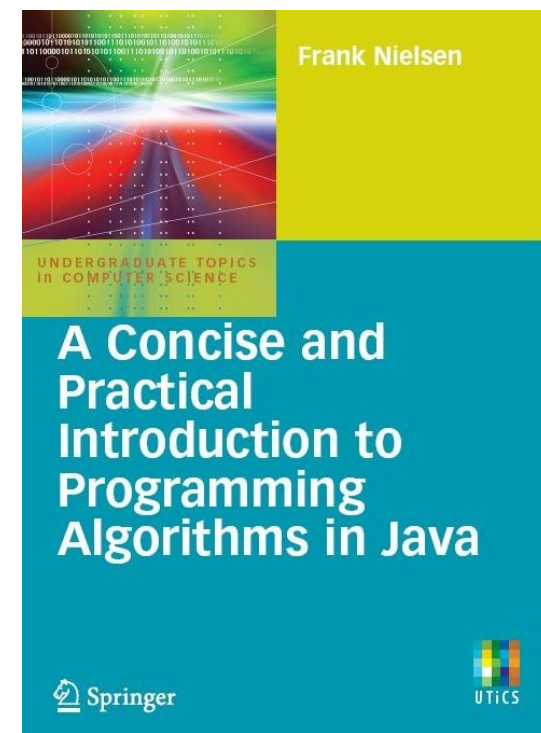


A Concise and Practical Introduction to Programming Algorithms in Java



Chapter 6: Searching and sorting

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Agenda

- Answering a few questions
- Searching (sequential, bisection)
- Sorting for faster searching
- Recursive sort
- Hashing

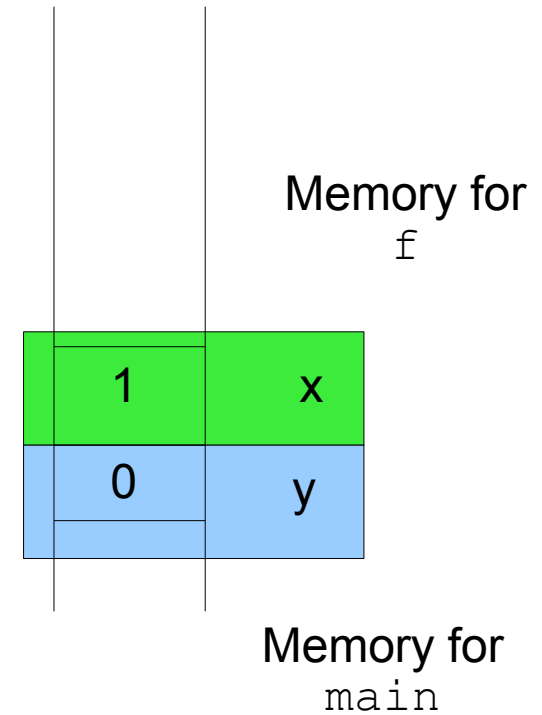
Swapping

Java is **pass-by-value** only.

Arrays and objects: **pass-by-reference** (=memory `value')

```
class FunctionCallSample
{
    public static void f(double x)
    {
x=1;
        System.out.println(x); //1
        return;
    }

    public static void main(String[] args)
    {
        int y=0;
f(y);
        System.out.println(y); //y is still 0
    }
}
```



Swapping & Strings

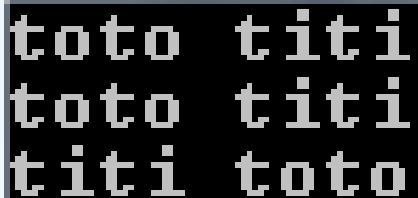
```
public static void Swap(String p, String q)
{
    String tmp;
    tmp=p;
    p=q;
    q=tmp;
}
```

```
public static void Swap(String [] t, int i, int j)
{
    String tmp;
    tmp=t[i];
    t[i]=t[j];
    t[j]=tmp;
}
```

```
String s1="toto";
String s2="titi";
String [] s=new String[2];
s[0]=new String("toto");
s[1]=new String("titi");

System.out.println(s1+" "+s2);
Swap(s1,s2);
System.out.println(s1+" "+s2);
Swap(s,0,1);
System.out.println(s[0]+" "+s[1]);
```

Result



```
toto titi
toto titi
titi toto
```

Swapping using wrapped integers

Let us **wrap** integers into a tailored object named `IntObj`

```
// Wrapping integers into an object
```

```
class IntObj {  
    private int x;  
  
    public IntObj(int x)  
    { this.x = x; }  
  
    public int getValue()  
    { return this.x; }  
  
    public void insertValue(int newx)  
    { this.x = newx; }  
}
```



```

public class SwapInt {

    // Pass-by-value (= pass-by-reference for objects)
    static void swap(IntObj p, IntObj q)
    {
        // interchange values inside objects
        int t = p.getValue();
        p.insertValue(q.getValue());
        q.insertValue(t);
    }

    public static void main(String[] args) {
        int a = 23, b = 12;

        System.out.println("Before| a:" + a + ", b: " + b);

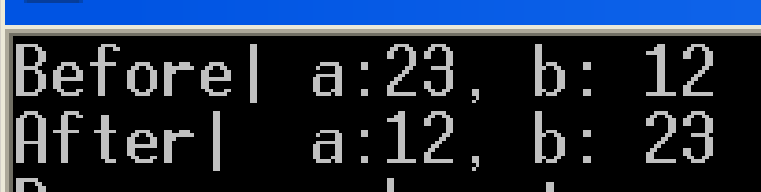
        IntObj aObj = new IntObj(a);
        IntObj bObj = new IntObj(b);

        swap(aObj, bObj);

        a = aObj.getValue();
        b = bObj.getValue();

        System.out.println("After|  a:" + a + ", b: " + b);
    }
}

```



```

Before| a:23, b: 12
After|  a:12, b: 23

```

Java wrapper class for **int** is **Integer** *BUT* it doesn't allow you to alter the data field inside.



