

M2 BIM – STRUCT – Lecture 1

Folding RNA *in silico*

Yann Ponty

AMIBio Team
CNRS LIX, École Polytechnique

Outline

Introduction

Dynamic programming 101

Why RNA?

RNA folding

RNA Structure(s)

Some representations of RNA structure

Some flavours of folding prediction

Thermodynamics vs Kinetics

Dynamic programming: Reminder

Free-energy minimization

Nussinov-style RNA folding

Foreword...

... or how to make a million bucks by giving change parsimoniously!!

Problem: You have access to unlimited amount of **1**, **20** and **50** cents coins. A client prefers to travel light, i.e. to **minimize the #coins**.

How to give **N** cents back in change without losing a customer?

Strategy #1: Start with *heaviest* coins, and then complete/fill-up with coins of *decreasing* value.

$$21 = ??$$

$$55$$

$$60$$

Foreword...

... or how to make a million bucks by giving change parsimoniously!!

Problem: You have access to unlimited amount of **1**, **20** and **50** cents coins. A client prefers to travel light, i.e. to **minimize the #coins**.

How to give **N** cents back in change without losing a customer?

Strategy #1: Start with *heaviest* coins, and then complete/fill-up with coins of *decreasing* value.

$$21 = \text{€}20 + \text{€}1$$

55??

60

Foreword...

... or how to make a million bucks by giving change parsimoniously!!

Problem: You have access to unlimited amount of **1**, **20** and **50** cents coins. A client prefers to travel light, i.e. to **minimize the #coins**.

How to give **N** cents back in change without losing a customer?

Strategy #1: Start with *heaviest* coins, and then complete/fill-up with coins of *decreasing* value.

$$21 = \text{€}20 + \text{€}1$$

$$55 = \text{€}50 + \text{€}1 + \text{€}1 + \text{€}1 + \text{€}1 + \text{€}1$$

60??

Foreword...

... or how to make a million bucks by giving change parsimoniously!!

Problem: You have access to unlimited amount of **1**, **20** and **50** cents coins. A client prefers to travel light, i.e. to **minimize the #coins**.

How to give **N** cents back in change without losing a customer?

Strategy #1: Start with *heaviest* coins, and then complete/fill-up with coins of *decreasing* value.

$$21 = \text{€}20 + \text{€}1$$

$$55 = \text{€}50 + \text{€}1 + \text{€}1 + \text{€}1 + \text{€}1 + \text{€}1$$

$$60 = \text{€}50 + \text{€}1 ??$$

Foreword...

... or how to make a million bucks by giving change parsimoniously!!

Problem: You have access to unlimited amount of **1**, **20** and **50** cents coins. A client prefers to travel light, i.e. to **minimize the #coins**.

How to give **N** cents back in change without losing a customer?

Strategy #1: Start with *heaviest* coins, and then complete/fill-up with coins of *decreasing* value.

$$21 = \text{€}20 + \text{€}1$$

$$55 = \text{€}50 + \text{€}20 + \text{€}10 + \text{€}10 + \text{€}10 + \text{€}1$$

$$\begin{aligned} 60 &= \text{€}50 + \text{€}20 + \text{€}10 ?? \\ &= \text{€}20 + \text{€}20 + \text{€}20 ! \end{aligned}$$

Problem *a priori* (?) non-solvable using such a *greedy* approach, as a (simpler) problem is already NP-complete (thus Efficient solution \Rightarrow 1M\$).

Foreword...

Strategy #2: Brute force enumeration $\rightarrow \#Coins^N$ (Ouch!)

Strategy #3: The following recurrence gives the minimal number of coins:

$$\text{Min}\#\text{Coins}(N) = \text{Min} \left\{ \begin{array}{l} \text{£1 coin} \rightarrow 1 + \text{Min}\#\text{Coins}(N - 1) \\ \text{£20 coin} \rightarrow 1 + \text{Min}\#\text{Coins}(N - 20) \\ \text{£50 coin} \rightarrow 1 + \text{Min}\#\text{Coins}(N - 50) \end{array} \right.$$

With some memory (N intermediate computations), the minimum number of coins can be obtained after $N \times \#Coins$ operations. An actual set of coins can be reconstructing by **tracing back** the choices performed at each stage, leading to the minimum.

Remark: We still haven't won the million, as N has **exponential value compared to the length of its encoding**, so the algorithm does not qualify as *efficient* (i.e. polynomial).

Still, this approach is much more efficient than a brute-force enumeration:
 \Rightarrow Dynamic programming.

Foreword...

Strategy #2: Brute force enumeration $\rightarrow \#Coins^N$ (Ouch!)

Strategy #3: The following recurrence gives the minimal number of coins:

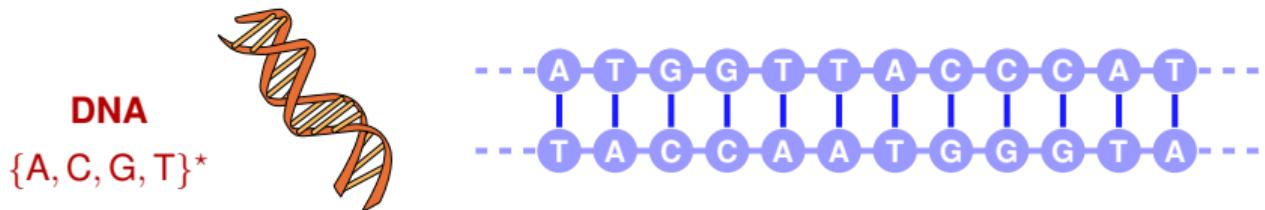
$$\text{Min}\#\text{Coins}(N) = \text{Min} \left\{ \begin{array}{l} \text{1 cent coin} \rightarrow 1 + \text{Min}\#\text{Coins}(N - 1) \\ \text{2 cent coin} \rightarrow 1 + \text{Min}\#\text{Coins}(N - 20) \\ \text{5 cent coin} \rightarrow 1 + \text{Min}\#\text{Coins}(N - 50) \end{array} \right.$$

With some memory (N intermediate computations), the minimum number of coins can be obtained after $N \times \#Coins$ operations. An actual set of coins can be reconstructing by **tracing back** the choices performed at each stage, leading to the minimum.

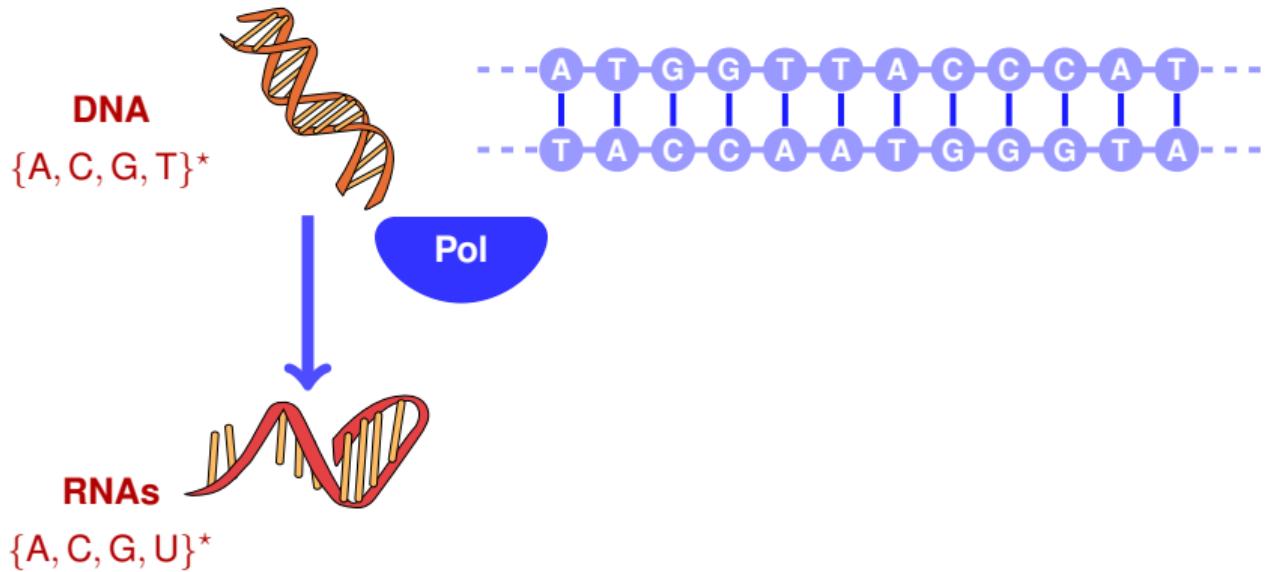
Remark: We still haven't won the million, as N has **exponential value compared to the length of its encoding**, so the algorithm does not qualify as *efficient* (i.e. polynomial).

Still, this approach is much more efficient than a brute-force enumeration:
 \Rightarrow Dynamic programming.

