M2 BIM – STRUCT - Lecture 2

Boltzmann equilibrium

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Paradigms in RNA structural bioinformatics



A – Kinetic Landscape

B – Evolution of concentrations

Continuous-time Markov chain

Given free-energy $E : \{A, C, G, U\}^* \times S \to \mathbb{R}$, at the Boltzmann equilibrium:

 $\mathbb{P}(S \mid w) \propto e^{-E(w,S)/RT}$

- Minimum Free-Energy (MFE): Relevant structure = Most stable/probable
- Partition function: Equilibrium properties of Boltzmann ensemble
- Kinetics: Finite-time evolution of concentrations/probabilities

Based on unambiguous decomposition of 2^{ary} structure into loops:

- Internal loops
- Bulges
- Terminal loops
- Multi loops
- Stackings

Free-energy Δ G of a loop depend on bases, assymmetry, dangles . . .

Experimentally determined + Interpolated for larger loops.



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MFold Unafold

- $E_H(i, j)$: Energy of terminal loop *enclosed by* (i, j) pair
- $E_{BI}(i, j)$: Energy of bulge or internal loop *enclosed by* (i, j) pair
- $E_S(i, j)$: Energy of stacking (i, j)/(i + 1, j 1)
- Penalty for multi loop (a), and occurrences of unpaired base (b) and helix (c) in multi loops.



DP recurrence

$$\begin{split} \mathcal{M}'_{i,j} &= \min \begin{cases} E_{\mathcal{H}}(i,j) \\ E_{S}(i,j) + \mathcal{M}'_{i+1,j-1} \\ \min_{i',j'}(E_{Bi}(i,i',j',j) + \mathcal{M}'_{i',j'}) \\ a + c + \min_{k}(\mathcal{M}_{i+1,k-1} + \mathcal{M}^{1}_{k,j-1}) \end{cases} \\ \mathcal{M}_{i,j} &= \min_{k} \left\{ \min(\mathcal{M}_{i,k-1}, b(k-1)) + \mathcal{M}^{1}_{k,j} \right\} \\ \mathcal{M}^{1}_{i,j} &= \min_{k} \left\{ b + \mathcal{M}^{1}_{i,j-1}, c + \mathcal{M}'_{i,j} \right\} \end{cases}$$

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Backtracking

Backtracking to reconstruct MFE structure:

Complexity:

For each min, $\mathcal{O}(n)$ potential contributors \Rightarrow Worst-case complexity in $\mathcal{O}(n^2)$ for naive backtrack. Keep best contributor for each Min \Rightarrow Backtracking in $\mathcal{O}(n)$

¹Using a trick/restriction for internal loops...

Backtracking to reconstruct MFE structure:

$$\mathcal{M}_{i,j}^{\prime} = \operatorname{Min}_{k-1} \left\{ \operatorname{min}_{i,k-1} \left\{ \operatorname{min}_{k,j}^{\prime} \left(\mathcal{E}_{Bl}(i,i',j',j) + \mathcal{M}_{i',j'}^{\prime} \right) - \left(\mathcal{E}_{s}(i,j) + \mathcal{H}_{i',j'}^{\prime} \right) \right) - \left(\mathcal{E}_{s}(i,j) + \mathcal{H}_{i',j'}^{\prime} \right) - \left(\mathcal{E}_{s}(i,j) + \mathcal{H}_{i',j'}^{\prime} \right) - \left(\mathcal{E}_{s}(i,j) + \mathcal{H}_{i',j'}^{\prime} \right) \right) - \left(\mathcal{E}_{s}(i,j) + \mathcal{H}_{i',j'}^{\prime} \right) - \left(\mathcal{E}_{s}(i,j) + \mathcal{H}_{i',j'}^{\prime} \right) \right) - \left(\mathcal{E}_{s}(i,j) + \mathcal{H}_{i',j'}^{$$

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$$\mathcal{M}'_{i,j} = \mathcal{M}_{i,j'} \{ \begin{array}{c} & & & & \\ & & & & \\ & & & \\ & & & & \\ & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\$$