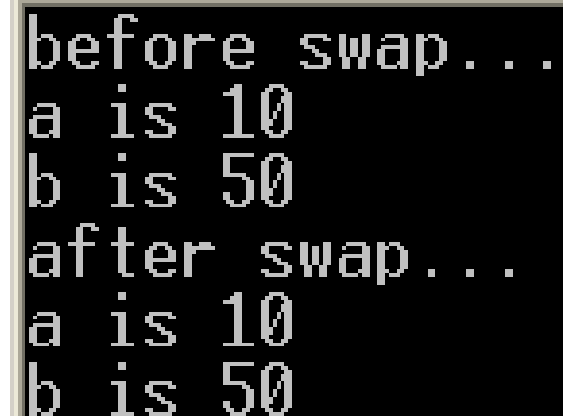
```

class SwapInteger
{
    public static void main(String args[])
    {
        Integer a, b;

        a = new Integer(10);
        b = new Integer(50);

        System.out.println("before swap...");
        System.out.println("a is " + a);
        System.out.println("b is " + b);
        swap(a, b); // References did not change (PASS-BY-VALUE)
        System.out.println("after swap...");
        System.out.println("a is " + a);
        System.out.println("b is " + b);
    }
}

```



```

before swap...
a is 10
b is 50
after swap...
a is 10
b is 50

```

```

public static void swap(Integer a, Integer b)
{
    Integer temp = a;
    a = b;
    b = temp;
}

```



Searching: Problem statement

- **Objects** are accessed via a corresponding **key**
- Each **object** stores its key and additional **fields**
- One seeks for **information** stored in an object from its key
(key= a handle)
- All objects are in the main memory (no external I/O)



More challenging problem:

Adding/removing or changing object attributes **dynamically**

Searching: Basic examples



Dictionary



**Object:
(word, definition)**



Key: word

```
class DictionaryEntry
{
    String word;
    String definition;
}
```



Dictionary

```
class DictionaryEntry
{
String word;
String definition;

DictionaryEntry(String w, String def)
{
this.word=new String(w); // Clone the strings
this.definition=new String(def);
}

}
```

```
class TestMyDictionary
{public static void main(String[] args){
    DictionaryEntry [] Dico =new DictionaryEntry[10];

    Dico[0]=new DictionaryEntry("INF311","Cours d'informatique");
    Dico[1]=new DictionaryEntry("ECO311","Cours d'economie");

}
}
```

Searching: Basic examples

- Objects denote people: Each object represents an individual
- The key (handle) is the lastname, firstname, ...
- Additional information fields are:
address, phone number, etc.



Firstname
Lastname
Phonenumber
int Age

```
class Individual {  
    String Firstname;  
    String Lastname;  
  
    String Address;  
    String Phonenumber;  
  
    int Age;  
    boolean Sex;  
}
```



Object

Searching: Basic examples

- Objects denote 2D point sets: Each object is a 2D point
- The key (handle) is the name, or x- or y- ordinate, etc.
- Additional information fields are:
color, name, etc.



x
y
name
color

```
class Color {int Red, Green, Blue;}  
  
class Point2D  
{  
  double x;  
  double y;  
  
  String name;  
  Color color;  
}
```

Sequential (linear) search

- Objects are stored in an **array** (of objects):
Any order is fine (= not necessarily sorted)
 - To seek for an object, we browse sequentially the array until we **find** the object, or that we report that it is **not inside**
 - Compare the **keys** of the query with object keys
-

```
class Person{
    String firstname, lastname; // key for searching
    String phone; // additional fields go here

    // Constructor
    Person(String l, String f, String p)
    {firstname=f; lastname=l; phone=p;}

};
```



```
static Person[] array=new Person[5];

static String LinearSearch(String lastname, String firstname)
{
    for(int i=0;i<array.length;i++)
    {
        if ((lastname.equals(array[i].lastname) &&
            (firstname.equals(array[i].firstname))))
            {return array[i].phone;}
    }

    return "Not found in my dictionary";
}
```

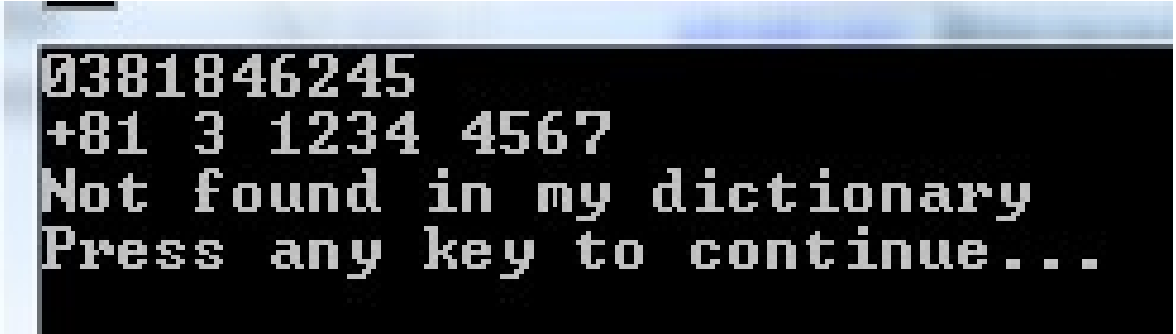
Array is class variable, here

```
class LinearSearch{
    static Person[] array=new Person[5];

    static String LinearSearch(String lastname, String firstname)
    {...}

    public static void main (String [] args)
    {
        array[0]=new Person("Nielsen","Frank","0169364089");
        array[1]=new Person("Nelson","Franck","04227745221");
        array[2]=new Person("Durand","Paul","0381846245");
        array[3]=new Person("Dupond","Jean","0256234512");
        array[4]=new Person("Tanaka","Ken","+81 3 1234 4567");

        System.out.println(LinearSearch("Durand","Paul"));
        System.out.println(LinearSearch("Tanaka","Ken"));
        System.out.println(LinearSearch("Durand","Jean"));
    }
}
```

