuit with *m* inputs and *n* outputs.

A state $|\psi\rangle$ is an element of \Bbbk^{2^m} :

$$\begin{aligned} &(i_0, i_1, i_2, \ldots) = (0, 0, 0, \ldots) & \rightsquigarrow & 1 \\ &(i_0, i_1, i_2, \ldots) = (1, 0, 0, \ldots) & \rightsquigarrow & 1 \\ &(i_0, i_1, i_2, \ldots) = (1, 1, 0, \ldots) & \rightsquigarrow & 0 \\ &\vdots & \rightsquigarrow & \vdots \end{aligned}$$

.

The hamiltonian of a boolean function

$$\mathbb{B}^m \rightarrow \mathbb{B}^n$$

is a linear function

$$k^{2^m} \rightarrow k^{2^n}$$

represented by a $2^m \times 2^n$ -matrix.

For instance NOT:

Operators

We define

$$a = \begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix}$$

$$a \left| 0 \right\rangle = \left| 1 \right\rangle \quad a \left| 1 \right\rangle = 0$$

the destruction operator

$$a^* \quad = \quad \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$$

$$a \left| 0 \right\rangle = 0 \quad a \left| 1 \right\rangle = \left| 0 \right\rangle$$

the number operator

$$N_a = aa^* = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$$
 $N_a |0\rangle = 0 |0\rangle$ $N_a |1\rangle = 1 |1\rangle$

Satisfying

$$aa^* + a^*a = 1$$

The hamiltonian associated to NOT is

 $a + a^*$

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$$a + a^*$$

The hamiltonian associated to CNOT is

$$NOT_bN_a + (I - N_a) = (b + b^*)aa^* + a^*a$$

Namely,

if a=1 then { b=not b }
if a=0 then {}

Traces (= feedback) correspond to taking exponentials.

THE GEOMETRY OF INTERACTION

Hamiltonian for programs

- Hamiltonian of a program
- Gol, execution formula (cf Abramsky, Ian Mackie)

STOCHASTIC MECHANICS

TODO: Baez stochastic

Dictionary

physics	programming
physical system	program
tangent bundle TQ	control flow graph (+ states)
speed \dot{q}	instruction

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🔋 Charles H Bennett.

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