

Directed Algebraic Topology and Concurrency

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Modeling and Analysing Interaction between Systems Laboratory

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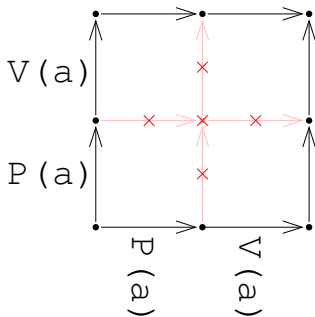
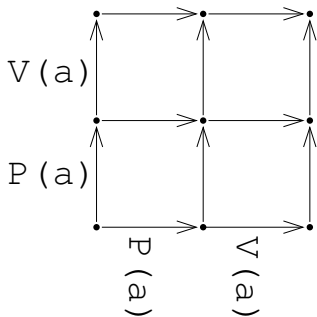
Underlying graph and Category of paths I

graph : 1-dimensional pre-simplicial set

$$\begin{array}{ccc} \begin{array}{c} A \\ \downarrow s \quad \downarrow t \\ V \end{array} & \begin{array}{ccc} A & \xrightarrow{\phi_1} & A' \\ \downarrow s \quad \downarrow t & & \downarrow s' \quad \downarrow t' \\ V & \xrightarrow{\phi_0} & V' \end{array} & \begin{array}{ccccc} A & \xrightarrow{\phi_1} & A' & \xrightarrow{\psi_1} & A'' \\ \downarrow s \quad \downarrow t & & \downarrow s' \quad \downarrow t' & & \downarrow s'' \quad \downarrow t'' \\ V & \xrightarrow{\phi_0} & V' & \xrightarrow{\psi_0} & V'' \end{array} \end{array}$$

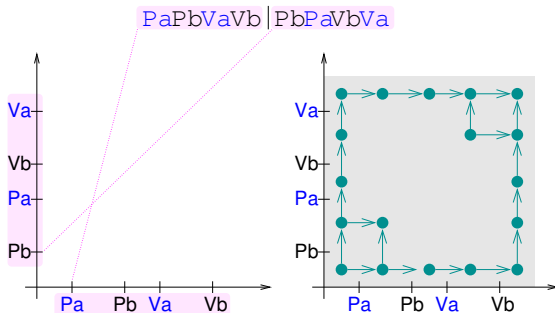
An example of model of a multi-task program

from *Edsger Wybe Dijkstra* "Pakken/Vrijlaten" language



Another example

from *Edsger Wybe Dijkstra* "Pakken/Vrijlaten" language

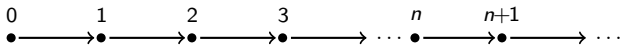


18 states and 20 arrows

Underlying graph and Category of paths II

adjunction between Cat and Grph

Path : morphism of graph from \mathbb{I}_n to Γ



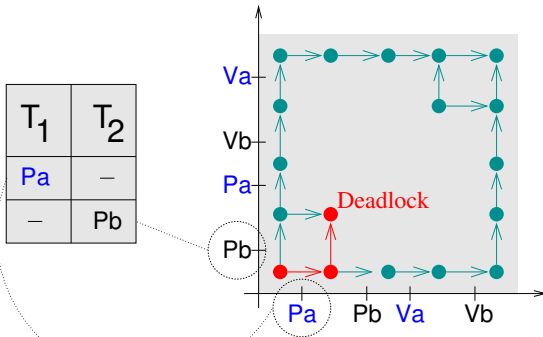
Forgetful functor $U : \text{Cat} \longrightarrow \text{Grph}$

“Category of paths” functor $F : \text{Grph} \longrightarrow \text{Cat}$

$$F \dashv U$$

A potential execution

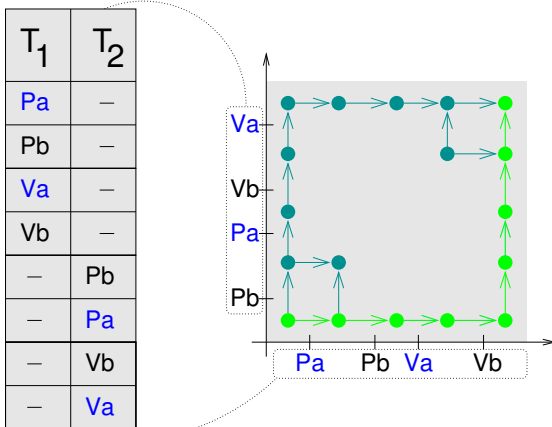
program $T_1 = PaPbVaVb \mid T_2 = PbPaVbVa$



