# Mathematical Programming: Modelling and Applications 

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## Outline

(1) Kissing number problem

## Kissing number problem

## Some examples

$$
n=6, K=2 \quad n=12, K=3
$$

## more dimensions



| $n$ | $\tau$ (lattice) | $\tau$ (nonlattice) |
| ---: | ---: | ---: |
| 0 | 0 |  |
| 1 | 2 |  |
| 2 | 6 |  |
| 3 | 12 |  |
| 4 | 24 |  |
| 5 | 40 |  |
| 6 | 72 |  |
| 7 | 126 |  |
| 8 | 240 |  |
| 9 | 272 | $(306)^{*}$ |
| 10 | 336 | $(500)^{*}$ |
| 11 | 438 | $(582)^{*}$ |
| 12 | 756 | $(840)^{*}$ |
| 13 | 918 | $(1130)^{*}$ |
| 14 | 1422 | $(1582)^{*}$ |
| 15 | 2340 |  |
| 16 | 4320 |  |
| 17 | 5346 |  |
| 18 | 7398 |  |
| 19 | 10668 |  |
| 20 | 17400 |  |
| 21 | 27720 |  |
| 22 | 49896 |  |

## Kissing number problem: exercise 1

Use AMPL to implement a MINLP formulation of the Kissing Number Problem (as an optimization problem).

Hint 1: .dat files are given for small/medium instances.
Hint 2: look at Maculan, Michelon, Smith' model (1995).
Hint 3: you can use additional variables to reformulate the binary products.
Hint 4: you can use the knpcheck.run script (which is given) to check your solution (in other words, you can include it at the bottom of your .run file).
Hint 5: use Baron.
Set a threshold (for example, CPU Time limit 100 s., i.e. maxtime $=100$ in the options of your .run script) and test these instances:
(a) $\mathrm{n}=12$ and $\mathrm{k}=3$
(b) $\mathrm{n}=7$ and $\mathrm{k}=2$

In theory, your program should be able to accomplish (a) but not (b).
Then, download knp.mod and knp.run. Your implementation should be quite similar.

## Kissing number problem: exercise 2

Use AMPL to implement the Kissing Number Problem (as a decision problem).
Hint: look at Kucherenko, Belotti, Liberti, Maculan' model (2007).
Set again a threshold (for example, CPU Time limit 100 s., i.e. maxtime=100 in the options of your run script) and test again these instances:
(a) $\mathrm{n}=12$ and $\mathrm{k}=3$
(b) $\mathrm{n}=7$ and $\mathrm{k}=2$

In theory, the program should not be able to accomplish neither (a) nor (b). Then, download knpfeas 18.mod and knpfeas18.run. Your implementation should be quite similar.

Change the constraint on $\alpha$, to see if something changes.

## Reminder: SDP

Semidefinite Programming (SDP) refers to optimization problems where decision variables are a symmetric matrix that is required to be semidefinite positive. The standard form is

$$
\left.\begin{array}{rrll}
\min & C \bullet X & &  \tag{1}\\
\forall i \leq m & A^{i} \bullet X & =b^{i} \\
& X & \succeq & 0,
\end{array}\right\}
$$

where:

- $C$ and $A^{i}$ (for $i \leq m$ ) are $n \times n$ symmetric matrices
- $X$ is an $n \times n$ matrix of decision variables,
- $b^{i} \in \mathbb{R}$ for all $i \leq m$
- for two $n \times n$ matrices $L=\left(\lambda_{i j}\right), M=\left(\mu_{i j}\right)$ we have $L \bullet M=\sum_{i, j \leq n} \lambda_{i j} \mu_{i j}$
- $X \succeq 0$, i.e. $X$ is positive semidefinite


## Reminder: SDP

To make Eq. (1) clearer, write out the componentwise product • of the matrices:
$C=\left(c_{j h}\right), A^{i}=\left(a_{j h}^{i}\right)$ and $X=\left(x_{j h}\right)$ :

$$
\forall i \leq m \min _{\sum_{j, h \leq n} c_{j h} x_{j h}}^{\sum_{j, h \leq n} a_{j h}^{i} x_{j h}=b^{i} .}
$$

This is just an LP subject to a semidefinite constraint $X \succeq 0$. What does $X \succeq 0$ mean? $X$, square and symmetric matrix, with real values is $\operatorname{PSD}$ if $\forall v \in \mathbb{R}^{n}$, we have that the scalar $v^{T} X v \geq 0$.
This is the same as requiring the decision variables $x_{j h}$ to take values such that, when they are arranged in a $n \times n$ array, the resulting matrix $A$ has non-negative eigenvalues.

## Kissing number problem: exercise 3

Use Octave to implement an SDP relaxation of the KNP (decision version, [KBLM07])

Then, download SDP_KBLM07.m (your code should not be very different).
Test some instances:
(b) $\mathrm{n}=7$ and $\mathrm{k}=2$
(c) $\mathrm{n}=8$ and $\mathrm{k}=2$
(d) $\mathrm{n}=9$ and $\mathrm{k}=2$
(e) ...

Can you explain this behaviour, with growing $n$ ?

## Kissing number problem: exercise 4

Apparently, this first version of a SDP relaxation of KNP (i.e. SDP_KBLM07.m) is feasible independently from $n$. This makes it not really useful (cf. Uselessness Theorem in Liberti's slides)

Look for an improved SDP relaxation that deals with this drawback. Initially, focus on 2D case only $(k=2)$.

Hint: add a constraint on the dimension of the sdp matrix (i.e. decision variable). Hint 2: for example, estimate the perimeter of the central sphere "consumed" by the placements.

## Kissing number problem: exercise 5

Extend your code to deal with 3D case $(k=3)$.
Hint: this would be an approximation.

## Kissing number problem: Upper bounds

Download delsartebnd.m, delsartelp.m and gegenbauer.m and test it in Octave in the following way:
[b,r] = delsartebnd (3, 60, 10);
Similarily, Download pfenderbnd.m and pfenderlp.m and test it in Octave in the following way:
[p,s] = pfenderbnd (3, 60, 10);
Your outcome should be, respectively:
delsartebnd: bound positivity in [0,0.000100], accepting pfenderbnd: bound positivity in [ $0,0.000100$ ], accepting
Can you explain what this means?

## Kissing number problem: scripting

In the end, look in these scripts:

- how to optimize unconstrained functions within Octave using sqp() (examples in dgpnlp.m, delsartebnd.m and pfenderbnd.m).
- how to call glpk() within octave (examples in delsartelp.m and pfenderlp.m).
- how to compute Gegenbauer polynomial coefficients.

