Mathematical Programming: Modelling and Applications

Leo Liberti

LIX, École Polytechnique

liberti@lix.polytechnique.fr

January 2018

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Outline

1 Kissing number problem

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Kissing number problem

Some examples

$$n = 6, K = 2$$
 $n = 12, K = 3$

more dimensions



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Use AMPL to implement a MINLP formulation of the Kissing Number Problem (as an **optimization problem**).

Hint 1: .dat files are given for small/medium instances. Hint 2: look at Maculan, Michelon, Smith' model (1995). Hint 3: you can use additional variables to reformulate the binary products. Hint 4: you can use the knpcheck.run script (which is given) to check your solution (in other words, you can include it at the bottom of your .run file). Hint 5: use Baron.

Set a threshold (for example, CPU Time limit 100 s., i.e. maxtime=100 in the options of your .run script) and test these instances:

(a) n=12 and k=3

(b) n=7 and k=2

In theory, your program should be able to accomplish (a) but not (b).

Then, download knp.mod and knp.run. Your implementation should be quite similar.

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Use AMPL to implement the Kissing Number Problem (as a decision problem).

Hint: look at Kucherenko, Belotti, Liberti, Maculan' model (2007).

Set again a threshold (for example, CPU Time limit 100 s., i.e. maxtime=100 in the options of your .run script) and test again these instances: (a) n=12 and k=3 (b) n=7 and k=2

In theory, the program should not be able to accomplish neither (a) nor (b). Then, download knpfeas18.mod and knpfeas18.run. Your implementation should be quite similar.

Change the constraint on α , to see if something changes.

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Reminder: SDP

SEMIDEFINITE PROGRAMMING (SDP) refers to optimization problems where decision variables are a symmetric matrix that is required to be semidefinite positive. The standard form is

$$\begin{array}{ccc} \min & C \bullet X \\ \forall i \leq m & A^i \bullet X &= b^i \\ & X \succeq 0, \end{array} \right\}$$
(1)

where:

- *C* and A^i (for $i \le m$) are $n \times n$ symmetric matrices
- X is an $n \times n$ matrix of decision variables,
- $b^i \in \mathbb{R}$ for all $i \leq m$
- for two $n \times n$ matrices $L = (\lambda_{ij}), M = (\mu_{ij})$ we have $L \bullet M = \sum_{i,j \le n} \lambda_{ij} \mu_{ij}$
- $X \succeq 0$, i.e. X is positive semidefinite

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Reminder: SDP

To make Eq. (1) clearer, write out the componentwise product • of the matrices: $C = (c_{jh}), A^i = (a^i_{jh})$ and $X = (x_{jh})$:

$$\min \sum_{j,h \le n} c_{jh} x_{jh}$$

 $orall i \le m \sum_{j,h \le n} a^i_{jh} x_{jh} = b^i.$

This is just an LP subject to a semidefinite constraint $X \succeq 0$. What does $X \succeq 0$ mean? *X*, square and symmetric matrix, with real values is PSD if $\forall v \in \mathbb{R}^n$, we have that the scalar $v^T X v \ge 0$.

This is the same as requiring the decision variables x_{jh} to take values such that, when they are arranged in a $n \times n$ array, the resulting matrix A has non-negative eigenvalues.

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Use Octave to implement an SDP relaxation of the KNP (decision version, [KBLM07])

Then, download SDP_KBLM07.m (your code should not be very different). Test some instances:

(b) n=7 and k=2 (c) n=8 and k=2 (d) n=9 and k=2 (e) ...

Can you explain this behaviour, with growing n?

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Apparently, this first version of a SDP relaxation of KNP (i.e. $SDP_KBLM07.m$) is feasible independently from *n*. This makes it not really useful (cf. *Uselessness Theorem* in Liberti's slides)

Look for an improved SDP relaxation that deals with this drawback. Initially, focus on **2D case** only (k = 2).

Hint: add a constraint on the dimension of the sdp matrix (i.e. decision variable). Hint 2: for example, estimate the perimeter of the central sphere "consumed" by the placements.

Extend your code to deal with 3D case (k = 3). Hint: this would be an approximation.

Kissing number problem: Upper bounds

Download delsartebnd.m, delsartelp.m and gegenbauer.m and test it in Octave in the following way: [b,r] = delsartebnd(3, 60, 10);

Similarily, Download pfenderbnd.m and pfenderlp.m and test it in Octave in the following way: [p,s] = pfenderbnd(3, 60, 10);

Your outcome should be, respectively: delsartebnd: bound positivity in [0,0.000100], accepting pfenderbnd: bound positivity in [0,0.000100], accepting

Can you explain what this means?

Kissing number problem: scripting

In the end, look in these scripts:

- how to optimize unconstrained functions within Octave using sqp() (examples in dgpnlp.m, delsartebnd.m and pfenderbnd.m).
- how to call glpk() within octave (examples in delsartelp.m and pfenderlp.m).
- how to compute Gegenbauer polynomial coefficients.