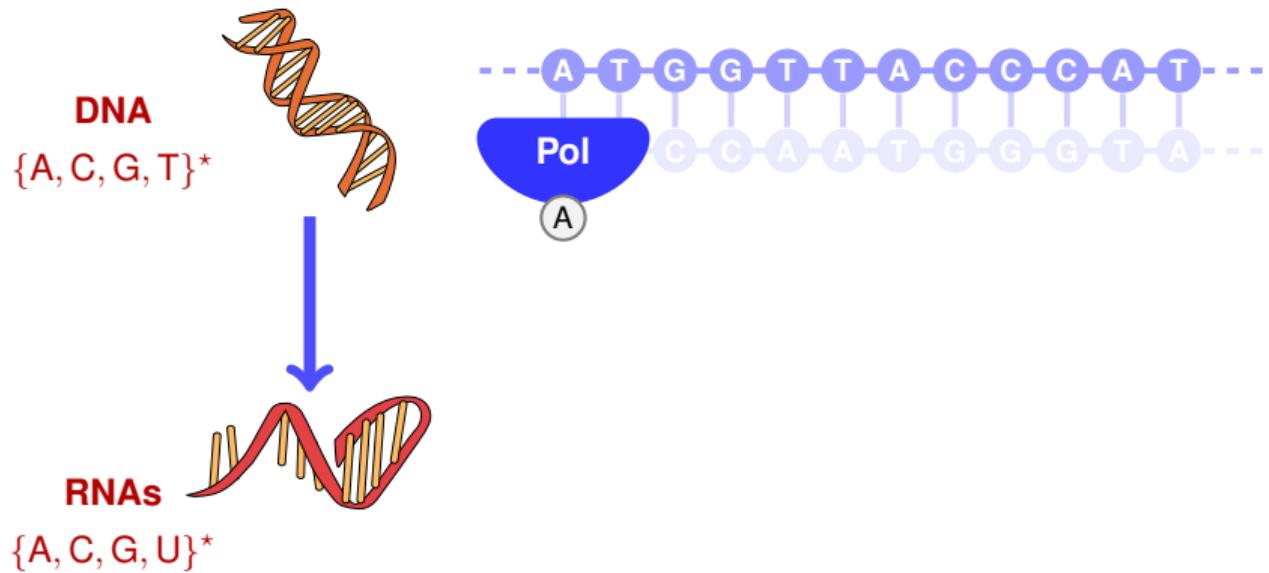
Fundamental dogma of molecular biology



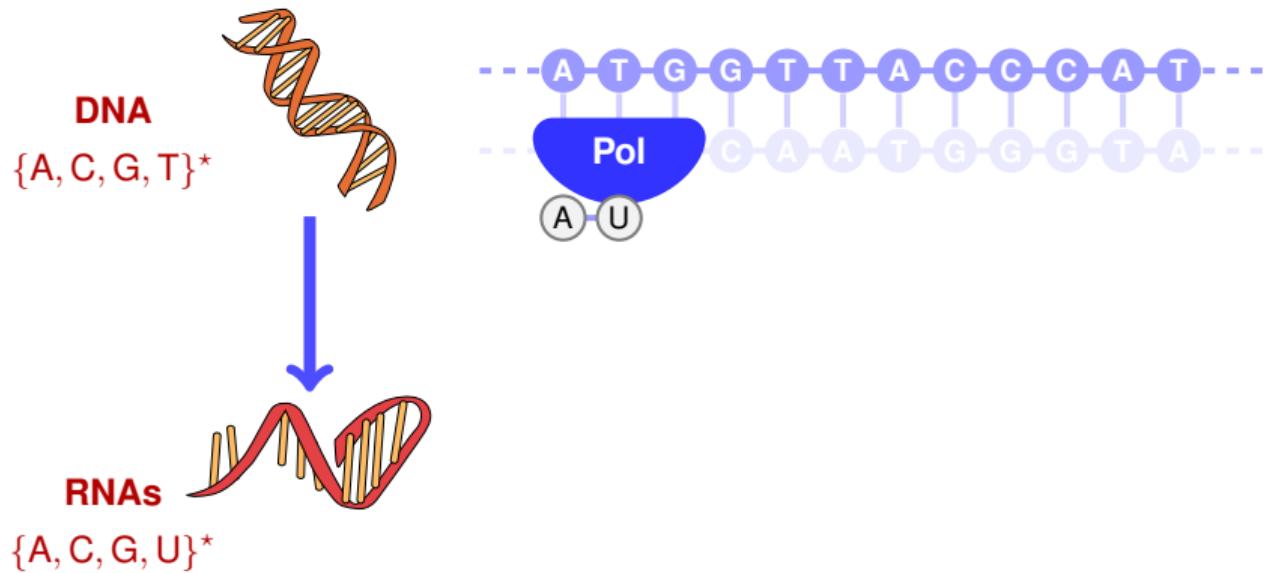
Fundamental dogma of molecular biology



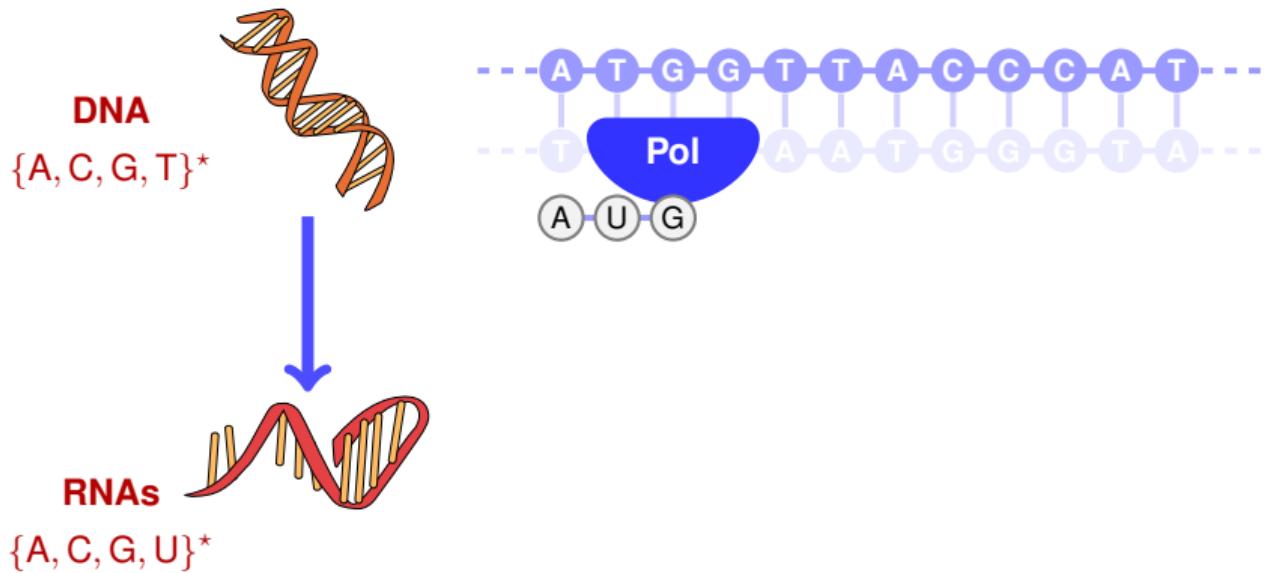
Fundamental dogma of molecular biology



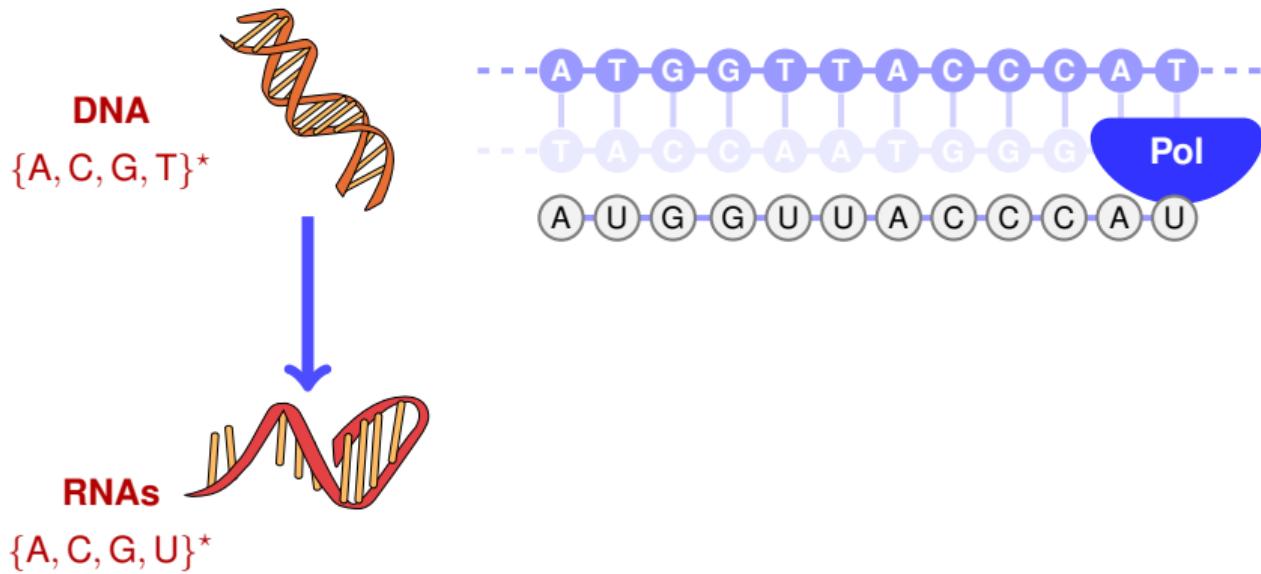
Fundamental dogma of molecular biology



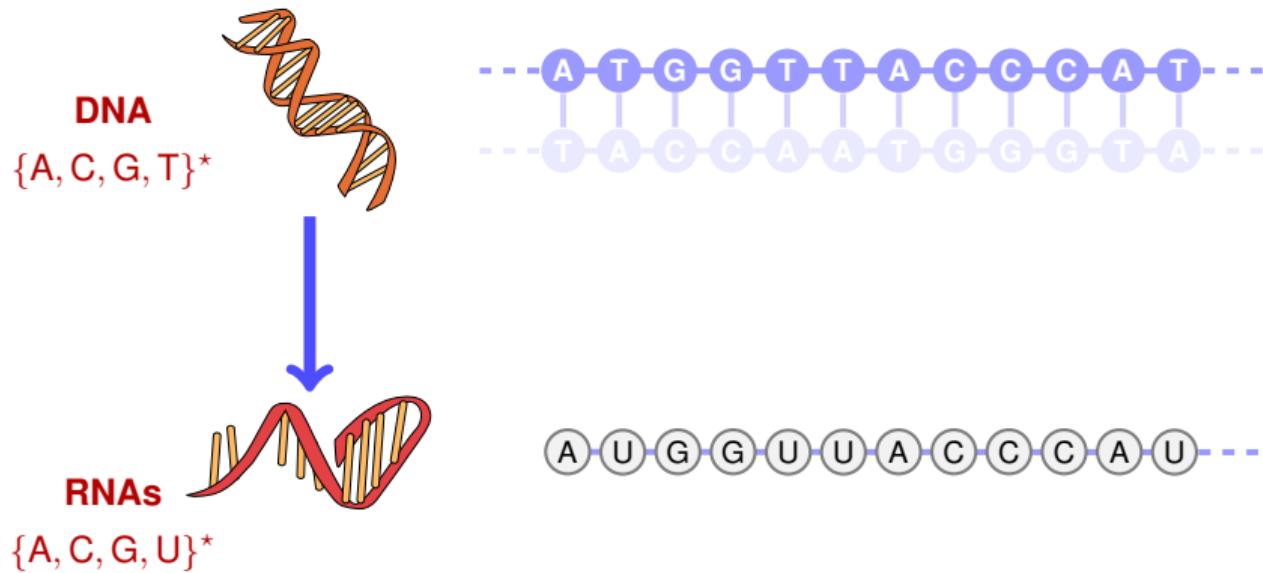
Fundamental dogma of molecular biology



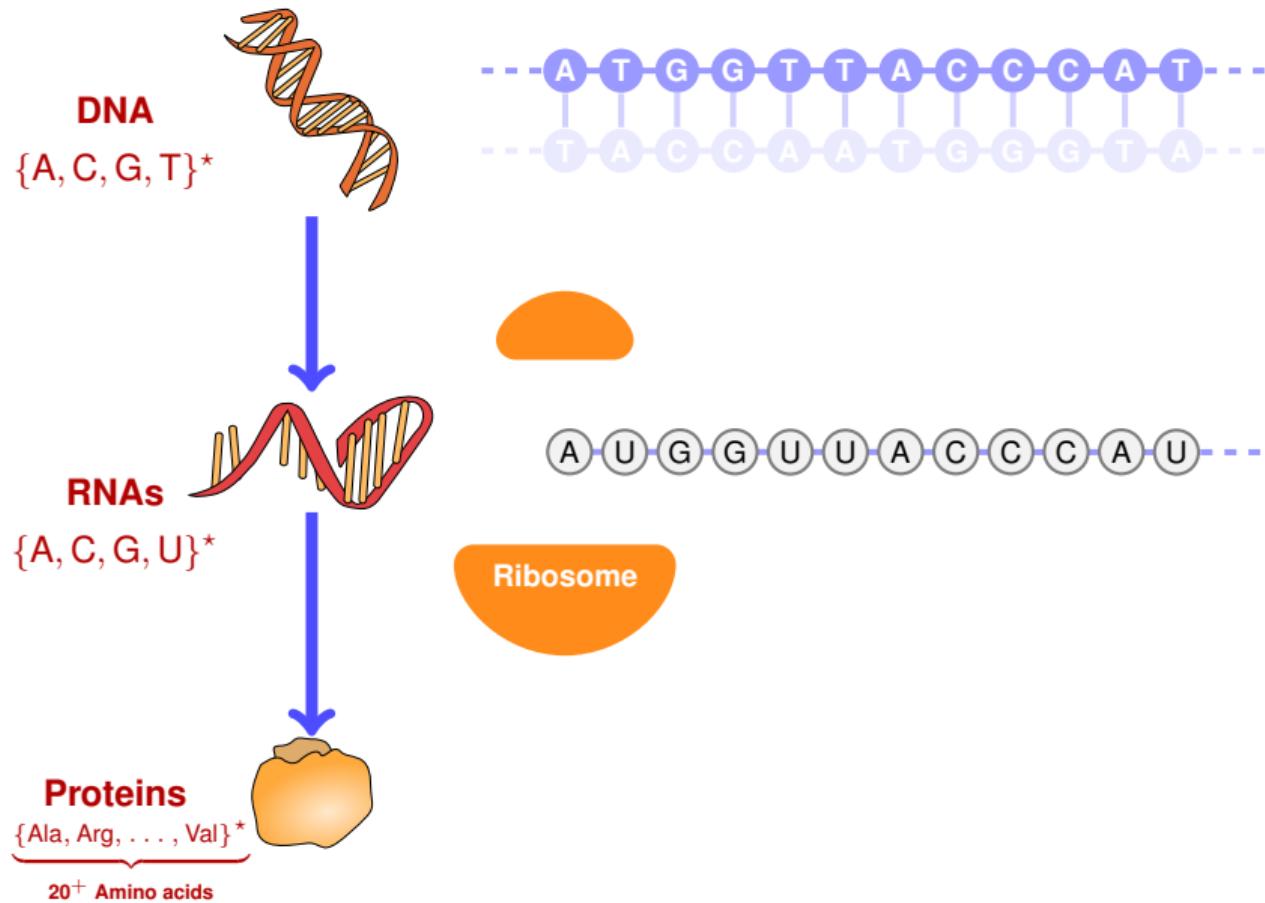
Fundamental dogma of molecular biology



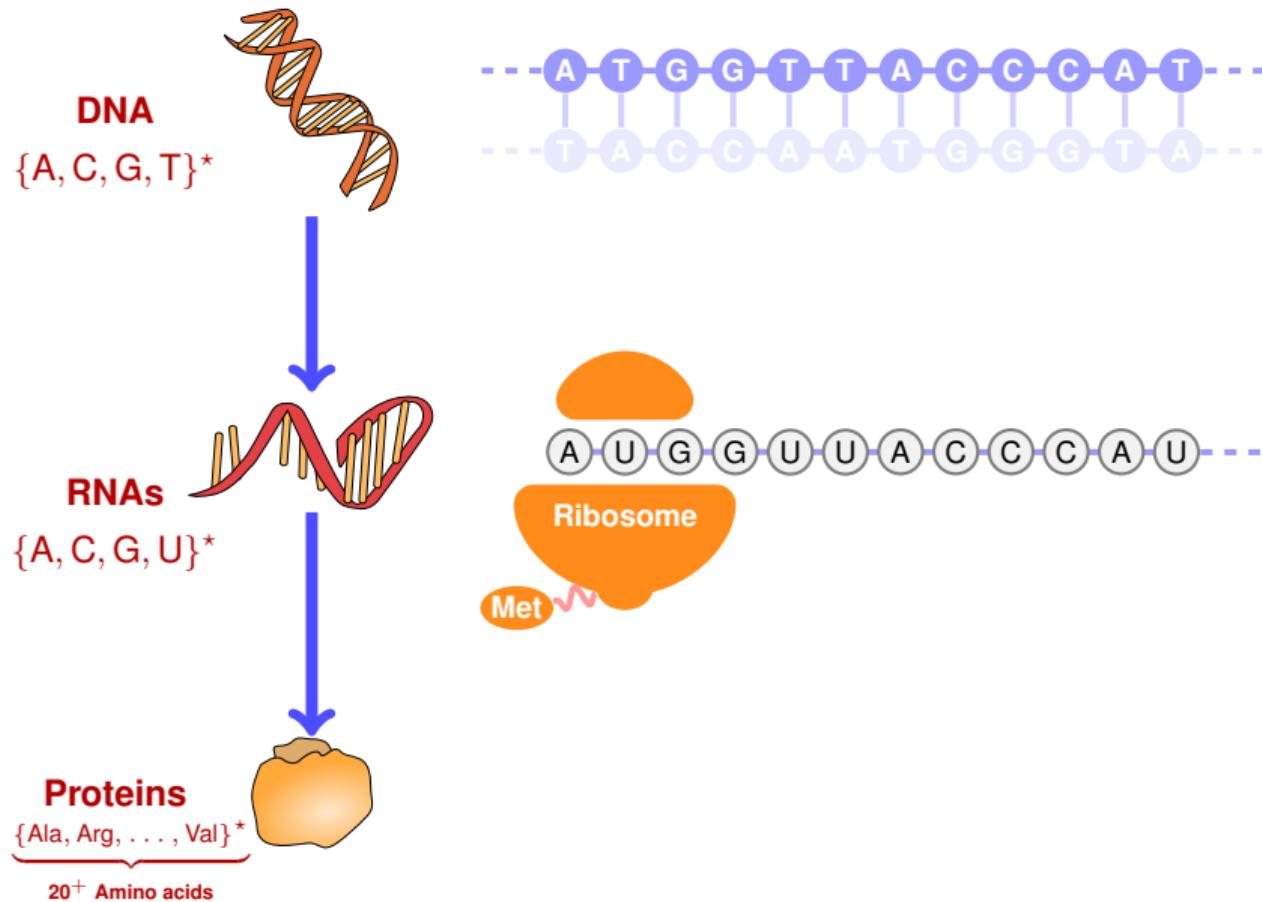
Fundamental dogma of molecular biology



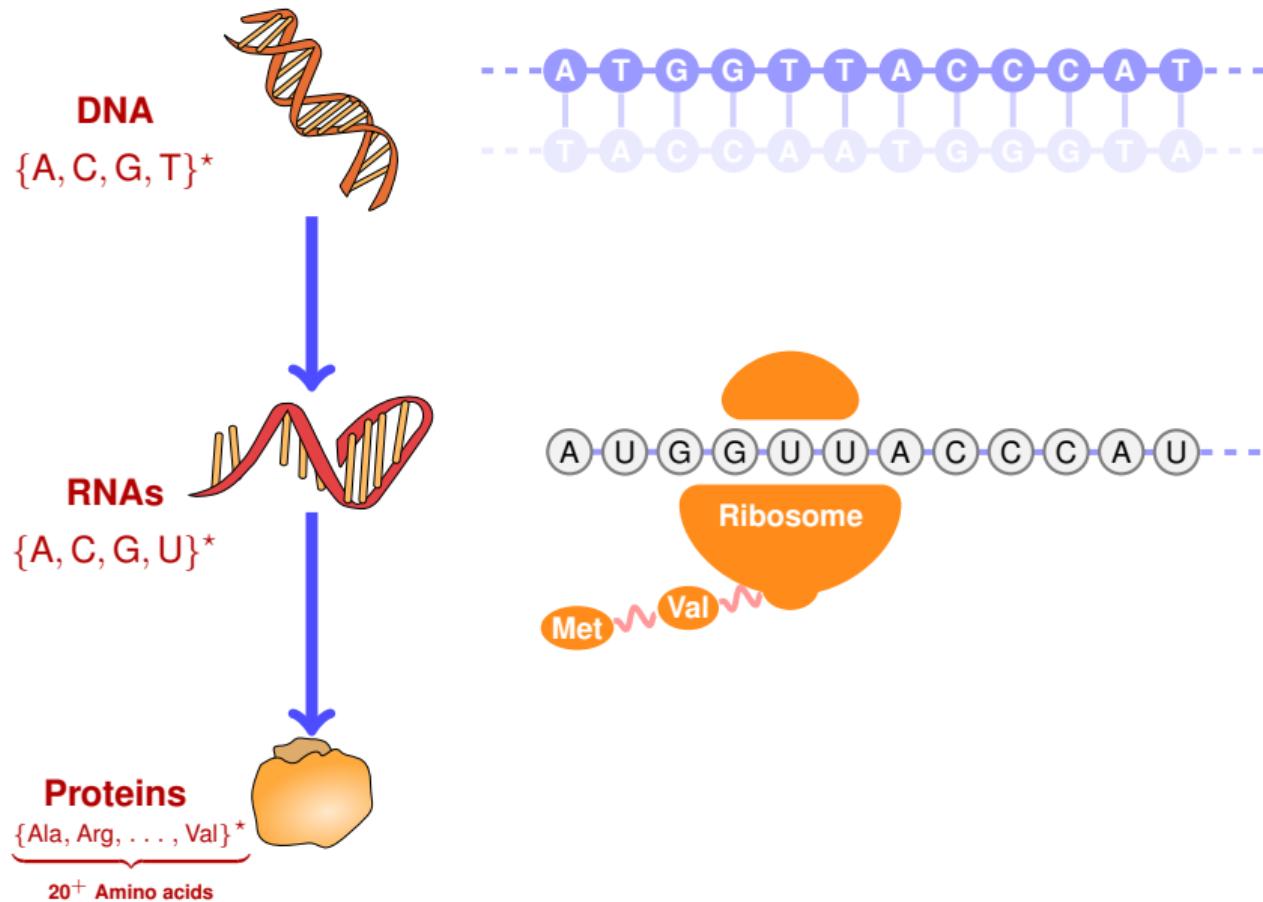