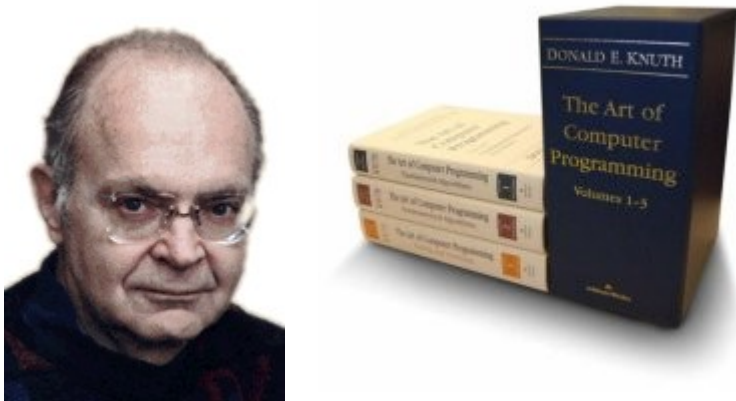


```
0381846245
+81 3 1234 4567
Not found in my dictionary
Press any key to continue...
```



Complexity of sequential search

- In the **worst case**, we need to **scan the full array**
- On **average**, we scan **half the array size** (if object is inside)



1	a[0]
2	a[1]
4	a[2]
8	a[3]
16	a[4]

Donald Knuth. The Art of Computer Programming, Volume 3: Sorting and Searching, Third Edition. Addison-Wesley, 1997. ISBN 0-201-89685-0. Section 6.1: Sequential Searching, pp.396–408.

Adding an element (object)

- To add dynamically an element, we first build a bigger array
- Then we use an index for storing the position of the last stored element
- When added an element, we increment the position of the 'end of array'
- Of course, we should check whether the array there is overflow or not
- Notice that it is difficult to erase elements...

Adding an element to the corpus

```
public static final int MAX_ELEMENTS=100;
static int nbelements=0;
static Person[] array=new Person[MAX_ELEMENTS];

static String LinearSearch(String lastname, String firstname)
{
    for(int i=0;i<nbelements;i++)
    {
        if ((lastname.equals(array[i].lastname) &&
            (firstname.equals(array[i].firstname))))
            return array[i].phone;
    }
    return "Not found in my dictionary";
}

static void AddElement(Person obj)
{if (nbelements<MAX_ELEMENTS)
    array[nbelements++]=obj; // At most MAX_ELEMENTS-1 here
    // nbelements is at most equal to MAX_ELEMENTS now
}
```



Adding elements to the corpus

```
public static void main (String [] args)
{
    AddElement(new Person("Nielsen", "Frank", "0169334089"));
    AddElement(new Person("Nelson", "Franck", "04227745221"));
    AddElement(new Person("Durand", "Paul", "0381846245"));
    AddElement(new Person("Dupond", "Jean", "0256234512"));
    AddElement(new Person("Tanaka", "Ken", "+81 3 1234 4567"));

    System.out.println(LinearSearch("Durand", "Paul"));
    System.out.println(LinearSearch("Tanaka", "Ken"));
    System.out.println(LinearSearch("Durand", "Jean"));
    AddElement(new Person("Durand", "Jean", "0199989796"));
    System.out.println(LinearSearch("Durand", "Jean"));

}
```



```
0381846245
+81 3 1234 4567
Not found in my dictionary
0199989796
```

Check if a person already belongs to the array before adding it
Under or oversaturated memory
(not convenient for non-predictable array sizes)

Dichotomic search (in sorted arrays)

- Assume all elements (keys) are **totally sorted** in the array
- For example, in increasing order
$$\text{array}[0] < \text{array}[1] < \dots <$$

(any lexicographic order would do)
- To find information, use a **dichotomic search** on the key
- Compare the key of the **middle element** with the query key:
 - If identical, we are done since we found it
 - If the key is greater, we search in the rightmost array part
 - If the key is less, we search in the leftmost array part
 - If the range [left,right] is ≤ 1 and element not inside then...
.....conclude it is not inside the array



Dichotomic search/Bisection search

```
static int Dichotomy(int [] array, int left, int right, int key)
{
    if (left>right) return -1;

    int m=(left+right)/2; // !!! Euclidean division !!!

    if (array[m]==key) return m;
    else
    {
        if (array[m]<key) return Dichotomy(array,m+1, right, key);
        else return Dichotomy(array,left,m-1, key);
    }
}
```

Think recursion!

```
static int DichotomicSearch(int [] array, int key)
{
    return Dichotomy(array,0,array.length-1, key);
}
```

Result is the index of the element position
Does it always terminate? Why?



Dichotomic search/Bisection search

```
public static void main (String[] args)
{
    int [] v={1, 6, 9, 12 , 45, 67, 76, 80, 95};

    System.out.println("Seeking for element 6: Position "+DichotomicSearch(v,6));
    System.out.println("Seeking for element 80: Position "+DichotomicSearch(v,80));
    System.out.println("Seeking for element 33: Position "+DichotomicSearch(v,33));
}
```

1, 6, 9, 12 , 45, 67, 76, 80, 95

9 elements [0,8] m=4 6<45

1, 6, 9, 12

4 elements [0,3] m=1 6=6 Found it return 1

```
Seeking for element 6: Position 1
Seeking for element 80: Position 7
Seeking for element 33: Position -1
```



Complexity of bisection search

We halve by 2 the array range at every recursion call

$n \Rightarrow n/2 \Rightarrow (n/2)/2 \Rightarrow ((n/2)/2)/2 \Rightarrow \dots \text{etc} \dots \Rightarrow 1 \text{ or } 0$

How many recursion calls?

$$2^k = n \Rightarrow k = \log(n)$$

(Execution stack (memory) is also of order $\log n$)

Big O(.)-notation of **ALGORITHMS** & complexity

Requires a **logarithmic number of operations**: $O(\log n)$

$t(1024) = 10$,

$t(2048) = 11$,

$t(65536) = 16$

$\Rightarrow$ (exponentially more) Efficient compare to sequential search

!!! SORTING HELPS A LOT !!!



```

static int Dichotomy(int [] array, int left, int right, int key)
{
    if (left>right)
        {throw new RuntimeException("Terminal state reached");
          //return -1;
         }

    int m=(left+right)/2;

    if (array[m]==key)
        {throw new RuntimeException("Terminal state reached");
          //return m;
         }
    else
    {
        if (array[m]<key) return Dichotomy(array,m+1, right, key);
        else return Dichotomy(array,left,m-1, key);
    }
}

