
Public-Key Cryptosystem Based on Isogenies

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Quantum Computer

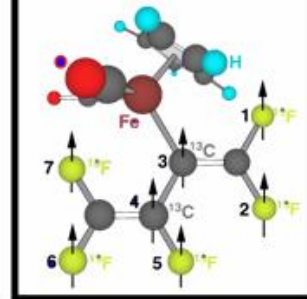
Public-key
cryptosystems



Shor's
algorithm



Quantum
computer



Breaking
with polynomial
complexity



Basic conceptions

- Non-supersingular elliptic curves over a finite field F_p : $Y^2 = X^3 + aX + b$; $j \neq 0, 1728$
- $\pi^2 - t\pi + p = 0$ - a Frobenius equation
- $D_\pi = t^2 - 4p$ - a Frobenius discriminant
- Isogenous elliptic curves
- Isogeny degree
- Isogeny kernel
- Modular equation: $\Phi_l(S, T) = 0$

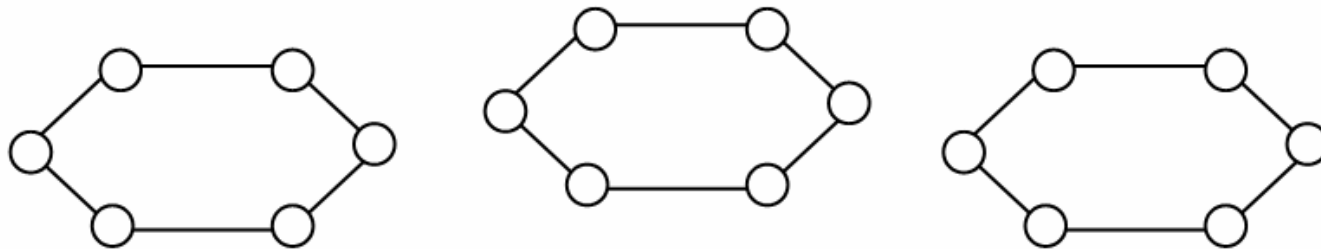
Branchless Cycles

- Elkies criterion: for an elliptic curve given, if

$$\left(\frac{D_\pi}{l}\right) = 1,$$

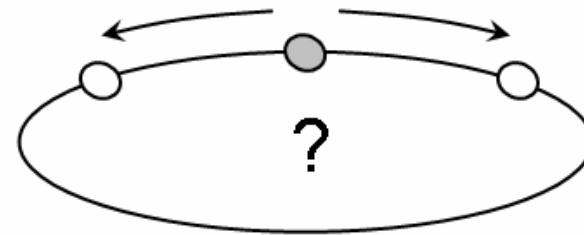
then there are two l -isogenous elliptic curves over F_p

- Isogenies of an Elkies degree form branchless cycles:



Direction Determination

- Frobenius equation for points of order l :
 $\pi^2 - t\pi + p = 0 \pmod{l}$



- $\left(\frac{t^2 - 4p}{l}\right) = 1 \Rightarrow$ there are 2 roots: π_1, π_2 over F_l -
– the Frobenius eigenvalues
- Action of the Frobenius endomorphism on an isogeny kernel is equivalent to multiplication of points by an eigenvalue [Elkies 1998]:
 $(X^p, Y^p) = \pi \cdot (X, Y)$ in $F_p[X, Y] / (Y^2 - X^3 - aX - b, H(X))$

Directed Step

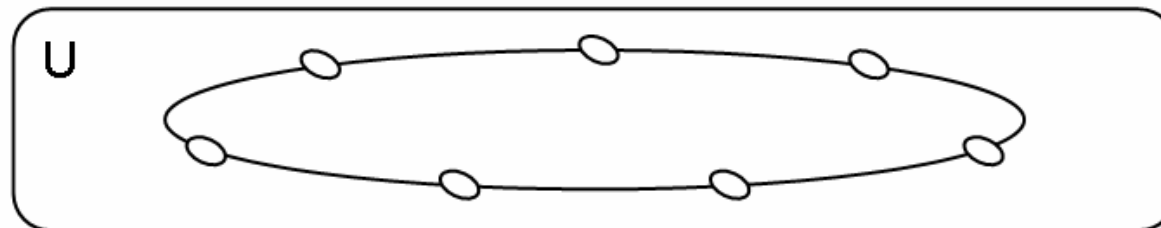
Input: field F_p , curve E , degree l , direction π