Deadlock

Anoter potential execution

program $T_1 = PaPbVaVb \mid T_2 = PbPaVbVa$

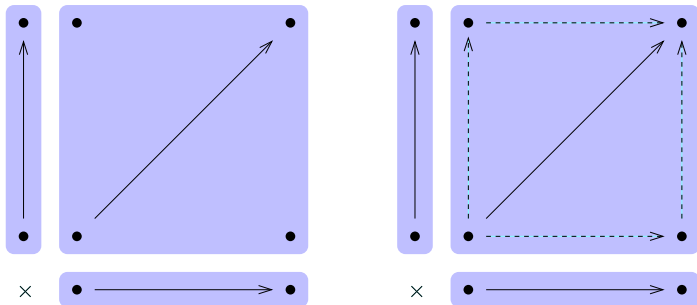


Termination

Underlying graph and Category of paths III

Cartesian products in Grph

$$F(\Gamma \times \Gamma') \not\cong F(\Gamma) \times F(\Gamma')$$



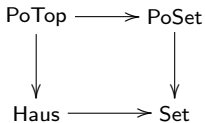
Transitions Systems, CCS/ π -calculus, Mazurkiewicz Traces ...

Partially ordered spaces

The category PoTop

Pospace $\vec{X} : \begin{cases} X & \text{topological space} \\ \sqsubseteq & \text{closed in } X \times X \end{cases}$

morphism f from \vec{X} to \vec{X}' : **continuous** and **order preserving** maps



Theorem

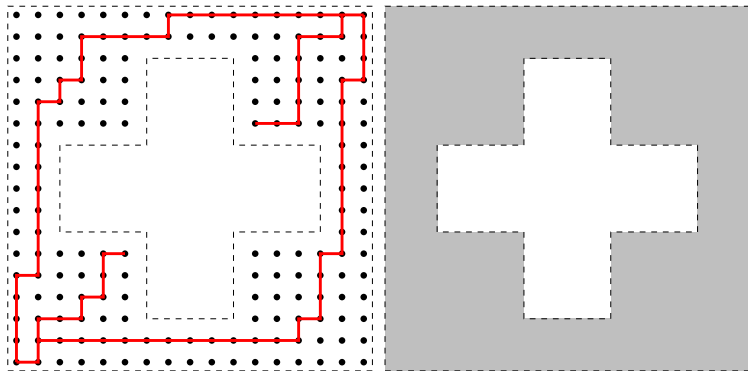
- *The directed compact unit segment is exponentiable in PoTop*
- *PoTop is complete and **cocomplete***
- *PoTop is symmetric monoidal closed*
- *CGPoTop is a Cartesian closed reflective subcategory of PoTop*
- *CPoTop is a (complete and cocomplete) Cartesian closed reflective subcategory of CGPoTop cogenerated by the directed compact unit segment*
- *PoTop has no loop*

Partially ordered spaces

examples

- Real line with standard order and topology : $\overrightarrow{\mathbb{R}}$
- Subset of a pospace (in particular $\overrightarrow{[0, 1]}$)
- Geometric realization of a graph
- Cartesian Product
- Closed subsets of a metric space together with inclusion

Size reduction



Graph $\Gamma_{\vec{X}}$

of paths on a pospace \vec{X}

- **paths** on \vec{X} : morphisms from $\overline{[0, 1]}$ to \vec{X}
- **arrows** of $\Gamma_{\vec{X}}$: paths on \vec{X}
- **source** and **target** of a path γ on, \vec{X} : $\gamma(0)$ and $\gamma(1)$