```

```

Exception in thread "main" java.lang.RuntimeException: Terminal state reached
    at DichotomySearch.Dichotomy(dichotomysearch.java:14)
    at DichotomySearch.Dichotomy(dichotomysearch.java:20)
    at DichotomySearch.DichotomicSearch(dichotomysearch.java:27)
    at DichotomySearch.main(dichotomysearch.java:34)

```


Sorting: A fundamental procedure

- Given an (unsorted) array of elements
- We are asked to get a **sorted array** in, say **increasing order**
- The only primitive operations (basic operations) are
 - **comparisons** and
 - element **swappings**

```
static boolean GreaterThan(int a, int b)
{return (a>b);}
```

```
static void swap (int [] array, int i, int j)
{
int tmp=array[i];
array[i]=array[j];
array[j]=tmp;
}
```



Sorting: Sorting by **selection**

Many different sorting techniques (endless world)

- First, seek for the **smallest element** of the array
(=SELECTION)
- Put it at the front (using a swapping operation)
- Iterate for the remaining part:
`array[1], ..., array[n-1]`

=> we seek for the smallest elements for all (sub)arrays
`array[j], array[j+1], ..., array[n-1]`



Sorting by **selection**

Pseudo-code: 2 NESTED LOOPS

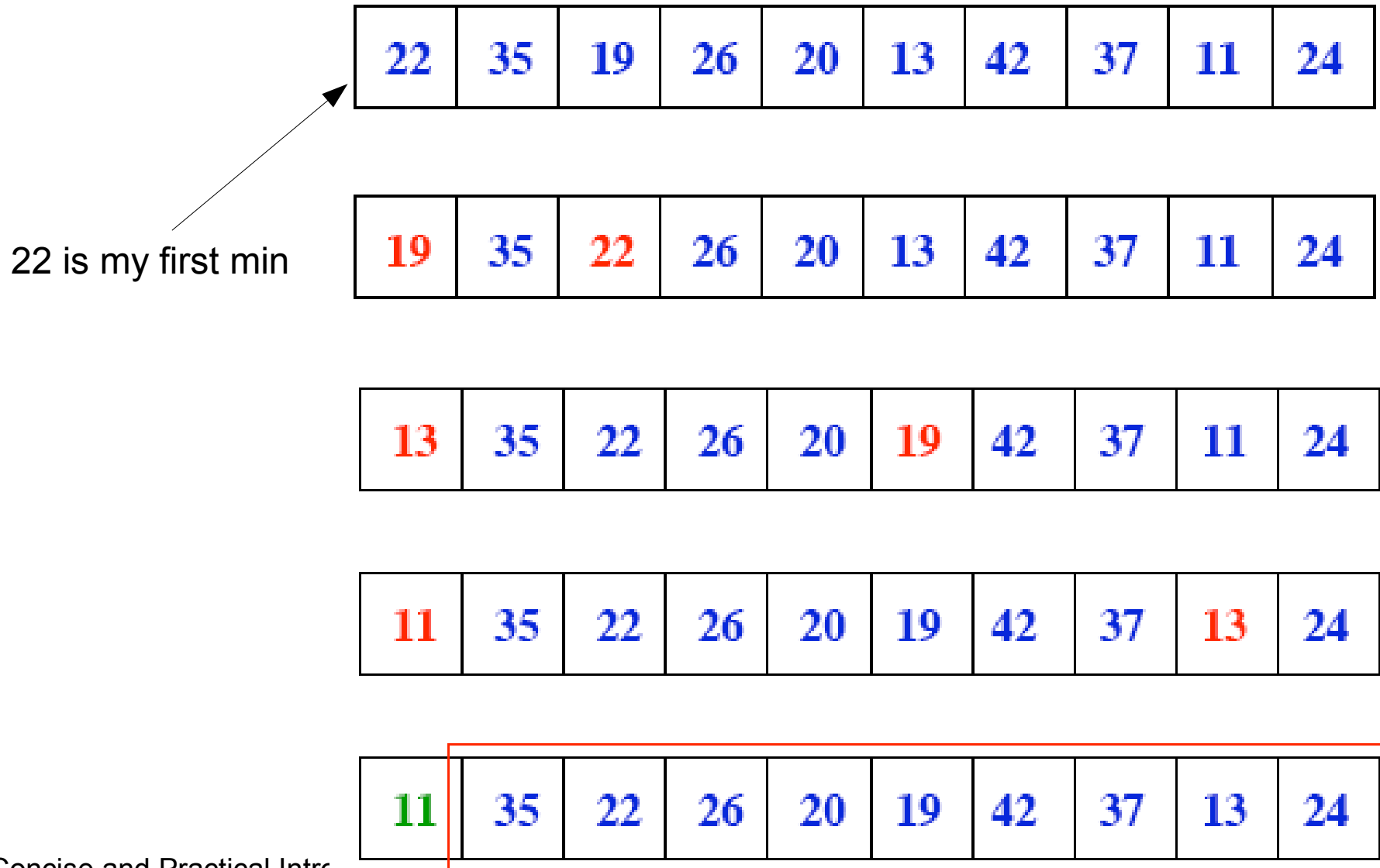
```
for i ← 0 to n-2 do
  min ← i
  for j ← (i + 1) to n-1 do
    if A[j] < A[min]
      min ← j
  swap A[i] and A[min]
```

```

for i ← 0 to n-2 do
  min ← i
  for j ← (i + 1) to n-1 do
    if A[j] < A[min]
      min ← j
  swap A[i] and A[min]

```

The inner loop (i=0)