Fundamental dogma of molecular biology



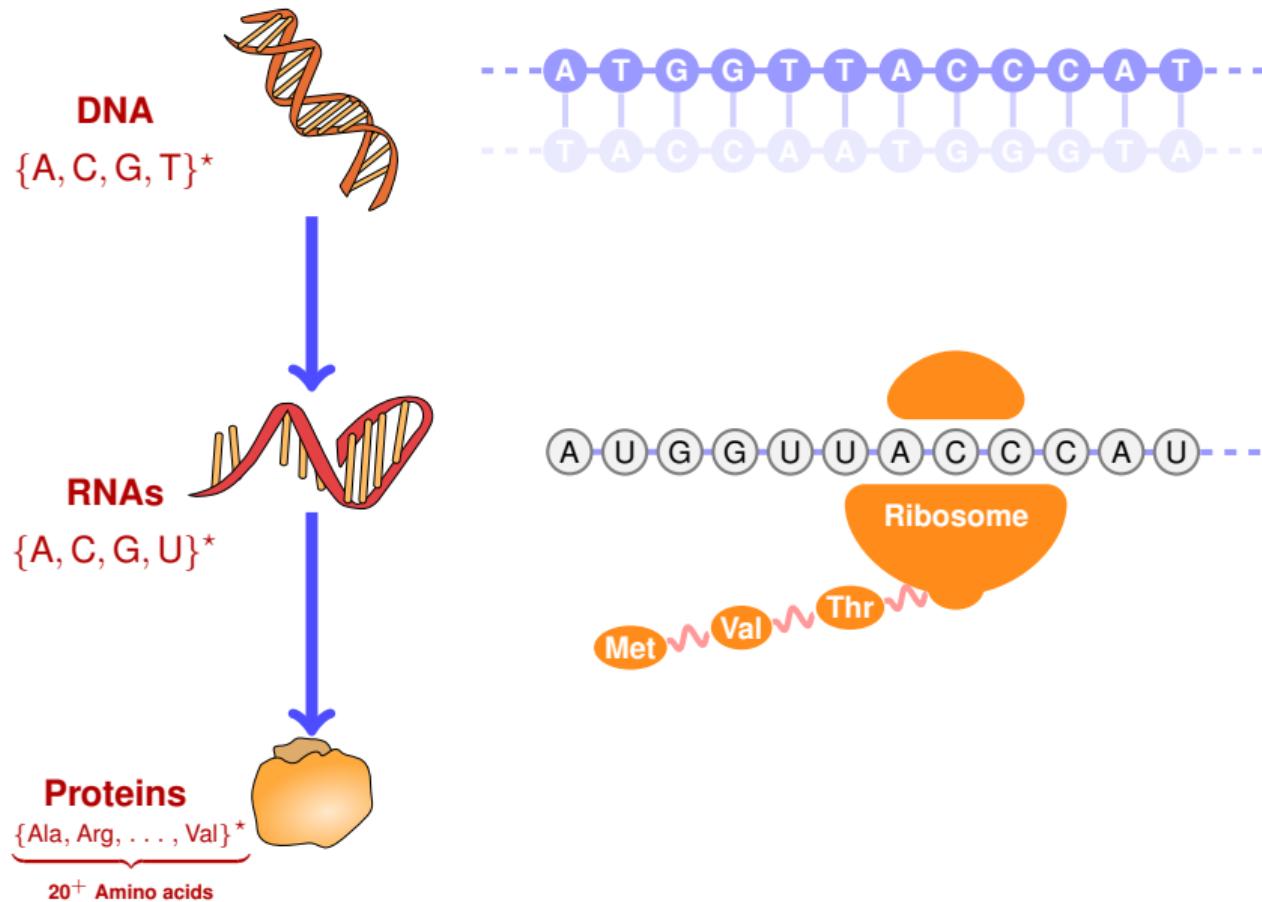
Fundamental dogma of molecular biology



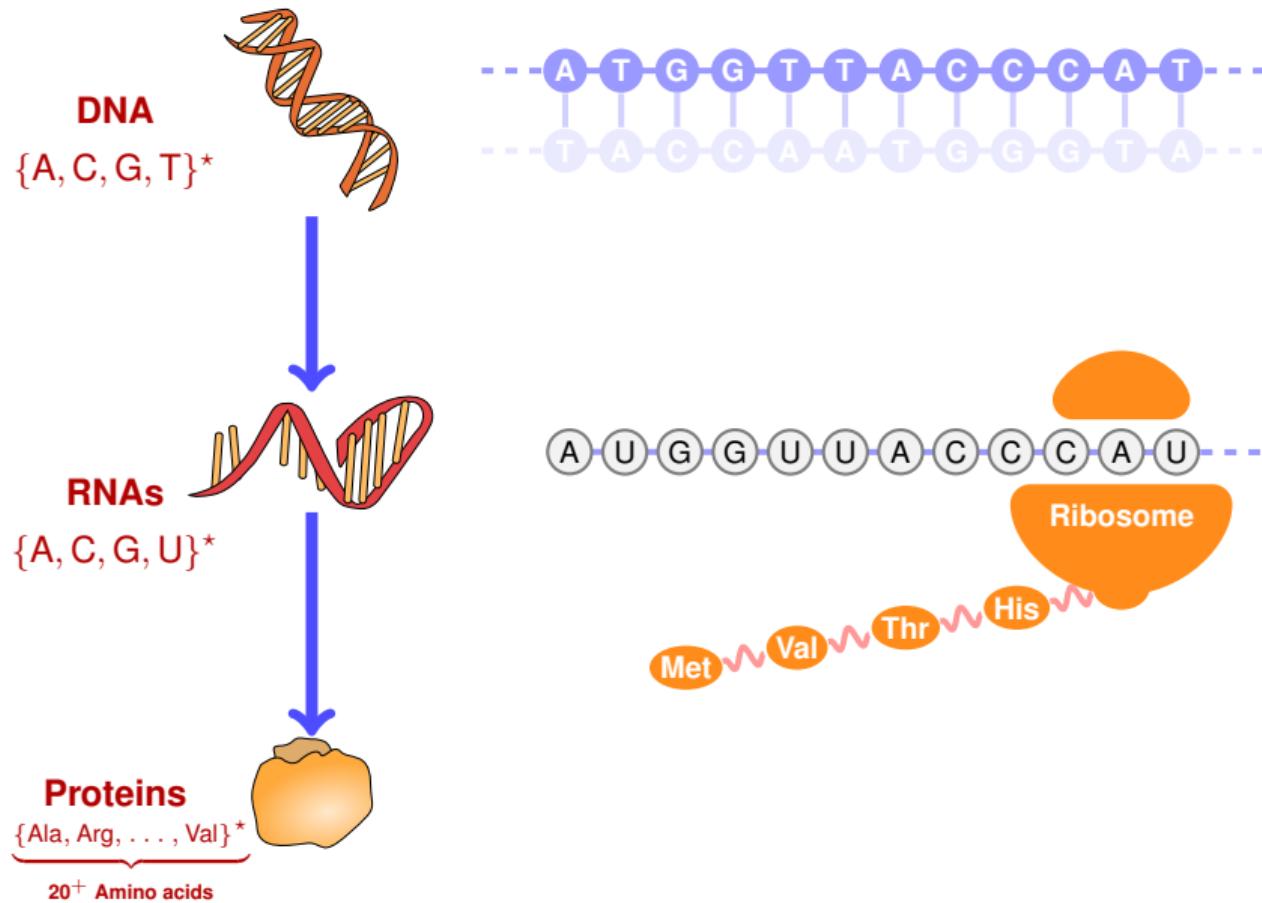
Fundamental dogma of molecular biology



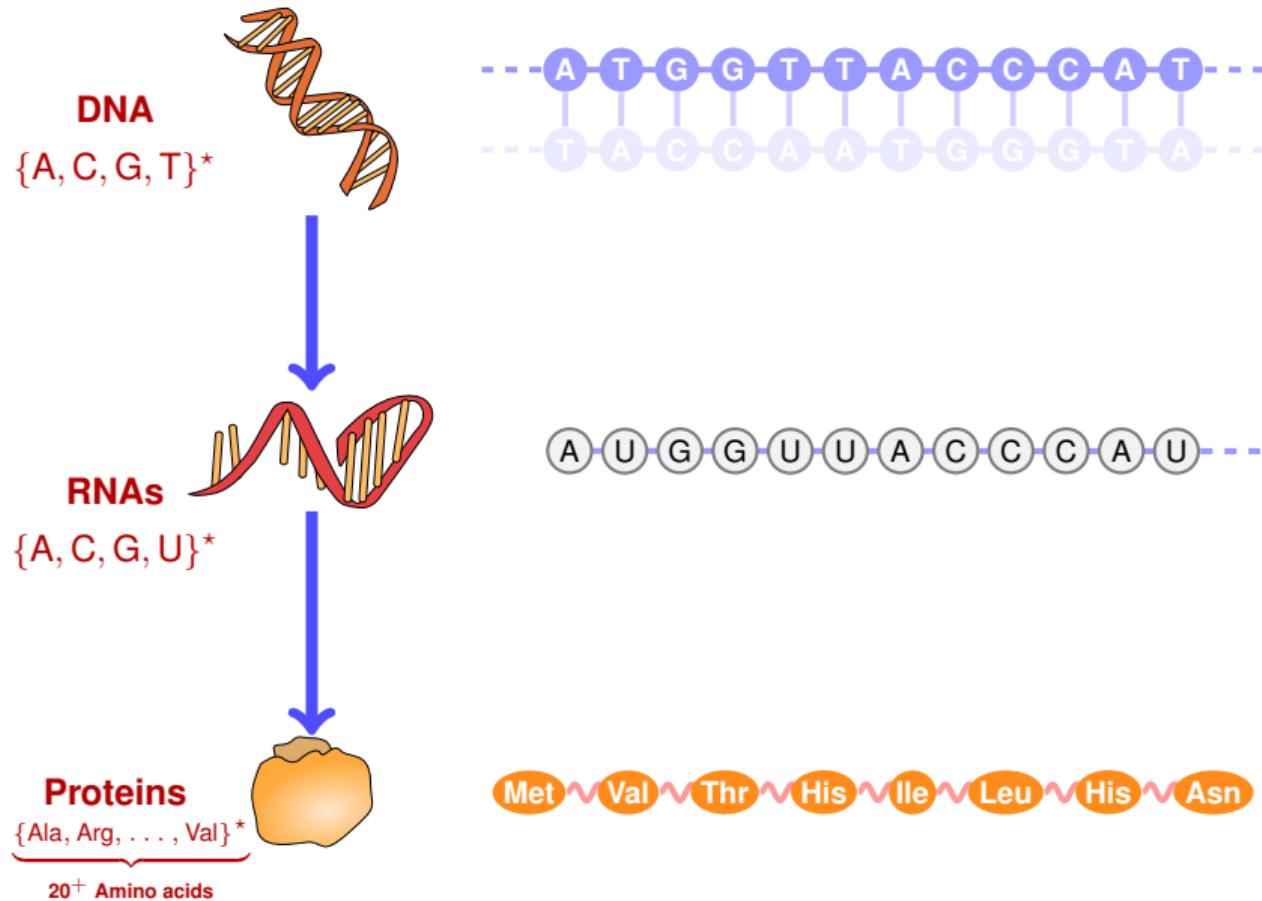
Fundamental dogma of molecular biology



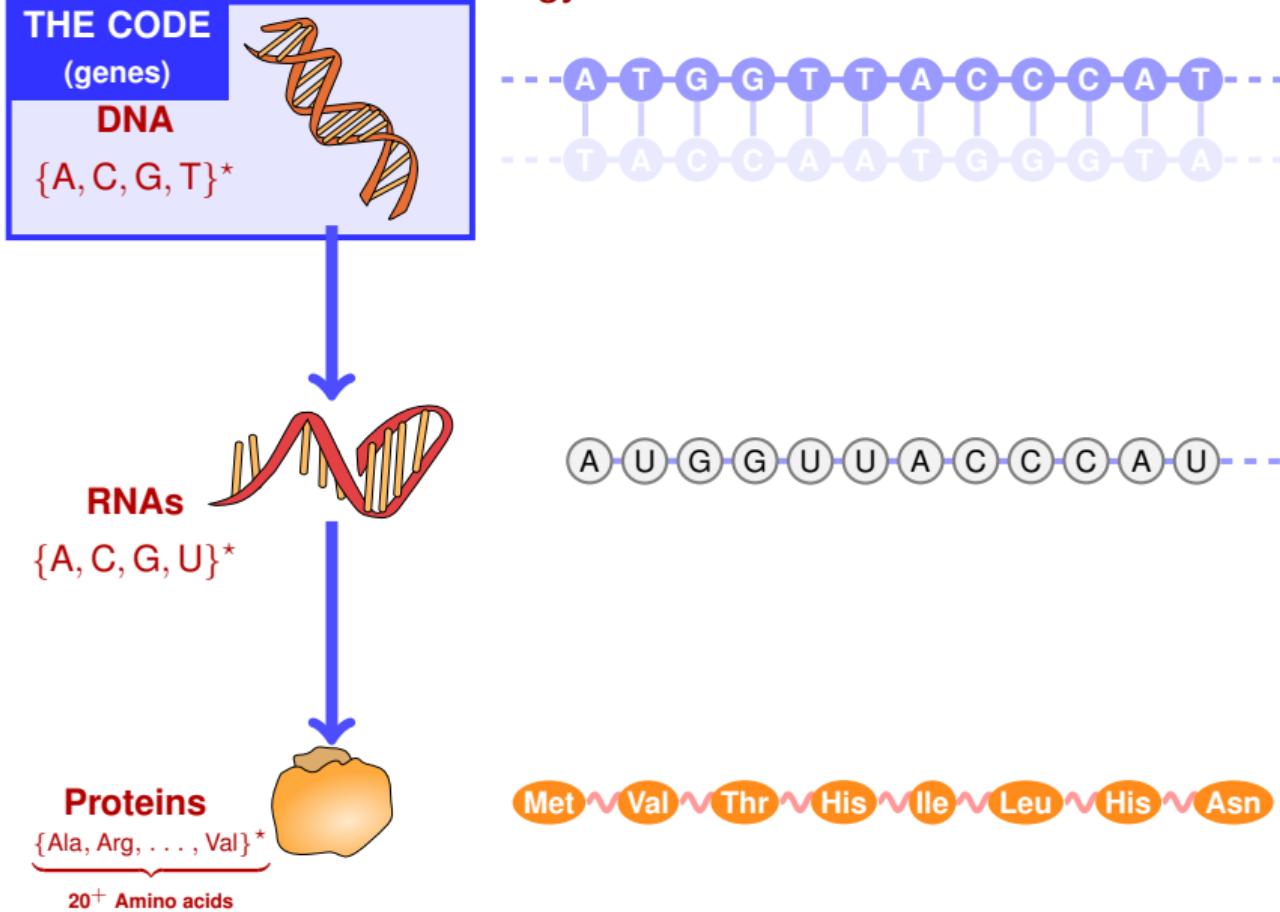
Fundamental dogma of molecular biology



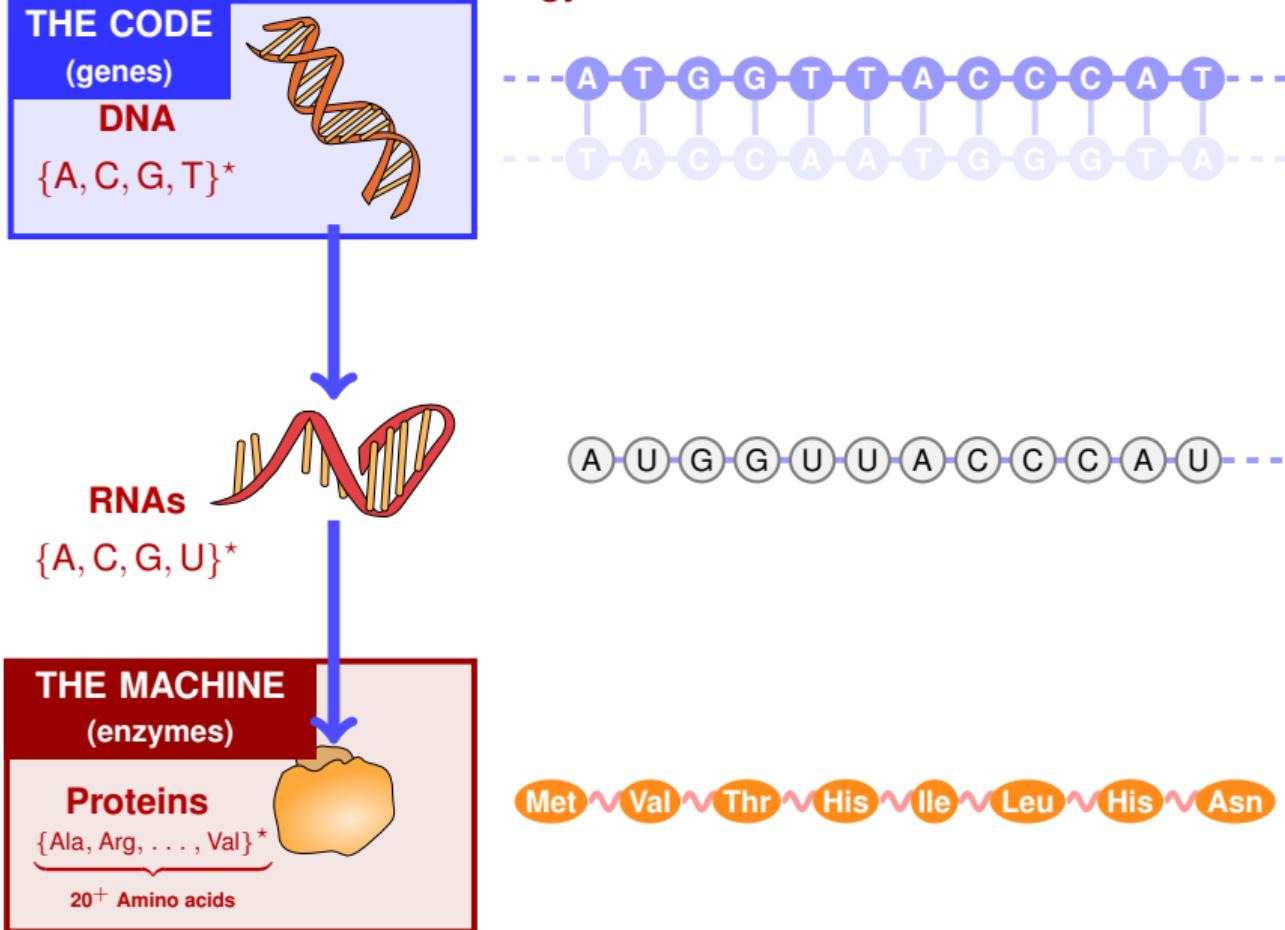
Fundamental dogma of molecular biology



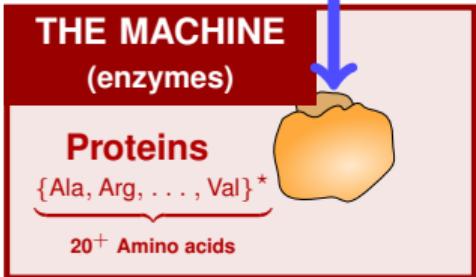
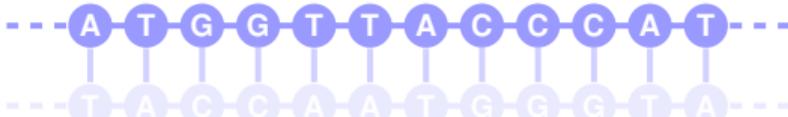
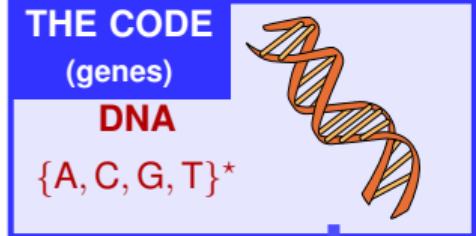
Fundamental dogma of molecular biology



