

Combinatorial Optimization in Bioinfo

Lecture 2 – Folding RNA *in silico*

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1 Introduction

- Dynamic programming 101
- Dynamic programming framework

2 Variations on RNA folding

- Why RNA?
- RNA folding
- RNA Structure(s)
- Some representations of RNA structure
- Thermodynamics vs Kinetics

3 Free-energy minimization

- Nussinov-style RNA folding
- Turner energy model
- MFold/Unafold
- Performances and the comparative approach
- Towards a 3D ab-initio prediction

4 Boltzmann ensemble

- Nussinov: Minimisation \Rightarrow Counting
- Computing the partition function
- Statistical sampling

...or how to make a million bucks by giving change parsimoniously!!

Problem: You have access to unlimited amount of **1**, **20** and **50** cents coins.
A client prefers to travel light, i.e. to **minimize the #coins**.
How to give **N** cents back in change without losing a customer?

Strategy #1: Start with *heaviest* coins, and then complete/fill-up with coins of *decreasing* value.

21 =??

55

60

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$$55 = \text{€50} + \text{€5} + \text{€5} + \text{€5} + \text{€5} + \text{€5}$$

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$$21 = 20 + 1$$

$55 =$  $+$  $+$  $+$  $+$  $+$ 

60 =  +  +  +  +  +  +  +  +  +  +  +  ??

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$$55 = \text{€50} + \text{€1} + \text{€1} + \text{€1} + \text{€1} + \text{€1}$$

$$\begin{aligned} 60 &= \text{€50} + \text{€1} + \text{€1} + \text{€1} + \text{€1} + \text{€1} + \text{€1} + \text{€1} + \text{€1} + \text{€1} + \text{€1} ?? \\ &= \text{€20} + \text{€20} + \text{€20} ! \end{aligned}$$

Problem *a priori* (!) non-solvable using such a *greedy* approach, as a (simpler) problem is already NP-complete (thus Efficient solution \Rightarrow 1M\$).

Strategy #2: Brute force enumeration $\rightarrow \#Coins^N$ (Ouch!)

Strategy #3: The following recurrence gives the minimal number of coins:

$$Min\#Coins(N) = \text{Min} \left\{ \begin{array}{ll} \img alt="1 Euro coin" data-bbox="441 258 471 301" & \rightarrow 1 + Min\#Coins(N - 1) \\ \img alt="2 Euro coin" data-bbox="434 311 478 364" & \rightarrow 1 + Min\#Coins(N - 20) \\ \img alt="5 Euro coin" data-bbox="434 374 478 438" & \rightarrow 1 + Min\#Coins(N - 50) \end{array} \right.$$

With some memory (N intermediate computations), the minimum number of coins can be obtained after $N \times \#Coins$ operations. An actual set of coins can be reconstructing by **tracing back** the choices performed at each stage, leading to the minimum.

Remark: We still haven't won the million, as N has **exponential value compared to the length of its encoding**, so the algorithm does not qualify as *efficient* (i.e. polynomial).

Still, this approach is much more efficient than a brute-force enumeration:
 \Rightarrow Dynamic programming.

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Dynamic programming = General optimization technique.

Prerequisite: Optimal solution for problem P can be derived from solutions to strict sub-problems of P .

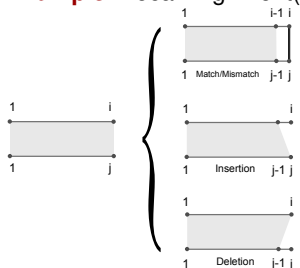
Bioinformatics :

Discrete solution space (alignments, structures...)

+ Additively-inherited objective function (cost, log-odd score, energy...)

⇒ Efficient dynamic programming scheme

Example: Local Alignment(Smith/Waterman)



$$\begin{aligned}
 W(i, 0) &= 0 \\
 W(0, j) &= 0 \\
 W(i, j) &= \max \begin{cases} W(i-1, j-1) + m_{i,j} \\ W(i-1, j) + p_i \\ W(i, j-1) + p_d \end{cases}
 \end{aligned}$$

Dynamic programming scheme defines a space of (sub)problems and a recurrence that relates the score of a problem to that of smaller problems.

Given a scheme, two steps :

- ▶ **Matrix filling:** Computation and tabulation of best scores (Computed from smaller problems to larger ones).
- ▶ **Traceback:** Reconstruct best solution from contributing subproblems.

Complexity of algorithm depends on:

- ▶ **Cardinality** of sub-problem space
- ▶ **Number of alternatives** considers at each step (#Terms in recurrence)

Smith&Waterman example:

- ▶ $i: 1 \rightarrow n + 1 \Rightarrow \Theta(n)$
 - ▶ $j: 1 \rightarrow m + 1 \Rightarrow \Theta(m)$
 - ▶ 3 operations at each step
- $\Rightarrow \Theta(m.n)$ time/memory

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i-1, j-1) + m_{i,j} \\ W(i-1, j) + p_i \\ W(i, j-1) + p_d \end{cases}$$

Necessary properties:

- ▶ **Correctness:** \forall sub-problem, the computed value must indeed maximize the objective function .

Proofs usually inductive, and quite technical, but very systematic.

Desirable properties of DP schemes:

- ▶ **Completeness** of space of solutions **generated by** decomposition.
Algorithmic tricks, by *cutting branches*, may violate this property.
- ▶ **Unambiguity:** Each solution is **generated** at most once.

\Rightarrow Under these properties, one can **enumerate** solution space.

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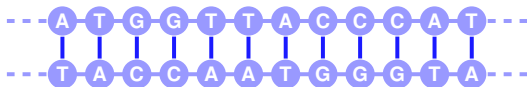
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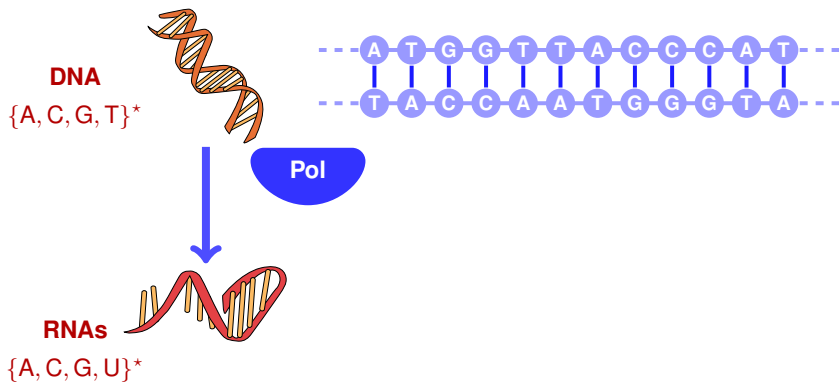
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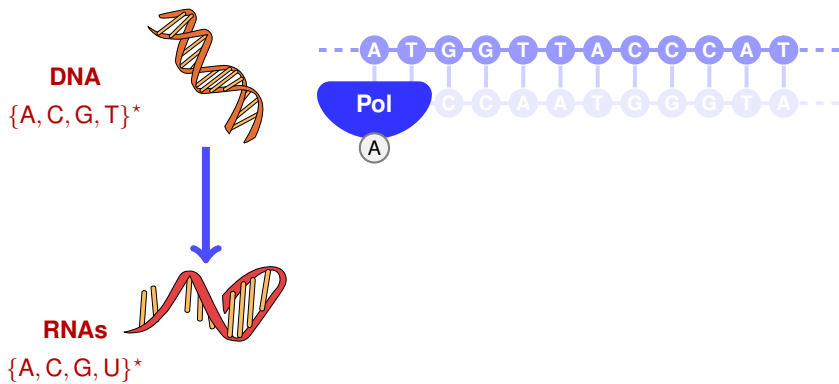
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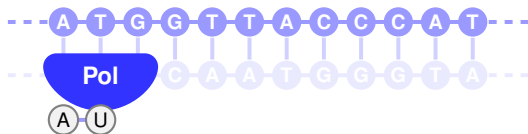
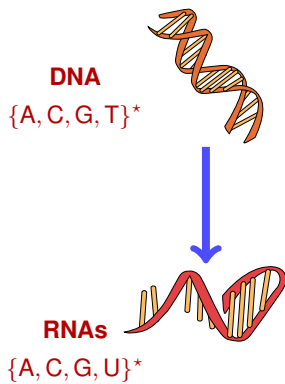
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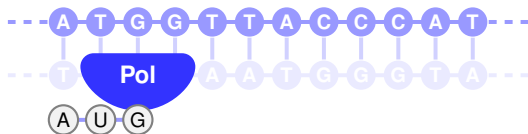
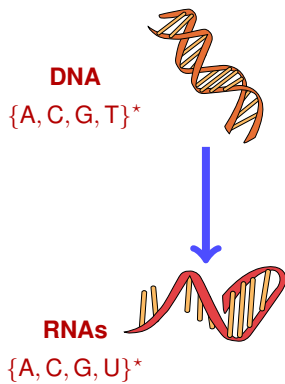
DNA
{A, C, G, T}^{*}

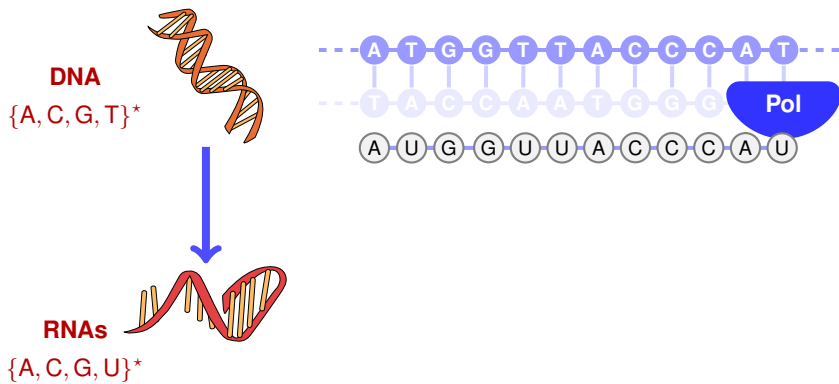


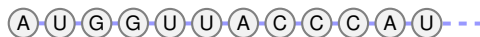
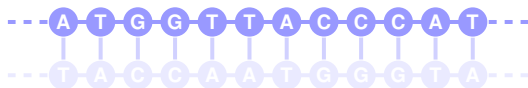
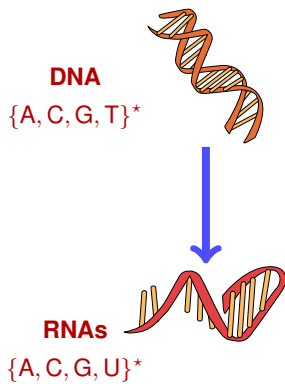