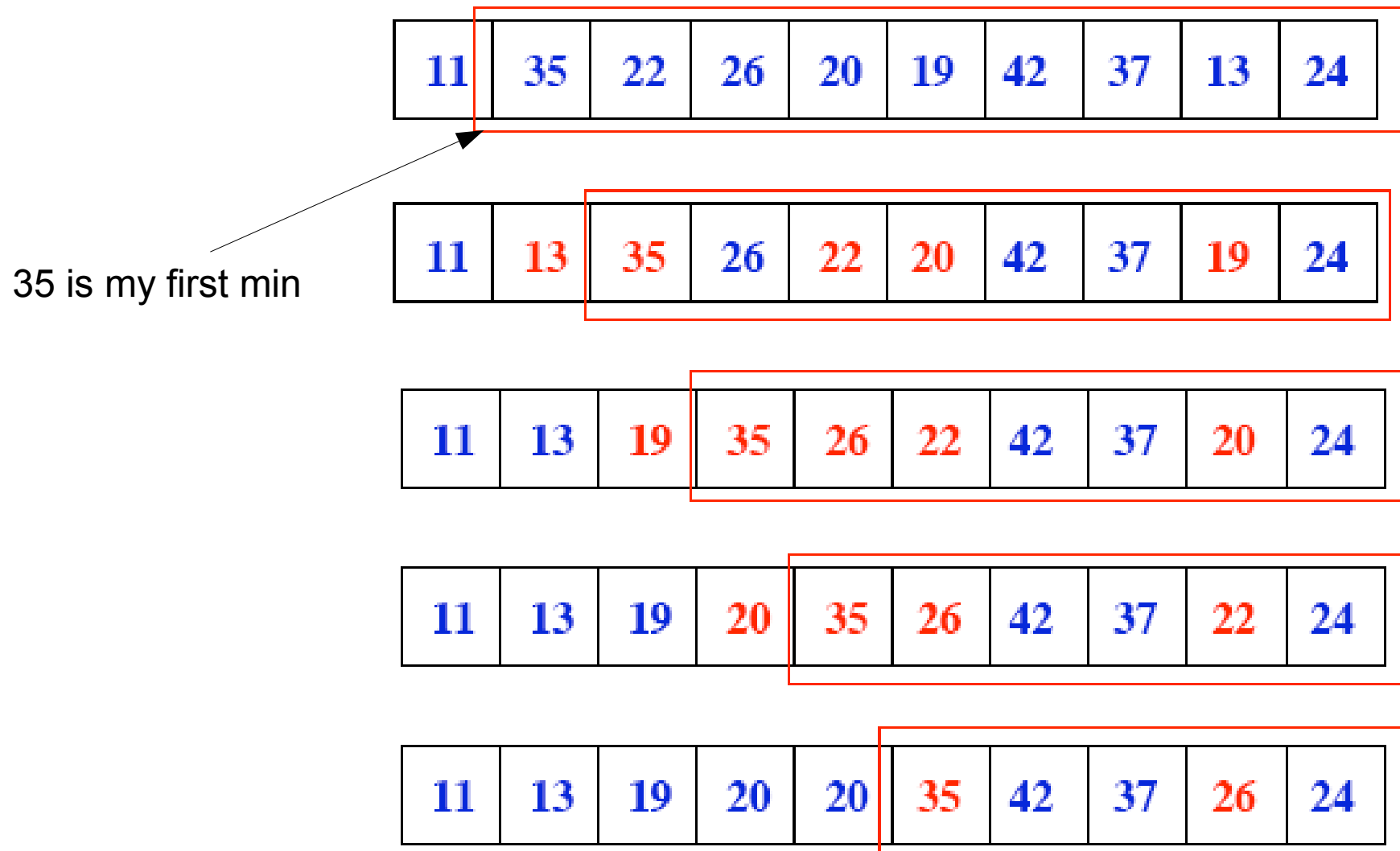


```

for i ← 0 to n-2 do
  min ← i
  for j ← (i + 1) to n-1 do
    if A[j] < A[min]
      min ← j
  swap A[i] and A[min]

```

The inner loop (i=1)



```

for i ← 0 to n-2 do
  min ← i
  for j ← (i + 1) to n-1 do
    if A[j] < A[min]
      min ← j
  swap A[i] and A[min]

```

The outer loop:
min of respective subarrays

11	13	19	20	22	24	42	37	35	26
----	----	----	----	----	----	----	----	----	----

11	13	19	20	22	24	26	42	37	35
----	----	----	----	----	----	----	----	----	----

11	13	19	20	22	24	26	35	42	37
----	----	----	----	----	----	----	----	----	----

11	13	19	20	22	24	26	35	37	42
----	----	----	----	----	----	----	----	----	----

Sorting by selection

Two nested loops:

- Select and put the smallest element in front of the array (inner loop)
- Repeat for all subarrays (at the right)

```
static boolean GreaterThan(int a, int b)
{return (a>b);}
```

```
static void swap (int [] array, int i, int j)
{int tmp=array[i];array[i]=array[j];array[j]=tmp;}
```

```
static void SelectionSort(int [] array)
{
int n=array.length;

for(int i=0;i<n-1;i++) {
    for(int j=i+1;j<n;j++) {
        if (GreaterThan(array[i],array[j]))
            swap(array,i,j);
    }
}
```



Sorting by selection

```
public static void main(String[] args)
{
    int [] array={22,35,19,26,20,13,42,37,11,24};

    SelectionSort(array);

    for(int i=0;i<array.length;i++)
        System.out.print(array[i]+" ");
    System.out.println("");
}
```



11 13 19 20 22 24 26 35 37 42

From sorting integers to objects

Generic algorithms for sorting any kind (=type) of elements

- Array `array` contains objects of the same type
- Define a predicate that returns whether an object is greater than another one or not
- Adjust the swap procedure

In Java 1.5, there are generics, beyond the scope of this course



For example, sorting events

```
class EventObject
{
    int year, month, day;
    EventObject(int y, int m, int d)
    {year=y; month=m; day=d; }

    static void Display(EventObject obj)
    {System.out.println(obj.year+"/"+obj.month+"/"+obj.day); }
}
```

```
static boolean GreaterThan(EventObject a, EventObject b)
{return (( a.year>b.year) ||
        ((a.year==b.year) && (a.month>b.month)) ||
        ((a.year==b.year) && (a.month==b.month) && (a.day>b.day)) ) ;}

static void swap (EventObject [] array, int i, int j)
{
    EventObject tmp=array[i];
    array[i]=array[j];
    array[j]=tmp;
}
```



sorting events

```
static void SelectionSort(EventObject [] array)
{
    int n=array.length;

    for(int i=0;i<n-1;i++)
        for(int j=i+1;j<n;j++)
            if (GreaterThan(array[i],array[j]))
                swap(array,i,j);
}

public static void main(String[] args)
{
    EventObject [] array=new EventObject[5];

    array[0]=new EventObject(2008,06,01);
    array[1]=new EventObject(2005,04,03);
    array[2]=new EventObject(2005,05,27);
    array[3]=new EventObject(2005,04,01);
    array[4]=new EventObject(2005,04,15);

    SelectionSort(array);
    for(int i=0;i<array.length;i++)
        EventObject.Display(array[i]);
        System.out.println("");
}
```

2005/4/1
2005/4/3
2005/4/15
2005/5/27
2008/6/1



Worst-case complexity of selection sort

Quadratic time: $O(n^2)$

= number of elementary operations

= number of comparisons+swap operations

...Try sorting a **reversed** sorted array...