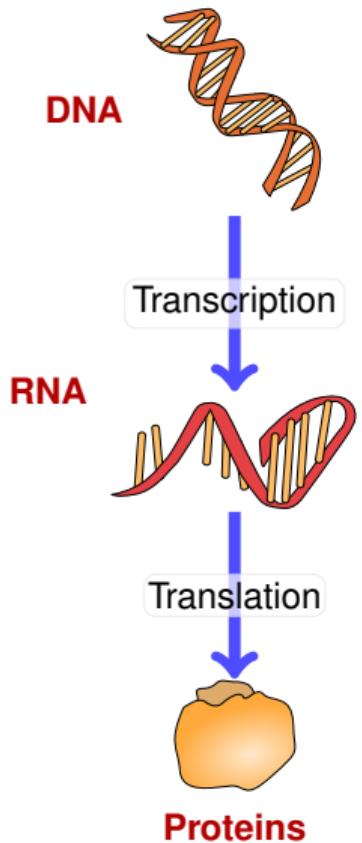
Fundamental dogma of molecular biology



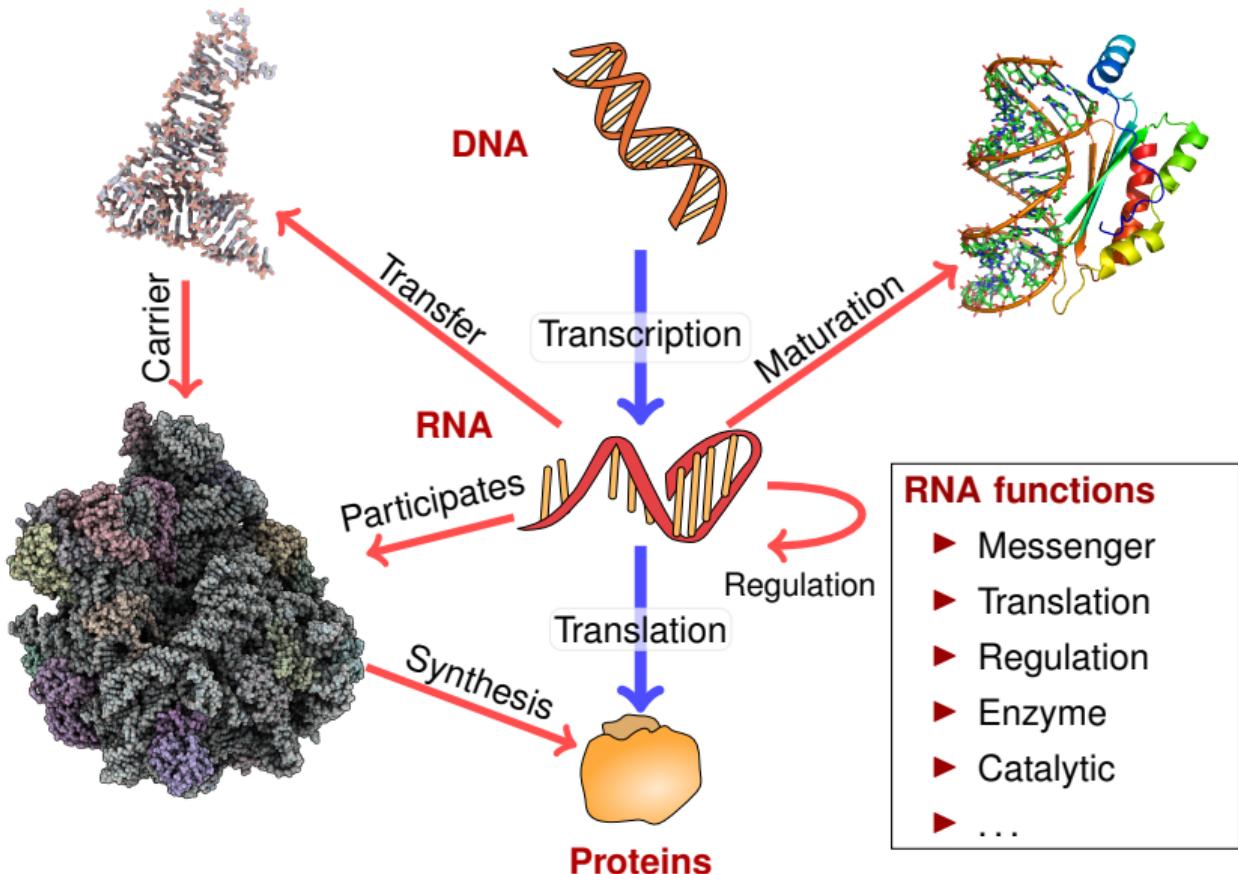
Fundamental dogma of molecular biology



Fundamental dogma of molecular biology



Fundamental dogma of molecular biology



RiboNucleic Acids (RNAs) in Human biology/health: Friends **and** Foes!

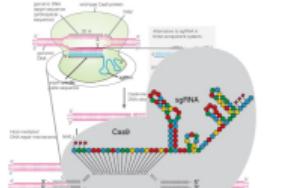


RiboNucleic Acids (RNAs) in Human biology/health: Friends **and** Foes!

Targeting system for DNA Editing

CRISPR therapies

Sickle-cell anemia, β -thalassamia, Leber congenital amaurosis (LCA), cancers...



Hendel et al, 2015; Agrotis & Ketteler, 2015

RiboNucleic Acids (RNAs)



Encodes proteins

mRNA Vaccines

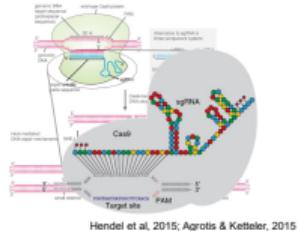
COVID-19, Malaria (Zika, CMV, Cancers?)

RiboNucleic Acids (RNAs) in Human biology/health: Friends **and** Foes!

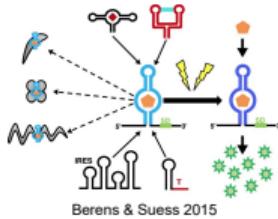
Targeting system for DNA Editing

CRISPR therapies

Sickle-cell anemia, β -thalassamia, Leber congenital amaurosis (LCA), cancers...



Hendel et al, 2015; Agrotis & Ketteler, 2015



Sensor of metabolites
Riboswitches

RiboNucleic Acids (RNAs)



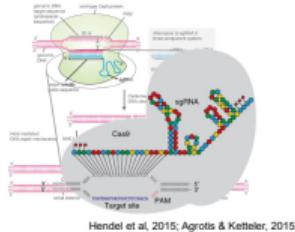
Encodes proteins
mRNA Vaccines
COVID-19, Malaria (Zika, CMV, Cancers?)

RiboNucleic Acids (RNAs) in Human biology/health: Friends **and** Foes!

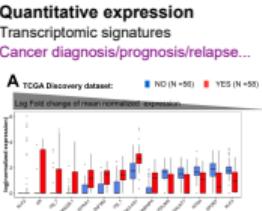
Targeting system for DNA Editing

CRISPR therapies

Sickle-cell anemia, β -thalassamia, Leber congenital amaurosis (LCA), cancers...



Sensor of metabolites
Riboswitches



[NGuyen et al, 2021]

RiboNucleic Acids (RNAs)



Encodes proteins

mRNA Vaccines

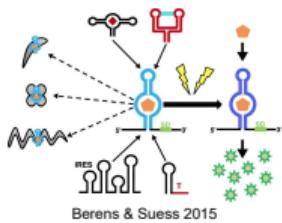
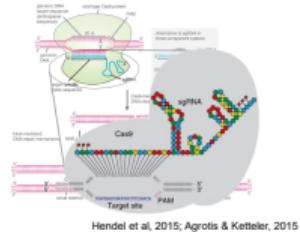
COVID-19, Malaria (Zika, CMV, Cancers?)

RiboNucleic Acids (RNAs) in Human biology/health: Friends and Foes!

Targeting system for DNA Editing

CRISPR therapies

Sickle-cell anemia, β -thalassamia, Leber congenital amaurosis (LCA), cancers...



Sensor of metabolites Riboswitches

RiboNucleic Acids (RNAs)

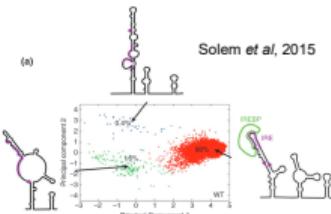
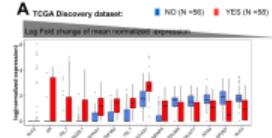


Encodes proteins
mRNA Vaccines
COVID-19, Malaria (Zika, CMV, Cancers?)

Quantitative expression

Transcriptomic signatures

Cancer diagnosis/prognosis/relapse...



Non-coding mutations

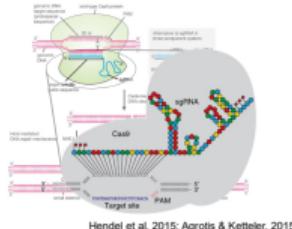
lncRNAs, miRNAs, structure-associated (RiboSnitches)
 β -thalassemia, duchenne muscular dystrophy,
Cystic fibrosis, Rett syndrome...

RiboNucleic Acids (RNAs) in Human biology/health: Friends **and** Foes!

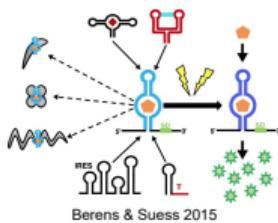
Targeting system for DNA Editing

CRISPR therapies

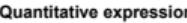
Sickle-cell anemia, β -thalassamia, Leber congenital amaurosis (LCA), cancers...



Hendel et al. 2015; Agrotis & Ketteler, 2015

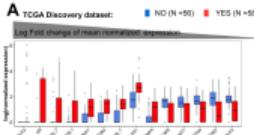


Sensor of metabolites
Riboswitches



Transcriptomic signatures

Cancer diagnosis/prognosis/relapse...



[Nguyen et al., 2021]

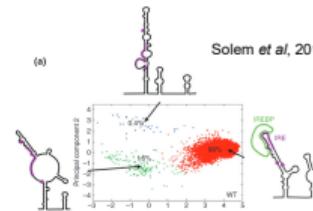
RiboNucleic Acids (RNAs)



Encodes proteins

mRNA Vaccines

COVID-19, Malaria (Zika, CMV, Cancers?)



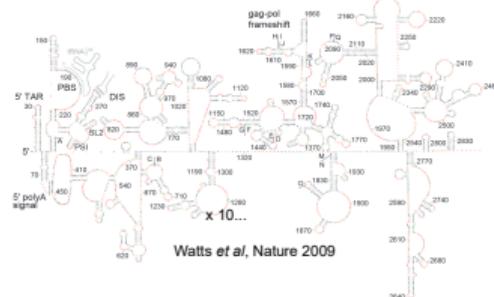
Solem et al. 2015

Non-coding mutations

lncRNAs, miRNAs, structure-associated (RiboSnitches)
β-thalassemia, duchenne muscular dystrophy,
Cystic fibrosis, Rett syndrome...

Genomic material for Human pathogens

HIV-1, SARS-CoV 2, HCoVs, MERS



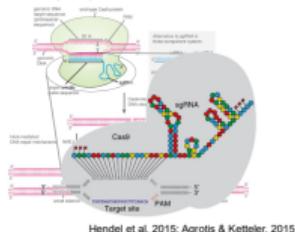
Watts et al. Nature 2009

RiboNucleic Acids (RNAs) in Human biology/health: Friends **and** Foes!

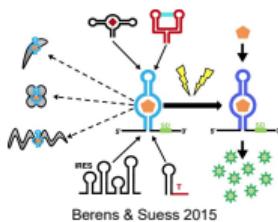
Targeting system for DNA Editing

CRISPR therapies

Sickle-cell anemia, β -thalassamia, Leber congenital amaurosis (LCA), cancers...

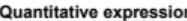


Hendel et al. 2015; Agrotis & Ketteler, 2015



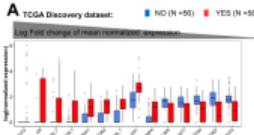
Sensor of metabolites

Riboswitches



Transcriptomic signatures

Cancer diagnosis/prognosis/relapse...



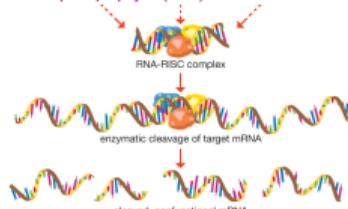
[Nguyen et al., 2021]

Regulation of gene expression

RNAi therapies (FDA approved)

Primary hyperoxaluria type 1 (PH1).

Hereditary transthyretin amyloidosis (ATTRv), Acute hepatic porphyria (AHP)



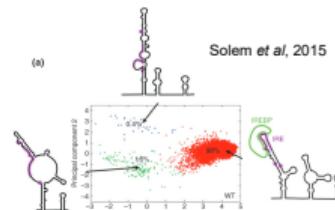
Encyclopaedia Britannica, Inc. 2013



Enriched proteins

Encodes proteins
mRNA messengers

mRNA Vaccines

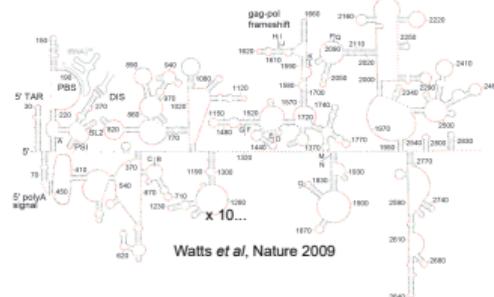


Non-coding mutations

lncRNAs, miRNAs, structure-associated (RiboSnitches)
β-thalassemia, duchenne muscular dystrophy,
Cystic fibrosis, Rett syndrome...

Genomic material for Human pathogens

HIV-1, SARS-CoV 2, HCoVs, MERS



1870

RiboNucleic Acids (RNAs) in Human biology/health: Friends **and** Foes!

Targeting system for DNA Editing

CRISPR therapies

Sickle-cell anemia, β-thalassamia, Leber congenital
amaurosis (LCA)



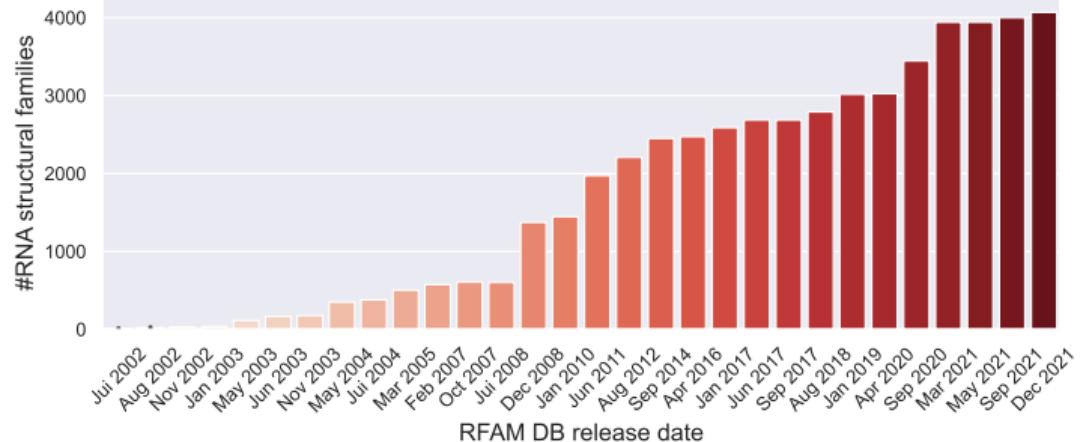
Quantitative expression

Transcriptomic signatures

Cancer diagnosis/prognosis/relapse...

(a)

Solem et al., 2015



RNA functional diversity is (largely) enabled by deep structural diversity

cleaved, nonfunctional mRNA

Encyclopaedia Britannica, Inc. 2013

mRNA Vaccines

COVID-19, Malaria (Zika, CMV, Cancers?)

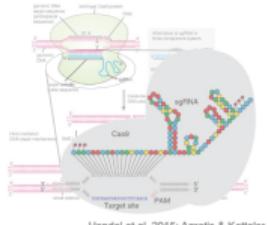
Watts et al., Nature 2009

RiboNucleic Acids (RNAs) in Human biology/health: Friends **and** Foes!

Targeting system for DNA Editing

CRISPR therapies

Sickle-cell anemia, β -thalassamia, Leber congenital amaurosis (LCA), cancers...



Hendel et al. 2015; Acosta & Ketteler 2015

Rational design

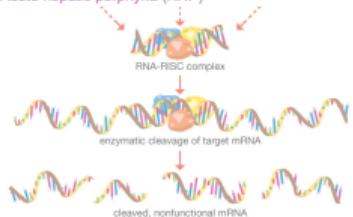
Regulation of gene expression

RNAi therapies (FDA approved)

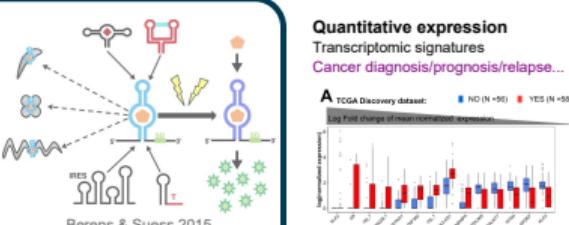
Primary hyperoxaluria type 1 (PH1):

Heredity, inheritance, and penetrance

Acute hepatic porphyria (AHP)

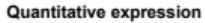


Encyclopaedia Britannica, Inc. 2013



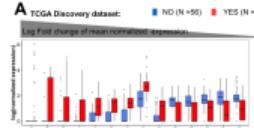
Sensor of metabolites
Riboswitches

RiboNucleic Acids (RNAs)



Transcriptomic signatures

Cancer diagnosis/prognosis/relapse...



[NGuyen et al., 2021]

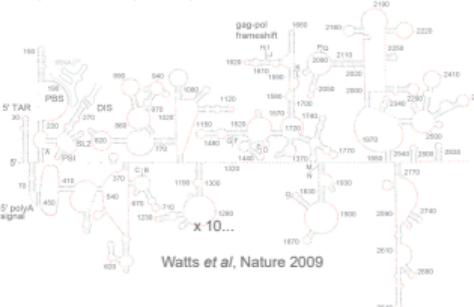
Non-coding mutations

NON CODING RNAs
lncRNAs, miRNAs, structure-associated (RiboSnitches)
 β -thalassemia, duchenne muscular dystrophy,
Cystic fibrosis, Rett syndrome...

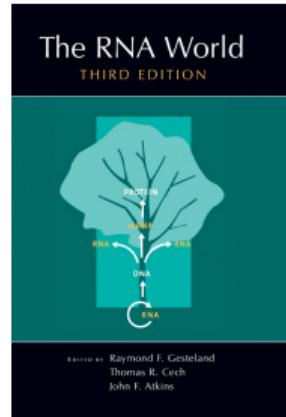
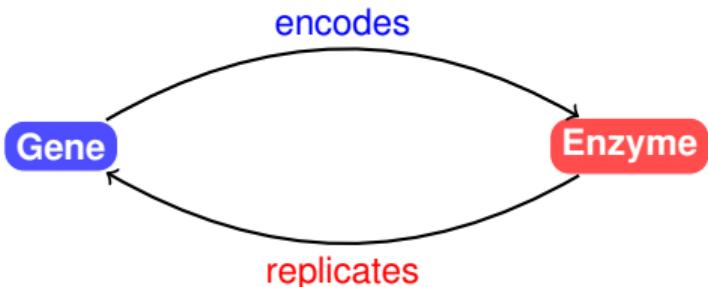
(2D) Structure Modeling

Genomic material for Human pathogens

HIV-1, SARS-CoV 2, HCoVs, MERS



RNA world: Resolving the *chicken vs egg* paradox at the origin of life...