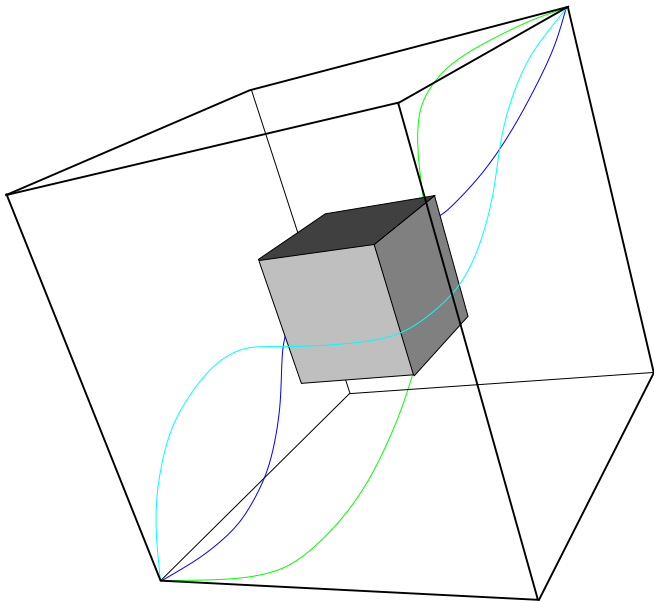
Image of a dipath

Singular facts about pospaces

- 1 The image of a dipath α on a pospace \vec{X} is either isomorphic (in PoSpc) to $\overrightarrow{[0, 1]}$ or $\{*\}$ (hence **no** directed *Peano* curve).
- 2 Two dipaths **sharing the same image** are **dihomotopic**.

Some paths around a cubic hole

$P(a).V(a) \mid P(a).V(a) \mid P(a).V(a)$ with $\alpha_a = 3$



Two “concatenations”

paths on $\Gamma_{\vec{X}}$ vs paths on \vec{X}

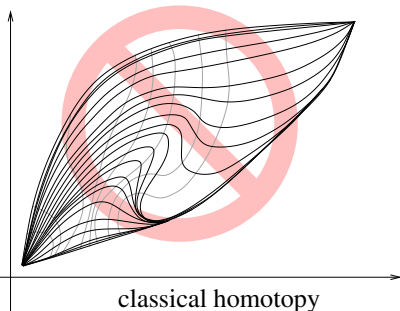
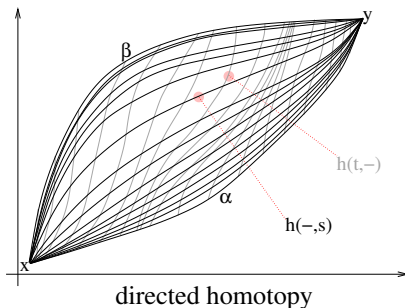
- Composition on $F(\Gamma_{\vec{X}})$ denoted by \circ
- Given $\gamma = (\gamma_n, \dots, \gamma_1)$ a path on $\Gamma_{\vec{X}}$, we define the following path on \vec{X}

$$(\nu(\gamma))(t) = \begin{cases} \gamma_k(nt - k) & \text{si } t \in [\frac{k}{n}, \frac{k+1}{n}[\text{ et } k < n - 1 \\ \gamma_n(nt - n + 1) & \text{si } t \in [\frac{n-1}{n}, 1] \end{cases}$$

Directed homotopy

what it is and looks like

Morphism h from $\overrightarrow{[0, 1]^2}$ to \overrightarrow{X} such that $U(h)$ is a homotopy from γ to δ