16, 15, 14, 13, 12, 11, 10, 9, 8, 7, 6, 5, 4, 3, 2, 1
1, 16, 15, 14, 13, 12, 11, 10, 9, 8, 7, 6, 5, 4, 3, 2
1, 2, 16, 15, 14, 13, 12, 11, 10, 9, 8, 7, 6, 5, 4, 3

...

$nb = 2 * \text{array.length} * (\text{array.length} - 1) / 2$
(2 times n choose 2)



```

class SelectionSort
{
    static int nboperations;

    static boolean GreaterThan(int a, int b)
    {nboperations++; return (a>b);}

    static void swap (int [] array, int i, int j)
    {
        nboperations++;
        int tmp=array[i];array[i]=array[j];array[j]=tmp;}

    public static void main(String[] args)
    {
        int [] array={16,15,14,13,12,11,10,9,8,7,6,5,4,3,2,1};
        nboperations=0;

        SelectionSort(array);

        System.out.println("Number of operations:"+nboperations);

        int nb=2*array.length*(array.length-1)/2;
        System.out.println("Number of operations:"+nb);
    }
}

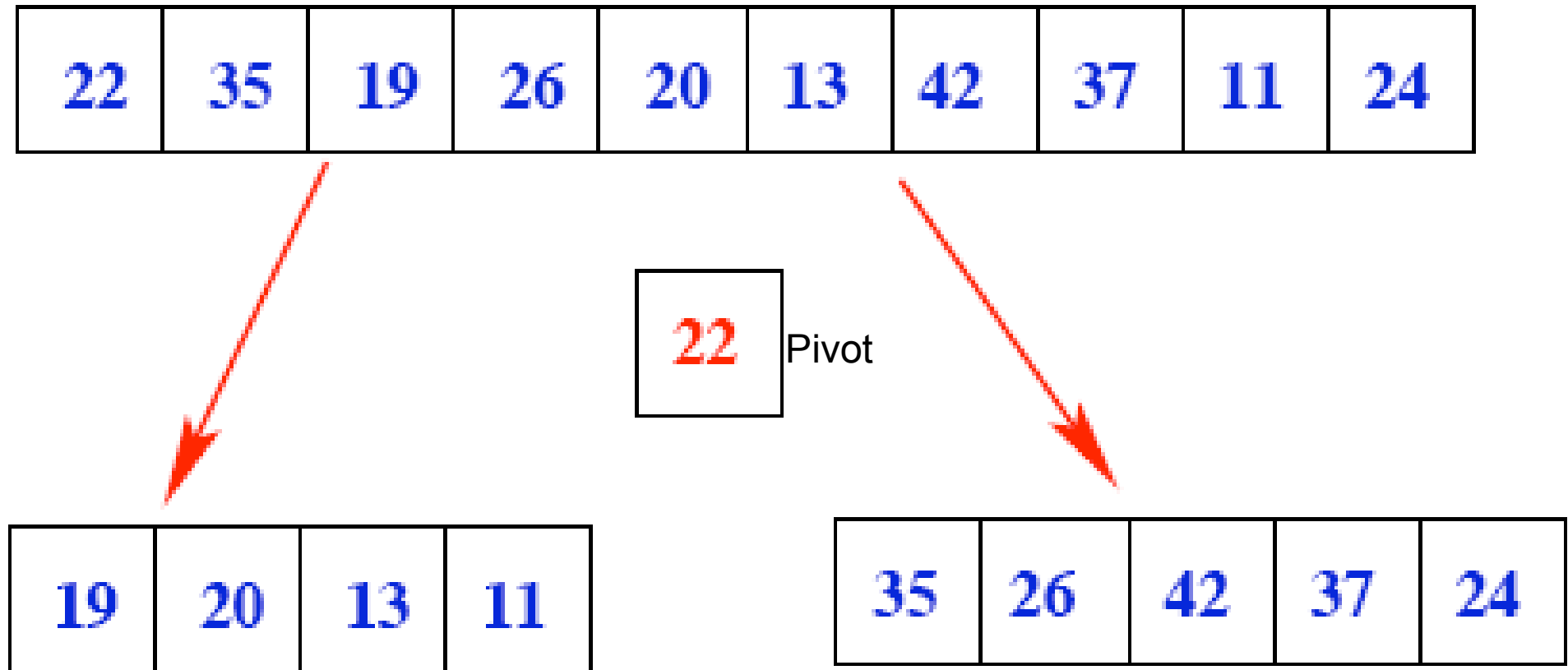
```