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$$\mathcal{M}_{i,j} = \operatorname{Min}_{k} \left\{ \min(\mathcal{M}_{i,k-1}, b(k-1)) + \mathcal{M}^{1}_{k,j} \right\}$$
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$$\mathcal{M}_{i,j} \leftarrow = -\operatorname{Min}_{k} \left\{ \min_{i=1}^{m} (\mathcal{M}_{i,k=1}, \mathcal{B}(k-1)) + \mathcal{M}^{1}_{k,j} \right\}$$
$$\mathcal{M}_{i,j}^{1} \leftarrow = -\operatorname{Min}_{k} \left\{ -b + \mathcal{M}^{1}_{-i;j=1;} c + \widetilde{\mathcal{M}}'_{i,j} \right\}$$

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Outline

Turner energy mode MFold/Unafold

Boltzmann ensemble

Nussinov: Minimisation \Rightarrow Counting Computing the partition function Statistical sampling

The canonical Boltzmann Ensemble

RNA *breathes* \Rightarrow There is no more than a single conformation.

New paradigm

The conformations of an RNA coexist in the Boltzmann distribution.



Consequence: The MFE probability can be arbitrarily small.

 \Rightarrow To understand how RNA acts, one must account for the set of alternative structures.

In particular, structurally close structures may *ally*, and become the most realistic candidate in the search for a functional conformation.

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Boltzmann Distribution: Definition

For each structure *S* compatible with an RNA ω , the Boltzmann distribution associates a Boltzmann factor $\mathcal{B}_{S,\omega} = e^{\frac{-E_{S,\omega}}{RT}}$, where:

- $E_{S,\omega}$ is the free-energy *S* (kCal.mol⁻¹)
- ► *T* is the temperature (K)
- *R* is the perfect gaz constant (1.986.10⁻³ kCal.K⁻¹.mol⁻¹)

To obtain a distribution, one simply renormalizes by the partition function

$${\mathcal{Z}}_\omega = \sum_{{\mathcal{S}} \in {\mathcal{S}}_\omega} {oldsymbol{e}}^{-{oldsymbol{\mathcal{E}}}_{S,\omega}}$$

where S_{ω} is the set of conformations that are compatibles with ω .

The Boltzmann probability of a structure *S* is simply given by

$$P_{S,\omega} = rac{e^{rac{-E_{S,\omega}}{RT}}}{\mathcal{Z}_{\omega}}$$



$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \begin{cases} j & i \text{ unpaired} \\ \min_{k=i+\theta+1} \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{cases}$$

Ambiguity? Consider *i*: Either unpaired, or paired to *k*. Sets of structures generated in these two cases are clearly disjoint. (also holds for various values of k) \Rightarrow Unambiguous decomposition

Completeness?True, since scheme explores every possible outcome for *i*. + Induction on interval length \Rightarrow Complete decomposition



Recurrence for minimal free-energy of a fold :

$$\begin{aligned} &\mathcal{N}_{i,t} &= 0, \quad \forall t \in [i, i+\theta] \\ &\mathcal{N}_{i,j} &= \min \left\{ \begin{array}{ll} \mathcal{N}_{i+1,j} & (i \text{ unpaired}) \\ \min_{k=i+\theta+1}^{j} \mathcal{E}_{i,k} + \mathcal{N}_{i+1,k-1} + \mathcal{N}_{k+1,j} & (i \text{ comp. with } k) \end{array} \right. \end{aligned}$$

Recurrence for counting compatible structures :

$$C_{i,t} = 1, \quad \forall t \in [i, i + \theta]$$

$$C_{i,j} = \sum \begin{cases} C_{i+1,j} & (i \text{ unpaired}) \\ \sum_{k=i+\theta+1}^{j} 1 \times C_{i+1,k-1} \times C_{k+1,j} & (i \text{ comp. with } k) \end{cases}$$

Decomposition matters, and the rest (MFE, count...) follows!