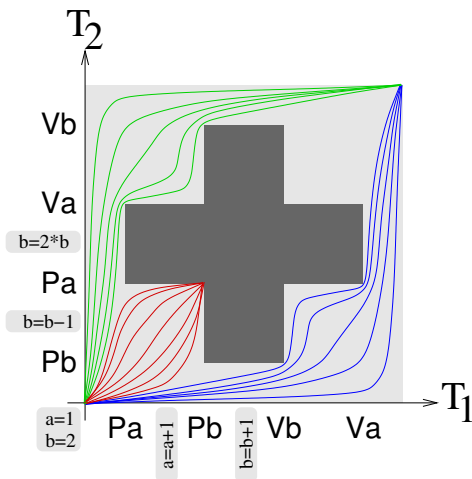
Directed homotopy

an example

T1 gets a and b before T2 $\Rightarrow a=2$ and $b=4$

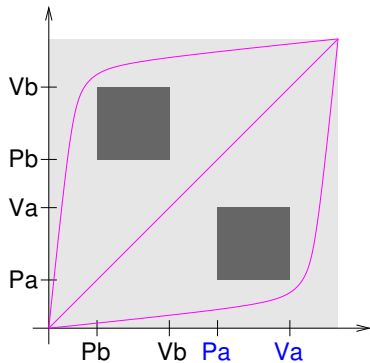
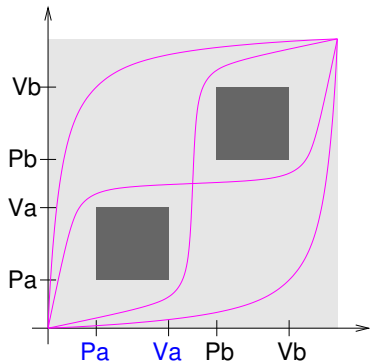
T2 gets b and a before T1 $\Rightarrow a=2$ and $b=3$

Each of T1 and T2 gets a resource
 \Rightarrow Deadlock with $a=2$ and $b=1$



A subtlety

directed homotopy is not classical homotopy



Loop-Free and One-Way categories

André Haefliger (scwol) / Dmitry Kozlov (acyclic) and Colin McLarty

Loop-Free category or small categories without loops (LfCat) :

$$\mathcal{C}[x, x] = \{\text{id}_x\} \text{ and } (\mathcal{C}[x, y] \times \mathcal{C}[y, x] \neq \emptyset \implies x = y)$$

One-Way category (OwCat) :

$$\mathcal{C}[x, x] = \{\text{id}_x\}$$

\mathcal{C} one-way $\iff \text{sk}(\mathcal{C})$ is loop-free

$$\text{LfCat} \hookrightarrow \text{OwCat} \hookrightarrow \text{Cat}$$

Let \sim be the congruence over $F(\Gamma_{\vec{X}})$ generated by

$$\left\{ ((\gamma_n, \dots, \gamma_1), (\delta_p, \dots, \delta_1)) \mid \text{there is a dihomotopy from } \nu(\gamma) \text{ to } \nu(\delta) \right\}$$

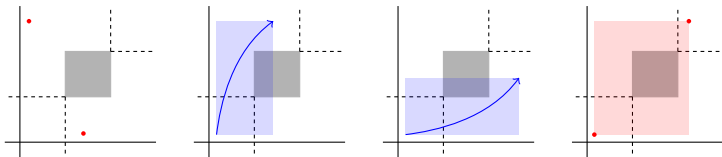
The fundamental category $\vec{\pi}_1(\vec{X})$ is $F(\Gamma_{\vec{X}})/\sim$ and we have

$$\vec{\pi}_1(\vec{X} \times \vec{Y}) \cong \vec{\pi}_1(\vec{X}) \times \vec{\pi}_1(\vec{Y})$$

$\vec{\pi}_1(\vec{X})$ is loop-free

van Kampen theorem

A detailed example

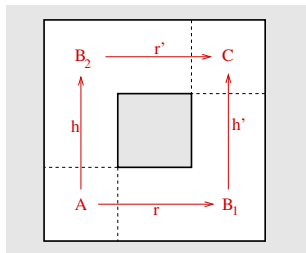


A detailed example

square with centered hole

$x \in$	$y \in$	$\vec{\pi}_1(\vec{X})[x, y]$
A	A	$\{\sigma_{x,y}\}$
B_1	B_1	$\{\sigma_{x,y}\}$
B_2	B_2	$\{\sigma_{x,y}\}$
C	C	$\{\sigma_{x,y}\}$
A	B_1	$\{r_{x,y}\}$
A	B_2	$\{h_{x,y}\}$
B_1	C	$\{h'_{x,y}\}$
B_2	C	$\{r'_{x,y}\}$
B_1	B_2	\emptyset
B_2	B_1	\emptyset
A	C	$\{u_{x,y}, d_{x,y}\}$

With $r'_{y,z} \circ h_{x,y} = u_{x,z}$, $h'_{y,z} \circ r_{x,y} = d_{x,z}$
 and 3 points x, y, z of the square
 such that $x \sqsubseteq y \sqsubseteq z$; if $x \not\sqsubseteq y$ then
 $\vec{\pi}_1(\vec{X})[x, y] = \emptyset$.