Algorithm:

- Find a root j_1 of $\Phi_l(j, T) = 0$ over F_p
- Compute an isogenous elliptic curve E_1
- Compute the polynomial $H_1(X)$ which determines the isogeny kernel
[Müller 1995]
- Check whether $(X^p, Y^p) = \pi \cdot (X, Y)$
in $F_p[X, Y] / (Y^2 - X^3 - aX - b, H_1(X))$
If not, then compute E_2 using the root j_2

Cycle of Prime Length

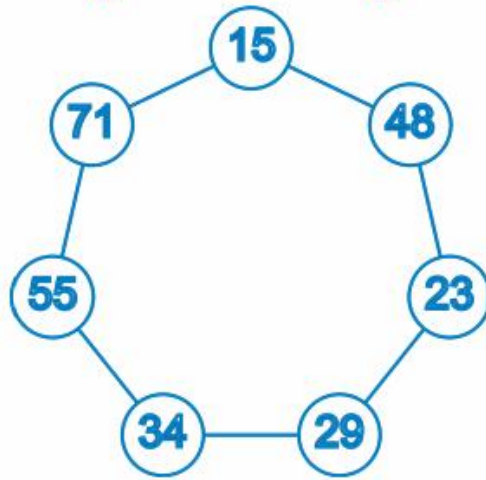
- U - a set of isogenous elliptic curves over F_p
- $\#U = H(D_\pi)$ - a class number [Schoof 1987]
- Practical observation:
 $\#U$ is prime \Rightarrow single isogeny cycle



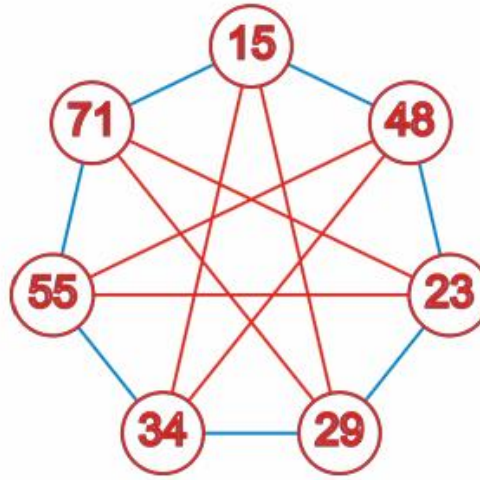
Isogeny Star

Example over F_{83} :