$$\begin{aligned} \mathcal{Z}_{i,t} &= 1, \quad \forall t \in [i, i+\theta] \\ \mathcal{Z}_{i,j} &= \sum \begin{cases} \mathcal{Z}_{i+1,j} \\ \sum_{k=i+\theta+1}^{j} 1 \times \mathcal{Z}_{i+1,k-1} \times \mathcal{Z}_{k+1,j} \end{cases} \end{aligned}$$

$$\begin{aligned} \mathcal{Z}_{i,t} &= 1, \quad \forall t \in [i, i + \theta] \\ \mathcal{Z}_{i,j} &= \sum \begin{cases} \mathcal{Z}_{i+1,j} \\ \sum_{k=l+\theta+1}^{j} e^{-\frac{e_{\text{bp}}(i,k)}{RI}} \times \mathcal{Z}_{i+1,k-1} \times \mathcal{Z}_{k+1,j} \end{aligned}$$



$$\begin{aligned} \mathcal{M}'_{i,j} &= \operatorname{Min} \begin{cases} & E_{H}(i,j) \\ & E_{S}(i,j) + \mathcal{M}'_{i+1,j-1} \\ & \operatorname{Min}(E_{Bi}(i,i',j',j) + \mathcal{M}'_{i',j'}) \\ & a + c + \operatorname{Min}(\mathcal{M}_{i+1,k-1} + \mathcal{M}^{1}_{k,j-1}) \\ & \mathcal{M}_{i,j} &= \operatorname{Min} \left\{ \operatorname{Min}(\mathcal{M}_{i,k-1}, b(k-1)) + \mathcal{M}^{1}_{k,j} \right\} \\ & \mathcal{M}^{1}_{i,j} &= \operatorname{Min} \left\{ b + \mathcal{M}^{1}_{i,j-1}, c + \mathcal{M}'_{i,j} \right\} \end{aligned}$$



$$\mathcal{M}'_{i,j} = \operatorname{Min} \begin{cases} e^{\frac{-\mathcal{E}_{H}(i,j)}{\mathcal{H}'}} \\ e^{\frac{-\mathcal{E}_{H}(i,j)}{\mathcal{H}'}} + \mathcal{M}'_{i+1,j-1} \\ \operatorname{Min} \left(e^{\frac{-\mathcal{E}_{H}(i,j)}{\mathcal{H}'}} + \mathcal{M}'_{i',j'} \right) \\ e^{\frac{-(\phi+\phi)}{\mathcal{H}'}} + \operatorname{Min} \left(\mathcal{M}_{i+1,k-1} + \mathcal{M}^{1}_{k,j-1} \right) \\ \mathcal{M}_{i,j} = \operatorname{Min} \left\{ \operatorname{Min} \left(\mathcal{M}_{i,k-1}, e^{\frac{-\phi(k-1)}{\mathcal{H}'}} \right) + \mathcal{M}^{1}_{k,j} \right\} \\ \mathcal{M}^{1}_{i,j} = \operatorname{Min} \left\{ e^{\frac{-\phi}{\mathcal{H}'}} + \mathcal{M}^{1}_{i,j-1}, e^{\frac{-\phi}{\mathcal{H}'}} + \mathcal{M}'_{i,j} \right\} \end{cases}$$



$$\mathcal{M}'_{i,j} = \operatorname{Min} \begin{cases} e^{-\frac{E_{H}(i,j)}{H^{1}}} \\ e^{-\frac{E_{H}(i,j)}{H^{1}}} \mathcal{M}'_{i+1,j-1} \\ \operatorname{Min} \left(e^{-\frac{E_{H}(i,l',j',j)}{H^{1}}} \mathcal{M}'_{i',j'} \right) \\ e^{-\frac{(a+c)}{H^{1}}} \operatorname{Min} \left(\mathcal{M}_{i+1,k-1} \mathcal{M}^{1}_{k,j-1} \right) \\ \mathcal{M}_{i,j} = \operatorname{Min} \left\{ \operatorname{Min} \left(\mathcal{M}_{i,k-1}, e^{-\frac{a(k-1)}{H^{1}}} \right) \mathcal{M}^{1}_{k,j} \right\} \\ \mathcal{M}^{1}_{i,j} = \operatorname{Min} \left\{ e^{\frac{-b}{H^{1}}} \mathcal{M}^{1}_{i,j-1}, e^{\frac{-c}{H^{1}}} \mathcal{M}'_{i,j} \right\} \end{cases}$$