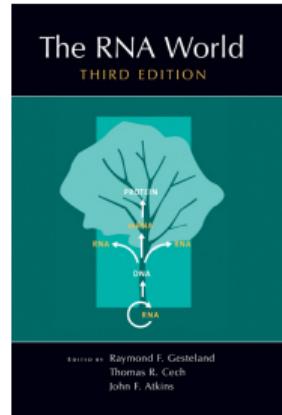
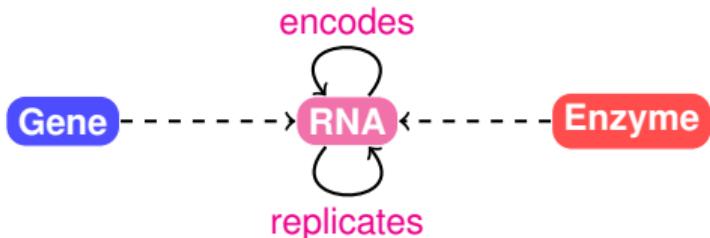


A gene big enough to specify an enzyme would be too big to replicate accurately without the aid of an enzyme of the very kind that it is trying to specify. So the system *apparently cannot get started*.

[...] This is the *RNA World*. To see how plausible it is, we need to look at why proteins are good at being enzymes but bad at being replicators; at why DNA is good at replicating but bad at being an enzyme; and finally why *RNA might just be good enough at both roles to break out of the Catch-22*.

R. Dawkins. *The Ancestor's Tale: A Pilgrimage to the Dawn of Evolution*

RNA world: Resolving the *chicken vs egg* paradox at the origin of life...



A gene big enough to specify an enzyme would be too big to replicate accurately without the aid of an enzyme of the very kind that it is trying to specify. So the system *apparently cannot get started*.

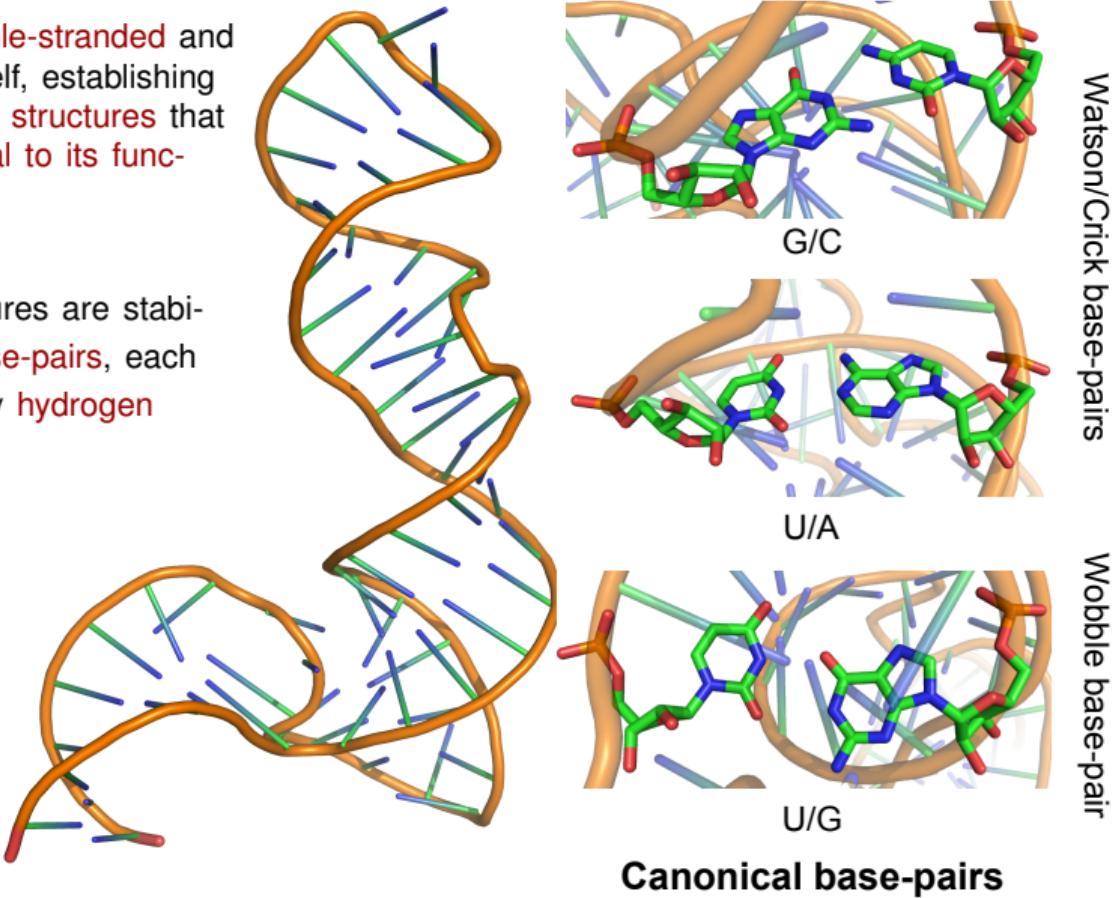
[...] This is the **RNA World**. To see how plausible it is, we need to look at why proteins are good at being enzymes but bad at being replicators; at why DNA is good at replicating but bad at being an enzyme; and finally why *RNA might just be good enough at both roles to break out of the Catch-22*.

R. Dawkins. *The Ancestor's Tale: A Pilgrimage to the Dawn of Evolution*

RNA folding

RNA is **single-stranded** and folds on itself, establishing complex 3D structures that are **essential** to its function(s).

RNA structures are stabilized by **base-pairs**, each mediated by hydrogen bonds.

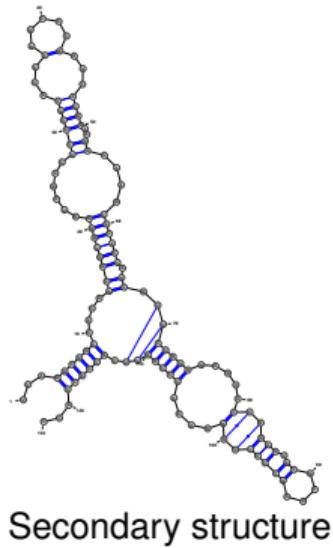


RNA Structure(s)

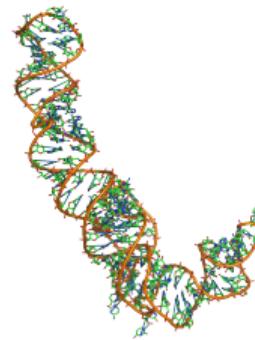
Three¹ levels of representation:

UUAGGCAGGCCACAGC
GGUGGGGUUGCCUCC
CGUACCCAUCCCGAA
CACGGAAGAUAGCC
CACCAAGCGUUCGGG
GAGUACUGGAGUGCG
CGAGCCUCUGGGAAA
CCGGUUCGCCGCCA
CC

Primary structure



Secondary structure



Tertiary structure

Source: 5s rRNA (PDB 1K73:B)

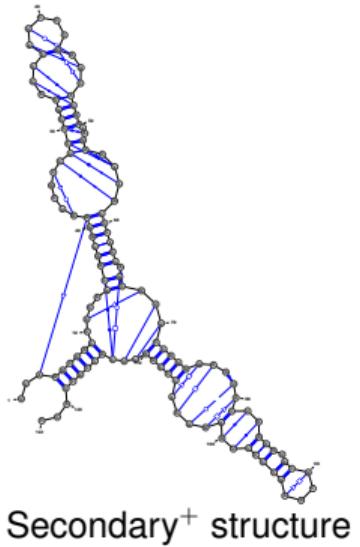
¹Well, mostly...

RNA Structure(s)

Three¹ levels of representation:

UUAGGCAGGCCACAGC
GGUGGGGUUGCCUCC
CGUACCCAUCCCGAA
CACGGAAGAUAGCC
CACCAAGCGUUCGGG
GAGUACUGGAGUGCG
CGAGCCUCUGGGAAA
CCGGUUCGCCGCCA
CC

Primary structure



Secondary⁺ structure



Tertiary structure

Source: 5s rRNA (PDB 1K73:B)

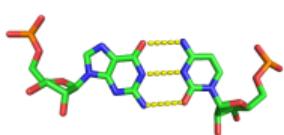
¹Well, mostly...

Ignored by secondary structure

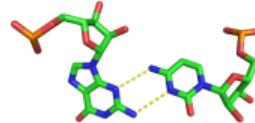
► Non-canonical base-pairs

Any base-pair **other than** {(A-U), (C-G), (G-U)}

Or interacting on non-standard edge (\neq WC/WC-Cis) [LW01].

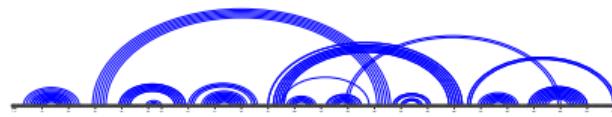


Canonique CG pair(WC/WC-Cis)



Non-canonical CG pair (Sugar/WC-Trans)

► Pseudoknots (PKs)



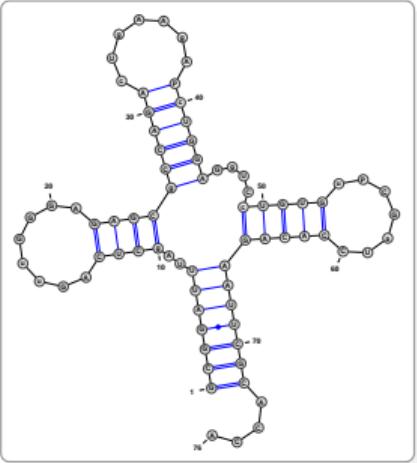
Pseudoknotted structure of group I ribozyme (PDBID: 1Y0Q:A)

Considering PKs may lead to better predictions, **but**:

- Some PK conformations are simply unfeasible;
- Folding *in silico* with general pseudoknots is NP-complete [LP00];

Still, folding on restricted classes of conformations seems promising [CDR⁺04].

Various representations for a versatile biomolecule



Outer-planar graphs

Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*

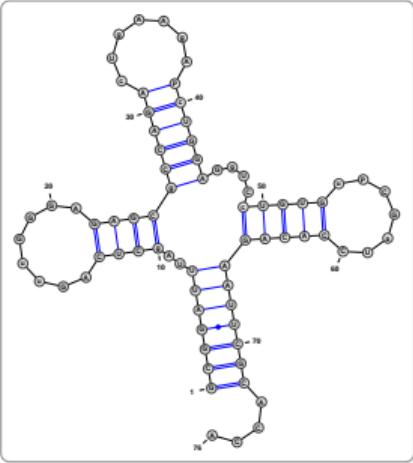
Supporting intuitions

Different representations

Common combinatorial structure

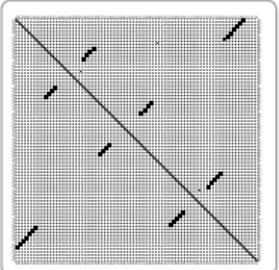
* Additional steric constraints

Various representations for a versatile biomolecule



Outer-planar graphs

Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*



Dot plots

Adjacency matrices*

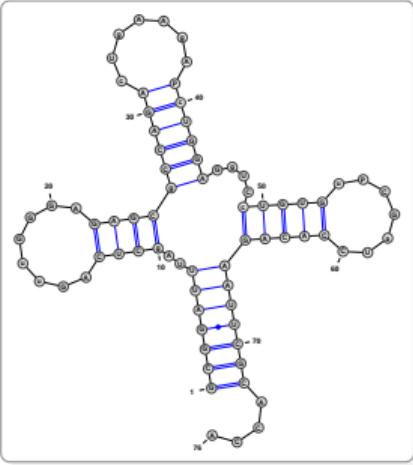
Supporting intuitions

Different representations

Common combinatorial structure

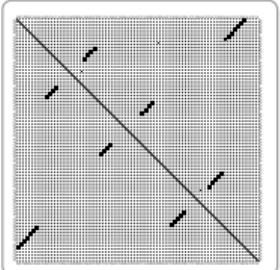
* Additional steric constraints

Various representations for a versatile biomolecule

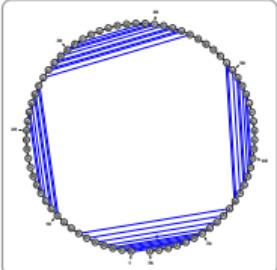


Outer-planar graphs

Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*



Dot plots
Adjacency matrices*



Non-crossing arc diagrams*

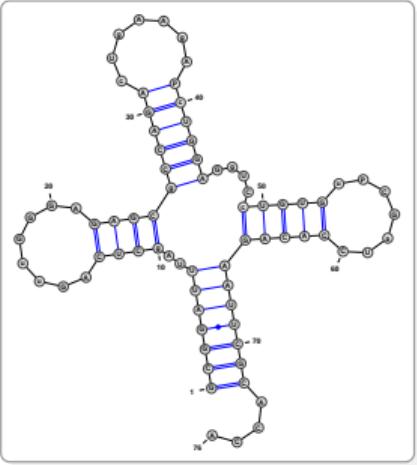
Supporting intuitions

Different representations

Common combinatorial structure

* Additional steric constraints

Various representations for a versatile biomolecule

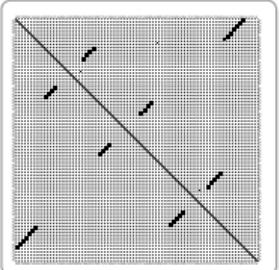


((((((...(((.....)))) (((((.....)))))).... (((((.....))))))))....

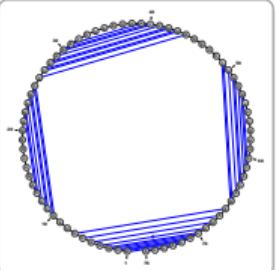
Motzkin words*

Outer-planar graphs

Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*



Dot plots
Adjacency matrices*



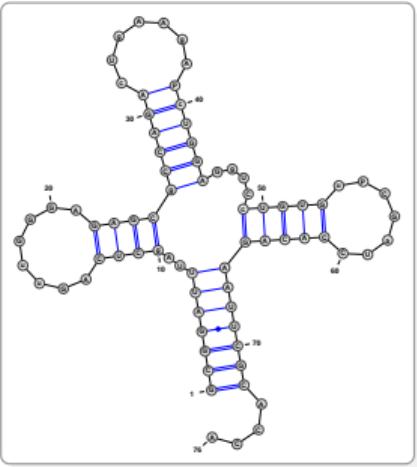
Non-crossing arc diagrams*

Supporting intuitions

Different representations
Common combinatorial structure

* Additional steric constraints

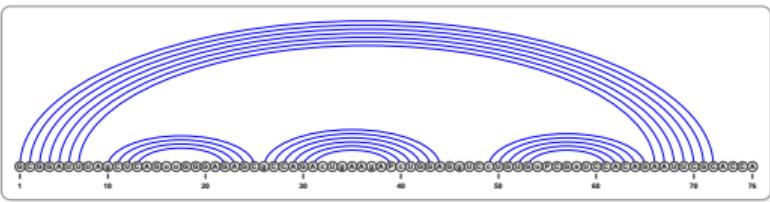
Various representations for a versatile biomolecule



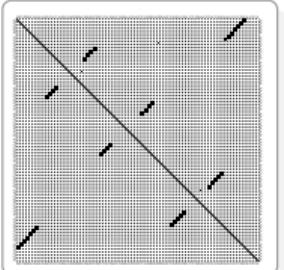
Outer-planar graphs
Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*

((((((...(((((.....)))) (((((.....))))))).... (((((.....))))))))....

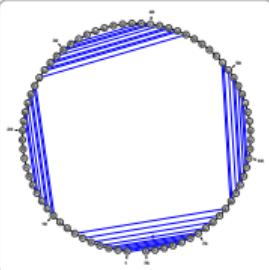
Motzkin words*



Non-crossing arc-annotated sequences*



Dot plots
Adjacency matrices*



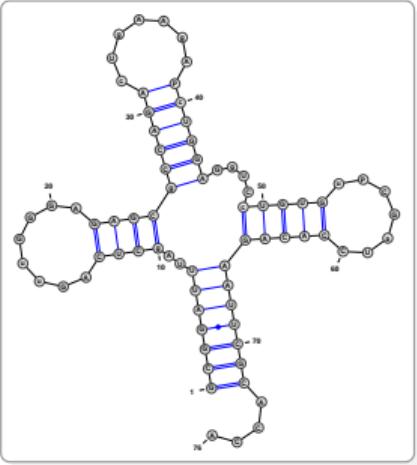
Non-crossing arc diagrams*

Supporting intuitions

Different representations
Common combinatorial structure

* Additional steric constraints

Various representations for a versatile biomolecule

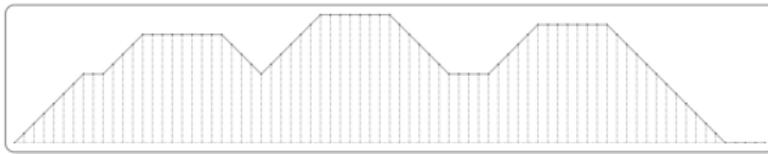


Outer-planar graphs

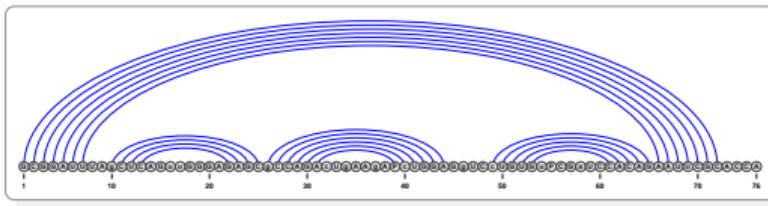
Hamiltonian-path, $\Delta(G) \leq 3$, 2-connected*

.....(.....,.....,.....,.....,.....,.....,.....,.....,.....,.....,.....,.....,.....,.....,.....,.....,.....)

