Mathematical Programming: Modelling and Applications

October 2009

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Writing the mathematical model

Preliminary observations

- Each batch must be dyed in the first bath, then in the second bath,....

The *k*-th bath must finish before that the (k+1)-th bath can start.

- The time for dyeing in the last bath must be bounded by the completion time for last batch.
- Two different batches cannot be dyed in the same bath at the same time.
- If the bath *i* is to be dyed before batch *j*, then *j* cannot be dyed before *i* (disjunction).

Mathematical model

• Sets

- *L* : set of fabric batches $L = \{1, ..., l\}$
- *V* : set of dyeing baths $V = \{1, ..., v\}$

Parameters

 $\forall i \in L, k \in V$ s_{ik} : time necessary to dye batch *i* in bath *k*

M: upper bound to completion time of last bath

Variables

T: completion time for last batch

 t_{ik} : starting time for dyeing batch *i* in bath *k*

 y_{ijk} = 1 if batch *i* is to be dyed before batch *j* in bath *k*, 0 otherwise.

Mathematical model

Objective function

we want to minimize the completion time of the last batch

min T

- Constraints
- sequential: $\forall i \in L, k \in V \setminus \{v\} \quad (t_{ik} + s_{ik} \leq t_{i(k+1)})$
- last bath: $\forall i \in L \quad (t_{iv} + s_{iv} \leq T)$
- non overlapping: $\forall i, j \in L, k \in V, i \neq j$ $(t_{ik} + s_{ik} \leq t_{jk} + M(1 y_{ijk}))$
- disjunction: $\forall i, j \in L, k \in V, i \neq j \quad (y_{ijk} + y_{jik} = 1)$