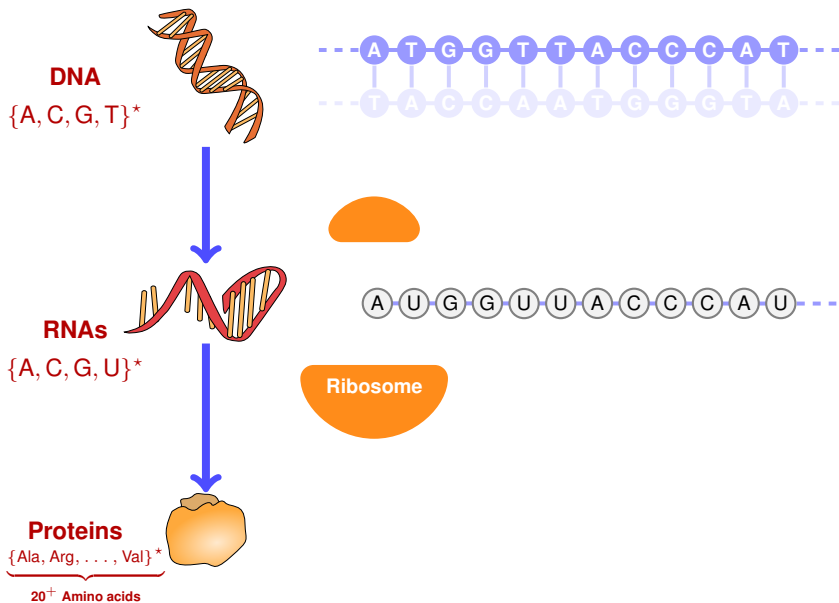


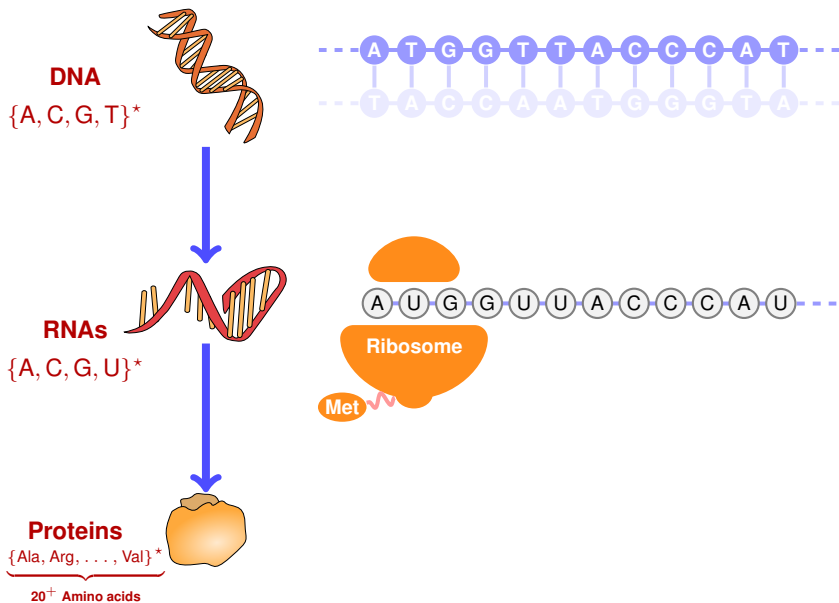


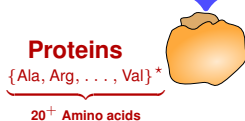
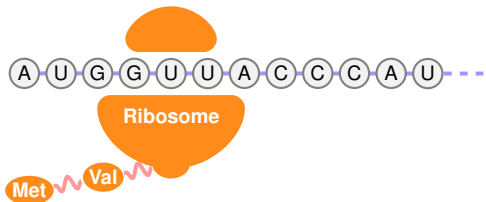
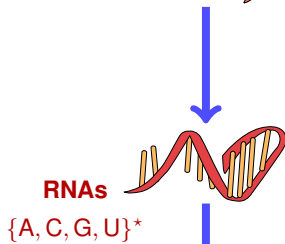
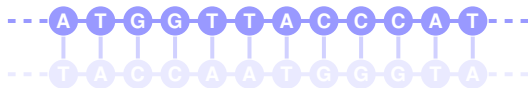


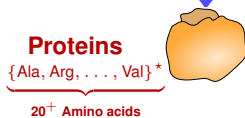
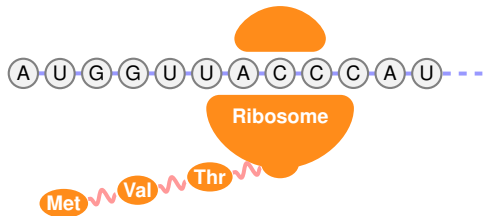
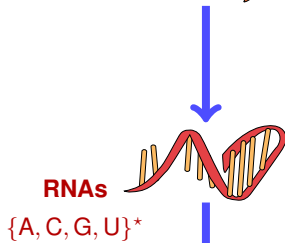
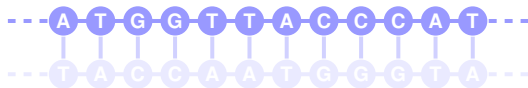


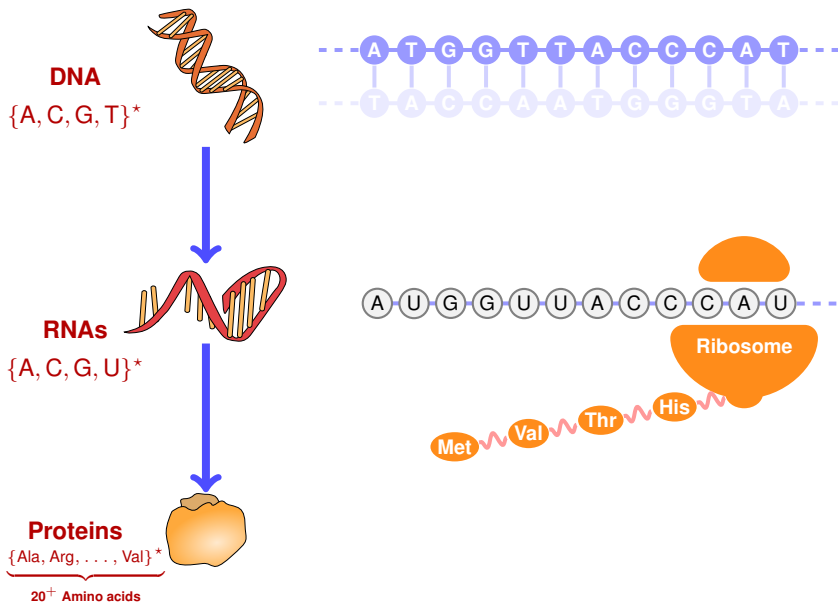


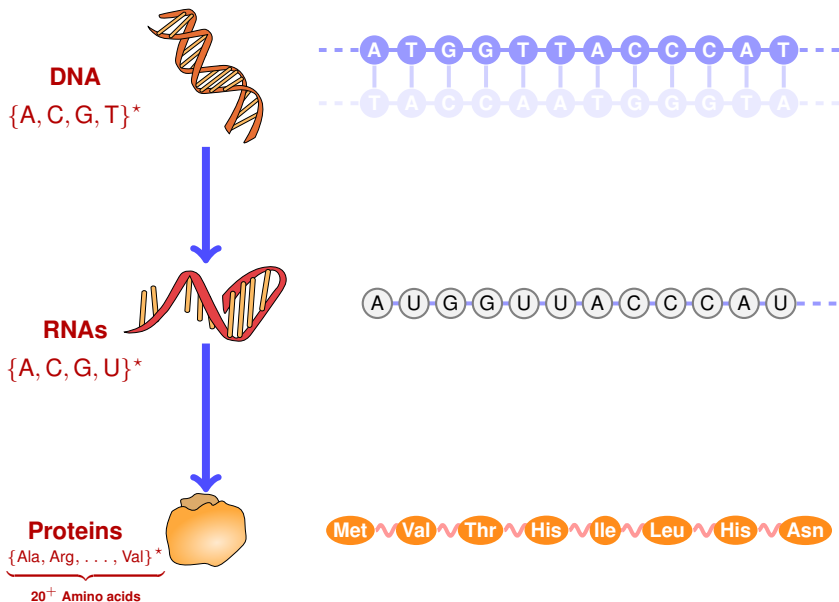


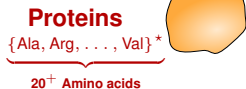
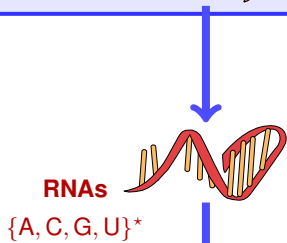
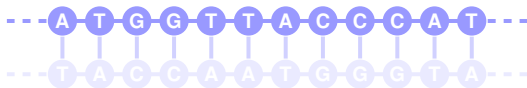
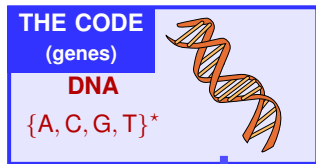


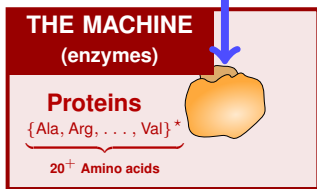
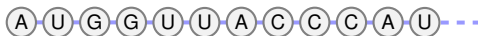
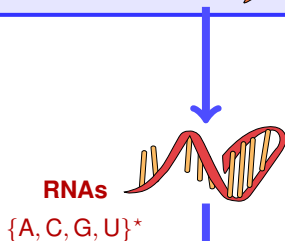
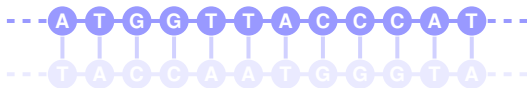
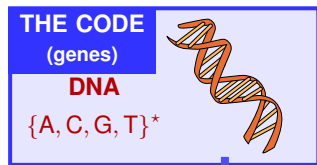