$$\begin{aligned} \mathcal{Z}'(i,j) &= \sum \begin{cases} e^{\frac{e}{RT}} \left(e^{\frac{e}{RT}} \right) \\ e^{\frac{-E_g(i,j)}{RT}} \mathcal{Z}'(i+1,j-1) \\ + \sum \left(e^{\frac{-E_g(i,j',j',j)}{RT}} \mathcal{Z}'(i',j') \right) \\ + e^{\frac{-(a+c)}{RT}} \sum \left(\mathcal{Z}(i+1,k-1)\mathcal{Z}^1(k,j-1) \right) \end{aligned} \\ \mathcal{Z}(i,j) &= \sum \left(\mathcal{Z}(i,k-1) + e^{\frac{-ck(k-1)}{RT}} \right) \mathcal{Z}^1(k,j) \\ \mathcal{Z}^1(i,j) &= e^{\frac{-b}{RT}} \mathcal{Z}^1(i,j-1) + e^{\frac{-cr}{RT}} \mathcal{Z}'(i,j) \end{aligned}$$

Partition function = Weighted count over compatible structures

$$\begin{aligned} \mathcal{Z}_{i,t} &= 1, \quad \forall t \in [i, i+\theta] \\ \mathcal{Z}_{i,j} &= \sum \left\{ \begin{array}{c} \mathcal{Z}_{i+1,j} \\ \sum_{k=i+\theta+1}^{j} e^{-\frac{\mathcal{E}_{bp}(i,k)}{RT}} \times \mathcal{Z}_{i+1,k-1} \times \mathcal{Z}_{k+1,j} \right. \end{aligned}$$

Validity of a partition function computation:

Completeness/Unambiguity of decomposition scheme

Correctness of Boltzmann factor Weight induced by backtrack = Product of derivations weights e^{-E/RT} → Weight products ⇔ Summing energy terms

$$e^{-E_{bp}(i,k)/RT} \times \mathcal{Z}_{i+1,k-1} \times \mathcal{Z}_{k+1,j} = \cdot \sum_{x} e^{-E(x)/RT} \cdot \sum_{y} e^{-E(y)/RT}$$
$$= \sum_{x,y} e^{-a/RT} \cdot e^{-E(x)/RT} \cdot e^{-E(y)/RT}$$
$$= \sum_{x,y} e^{-(E_{bp}(i,k)+E(x)+E(y))/RT}$$

Partition function = Weighted count over compatible structures

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$$^{-E/H} \rightarrow$$
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Statistical sampling of RNA 2ary structures

MFE (\Leftrightarrow Max probability) may be heavily dominated by a set \mathcal{B} of structurally similar suboptimal structures.

 \Rightarrow Functional conformation probably closer to ${\cal B}$ than to MFE.



Proof-of-concept: [DCL05]

- Sample structures within Boltzmann probability
- Cluster structures
- Build and return consensus structure of the heaviest cluster

 \Rightarrow Relative improvement for specificity (+17.6%) and sensitivity (+21.74%, except group II introns)

Problem

How to sample from the Boltzmann ensemble?

Goal [DL03]: From sequence ω , draw *S* with prob. $e^{-E_S/RT}/Z$

- 1. Draw uniform random number $r \in [0, \mathbb{Z}'(i, j))$
- 2. Subtract from *r* the contributions of $\mathcal{Z}'(i, j)$ until r < 0
- 3. Recurse over associated regions/matrices

$$\mathcal{Z}'(i,j) \in \left\{ \begin{array}{c} -- \Rightarrow e^{\frac{-E_{H}(i,j)}{RT}} + e^{\frac{-E_{S}(i,j)}{RT}} \mathcal{Z}'(i+1,j-1) \\ \Rightarrow e^{\frac{-(a+c)}{RT}} \sum \left(\mathcal{Z}(i+1,k-1) \mathcal{Z}^{1}(k,j-1) \right) \end{array} \right\}$$