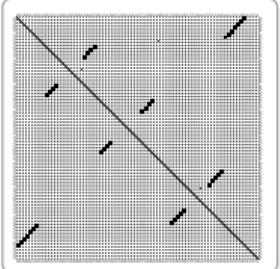
Motzkin words*



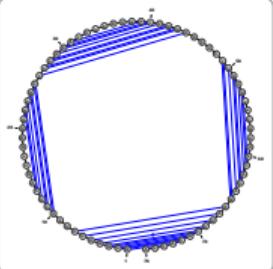
Positive 1D meanders* over $S = \{+1, -1, 0\}$



Non-crossing arc-annotated sequences*



Dot plots
Adjacency matrices*



Non-crossing arc diagrams*

Supporting intuitions

Different representations

Common combinatorial structure

* Additional steric constraints

Outline

Introduction

Dynamic programming 101

Why RNA?

RNA folding

RNA Structure(s)

Some representations of RNA structure

Some flavours of folding prediction

Thermodynamics vs Kinetics

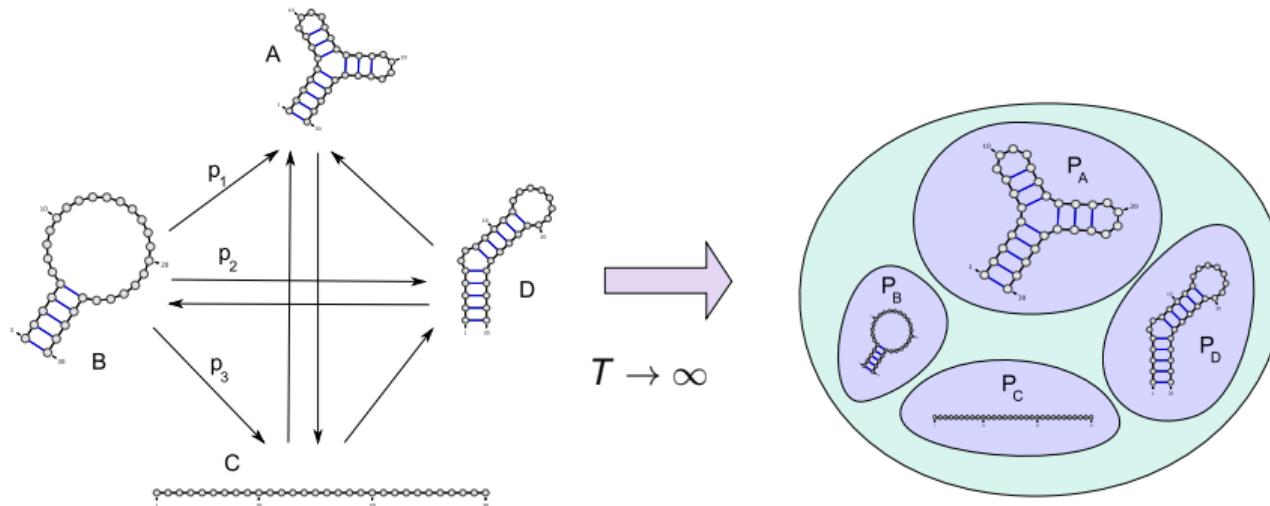
Dynamic programming: Reminder

Free-energy minimization

Nussinov-style RNA folding

Thermodynamics aparté

At the nanoscopic scale, RNA structure *fluctuates* (\approx Markov process).



Convergence towards a **stationary distribution** at the **Boltzmann equilibrium**, where the probability of a conformation only depends on its **free-energy**.

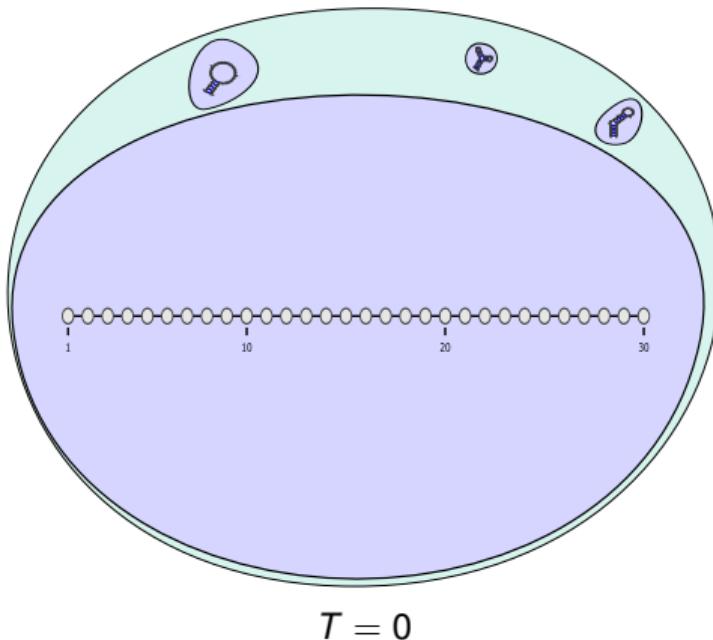
Corollary: Initial conformation does not matter.

Questions: For a given **conformation space** and **free-energy** model:

- Determine most stable (Minimum Free-Energy) structure at equilibrium;
- Compute average properties of Boltzmann ensemble;

Away from equilibrium

Transcription: RNA synthesized, supposedly without structure²

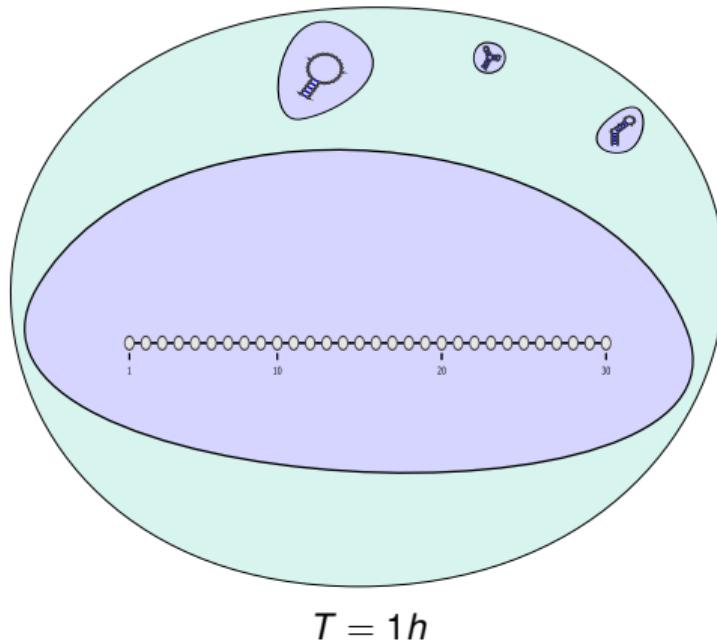


But most mRNAs are degrade before 7h (Org.: Souris [SSN⁺⁰⁹]).

²Except for co-transcriptional folding...

Away from equilibrium

Transcription: RNA synthesized, supposedly without structure²

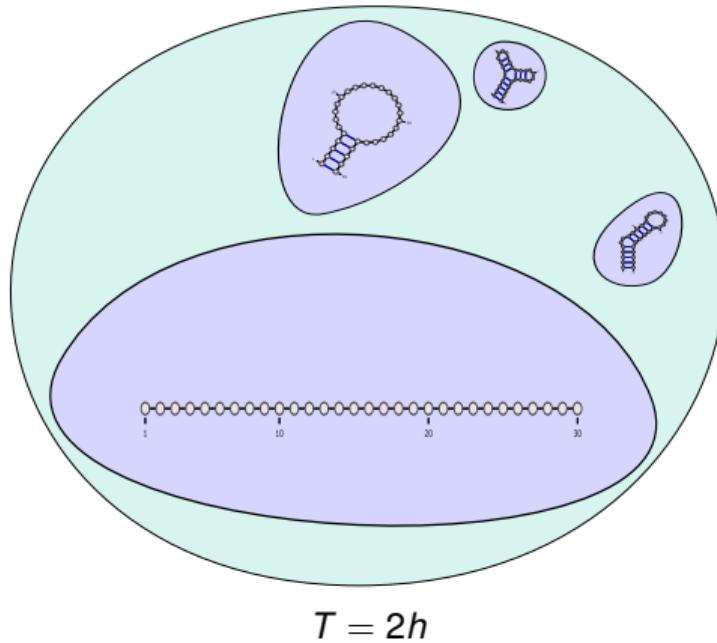


But most mRNAs are degrade before 7h (Org.: Souris [SSN⁺⁰⁹]).

²Except for co-transcriptional folding...

Away from equilibrium

Transcription: RNA synthesized, supposedly without structure²

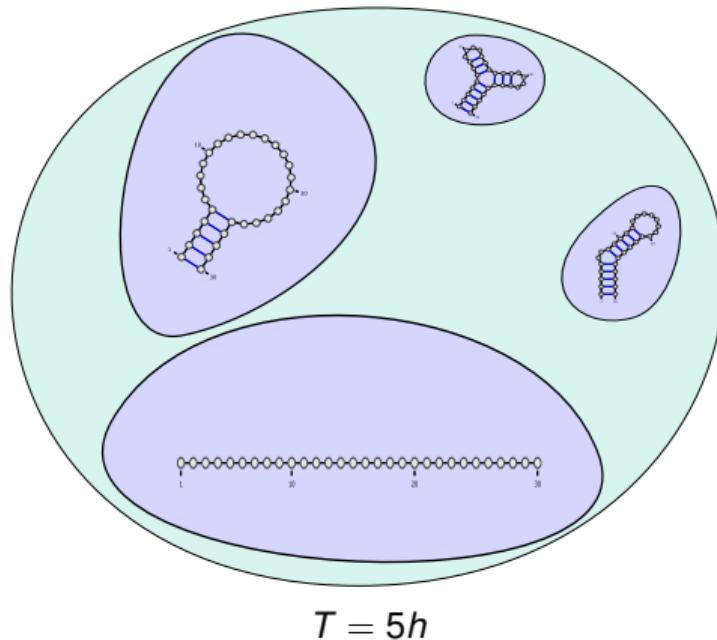


But most mRNAs are degrade before 7h (Org.: Souris [SSN⁺⁰⁹]).

²Except for co-transcriptional folding...

Away from equilibrium

Transcription: RNA synthesized, supposedly without structure²

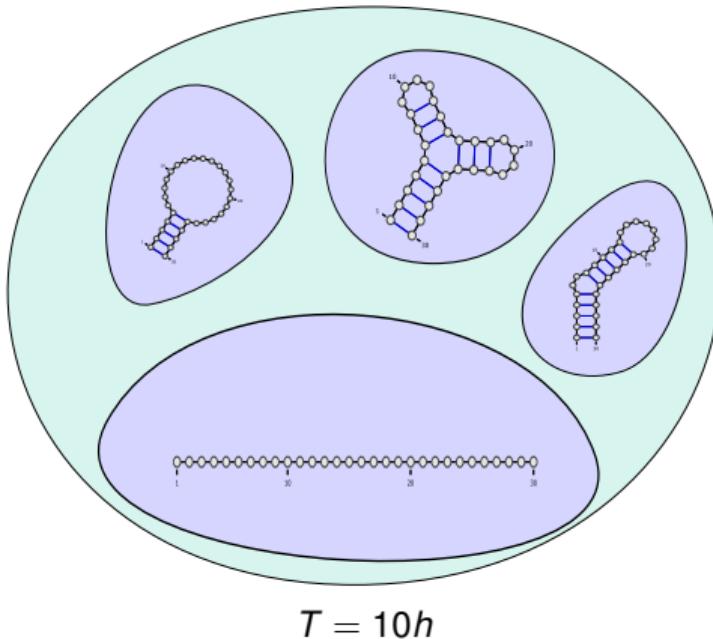


But most mRNAs are degrade before 7h (Org.: Souris [SSN⁺⁰⁹]).

²Except for co-transcriptional folding...

Away from equilibrium

Transcription: RNA synthesized, supposedly without structure²

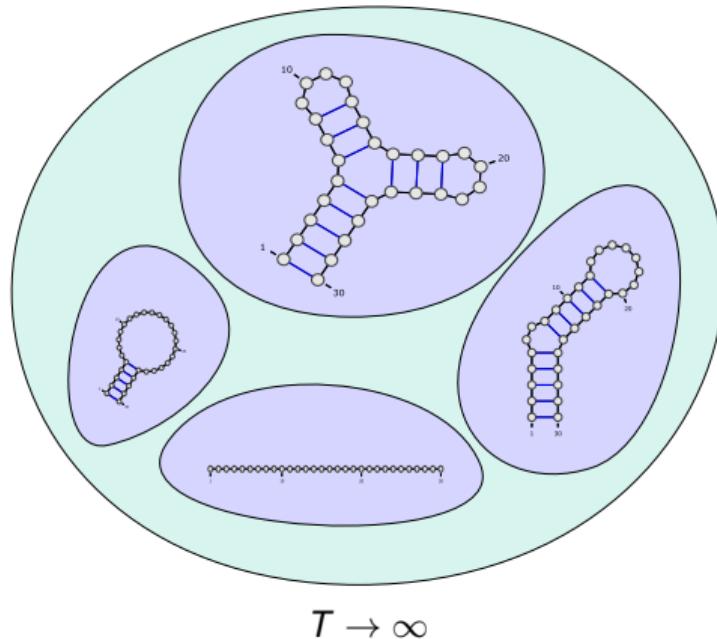


But most mRNAs are degrade before 7h (Org.: Souris [SSN⁺⁰⁹]).

²Except for co-transcriptional folding...

Away from equilibrium

Transcription: RNA synthesized, supposedly without structure²

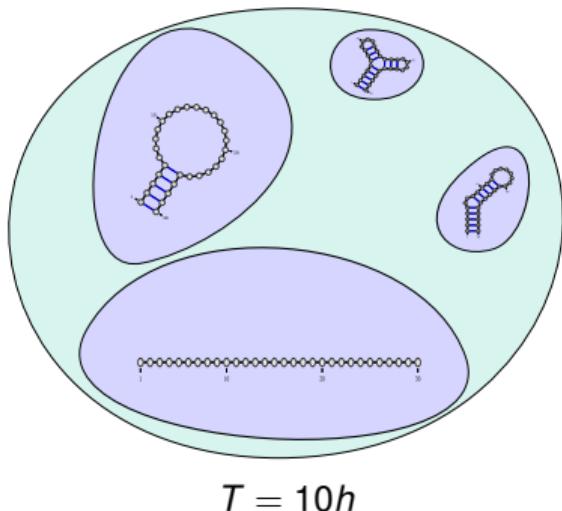


But most mRNAs are degrade before 7h (Org.: Souris [SSN⁺09]).

²Except for co-transcriptional folding...

Away from equilibrium

Transcription: RNA synthesized, supposedly without structure²



But most mRNAs are degrade before 7h (Org.: Souris [SSN⁺⁰⁹]).

- A. Determine most stable (Minimum Free-Energy) structure at equilibrium;
- B. Compute average properties of Boltzmann ensemble;
- C. Determine most likely structure at finite time T .

(c.f. H. Isambert through simulation, NP-complete deterministically [MTSC09])

²Except for co-transcriptional folding...

Dynamic programming: General principle

Dynamic programming = General optimization technique.

Prerequisite: Optimal solution for problem P can be derived from solutions to strict sub-problems.

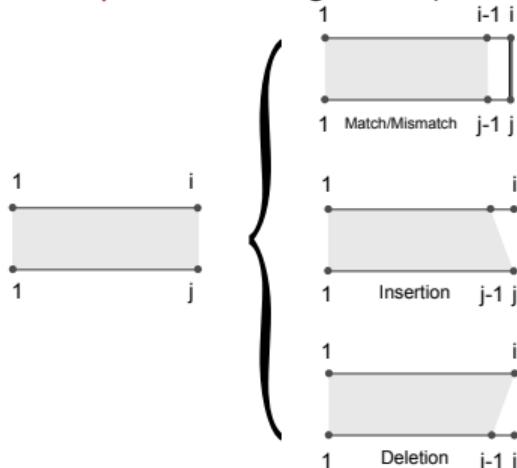
Bioinformatics :

Discrete solution space (alignments, structures...)

+ Additively-inherited objective function (cost, log-odd score, energy...)

⇒ Efficient dynamic programming scheme

Example: Local Alignment(Smith/Waterman)



$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Algorithmic details

Dynamic programming scheme defines a space of (sub)problems and a recurrence that relates the score of a problem to that of smaller problems.

Given a scheme, two steps :

- ▶ Matrix filling: Computation and tabulation of best scores (Computed from smaller problems to larger ones).
- ▶ Traceback: Reconstruct best solution from contributing subproblems.