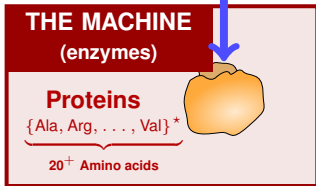
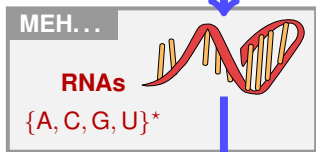
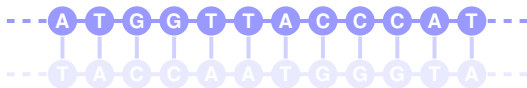
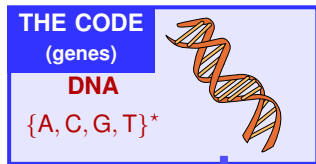


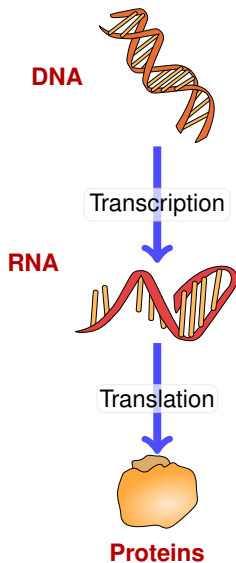


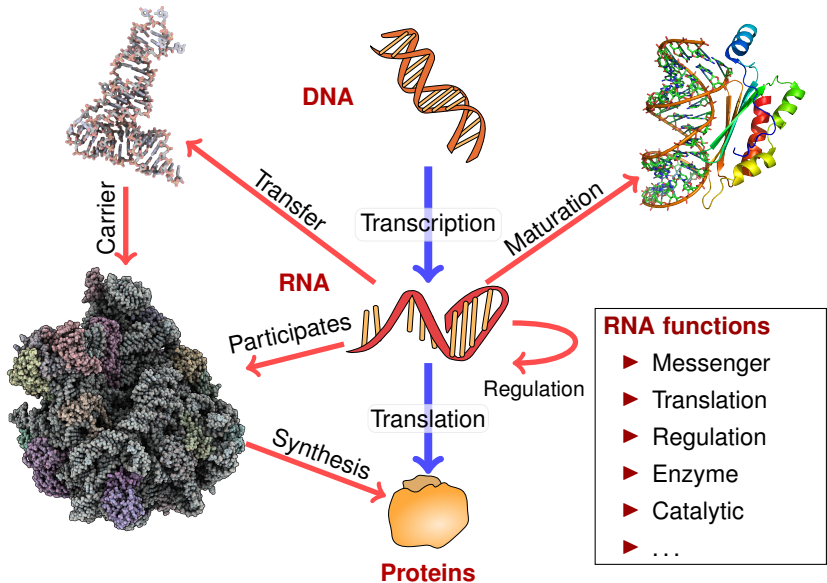


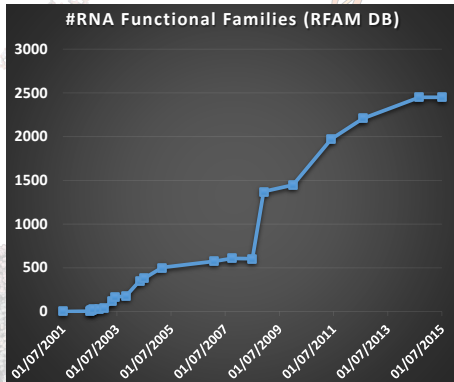








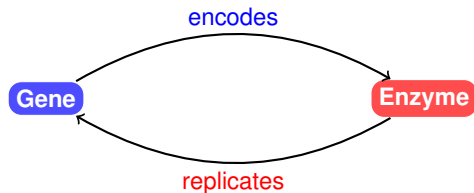




RNA functions

- ▶ Messenger
- ▶ Translation
- ▶ Regulation
- ▶ Enzyme
- ▶ Catalytic
- ▶ ...

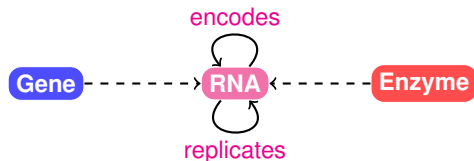
Proteins



A **gene** big enough to specify **an enzyme** would be too big to replicate accurately without the aid of **an enzyme** of the very kind that it is trying to specify. So the system *apparently cannot get started*.

[...] This is the **RNA World**. To see how plausible it is, we need to look at why proteins are good at being enzymes but bad at being replicators; at why DNA is good at replicating but bad at being an enzyme; and finally why **RNA might just be good enough at both roles to break out of the Catch-22**.

R. Dawkins. *The Ancestor's Tale: A Pilgrimage to the Dawn of Evolution*



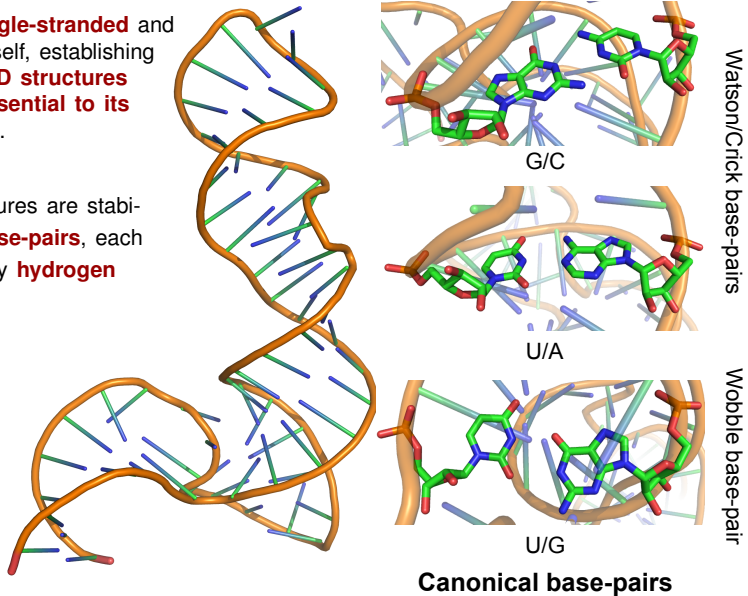
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RNA is **single-stranded** and **folds** on itself, establishing **complex 3D structures** that are **essential to its function(s)**.

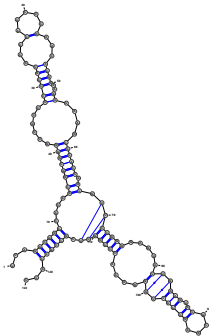
RNA structures are stabilized by **base-pairs**, each mediated by **hydrogen bonds**.



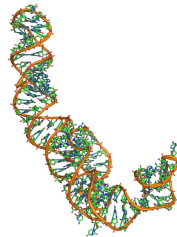
Three¹ levels of representation:

```
UUAGGCGGCCACAGC
GGUGGGGUUGCCUCC
CGUACCAUCCCGAA
CACGGAAGAUAGCC
CACCAGCGUCCGGG
GAGUACUGGAGUGCG
CGAGCCUCUGGAAA
CCCGGUUCGCCGCA
CC
```

Primary structure



Secondary structure



Tertiary structure

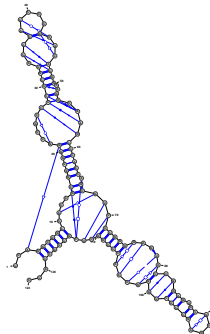
Source: 5s rRNA (PDB 1K73:B)

¹Well, mostly...