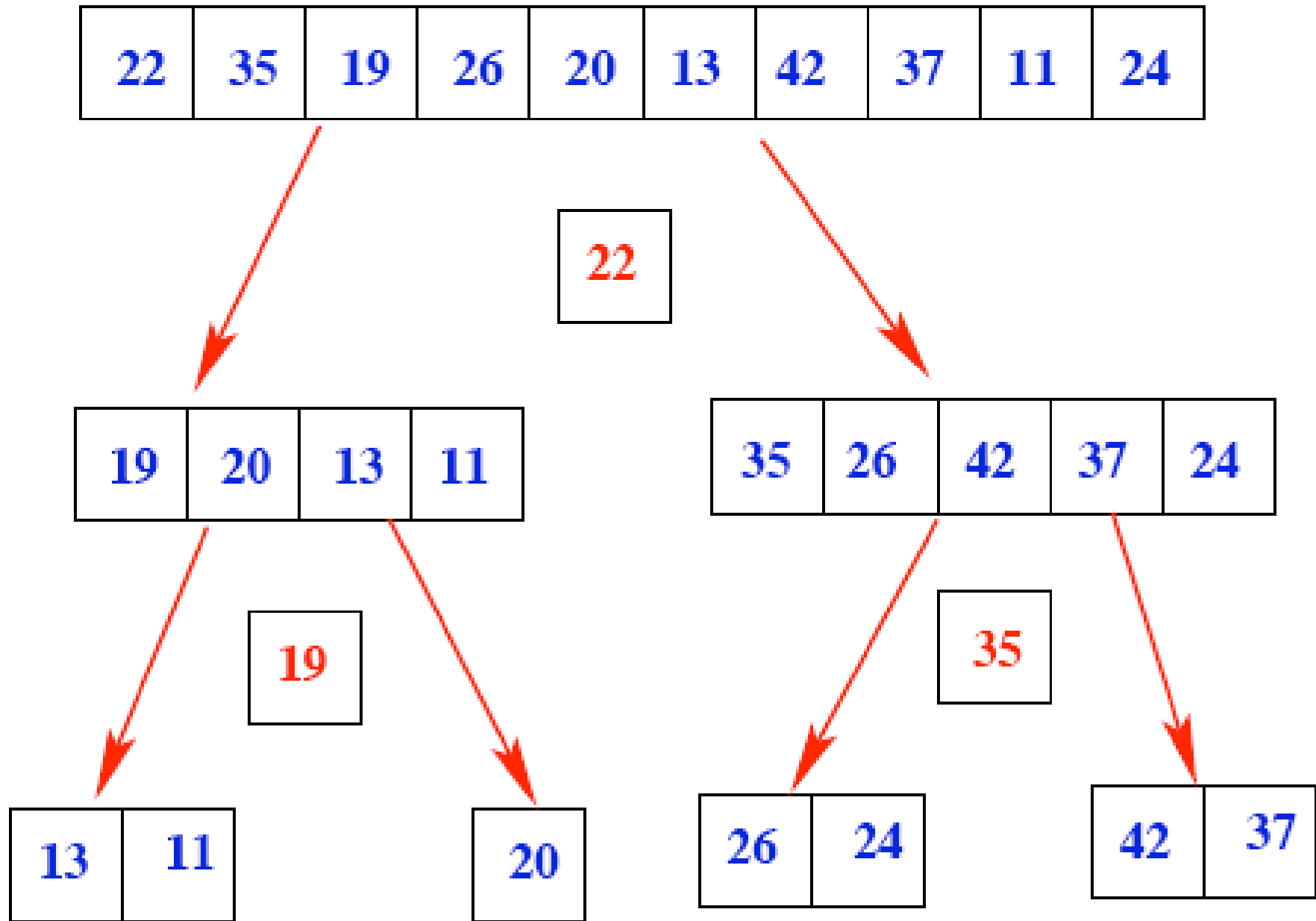
- Classify the elements according to `array[0]` (the **pivot**):
 - those *smaller than* the pivot: `arrayleft`
 - those *greater than* the pivot: `arrayright`
 - those *equal to* the pivot (in case of ties): `arraypivot`
- **Solve recursively** the sorting in `arrayleft` / `arrayright`, and recompose as

`arrayleft arraypivot arrayright`
- Note that a single element is sorted (= **terminal case**)

A recursive approach: **Quicksort**



Partition the array into 3 pieces




```

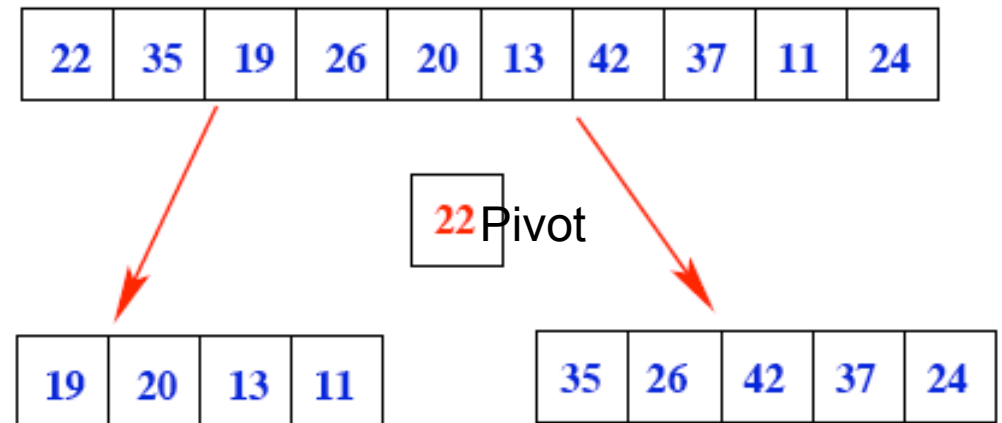
static int partition(int [] array, int left, int right)
{
    int m=left;// pivot

    for(int i=left+1; i<=right;i++)
    {
        if (GreaterThan(array[left],array[i]))
        {
            swap(array,i,m+1);
            m++;// we extend left side array
        }
    }

    // place pivot now
    swap(array,left,m);

    return m;
}

```



Sorting is **in-place**
(require no extra memory storage)

```
static void quicksort(int [] array, int left, int right)
{
    if (right>left)
    {
        int pivot=partition(array,left,right);
        quicksort(array,left,pivot-1);
        quicksort(array,pivot+1,right);
    }
}
```

```
static void QuickSort(int [] array)
{
    quicksort(array,0, array.length-1);
}
```



Quicksort: Analyzing the running time

- **Worst-case** running time is quadratic $O(n^2)$
- **Expected** running time is $O(n \log n)$
- Best case is linear $O(n)$ (all elements are identical)

<http://www.cs.ubc.ca/~harrison/Java/sorting-demo.html>

- **Lower bound** is $n \log n$: Any algorithm sorting n numbers require $n \log n$ comparison operations

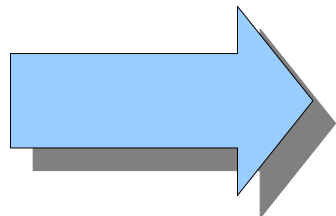
Las Vegas complexity analysis, tail bounds (how much deviation), etc...