Complexity of algorithm depends on:

- ▶ Cardinality of sub-problem space
- ▶ Number of alternatives considers at each step (#Terms in recurrence)

Smith&Waterman example:

- ▶ $i: 1 \rightarrow n + 1 \Rightarrow \Theta(n)$
- ▶ $j: 1 \rightarrow m + 1 \Rightarrow \Theta(m)$
- ▶ 3 operations at each step
 $\Rightarrow \Theta(m.n)$ time/memory

$$W(i, 0) = 0$$
$$W(0, j) = 0$$
$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_d = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

	A	C	A	C	A	C	T	A
A	0	0	0	0	0	0	0	0
G	0							
C	0							
A	0							
C	0							
A	0							
C	0							
A	0							

Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_d = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

	A	C	A	C	A	C	T	A
	0	0	0	0	0	0	0	0
A	0 → 2							
G	0							
C	0							
A	0							
C	0							
A	0							
C	0							
A	0							

Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_d = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

	A	C	A	C	A	C	T	A
	0	0	0	0	0	0	0	0
A	0	2 → 1						
G	0							
C	0							
A	0							
C	0							
A	0							
C	0							
A	0							

Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_d = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

	A	C	A	C	A	C	T	A
	0	0	0	0	0	0	0	0
A	0	2 → 1 → 2						
G	0							
C	0							
A	0							
C	0							
A	0							
C	0							
A	0							

Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

	A	C	A	C	A	C	T	A
	0	0	0	0	0	0	0	0
A	0	2 → 1	2 → 1					
G	0							
C	0							
A	0							
C	0							
A	0							
C	0							
A	0							

Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_d = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

	A	C	A	C	A	C	T	A
A	0	0	0	0	0	0	0	0
G	0							
C	0							
A	0							
C	0							
A	0							
C	0							
A	0							

Complete example

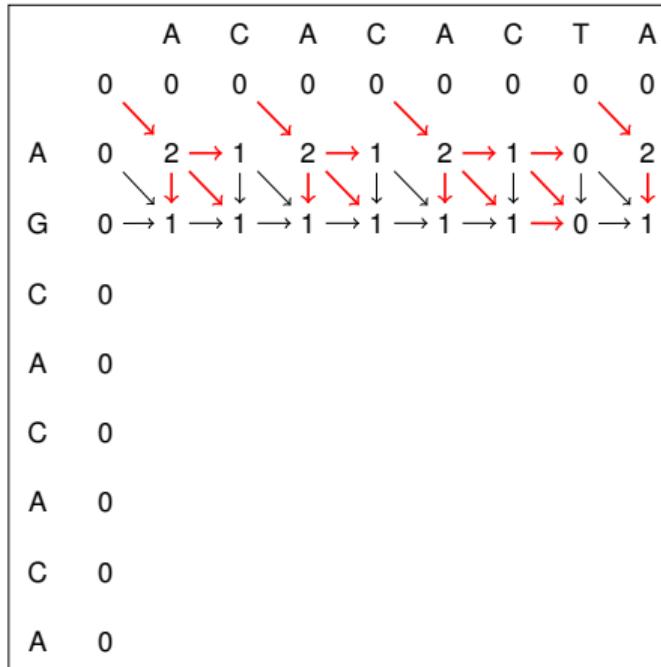
Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$



Complete example

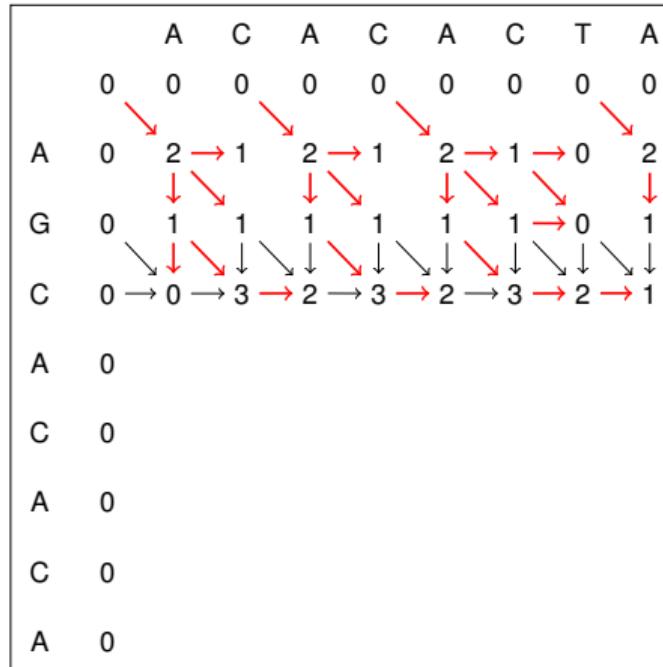
Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$



Complete example

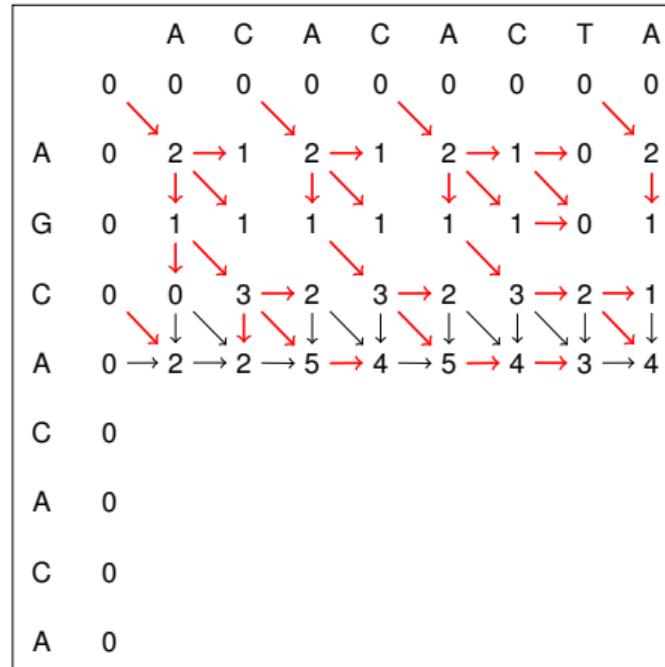
Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$



Complete example

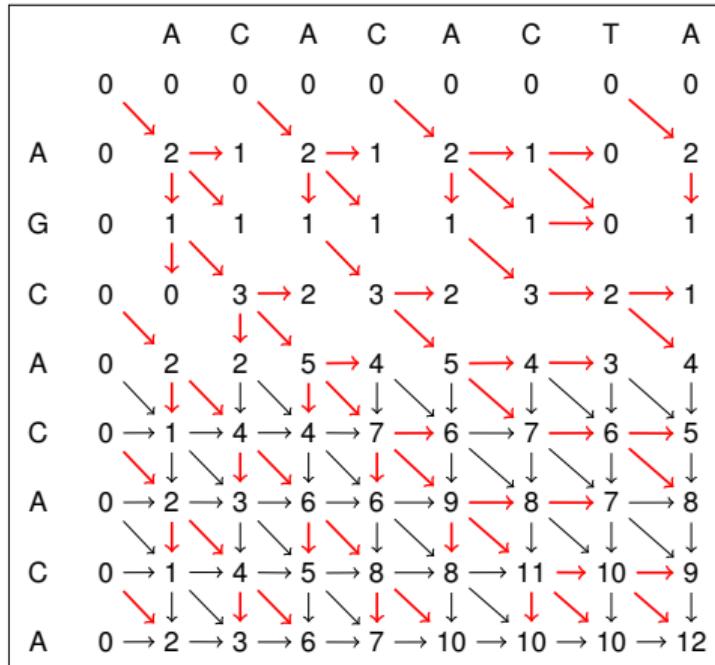
Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$



Complete example

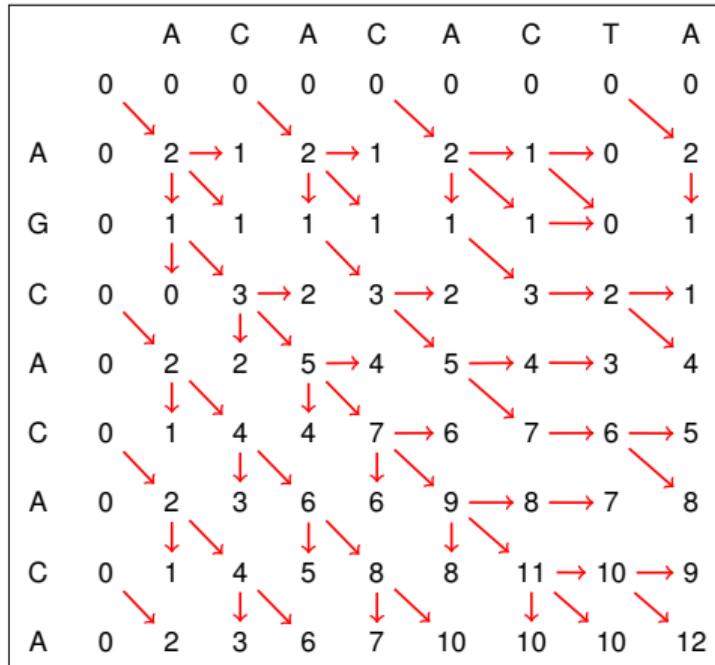
Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$



Complete example

Example: Local alignment of AGCACACCA and ACACACTA

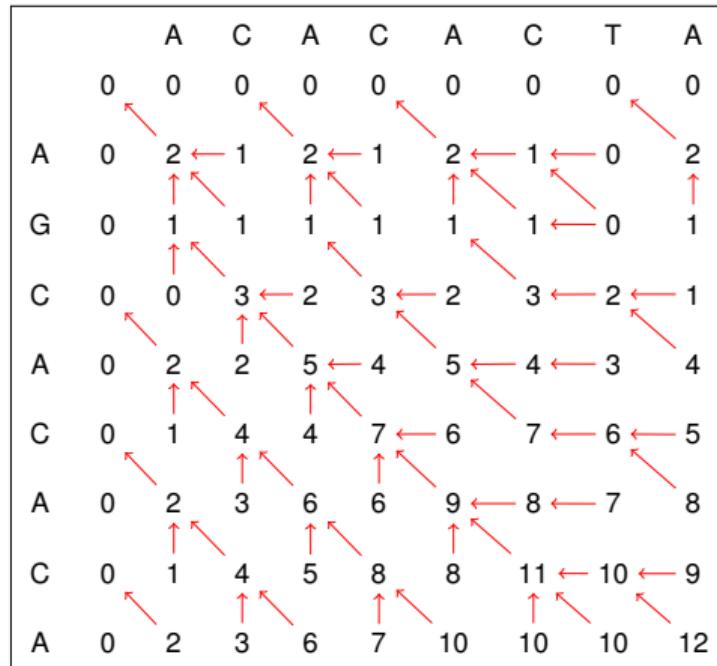
Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Best alignment



Complete example

Example: Local alignment of AGCACACCA and ACACACTA

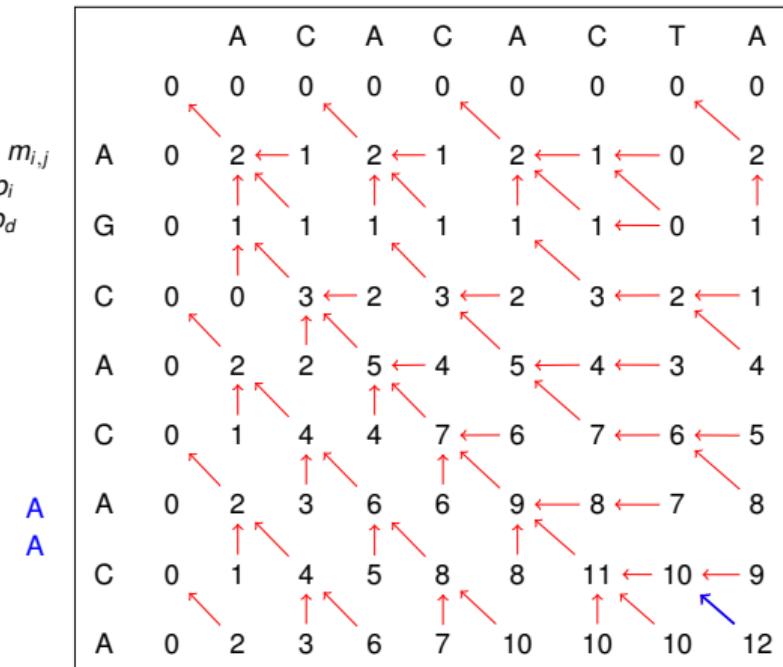
Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Best alignment



Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

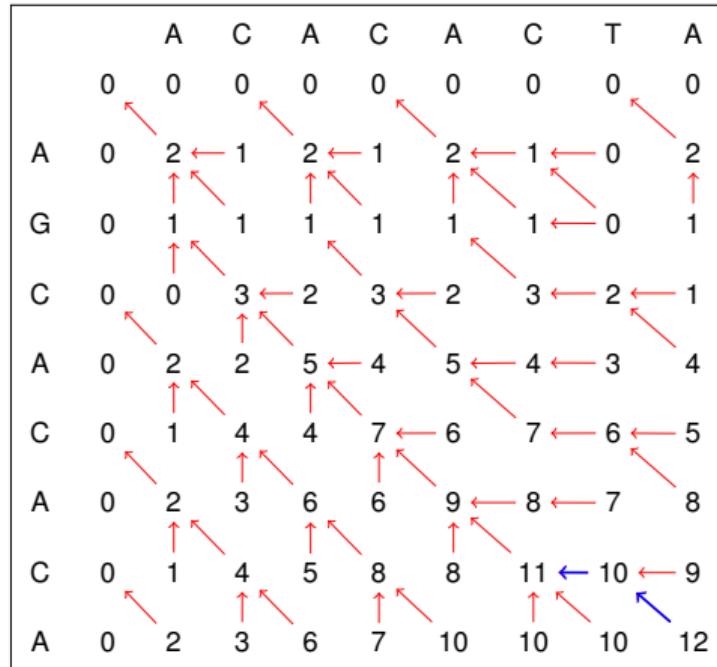
$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Best alignment

- T A



Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

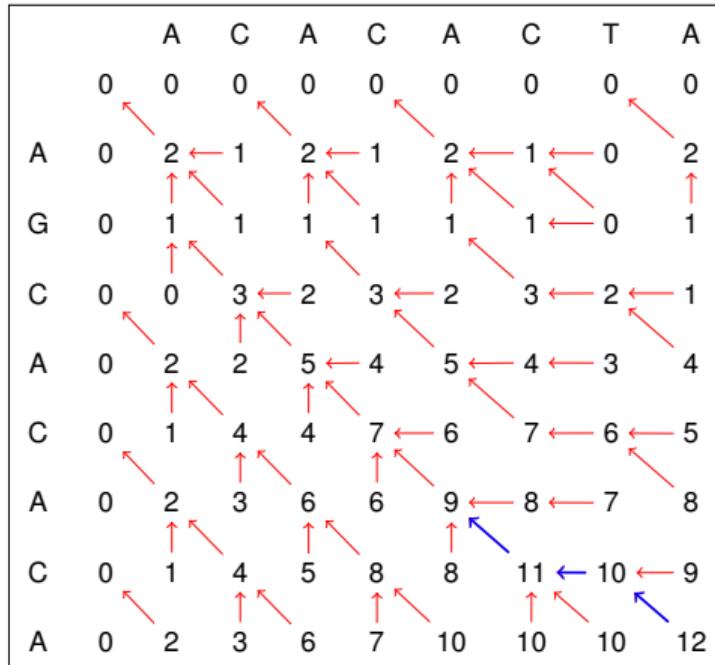
$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Best alignment

C - A
C T A



Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

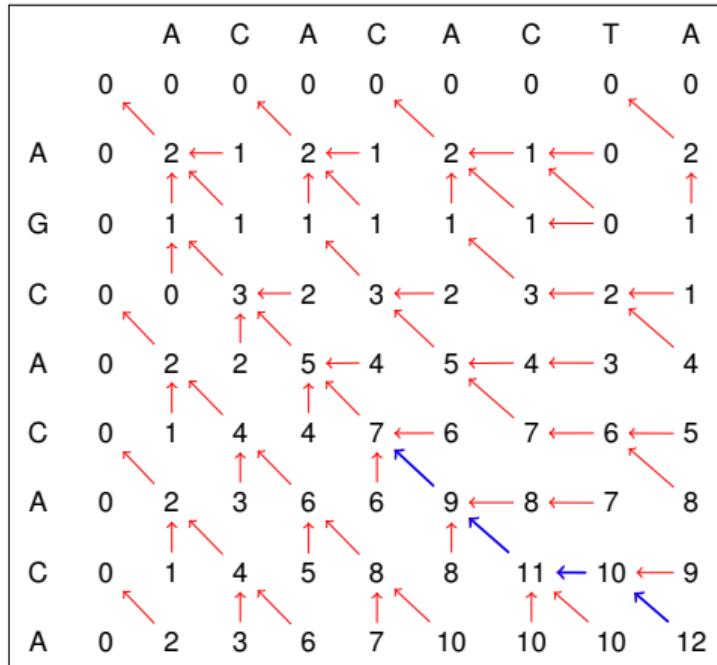
$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Best alignment

A C - A
A C T A



Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

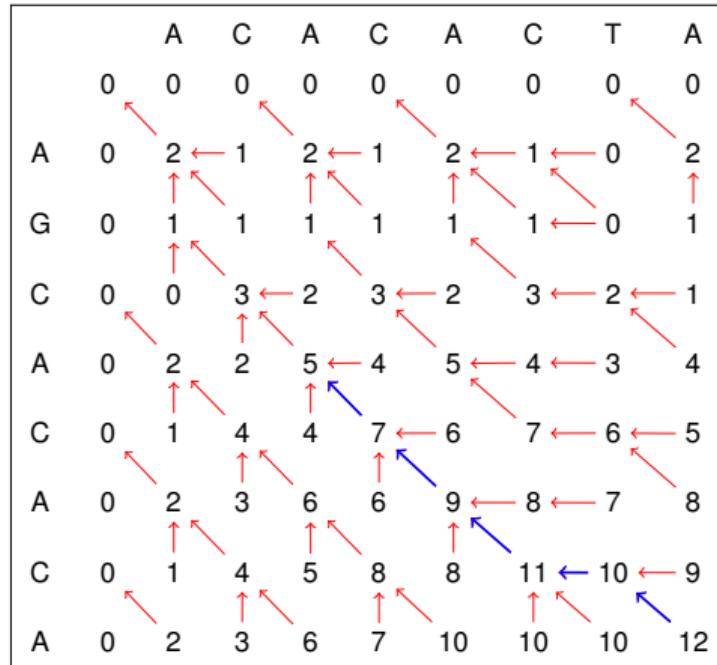
$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Best alignment