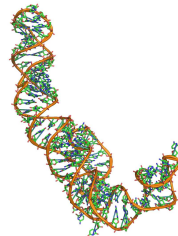
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CACGGAAGAUAGCC
CACCAGCGUCCGGG
GAGUACUGGAGUGCG
CGAGCCUCUGGGAAA
CCCGGUUCGCCGCCA
CC
```

Primary structure



Secondary⁺ structure



Tertiary structure

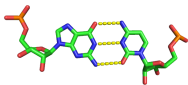
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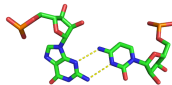
► **Non-canonical base-pairs**

Any base-pair **other than** {(A-U), (C-G), (G-U)}

Or interacting on non-standard edge (\neq WC/WC-Cis) [LW01].

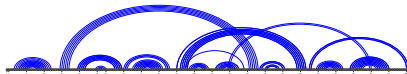


Canonique CG pair(WC/WC-Cis)



Non-canonique CG pair (Sugar/WC-Trans)

► **Pseudoknots (PKs)**

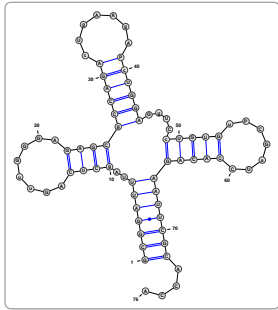


Pseudoknotted structure of group I ribozyme (PDBID: 1Y0Q:A)

Considering PKs may lead to better predictions, **but**:

- Some PK conformations are simply unfeasible;
- Folding *in silico* with general pseudoknots is NP-complete [LP00];

Still, folding on restricted classes of conformations seems promising [CDR⁺04].



Outer-planar graphs

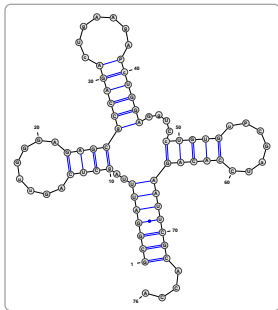
Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*

Supporting intuitions

Different representations

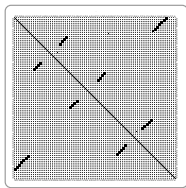
Common combinatorial structure

* Additional steric constraints



Outer-planar graphs

Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*



Dot plots

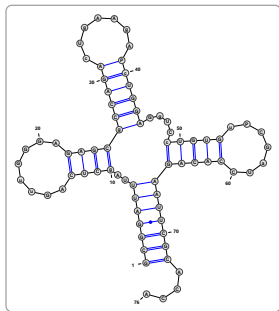
Adjacency matrices*

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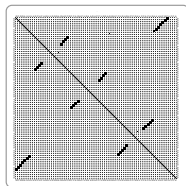
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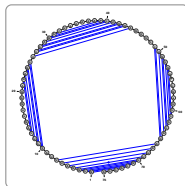
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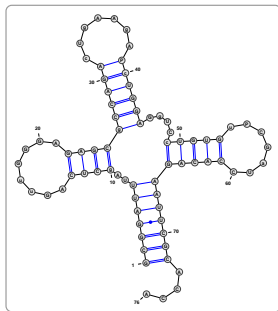
Non-crossing arc diagrams*

Supporting intuitions

Different representations

Common combinatorial structure

* Additional steric constraints

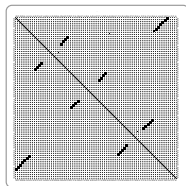


Outer-planar graphs

Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*

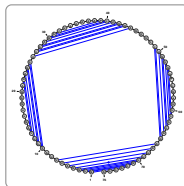
((((((((.....))))))(((.....)))))).....(((.....)))))).....

Motzkin words*



Dot plots

Adjacency matrices*



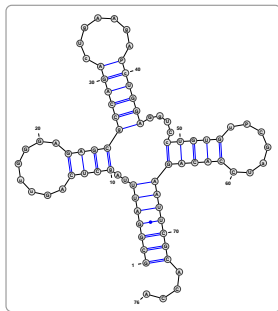
Non-crossing arc diagrams*

Supporting intuitions

Different representations

Common combinatorial structure

* Additional steric constraints

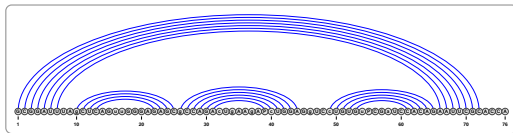


Outer-planar graphs

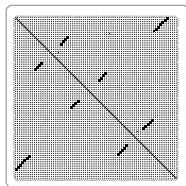
Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*

((((((...(((...))))))(((...))))...(((...))))))...))

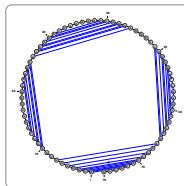
Motzkin words*



Non-crossing arc-annotated sequences*



Dot plots
Adjacency matrices*



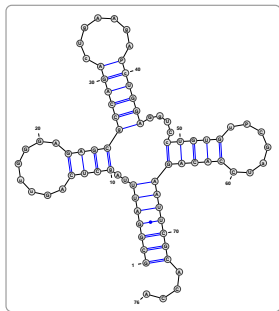
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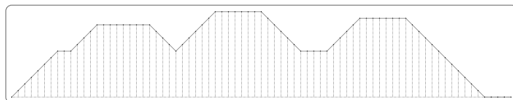


Outer-planar graphs

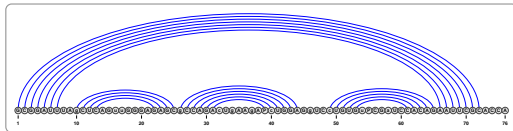
Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*

((((((...(((.....))))))((((((.....))))))....(((.....))))))....

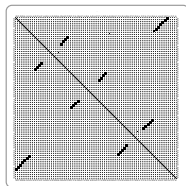
Motzkin words*



Positive 1D meanders* over $S = \{+1, -1, 0\}$

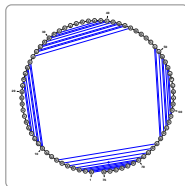


Non-crossing arc-annotated sequences*



Dot plots

Adjacency matrices*



Non-crossing arc diagrams*

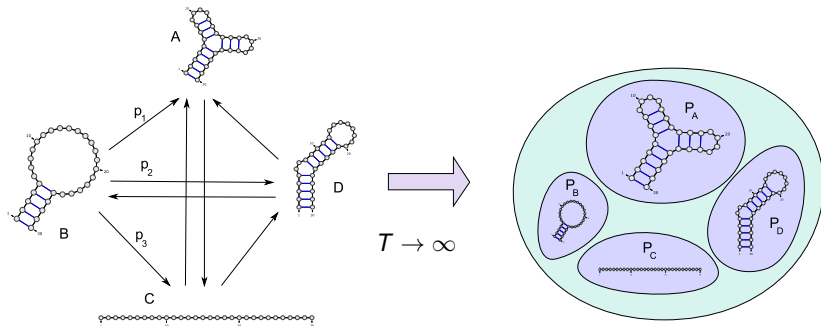
Supporting intuitions

Different representations

Common combinatorial structure

* Additional steric constraints

At the nanoscopic scale, RNA structure *fluctuates* (\approx Markov process).



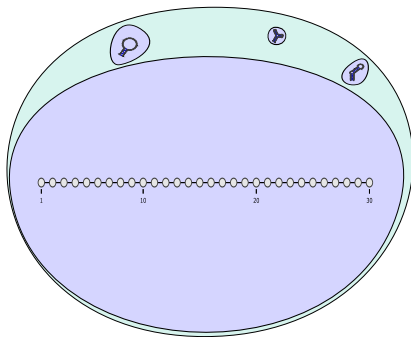
Convergence towards a **stationary distribution** at the **Boltzmann equilibrium**, where the probability of a conformation only depends on its **free-energy**.

Corollary: Initial conformation does not matter.

Questions: For a given **conformation space** and **free-energy** model:

- A.** Determine most stable (Minimum Free-Energy) structure at equilibrium;
- B.** Compute average properties of Boltzmann ensemble;

Transcription: RNA synthesized, supposedly without structure²

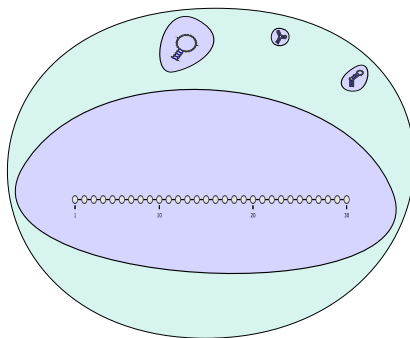


$$T = 0$$

But most mRNAs are degraded before 7h (Org.: Souris [SSN⁺09]).

²Except for co-transcriptional folding...

Transcription: RNA synthesized, supposedly without structure²

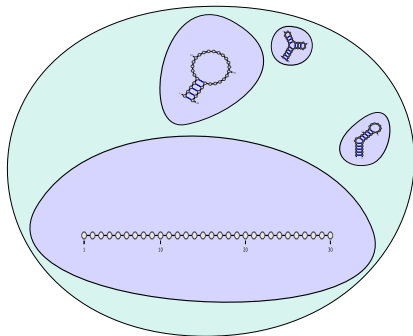


$$T = 1h$$

But most mRNAs are degraded before 7h (Orig.: Souris [SSN⁺09]).

²Except for co-transcriptional folding...

Transcription: RNA synthesized, supposedly without structure²

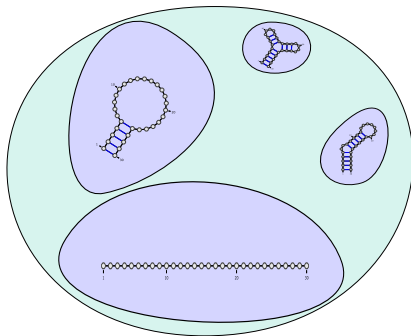


$$T = 2h$$

But most mRNAs are degraded before 7h (Orig.: Souris [SSN⁺09]).

²Except for co-transcriptional folding...

Transcription: RNA synthesized, supposedly without structure²

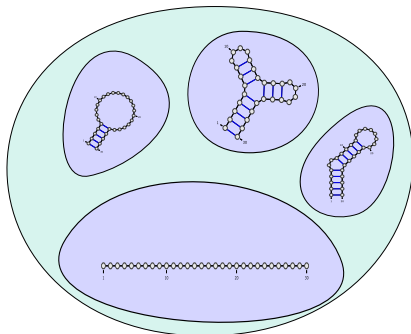


$$T = 5h$$

But most mRNAs are degraded before 7h (Orig.: Souris [SSN⁺09]).

²Except for co-transcriptional folding. . .

Transcription: RNA synthesized, supposedly without structure²

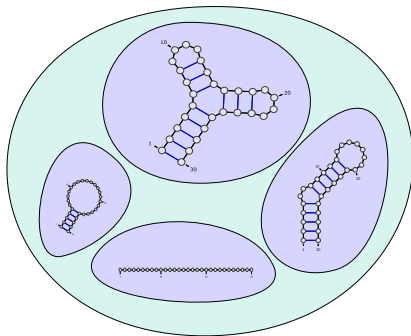


$$T = 10h$$

But most mRNAs are degraded before 7h (Orig.: Souris [SSN⁺09]).

²Except for co-transcriptional folding...

Transcription: RNA synthesized, supposedly without structure²

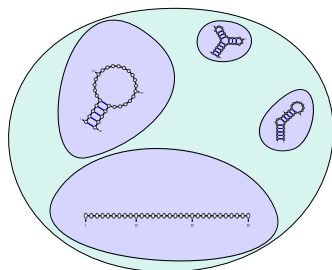


$$T \rightarrow \infty$$

But most mRNAs are degraded before 7h (Org.: Souris [SSN⁺09]).

²Except for co-transcriptional folding...

Transcription: RNA synthesized, supposedly without structure²



$$T = 10h$$

But most mRNAs are degraded before 7h (Org.: Souris [SSN⁺09]).

- A.** Determine most stable (Minimum Free-Energy) structure at equilibrium;
- B.** Compute average properties of Boltzmann ensemble;
- C. Determine most likely structure at finite time T .**
(c.f. H. Isambert through simulation, NP-complete deterministically [MTSC09])

²Except for co-transcriptional folding. . .

1 Introduction

- Dynamic programming 101
- Dynamic programming framework

2 Variations on RNA folding

- Why RNA?
- RNA folding
- RNA Structure(s)
- Some representations of RNA structure
- Thermodynamics vs Kinetics

3 Free-energy minimization

- Nussinov-style RNA folding
- Turner energy model
- MFold/Unafold
- Performances and the comparative approach
- Towards a 3D ab-initio prediction

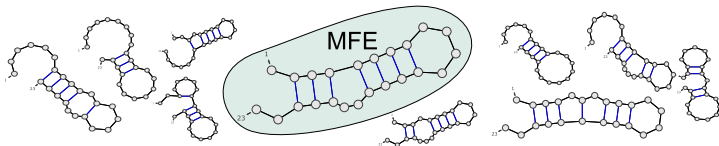
4 Boltzmann ensemble

- Nussinov: Minimisation \Rightarrow Counting
- Computing the partition function
- Statistical sampling

Problem A: Determine Minimum Free-Energy structure (MFE).

Ab initio folding prediction =

Predict RNA structure from its sequence ω only.



- **Conformations:** Set S_ω of secondary structures **compatible** (w.r.t. **base-pairing constraints**) with primary structure ω .
- **Free-Energy:** Function $E_{\omega, S}$ (KCal.mol^{-1}), **additive** on motifs occurring in any sequence/conformation couple (ω, S) .
- **Native structure:** Functional conformation of the biomolecule.

Remarks:

- Not necessarily unique (Kinetics, or bi-stable structures);
- In presence of PKs \rightarrow Ambiguous: Which is the native conformation?

Nussinov/Jacobson energy model (NJ)

Base-pair maximization (with a twist):

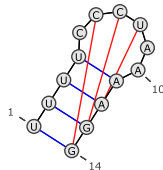
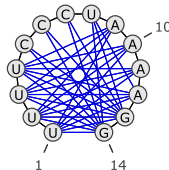
- ▶ Additive model on **independently contributing** base-pairs;
- ▶ **Canonical base-pairs** only: Watson/Crick (A/U,C/G) and Wobble (G/U)

$$\Rightarrow E_{\omega, S} = -\#Paires(S)$$