Hashing: Principle and techniques

- Hashing: **fundamental technique** in computer science
(see INF 421.a')
- Store object x in position $h(x)$ (`int`) in an array
- Problem occurs if two objects x and y
are stored on the same cell: **Collision**.
- Finding good hashing functions that minimize collisions, and adopting a good search policy in case of collisions are key problems of hashing.

```
int i;  
array[i]
```



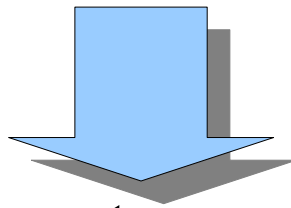
```
Object Obj=new Object();  
int i;  
i=h(Obj); // hashing function  
array[i]
```

Hashing functions

- Given a **universe X** of keys and for any x in X , we need to find an integer $h(x)$ **between 0 and m**
- Usually *easy to transform* the object into an integer:

For example, for strings just add the ASCII codes of characters

- The problem is then to transform
a set of n (sparse) integers

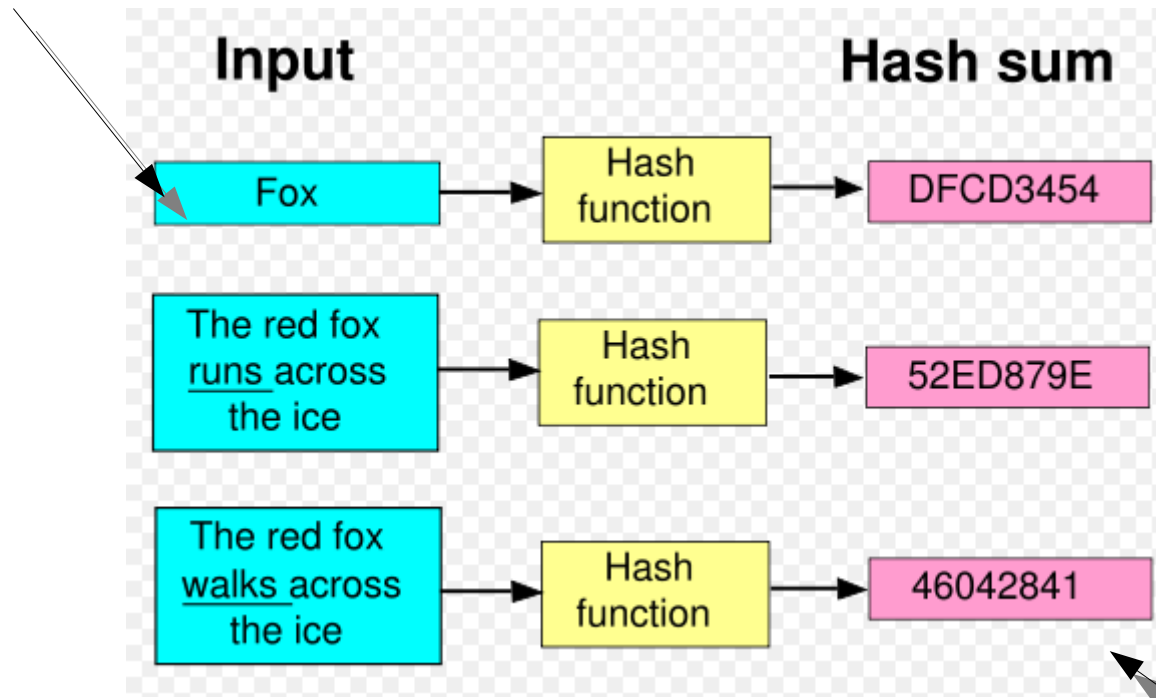


into a compact array of size $m \ll N$.

(<< means much less than)

Hashing strings

String lengths
may vary much



http://en.wikipedia.org/wiki/Hash_function

Constant-size
representation
(compact)



Hashing functions

Key idea is to take the modulo operation

$h(k) = k \bmod m$ where m is a **prime number**.

```
static int m=23;

static int String2Integer(String s)
{
    int result=0;

    for(int j=0;j<s.length();j++)
        result+=(int)s.charAt(j);

    return result;
}

// Note that m is a static variable
static int HashFunction(int l)
{return l%m;}
```



```

public static void main (String[] args)
{
String [] animals={"cat","dog","parrot","horse","fish",
"shark","pelican","tortoise", "whale", "lion",
"flamingo", "cow", "snake", "spider", "bee", "peacock",
"elephant", "butterfly"};

int i;
String [] HashTable=new String[m];

for(i=0;i<m;i++)
    HashTable[i]=new String("-->");

for(i=0;i<animals.length;i++)
    {int pos=HashFunction(String2Integer(animals[i]));
    HashTable[pos]+=(" "+animals[i]);
    }

for(i=0;i<m;i++)
    System.out.println("Position "+i+"\t"+HashTable[i]);

}

```



Position 0	--> whale
Position 1	--> snake
Position 2	-->
Position 3	-->
Position 4	-->
Position 5	-->
Position 6	-->
Position 7	--> cow
Position 8	--> shark
Position 9	-->
Position 10	-->
Position 11	-->
Position 12	--> fish
Position 13	--> cat
Position 14	-->
Position 15	--> dog tortoise
Position 16	--> horse
Position 17	--> flamingo
Position 18	-->
Position 19	--> pelican
Position 20	--> parrot lion
Position 21	-->
Position 22	-->



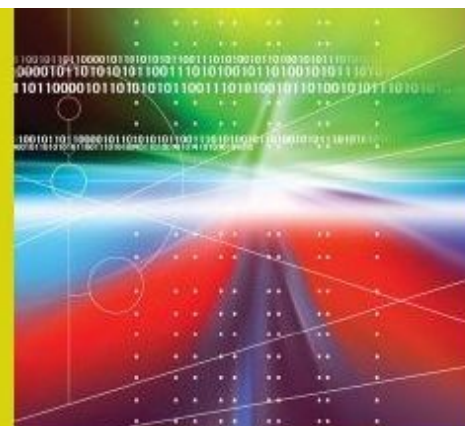


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Frank Nielsen

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