Yoneda morphism $\sigma \in \mathcal{C}[x, y]$

preserving the past and the future I

future if $\mathcal{C}[y, z] \neq \emptyset$, then $\mathcal{C}[y, z] \longrightarrow \mathcal{C}[x, z]$ is a bijection and

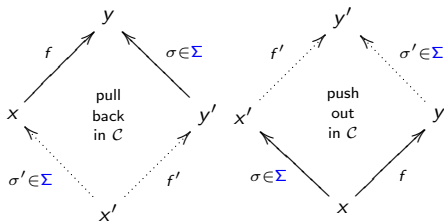
$$\gamma \longmapsto \gamma \circ \sigma$$

past if $\mathcal{C}[z, x] \neq \emptyset$, then $\mathcal{C}[z, x] \longrightarrow \mathcal{C}[z, y]$ is a bijection

$$\delta \longmapsto \sigma \circ \delta$$

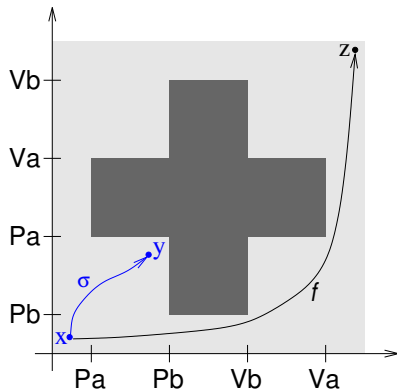
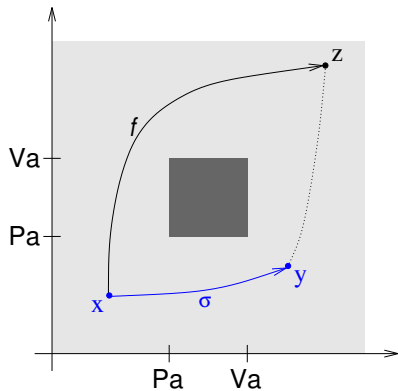
Yoneda system Σ of a small category \mathcal{C} preserving the past and the future II

- 1 Σ is stable under composition,
- 2 Σ contains all the isomorphisms of \mathcal{C} ,
- 3 all the elements of Σ are *Yoneda* morphisms and
- 4 Σ is stable under change and cochange of base.



Yoneda systems

Example



Structure of Σ -components

\mathcal{C} loop-free category and Σ Yoneda system over \mathcal{C}

Theorem

- 1 $\exists z \Sigma[x, z] \times \Sigma[y, z] \neq \emptyset$ iff $\exists z \Sigma[z, x] \times \Sigma[z, y] \neq \emptyset$
- 2 “ $\exists z \Sigma[x, z] \times \Sigma[y, z] \neq \emptyset$ ” defines an equivalence relation $x \sim y$
- 3 Given any \sim -equivalence class K , the full subcategory of \mathcal{C} whose set of objects is K is a non empty lattice
- 4 If $a \sim b$, then the following square is both a pullback and a pushout in \mathcal{C} .

$$\begin{array}{ccc} x & & a \longrightarrow a \vee b \\ \Sigma \uparrow & x \xrightarrow{\Sigma} z & \uparrow \quad \uparrow \\ z & & a \wedge b \longrightarrow b \\ \xrightarrow{\Sigma} y & \uparrow \Sigma & \end{array}$$

Locale of *Yoneda* systems

topology without point over a one-way category

Theorem

*The collection, ordered by inclusion, of the Yoneda systems of a one-way category, forms a **locale** whose maximum is denoted $\overline{\Sigma}$. Beside, its minimum is the collection of all isomorphisms of \mathcal{C} .*

Category of components

generalizing the set of arcwise components

The **category of components** of a loop-free category \mathcal{C} is the quotient $\mathcal{C}/\bar{\Sigma}$

Theorem

A loop-free category \mathcal{C} is a non empty lattice iff its category of components is $\{\}$*