Folding in NJ model \Leftrightarrow **Base-pair (weight)** maximization

Example:

UUUUCCCUAAAAGG



Variant: Weight each pair with $-\#Hydrogen\ bonds$

$$\Delta G(G \equiv C) = -3$$

$$\Delta G(A = U) = -2$$

$$\Delta G(G - U) = -1$$

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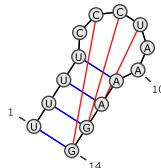
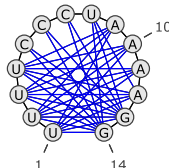
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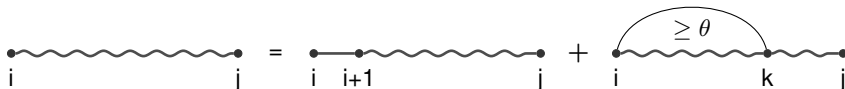


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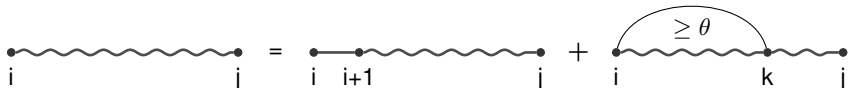
$$\Delta G(A = U) = -2$$

$$\Delta G(G - U) = -1$$



$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \begin{cases} N_{i+1,j} & i \text{ unpaired} \\ \min_{k=i+\theta+1}^j \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{cases}$$

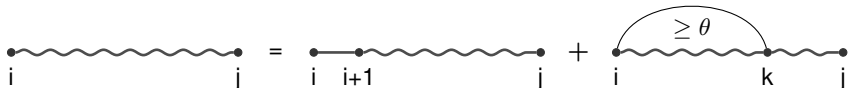


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Correctness. Goal = Show that MFE over interval $[i, j]$ is indeed found in $N_{i,j}$ after completing the computation. Proceed by induction:

- ▶ Assume that property holds for any $[i', j']$ such that $j' - i' < n$.
- ▶ Consider $[i, j], j - i = n$. Let $\text{MFE}_{i,j} :=$ Base-pairs of best struct. on $[i, j]$. Then first position i in $\text{MFE}_{i,j}$ is either:
 - ▶ **Unpaired:** $\text{MFE}_{i,j} = \text{MFE}_{i+1,j} \rightarrow \text{free-energy} = N_{i+1,j}$
 - ▶ **Paired to k :** $\text{MFE}_{i,j} = \{(i, k)\} \cup \text{MFE}_{i+1,k-1} \cup \text{MFE}_{k+1,j}$
 (Indeed, any BP between $[i+1, k-1]$ and $[k+1, j]$ would cross (i, k))
 $\rightarrow \text{free-energy} = \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j}$

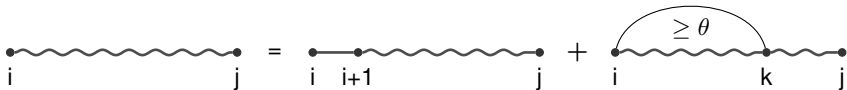


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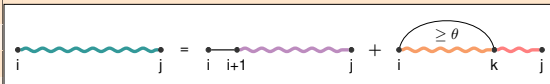
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	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A

C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



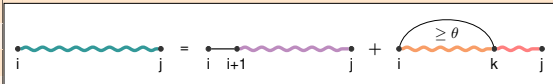
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A

C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0

Diagram illustrating the Nussinov/Jacobson RNA secondary structure prediction algorithm. It shows a sequence from index i to j . The first part is a blue wavy line from i to j . This is equal to the sum of two terms: 1) a black line from i to $i+1$ followed by a blue wavy line from $i+1$ to j , and 2) a term where a blue wavy line from i to k is paired with a red wavy line from k to j , with a curved arrow labeled $\geq \theta$ above the pair.

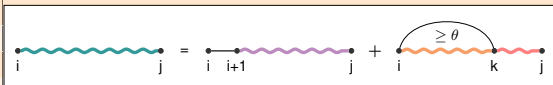
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A

C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0

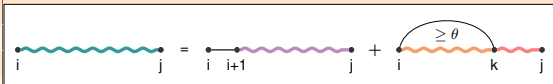


	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A

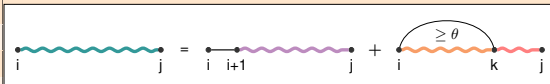
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
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U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



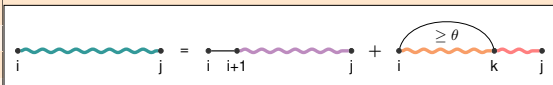
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



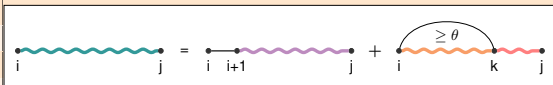
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	(.)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



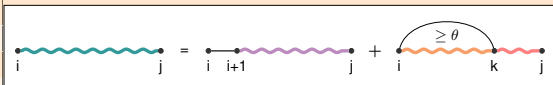
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	(.)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



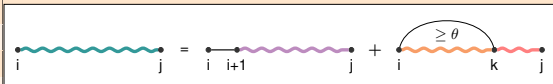
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



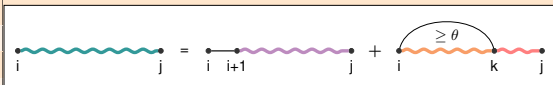
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



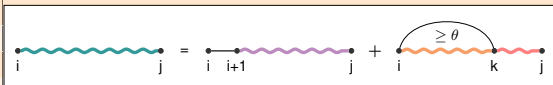
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	(.)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



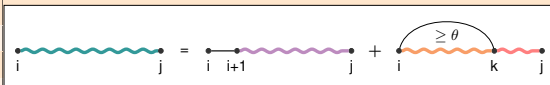
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	(.)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



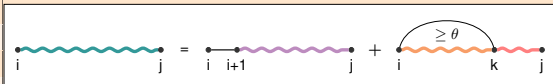
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	(.)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
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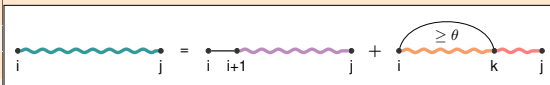
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	(.)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
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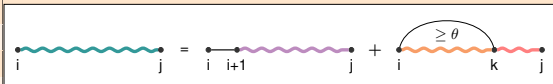
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	((.))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



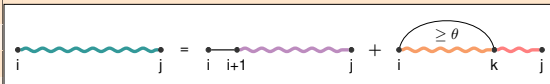
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	((.))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



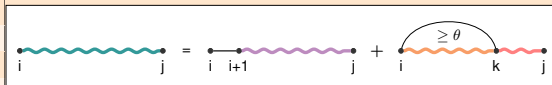
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	((.))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



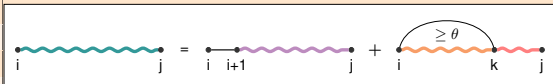
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
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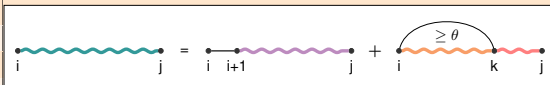
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



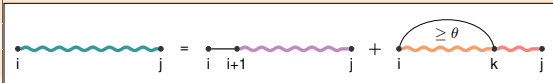
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0



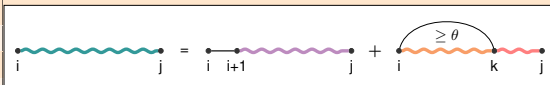
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
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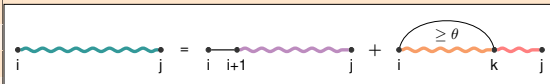
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G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
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C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
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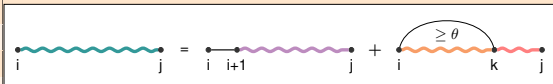
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G																	0	0
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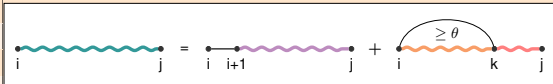
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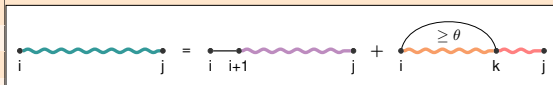
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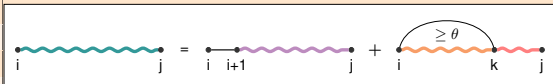
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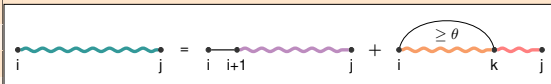
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U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
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G														0	0	0	0	0
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C																0	0	0
G																	0	0
A																		0



	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
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U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
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A																		0



	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)	.	((.	.	.))))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
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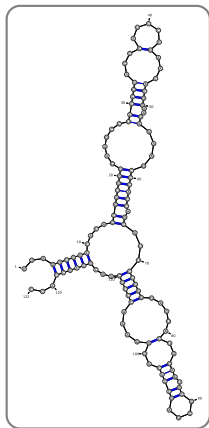
Based on **unambiguous** decomposition of 2^{ary} structure into **loops**:

- ▶ Internal loops
- ▶ Bulges
- ▶ Terminal loops
- ▶ Multi loops
- ▶ Stackings

Free-energy ΔG of a loop depend on
bases, asymmetry, dangles ...

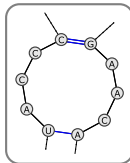
Experimentally determined
+ Interpolated for larger loops.

Improved results by taking stacking into account.



Based on **unambiguous** decomposition of 2^{ary} structure into **loops**:

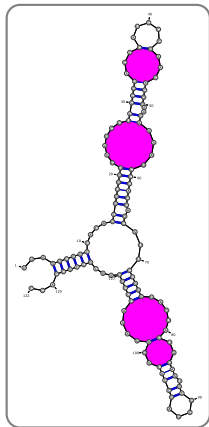
- ▶ Internal loops