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- 3. Recurse over associated regions/matrices

$$\mathcal{Z}'(i,j) = \sum \begin{cases} e^{\frac{-\mathcal{E}_{H}(i,j)}{RT}} + e^{\frac{-\mathcal{E}_{S}(i,j)}{RT}} \mathcal{Z}'(i+1,j-1) & \mathbb{A} \\ \sum \left(e^{\frac{-\mathcal{E}_{BI}(i,i',j',j)}{RT}} \mathcal{Z}'(i',j') \right) & \mathbb{B} \\ e^{\frac{-(a+c)}{RT}} \sum \left(\mathcal{Z}(i+1,k-1) \mathcal{Z}^{1}(k,j-1) \right) & \mathbb{C} \end{cases}$$

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$$A_{1} - A_{2} - B_{i} - B_{i+1} - \dots - B_{j-1} - B_{j} - C_{i} - C_{i+1} - \dots - C_{j-1} - C_{j} \end{cases}$$

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- **2**. Subtract from *r* the contributions of $\mathcal{Z}'(i, j)$ until r < 0
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$$\mathcal{Z}'(i,j) = \sum \begin{cases} e^{-\frac{E_{H}(i,j)}{RT}} + e^{-\frac{E_{S}(i,j)}{RT}} \mathcal{Z}'(i+1,j-1) & \mathbb{A} \\ \sum \left(e^{-\frac{E_{B}(i,i',j',j')}{RT}} \mathcal{Z}'(i',j') \right) & \mathbb{B} \\ e^{-\frac{(a+c)}{RT}} \sum \left(\mathcal{Z}(i+1,k-1)\mathcal{Z}^{1}(k,j-1) \right) & \mathbb{C} \\ \\ \mathbf{A}_{1} - \mathbf{A}_{2} - \frac{B_{i}}{B_{i}} - \frac{B_{i+1}}{B_{i+1}} - \dots - \frac{B_{j-1}}{B_{j}} - \frac{B_{j}}{C_{i}} - \frac{C_{i+1}}{C_{i+1}} - \dots - \frac{C_{j-1}}{C_{j}} - \frac{C_{j}}{C_{j}} \\ \end{cases}$$

Goal [DL03]: From sequence ω , draw *S* with prob. $e^{-E_S/RT}/Z$

- 1. Draw uniform random number $r \in [0, \mathbb{Z}'(i, j))$
- **2**. Subtract from *r* the contributions of $\mathcal{Z}'(i, j)$ until r < 0
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$$\mathcal{Z}'(i,j) = \sum \begin{cases} e^{-\frac{E_{H}(i,j)}{RT}} + e^{-\frac{E_{S}(i,j)}{RT}} \mathcal{Z}'(i+1,j-1) & \mathbb{A} \\ \sum \left(e^{-\frac{E_{B}(i,i',j',j')}{RT}} \mathcal{Z}'(i',j') \right) & \mathbb{B} \\ e^{-\frac{(a+c)}{RT}} \sum \left(\mathcal{Z}(i+1,k-1)\mathcal{Z}^{1}(k,j-1) \right) & \mathbb{C} \\ \end{bmatrix}$$

$$A_{1} - A_{2} - B_{i} - B_{i+1} - \dots - B_{j-1} - B_{j} - C_{i} - C_{i+1} - \dots - C_{j-1} - C_{j}$$

Goal [DL03]: From sequence ω , draw *S* with prob. $e^{-E_S/RT}/Z$

Principle: Choose derivation with prob. prop. to its contribution to part. fun. Precomputation: Compute part. fun. versions of matrices $(\mathcal{Z}, \mathcal{Z}', \mathcal{Z}^1)$. Stochastic backtrack:

- 1. Draw uniform random number $r \in [0, \mathbb{Z}'(i, j))$
- **2**. Subtract from *r* the contributions of $\mathcal{Z}'(i, j)$ until r < 0
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$$\mathcal{Z}'(i,j) = \sum \begin{cases} e^{\frac{-E_{H}(i,j)}{RT}} + e^{\frac{-E_{S}(i,j)}{RT}} \mathcal{Z}'(i+1,j-1) & \mathbb{A} \\ \sum \left(e^{\frac{-E_{BI}(i,i',j',j)}{RT}} \mathcal{Z}'(i',j') \right) & \mathbb{B} \\ e^{\frac{-(a+c)}{RT}} \sum \left(\mathcal{Z}(i+1,k-1) \mathcal{Z}^{1}(k,j-1) \right) & \mathbb{C} \end{cases}$$

Correctness: Each $S \in S_{\omega}$ uniquely generated (DP scheme unambiguity) Therefore the probability of generated *S* is

$$p_{S} = \frac{\mathcal{B}(E_{1})}{\mathcal{B}(S_{W})} \cdot \frac{\mathcal{B}(E_{2})}{\mathcal{B}(E_{1})} \cdot \frac{\mathcal{B}(E_{3})}{\mathcal{B}(E_{2})} \cdots \frac{\mathcal{B}(\{S\})}{\mathcal{B}(E_{m})}$$

Goal [DL03]: From sequence ω , draw *S* with prob. $e^{-E_S/RT}/Z$

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Correctness: Each $S \in S_{\omega}$ uniquely generated (DP scheme unambiguity) Therefore the probability of generated *S* is

$$p_{S} = \frac{1}{\mathcal{B}(\mathcal{S}_{W})} \cdot \frac{1}{1} \cdot \frac{1}{1} \dots \frac{\mathcal{B}(\{S\})}{1}$$

Goal [DL03]: From sequence ω , draw *S* with prob. $e^{-E_S/RT}/Z$

Principle: Choose derivation with prob. prop. to its contribution to part. fun. Precomputation: Compute part. fun. versions of matrices $(\mathcal{Z}, \mathcal{Z}', \mathcal{Z}^1)$. Stochastic backtrack:

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Correctness: Each $S \in S_{\omega}$ uniquely generated (DP scheme unambiguity) Therefore the probability of generated S is

$$p_{\mathcal{S}} = rac{\mathcal{B}(\{\mathcal{S}\})}{\mathcal{B}(\mathcal{S}_w)} = rac{e^{-E_{\mathcal{S}}/RT}}{\mathcal{Z}} = P_{\mathcal{S},\omega}$$

Complexity

Goal [DL03]: From sequence ω , draw *S* with prob. $e^{-E_S/RT}/Z$

Stochastic backtrack:

- 1. Draw uniform random number $r \in [0, \mathbb{Z}'(i, j))$
- **2**. Subtract from *r* the contributions of $\mathcal{Z}'(i, j)$ until r < 0
- 3. Recurse over associated regions/matrices

$$\mathcal{Z}'(i,j) \in \underbrace{\{--, -\}}_{RT} e^{-\frac{E_{H}(i,j)}{RT}} + e^{-\frac{E_{S}(i,j)}{RT}} \mathcal{Z}'(i+1,j-1)$$

$$(A) = \underbrace{\{--, -\}}_{RT} e^{-\frac{E_{H}(i,j)}{RT}} \mathcal{Z}'(i,j') = \underbrace{\{--, -\}}_{RT} e^{-\frac{E_{H}(i,j)}{RT}} e^{-\frac{E_{H}(i,j)}{RT}} \mathcal{Z}'(i,j') = \underbrace{\{--, -\}}_{RT} e^{-\frac{E_{H}(i,j)}{RT}} e^{-\frac{E_{H}$$

Average-case complexity in $\Theta(k \times n\sqrt{n})$ (homopolymer model) [Pon08]. Boustrophedon search $\Rightarrow O(k \times n \log n)$ worst-case [Pon08].

Complexity

Goal [DL03]: From sequence ω , draw *S* with prob. $e^{-E_S/RT}/Z$

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Average-case complexity in $\Theta(k \times n\sqrt{n})$ (homopolymer model) [Pon08]. Boustrophedon search $\Rightarrow O(k \times n \log n)$ worst-case [Pon08].

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