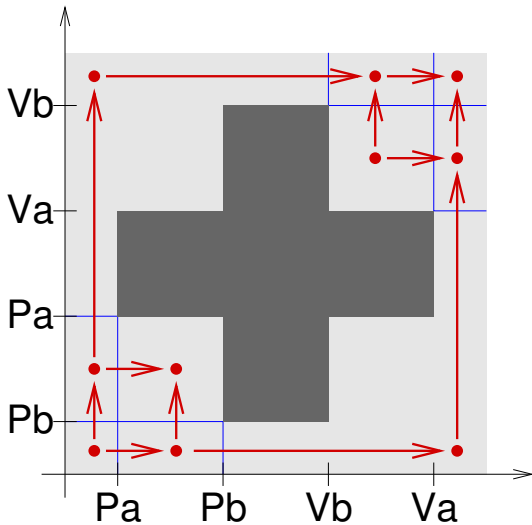
Fundamental theorem

\mathcal{C} one-way category and Σ Yoneda system over \mathcal{C}

Theorem

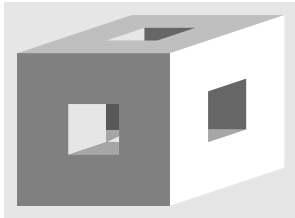
- 1 the collection Σ is *pure* in \mathcal{C} ($\beta \circ \alpha \in \Sigma \Rightarrow \beta, \alpha \in \Sigma$),
- 2 the category \mathcal{C}/Σ is *loop-free* and the category $\mathcal{C}[\Sigma^{-1}]$ is *one-way*
- 3 the categories $\mathcal{C}[\Sigma^{-1}]$ and \mathcal{C}/Σ are *equivalent* and
- 4 the category $\mathcal{C}[\Sigma^{-1}]$ is *fibered* over the base \mathcal{C}/Σ .

The category of components of the swiss flag

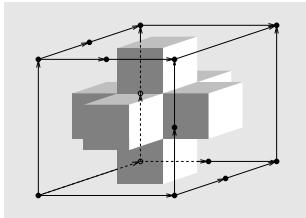


The category of components

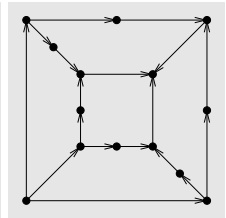
Menger sponge first iteration : $P(a) \cdot V(a) \mid P(a) \cdot V(a) \mid P(a) \cdot V(a)$ with $\alpha_a = 2$



Interior of the pospace



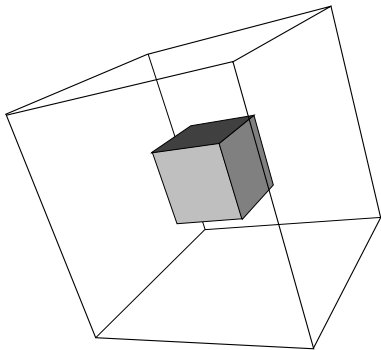
Category of components



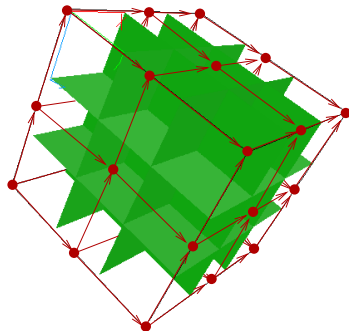
Flattened

The components category

of a 2-semaphore : $P(a) \cdot V(a) \mid P(a) \cdot V(a) \mid P(a) \cdot V(a)$ avec $\alpha_a = 3$



the pospace



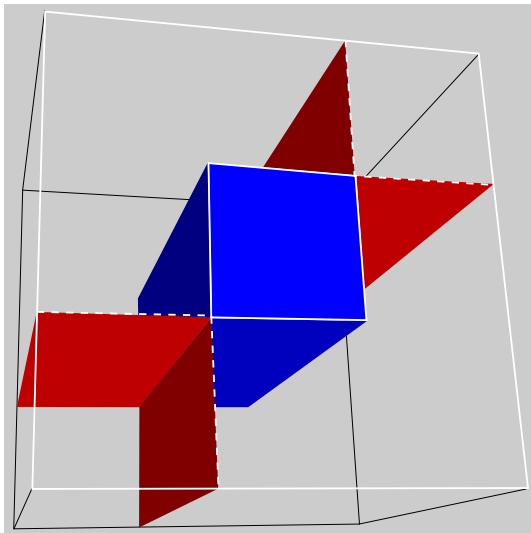
its category of components

Theorem (J. Hashimoto -T. Balabonski)

*The monoid of (isomorphism classes of) non empty, connected, finite, loop-free categories is countable and **free***

Example of product

parallel “independent” composition



The Directed Circle

Obects : S^1 Morphisms : $S^1 \times \mathbb{N} \times S^1$ Identities : $(x,0,x)$ for $x \in S^1$