- ▶ Bulges
- ▶ Terminal loops
- ▶ Multi loops
- ▶ Stackings



Free-energy ΔG of a loop depend on
bases, asymmetry, dangles ...

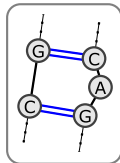
Experimentally determined
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Improved results by taking stacking into account.



Based on **unambiguous** decomposition of 2^{ary} structure into **loops**:

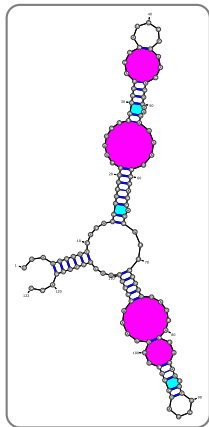
- ▶ Internal loops
- ▶ Bulges
- ▶ Terminal loops
- ▶ Multi loops
- ▶ Stackings



Free-energy ΔG of a loop depend on
bases, asymmetry, dangles ...

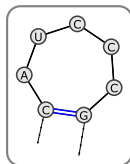
Experimentally determined
+ Interpolated for larger loops.

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Based on **unambiguous** decomposition of 2^{ary} structure into **loops**:

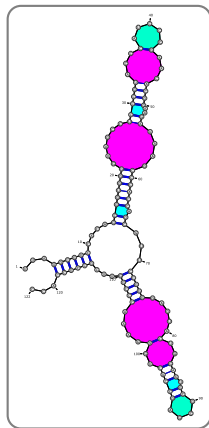
- ▶ Internal loops
- ▶ Bulges
- ▶ Terminal loops
- ▶ Multi loops
- ▶ Stackings



Free-energy ΔG of a loop depend on
bases, asymmetry, dangles ...

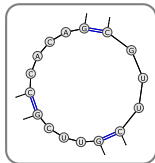
Experimentally determined
+ Interpolated for larger loops.

Improved results by taking stacking into account.



Based on **unambiguous** decomposition of 2^{ary} structure into **loops**:

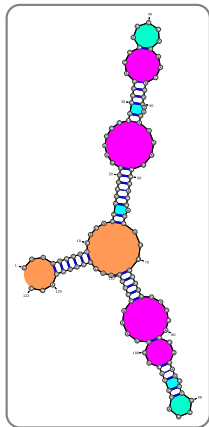
- ▶ Internal loops
- ▶ Bulges
- ▶ Terminal loops
- ▶ Multi loops
- ▶ Stackings



Free-energy ΔG of a loop depend on
bases, asymmetry, dangles ...

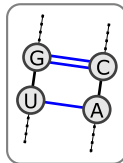
Experimentally determined
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Improved results by taking stacking into account.



Based on **unambiguous** decomposition of 2^{ary} structure into **loops**:

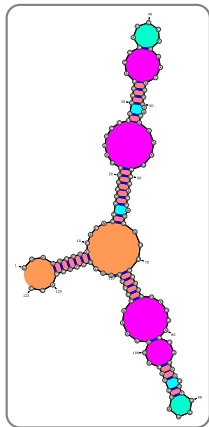
- ▶ Internal loops
- ▶ Bulges
- ▶ Terminal loops
- ▶ Multi loops
- ▶ Stackings

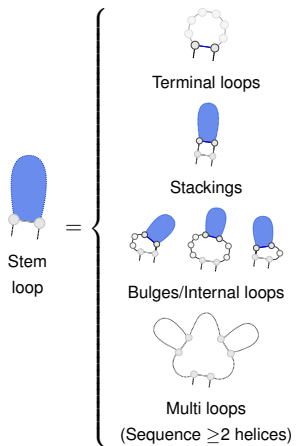


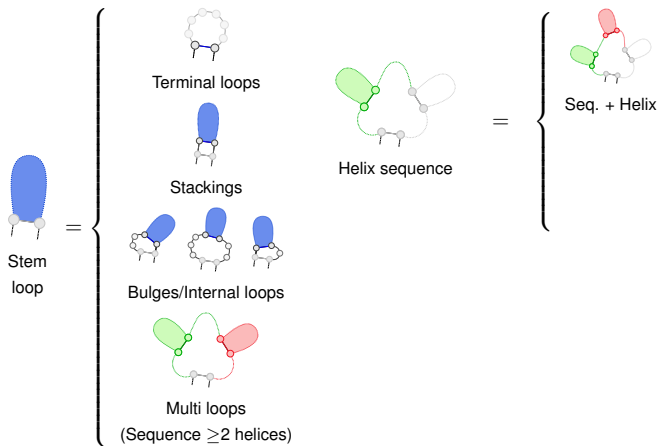
Free-energy ΔG of a loop depend on
bases, asymmetry, dangles ...

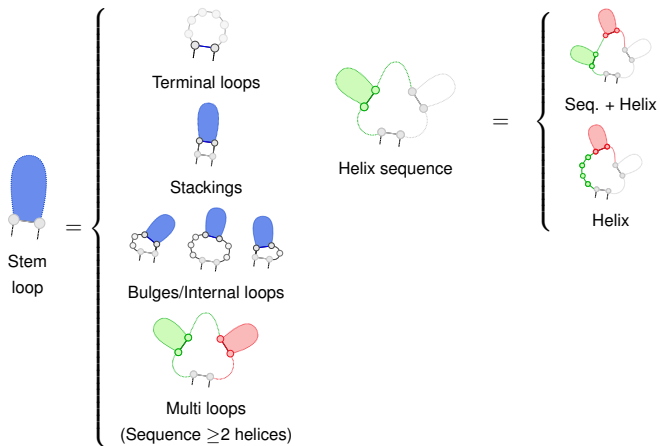
Experimentally determined
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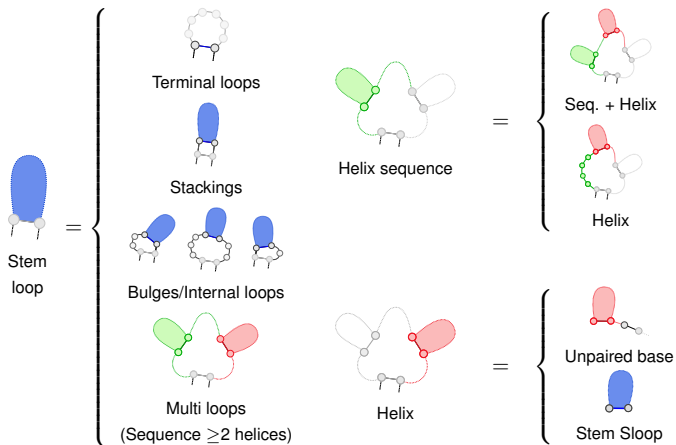
Improved results by taking stacking into account.



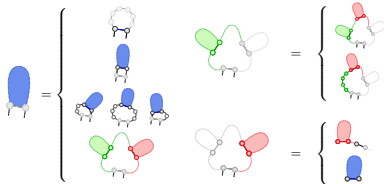








- ▶ $E_H(i, j)$: Energy of terminal loop *enclosed by* (i, j) pair
- ▶ $E_{BI}(i, j)$: Energy of bulge or internal loop *enclosed by* (i, j) pair
- ▶ $E_S(i, j)$: Energy of stacking $(i, j)/(i + 1, j - 1)$
- ▶ Penalty for multi loop (a), and occurrences of unpaired base (b) and helix (c) in multi loops.



DP recurrence

$$\begin{aligned}
 \mathcal{M}'_{i,j} &= \min \begin{cases} E_H(i, j) \\ E_S(i, j) + \mathcal{M}'_{i+1, j-1} \\ \text{Min}_{i', j'} (E_{BI}(i, i', j', j) + \mathcal{M}'_{i', j'}) \\ a + \text{Min}_k (\mathcal{M}_{i+1, k-1} + \mathcal{M}^1_{k, j-1}) \end{cases} \\
 \mathcal{M}_{i,j} &= \text{Min}_k \left\{ \min (\mathcal{M}_{i, k-1}, b(k-1)) + \mathcal{M}^1_{k, j} \right\} \\
 \mathcal{M}^1_{i,j} &= \text{Min}_k \left\{ b + \mathcal{M}^1_{i, j-1}, c + \mathcal{M}'_{i, j} \right\}
 \end{aligned}$$

Backtracking to reconstruct MFE structure:

$$\begin{aligned}
 \mathcal{M}'_{i,j} &= \text{Min} \left\{ \begin{array}{l} E_H(i,j) \\ E_S(i,j) + \mathcal{M}'_{i+1,j-1} \\ \text{Min}_{i',j'} (E_{BI}(i,i',j',j) + \mathcal{M}'_{i',j'}) \\ a + \text{Min}_k (\mathcal{M}_{i+1,k-1} + \mathcal{M}^1_{k,j-1}) \end{array} \right\} \\
 \mathcal{M}_{i,j} &= \text{Min}_k \left\{ \min (\mathcal{M}_{i,k-1}, b(k-1)) + \mathcal{M}^1_{k,j} \right\} \\
 \mathcal{M}^1_{i,j} &= \text{Min}_k \left\{ b + \mathcal{M}^1_{i,j-1}, c + \mathcal{M}'_{i,j} \right\}
 \end{aligned}$$

Complexity:

For each min, $\mathcal{O}(n)$ potential contributors

⇒ **Worst-case** complexity in $\mathcal{O}(n^2)$ for **naive backtrack**.

Keep best contributor for each Min ⇒ **Backtracking in $\mathcal{O}(n)$**

⇒ `Unafold` [MZ08]/`RNAfold` [HFS⁺94] compute the MFE for the Turner model in **overall**³ time/space complexities in $\mathcal{O}(n^3)/\mathcal{O}(n^2)$

³Using a trick/restriction for internal loops...

Backtracking to reconstruct MFE structure:

$$\begin{array}{c}
 \mathcal{M}'_{i,j} \leftarrow \text{Min} \left\{ \begin{array}{l} E_H(i,j) \\ E_S(i,j) + \mathcal{M}'_{i+1,j-1} \\ \text{Min}_{i',j'} (E_{BI}(i,i',j',j) + \mathcal{M}'_{i',j'}) \\ a + \text{Min}_k (\mathcal{M}_{i+1,k-1} + \mathcal{M}^1_{k,j-1}) \end{array} \right\} \\
 \mathcal{M}_{i,j} = \text{Min}_k \left\{ \min (\mathcal{M}_{i,k-1}, b(k-1)) + \mathcal{M}^1_{k,j} \right\} \\
 \mathcal{M}^1_{i,j} = \text{Min}_k \left\{ b + \mathcal{M}^1_{i,j-1}, c + \mathcal{M}'_{i,j} \right\}
 \end{array}$$

Complexity:

For each min, $\mathcal{O}(n)$ potential contributors

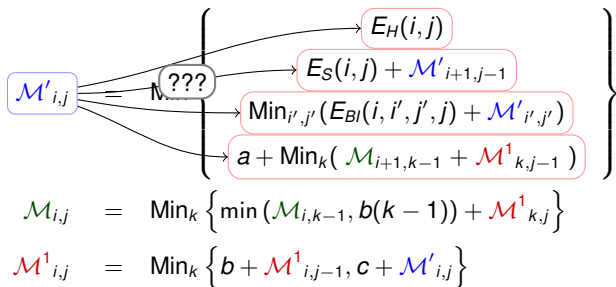
\Rightarrow **Worst-case** complexity in $\mathcal{O}(n^2)$ for **naive backtrack**.

Keep best contributor for each Min \Rightarrow **Backtracking in $\mathcal{O}(n)$**

\Rightarrow `Unafold` [MZ08]/`RNAfold` [HFS⁺94] compute the MFE for the Turner model in **overall**³ time/space complexities in $\mathcal{O}(n^3)/\mathcal{O}(n^2)$

³Using a trick/restriction for internal loops...

Backtracking to reconstruct MFE structure:



Complexity:

For each min, $\mathcal{O}(n)$ potential contributors

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Backtracking to reconstruct MFE structure:

$$\begin{aligned}
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 \mathcal{M}_{i,j} &= \text{Min}_k \left\{ \min (\mathcal{M}_{i,k-1}, b(k-1)) + \mathcal{M}^1_{k,j} \right\} \\
 \mathcal{M}^1_{i,j} &= \text{Min}_k \left\{ b + \mathcal{M}^1_{i,j-1}, c + \mathcal{M}'_{i,j} \right\}
 \end{aligned}$$

Complexity:

For each min, $\mathcal{O}(n)$ potential contributors

⇒ **Worst-case** complexity in $\mathcal{O}(n^2)$ for **naive backtrack**.

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Backtracking to reconstruct MFE structure:

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Starting from sequence, find conformation that minimizes free-energy.

Advantages:

- ▶ Mechanical nature allows the (in)validation of models
- ▶ Reasonable complexity $\mathcal{O}(n^3)/\mathcal{O}(n^2)$ time/space
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Limitations:

- ▶ Hard to include PKs
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- ▶ No cooperativity
- ▶ Limited performances

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Starting from homologous sequences, postulate common structure and find best possible tradeoff between folding & alignment.

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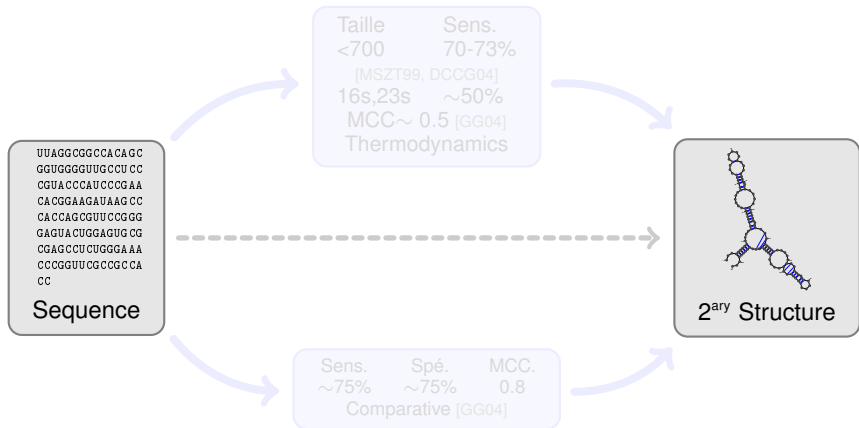
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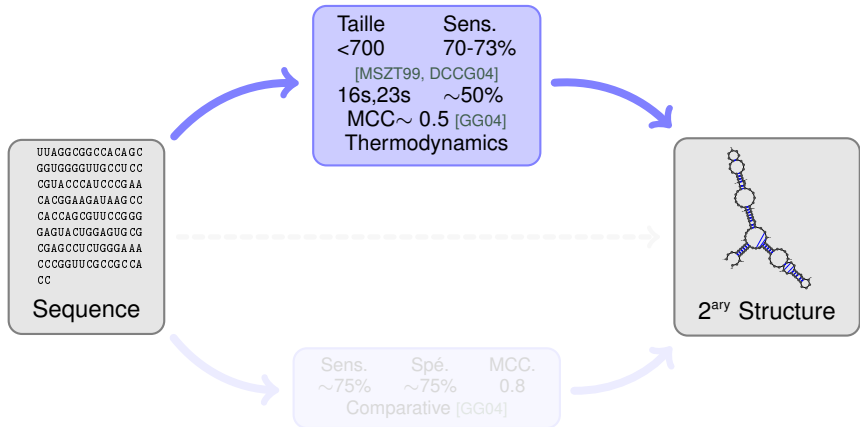
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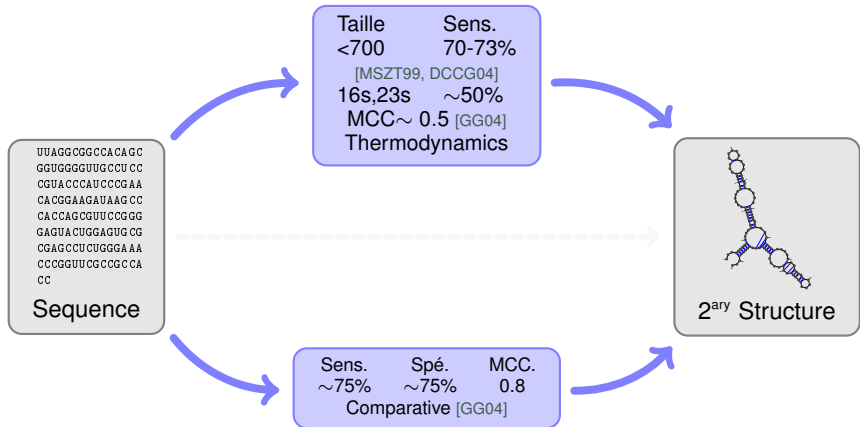
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Reminder:
$$MCC = \frac{t^+ t^- - f^+ f^-}{\sqrt{(t^+ + f^+)(t^+ + f^-)(t^- + f^+)(t^- + f^-)}}$$



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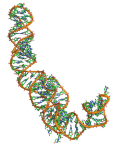
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Goal: From sequence to all-atom/coarse grain 3D models!!!

- ▶ Comparative models + Molecular dynamics: RNA2D3D [SYKB07]
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```
UUAGGCGGCCACAGC
GGUGGGGUUGCCUCC
CGUACCCAUCCGAA
CACGGAAGUAAGCC
CACCAGCGUUCGGG
GAGUACUGGAGUGCG
CGAGCCUCUGGGA AA
CCCGGUUCGCCGCA
CC
```