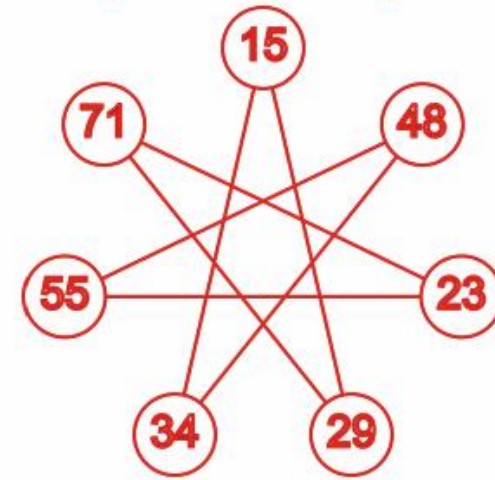
Isogenies of degree 3



Star



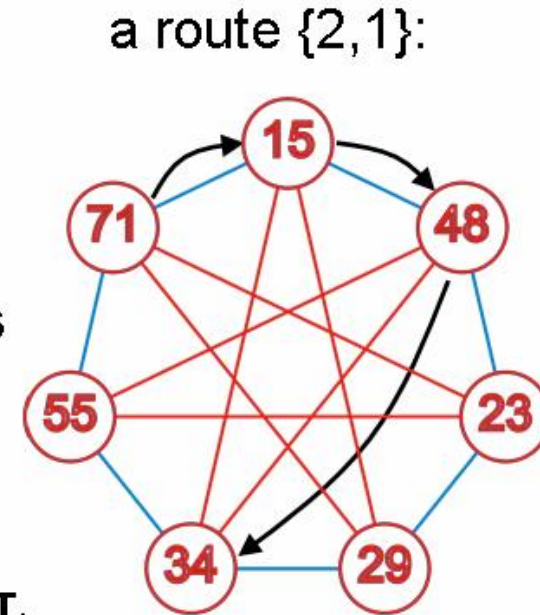
Isogenies of degree 5



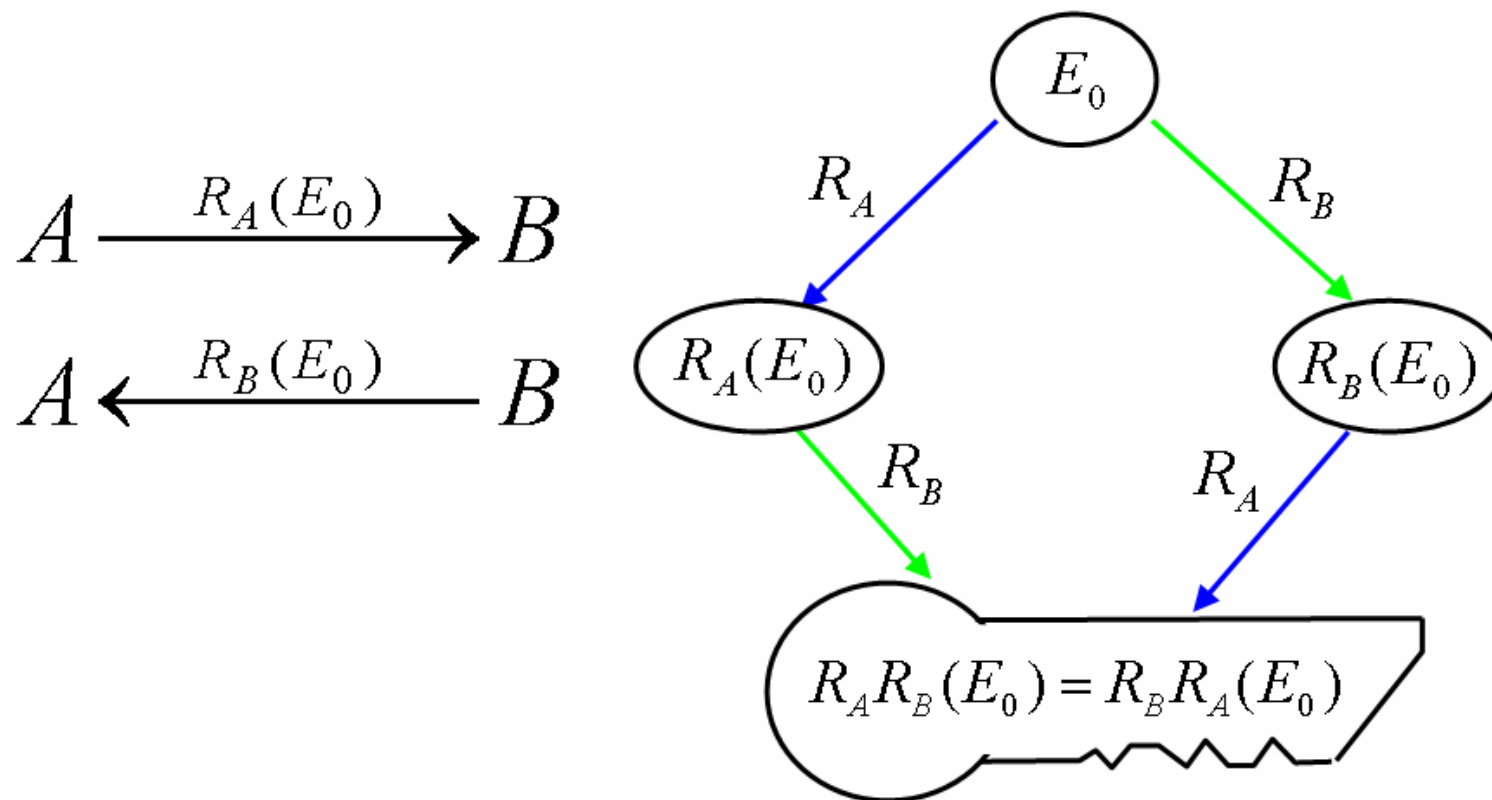
A graph of prime number of elliptic curves,
connected by isogenies of Elkies degrees