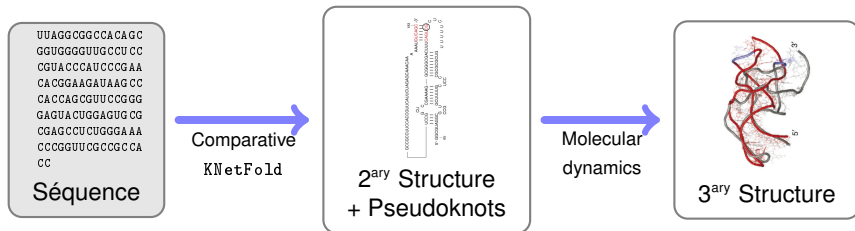
Séquence



3^{ary} Structure

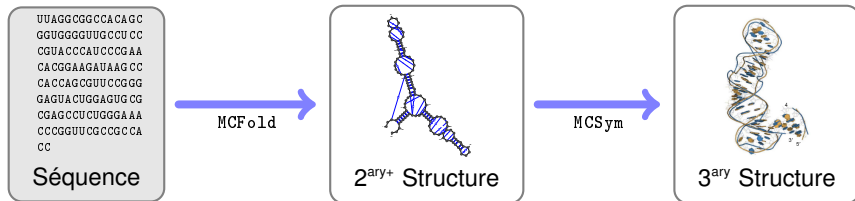
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1 Introduction

- Dynamic programming 101
- Dynamic programming framework

2 Variations on RNA folding

- Why RNA?
- RNA folding
- RNA Structure(s)
- Some representations of RNA structure
- Thermodynamics vs Kinetics

3 Free-energy minimization

- Nussinov-style RNA folding
- Turner energy model
- MFold/Unafold
- Performances and the comparative approach
- Towards a 3D ab-initio prediction

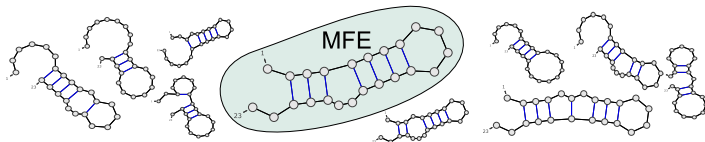
4 Boltzmann ensemble

- Nussinov: Minimisation \Rightarrow Counting
- Computing the partition function
- Statistical sampling

RNA *breathes* \Rightarrow There is no more than a single conformation.

New paradigm

The conformations of an RNA **coexist** in the **Boltzmann distribution**.



Consequence: The MFE probability can be arbitrarily small.

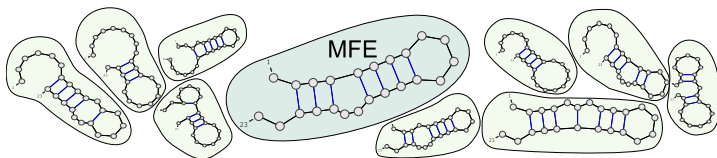
\Rightarrow To understand how RNA acts, one must account for the set of alternative structures.

In particular, structurally close structures may *ally*, and become the most realistic candidate in the search for a functional conformation.

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For each structure S compatible with an RNA ω , the Boltzmann distribution associates a **Boltzmann factor** $\mathcal{B}_{S,\omega} = e^{\frac{-E_{S,\omega}}{RT}}$, where:

- ▶ $E_{S,\omega}$ is the free-energy of S (kCal.mol⁻¹)
- ▶ T is the temperature (K)
- ▶ R is the perfect gas constant (1.986.10⁻³ kCal.K⁻¹.mol⁻¹)

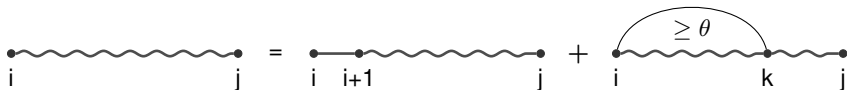
To obtain a distribution, one simply renormalizes by the **partition function**

$$\mathcal{Z}_\omega = \sum_{S \in \mathcal{S}_\omega} e^{\frac{-E_{S,\omega}}{RT}}$$

where \mathcal{S}_ω is the set of conformations that are compatibles with ω .

The **Boltzmann probability** of a structure S is simply given by

$$P_{S,\omega} = \frac{e^{\frac{-E_{S,\omega}}{RT}}}{\mathcal{Z}_\omega}.$$



$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \begin{cases} N_{i+1,j} & i \text{ unpaired} \\ \min_{k=i+\theta+1}^j \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{cases}$$

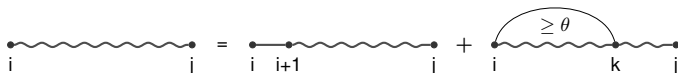
Ambiguity? Consider i : Either **unpaired**, or **paired** to k .

Sets of structures generated in these two cases are clearly disjoint.

(also holds for various values of k) \Rightarrow **Unambiguous** decomposition

Completeness? True, since scheme explores every possible outcome for i .

+ Induction on interval length \Rightarrow **Complete** decomposition



Recurrence for **minimal free-energy** of a fold :

$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \begin{cases} N_{i+1,j} & (i \text{ unpaired}) \\ \min_{k=i+\theta+1}^j E_{i,k} + N_{i+1,k-1} + N_{k+1,j} & (i \text{ comp. with } k) \end{cases}$$

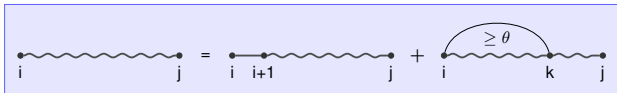
Recurrence for **counting compatible structures** :

$$C_{i,t} = 1, \quad \forall t \in [i, i + \theta]$$

$$C_{i,j} = \sum \begin{cases} C_{i+1,j} & (i \text{ unpaired}) \\ \sum_{k=i+\theta+1}^j 1 \times C_{i+1,k-1} \times C_{k+1,j} & (i \text{ comp. with } k) \end{cases}$$

Decomposition matters, and the rest (MFE, count. . .) follows!

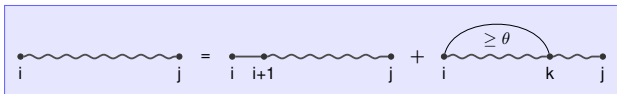
Partition function = Weighted count over compatible structures



$$\mathcal{Z}_{i,t} = 1, \quad \forall t \in [i, i + \theta]$$

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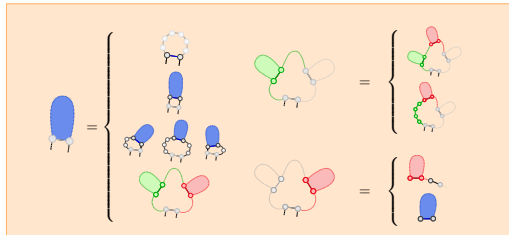
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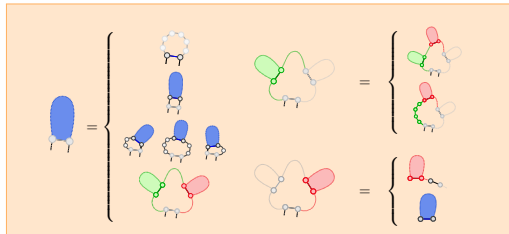
$$Z_{i,j} = \sum \left\{ \begin{array}{l} Z_{i+1,j} \\ \sum_{k=i+\theta+1}^j e^{\frac{-E_{bp}(i,k)}{RT}} \times Z_{i+1,k-1} \times Z_{k+1,j} \end{array} \right.$$

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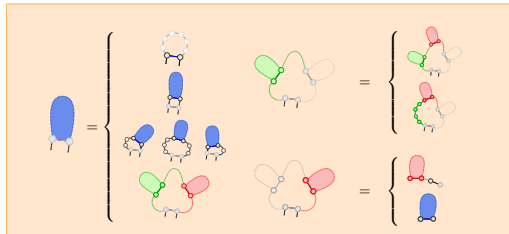
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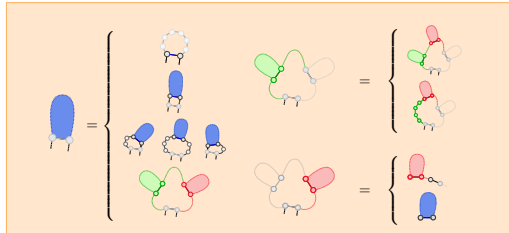
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Validity of a partition function computation:

- Completeness/Unambiguity of decomposition scheme
- Correctness of Boltzmann factor

Weight induced by backtrack = Product of derivations weights

$e^{-E/RT} \rightarrow$ Weight products \Leftrightarrow Summing energy terms

$$\begin{aligned} e^{-E_{bp}(i,k)/RT} \times \mathcal{Z}_{i+1,k-1} \times \mathcal{Z}_{k+1,j} &= \cdot \sum_x e^{-E(x)/RT} \cdot \sum_y e^{-E(y)/RT} \\ &= \sum_{x,y} e^{-a/RT} \cdot e^{-E(x)/RT} \cdot e^{-E(y)/RT} \\ &= \sum_{x,y} e^{-(E_{bp}(i,k)+E(x)+E(y))/RT} \end{aligned}$$

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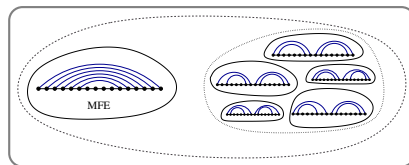
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MFE (\Leftrightarrow Max probability) may be **heavily dominated** by a set \mathcal{B} of **structurally similar** suboptimal structures.

\Rightarrow Functional conformation probably closer to \mathcal{B} than to MFE.



Proof-of-concept: [DCL05]

- ▶ Sample structures within Boltzmann probability
- ▶ Cluster structures
- ▶ Build and return consensus structure of the heaviest cluster

\Rightarrow Relative improvement for specificity (+17.6%) and sensitivity (+21.74%, except group II introns)

Problem

How to sample from the Boltzmann ensemble?

Goal [DL03]: From sequence ω , draw S with prob. $e^{-E_S/RT} / \mathcal{Z}$

Principle: Choose derivation with prob. prop. to its contribution to part. fun.

Precomputation: Compute part. fun. versions of matrices (\mathcal{Z} , \mathcal{Z}' , \mathcal{Z}^1).

Stochastic backtrack:

- 1 Draw uniform random number $r \in [0, \mathcal{Z}'(i, j))$
- 2 Subtract from r the contributions of $\mathcal{Z}'(i, j)$ until $r < 0$
- 3 Recurse over associated regions/matrices

$$\mathcal{Z}'(i, j) \equiv \boxed{???} \begin{cases} \rightarrow e^{\frac{-E_H(i, j)}{RT}} + e^{\frac{-E_S(i, j)}{RT}} \mathcal{Z}'(i+1, j-1) & \text{A} \\ \rightarrow \sum \left(e^{\frac{-E_{BL}(i, i', j', j)}{RT}} \mathcal{Z}'(i', j') \right) & \text{B} \\ \rightarrow e^{\frac{-(a)}{RT}} \sum (\mathcal{Z}(i+1, k-1) \mathcal{Z}^1(k, j-1)) & \text{C} \end{cases}$$

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$A_1 | A_2 | B_i | B_{i+1} | \dots | B_{j-1} | B_j | \boxed{C_i} | C_{i+1} | \dots | C_{j-1} | C_j$

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- 2 Subtract from r the contributions of $\mathcal{Z}'(i, j)$ until $r < 0$
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$A_1 | A_2 | B_i | B_{i+1} | \dots | B_{j-1} | B_j | C_i | C_{i+1} | \dots | C_{j-1} | C_j$

Goal [DL03]: From sequence ω , draw S with prob. $e^{-E_S/RT} / \mathcal{Z}$

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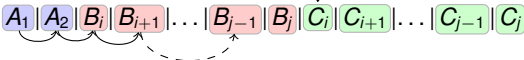
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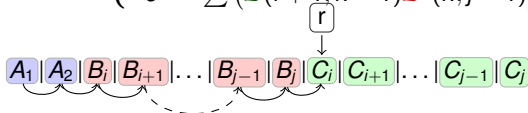
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Correctness: Each $S \in \mathcal{S}_\omega$ uniquely generated (DP scheme unambiguity)
Therefore the probability of generated S is

$$p_S = \frac{\mathcal{B}(E_1)}{\mathcal{B}(\mathcal{S}_\omega)} \cdot \frac{\mathcal{B}(E_2)}{\mathcal{B}(E_1)} \cdot \frac{\mathcal{B}(E_3)}{\mathcal{B}(E_2)} \cdots \frac{\mathcal{B}(\{S\})}{\mathcal{B}(E_m)}$$

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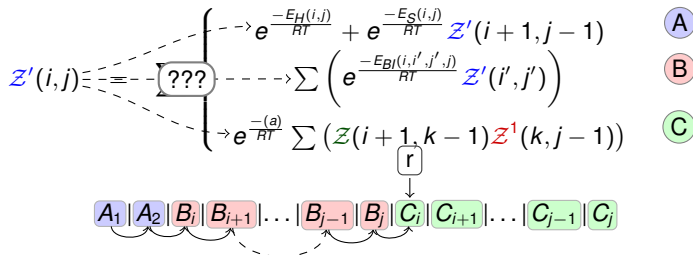
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Average-case complexity in $\Theta(k \times n\sqrt{n})$ (homopolymer model) [Pon08].

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After $\Theta(n)$ operations, recurse over region of length $n-1$
 \Rightarrow Worst-case complexity in $\mathcal{O}(k \times n^2)$ for k samples

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