Route on Star

- For given
 - F_p – a finite field
 - E – an elliptic curve in a star
 - $\{l_i\}$ – a set of isogeny degrees
 - $\{\pi_i\}$ – a set of positive directions
- A route is a set $R=\{r_i\}$,
where r_i is a number of steps
by l_i -isogeny in the direction π_i
- Routes are commutative: $R_A R_B = R_B R_A$



Key Agreement



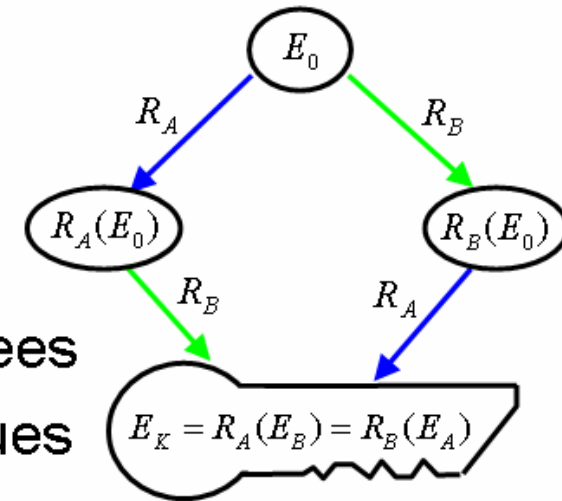
Key Agreement – Algorithm

Common parameters:

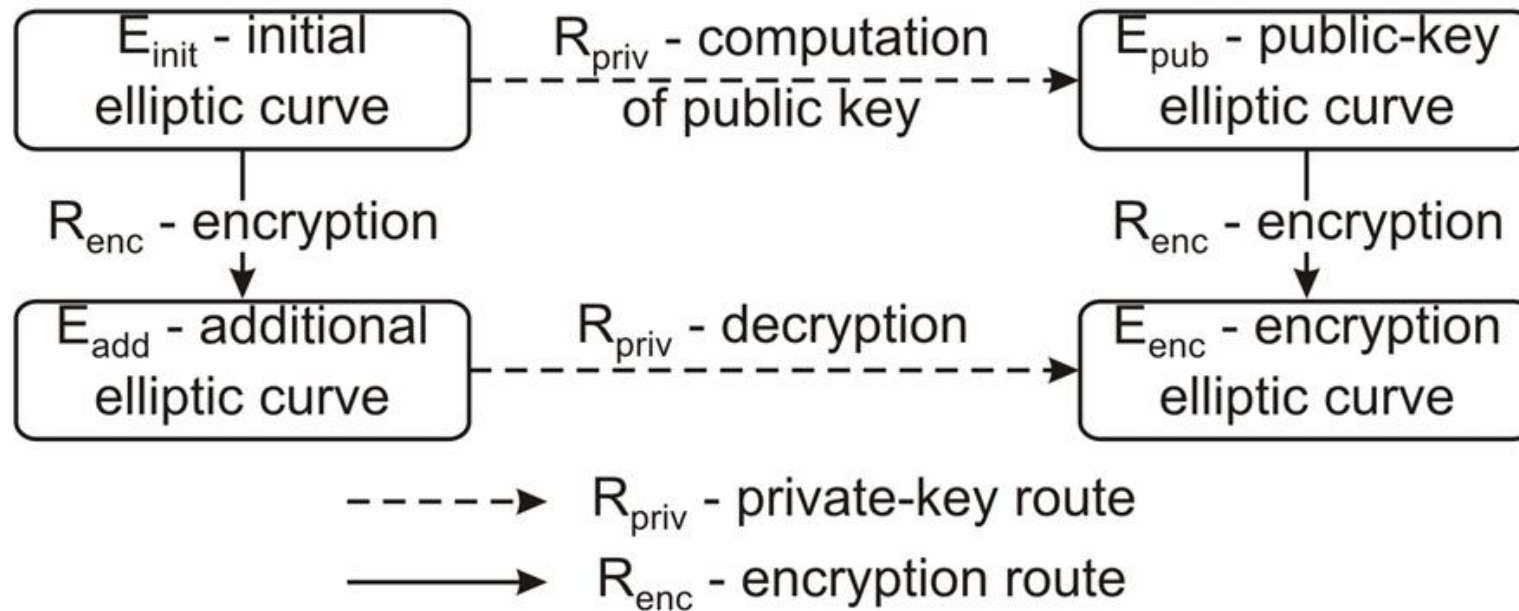
- F_p – a finite field
- E_0 – an initial elliptic curve
- $\{l_i\}$ – a set of Elkies isogeny degrees
- $\{\pi_i\}$ – a set of Frobenius eigenvalues

Algorithm:

- A randomly chooses a route R_A and sends $E_A = R_A(E_0)$
- B randomly chooses a route R_B and sends $E_B = R_B(E_0)$
- A computes $E_K = R_A(E_B)$, B computes $E_K = R_B(E_A)$
- Resulting key is the j-invariant of E_K



Public-Key Encryption



Security

- Problem of searching for a route between elliptic curves
- Solving methods on an $\#U$ -curves star:
 - Brute-force: $O(\#U)$ isogenous steps
 - Meet-in-the-middle: $O(\sqrt{\#U})$ isogenous steps
 - Others - ?

Quantum Computer Resistance

- An algorithm of a route search requires a subroutine, which calculates a chain of isogeny steps
- Calculation of an isogeny chain requires consecutive solving of modular equation $\Phi_l(j, T) = 0$, where j is being changed with every step
- Leads to exponential time of the algorithm

Complexity and Sizes

- Key agreement complexity:
 - $O(\log \#U)$ isogeny steps, or
 - $O(\log^4 p)$ field operations
- Consuming operations:
 - $X^p \bmod H(X)$
 - solving of $\Phi_l(j, T) = 0$
- For 2^{80} secrecy:
 - field characteristic: $p \sim 2^{320}$
 - star size $\sim 2^{160}$
 - number of isogeny degrees ~ 40
 - steps per degree: $0 \dots \pm 8$

Parameters Selection

- Obtaining a large prime $\#U$ is very complicated
- Hypothesis: $\#U$ must have a large prime divisor
- Choose $D_\pi = D f^2$, where f is a large prime conductor and $h(D)$ is small. Then [Cohen, 1996]

$$h_{D_\pi} = h_D \cdot \left(f - \left(\frac{D}{f} \right) \right) = h_D \cdot (f \pm 1)$$

Choose f such that $\frac{f \pm 1}{2}$ is prime

Test Implementation

- Mathematica 5.0
- $F_{2038074743}$
- Star of 55103 elliptic curves (prime), chosen by direct computation of a class number
- 6 isogeny degrees: {3, 5, 7, 11, 13, 17}
- 0...9 steps per each isogeny degree

The End

A. Rostovtsev and A. Stolbunov
Public-Key Cryptosystem Based on Isogenies
<http://eprint.iacr.org/2006/145>