C A C - A
C A C T A



Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

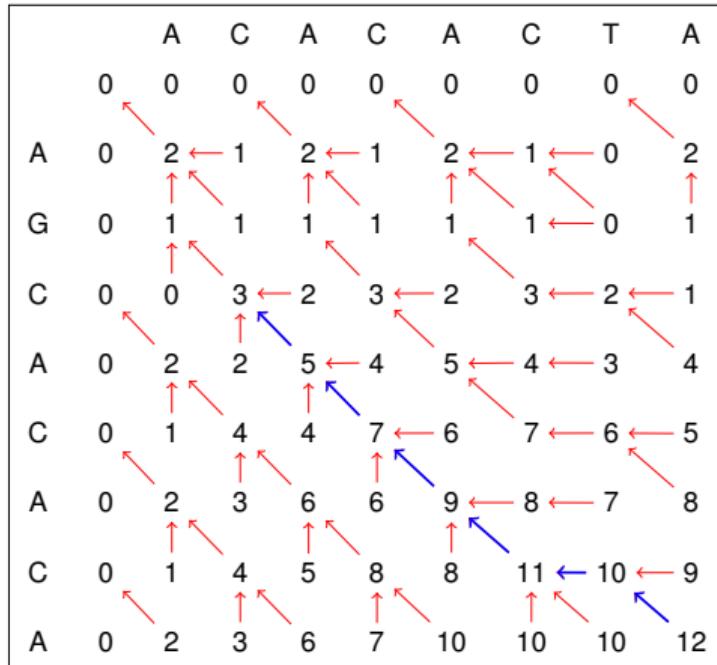
$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Best alignment

A	C	A	C	-	A
A	C	A	C	T	A



Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

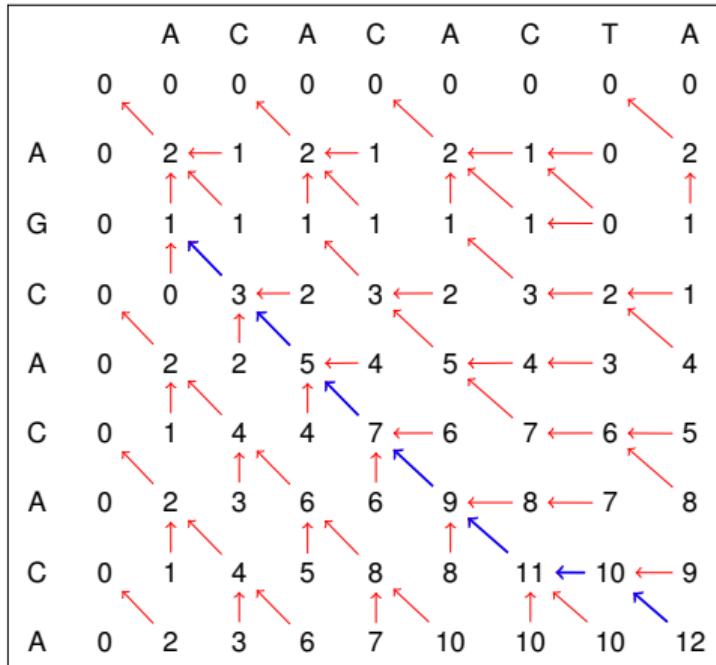
$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Best alignment

C	A	C	A	C	-	A
C	A	C	A	C	T	A



Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

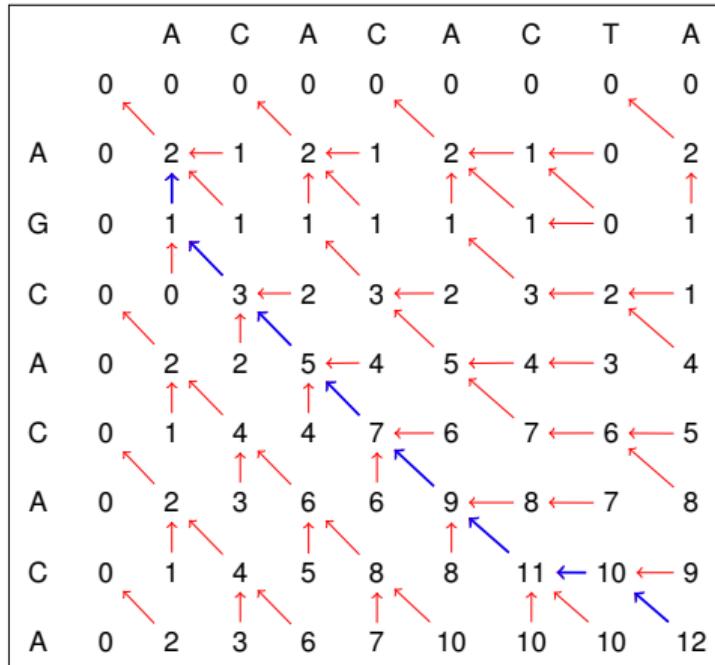
$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Best alignment

G	C	A	C	A	C	-	A
-	C	A	C	A	C	T	A



Complete example

Example: Local alignment of AGCACACCA and ACACACTA

Costs: Match $m_{i,j} = +2$, Insertion/Déletion $p_i = p_j = -1$

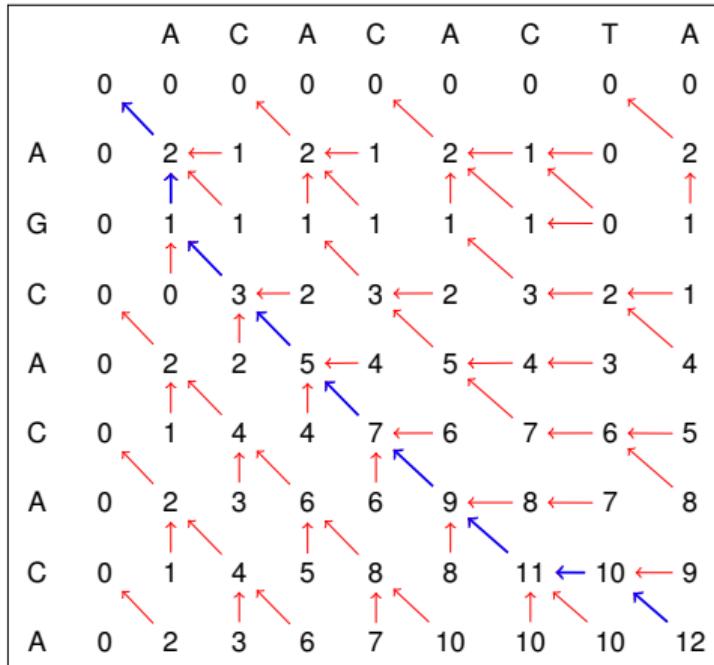
$$W(i, 0) = 0$$

$$W(0, j) = 0$$

$$W(i, j) = \max \begin{cases} W(i - 1, j - 1) + m_{i,j} \\ W(i - 1, j) + p_i \\ W(i, j - 1) + p_d \end{cases}$$

Best alignment

A	G	C	A	C	A	C	-	A
A	-	C	A	C	A	C	T	A



Properties of DP schemes

Necessary properties:

- ▶ **Correctness:** \forall sub-problem, the computed value must indeed maximize the objective function .

Proofs usually inductive, and quite technical, but very systematic.

Desirable properties of DP schemes:

- ▶ **Completeness** of space of solutions generated by decomposition.
Algorithmic tricks, by *cutting branches*, may violate this property.
- ▶ **Unambiguity:** Each solution is generated at most once.
⇒ Under these properties, one can enumerate solution space.

Outline

Introduction

Dynamic programming 101

Why RNA?

RNA folding

RNA Structure(s)

Some representations of RNA structure

Some flavours of folding prediction

Thermodynamics vs Kinetics

Dynamic programming: Reminder

Free-energy minimization

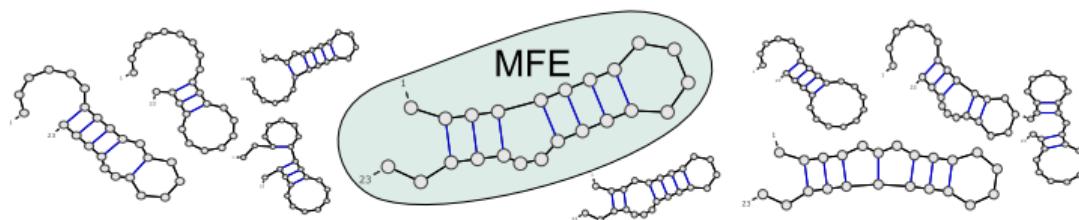
Nussinov-style RNA folding

Folding by minimizing free-energy

Problem A: Determine Minimum Free-Energy structure (MFE).

Ab initio folding prediction =

Predict RNA structure from its sequence ω only.



- ▶ **Conformations:** Set S_ω of secondary structures compatible (w.r.t. base-pairing constraints) with primary structure ω .
- ▶ **Free-Energy:** Function $E_{\omega, S}$ (KCal.mol⁻¹), additive on motifs occurring in any sequence/conformation couple (ω, S).
- ▶ **Native structure:** Functional conformation of the biomolecule.

Remarks:

- ▶ Not necessarily unique (Kinetics, or bi-stable structures);
- ▶ In presence of PKs → Ambiguous: Which is the native conformation?

Nussinov/Jacobson model

Nussinov/Jacobson energy model (NJ)

Base-pair maximization (with a twist):

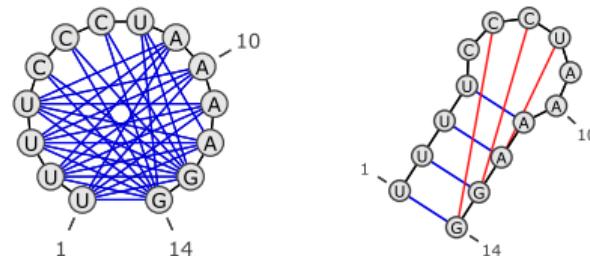
- ▶ Additive model on independently contributing base-pairs;
- ▶ Canonical base-pairs only: Watson/Crick (A/U,C/G) and Wobble (G/U)

$$\Rightarrow E_{\omega,S} = -\# \text{Paires}(S)$$

Folding in NJ model \Leftrightarrow Base-pair (weight) maximization

Example:

UUUUUCCCUAAAAGG



Variant: Weight each pair with $-\#$ Hydrogen bonds

$$\Delta G(G \equiv C) = -3$$

$$\Delta G(A = U) = -2$$

$$\Delta G(G - U) = -1$$

Nussinov/Jacobson model

Nussinov/Jacobson energy model (NJ)

Base-pair maximization (with a twist):

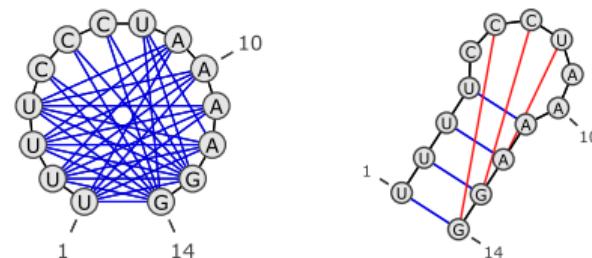
- ▶ Additive model on independently contributing base-pairs;
- ▶ Canonical base-pairs only: Watson/Crick (A/U,C/G) and Wobble (G/U)

$$\Rightarrow E_{\omega,S} = -\# \text{Paires}(S)$$

Folding in NJ model \Leftrightarrow Base-pair (weight) maximization

Example:

UUUUUCCCUAAAAGG



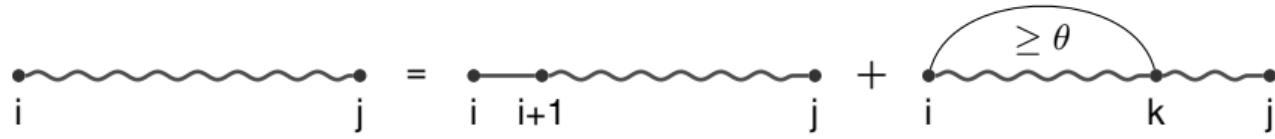
Variant: Weight each pair with $-\#$ Hydrogen bonds

$$\Delta G(G \equiv C) = -3$$

$$\Delta G(A = U) = -2$$

$$\Delta G(G - U) = -1$$

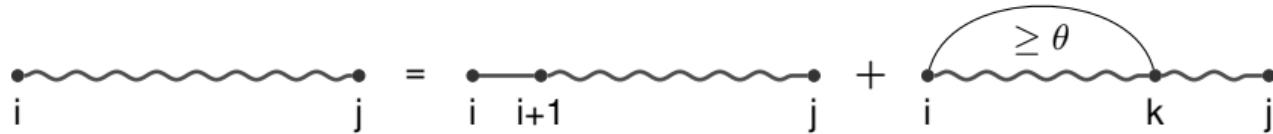
Nussinov/Jacobson DP scheme



$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \begin{cases} N_{i+1,j} & i \text{ unpaired} \\ \min_{k=i+\theta+1}^j \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{cases}$$

Nussinov/Jacobson DP scheme



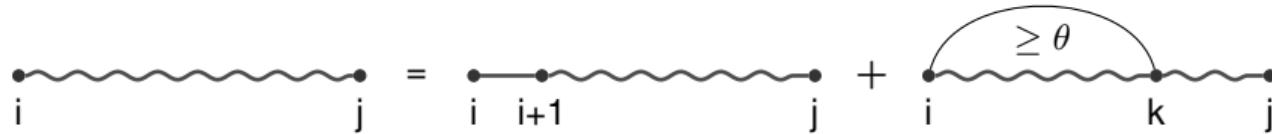
$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \begin{cases} N_{i+1,j} & i \text{ unpaired} \\ \min_{k=i+\theta+1}^j \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{cases}$$

Correctness. Goal = Show that MFE over interval $[i, j]$ is indeed found in $N_{i,j}$ after completing the computation. Proceed by induction:

- ▶ Assume that property holds for any $[i', j']$ such that $j' - i' < n$.
- ▶ Consider $[i, j], j - i = n$. Let $\text{MFE}_{i,j} := \text{Base-pairs of best struct. on } [i, j]$. Then first position i in $\text{MFE}_{i,j}$ is either:
 - ▶ Unpaired: $\text{MFE}_{i,j} = \text{MFE}_{i+1,j}$ → free-energy = $N_{i+1,j}$
 - ▶ Paired to k : $\text{MFE}_{i,j} = \{(i, k)\} \cup \text{MFE}_{i+1,k-1} \cup \text{MFE}_{k+1,j}$.
(Indeed, any BP between $[i+1, k-1]$ and $[k+1, j]$ would cross (i, k))→ free-energy = $\Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j}$

Nussinov/Jacobson DP scheme



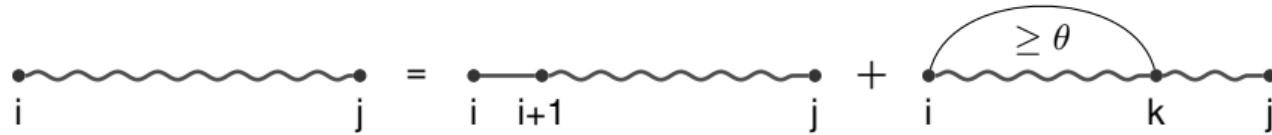
$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \begin{cases} N_{i+1,j} & i \text{ unpaired} \\ \min_{k=i+\theta+1}^j \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{cases}$$

Correctness. Goal = Show that MFE over interval $[i, j]$ is indeed found in $N_{i,j}$ after completing the computation. Proceed by induction:

- ▶ Assume that property holds for any $[i', j']$ such that $j' - i' < n$.
- ▶ Consider $[i, j], j - i = n$. Let $\text{MFE}_{i,j} :=$ Base-pairs of best struct. on $[i, j]$. Then first position i in $\text{MFE}_{i,j}$ is either:
 - ▶ **Unpaired:** $\text{MFE}_{i,j} = \text{MFE}_{i+1,j}$ $\rightarrow \text{free-energy} = N_{i+1,j}$
 - ▶ **Paired to k :** $\text{MFE}_{i,j} = \{(i, k)\} \cup \text{MFE}_{i+1,k-1} \cup \text{MFE}_{k+1,j}$.
(Indeed, any BP between $[i+1, k-1]$ and $[k+1, j]$ would cross (i, k)) $\rightarrow \text{free-energy} = \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j}$

Nussinov/Jacobson DP scheme



$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \min \begin{cases} N_{i+1,j} & i \text{ unpaired} \\ \min_{k=i+\theta+1}^j \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{cases}$$

Correctness. Goal = Show that MFE over interval $[i, j]$ is indeed found in $N_{i,j}$ after completing the computation. Proceed by induction:

- ▶ Assume that property holds for any $[i', j']$ such that $j' - i' < n$.
- ▶ Consider $[i, j], j - i = n$. Let $\text{MFE}_{i,j} :=$ Base-pairs of best struct. on $[i, j]$. Then first position i in $\text{MFE}_{i,j}$ is either:
 - ▶ **Unpaired:** $\text{MFE}_{i,j} = \text{MFE}_{i+1,j}$ $\rightarrow \text{free-energy} = N_{i+1,j}$
 - ▶ **Paired to k :** $\text{MFE}_{i,j} = \{(i, k)\} \cup \text{MFE}_{i+1,k-1} \cup \text{MFE}_{k+1,j}$.
(Indeed, any BP between $[i+1, k-1]$ and $[k+1, j]$ would cross (i, k)) $\rightarrow \text{free-energy} = \Delta G_{i,k} + N_{i+1,k-1} + N_{k+1,j}$

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	
A	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8		
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7		
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	



Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A

C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	2	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7	7	7	7
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7	7	7
C	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5	5	5	5
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	3	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0



Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
C
G	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	2	2	2	2	4	4	5	7	7	7	8	10		
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10		
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8		
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8		
U	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5			
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	i		j	=	i		j		+	i		k	j	$\geq \theta$	0	0	0	
G																0	0	0
A																0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
C
G	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	11	11
G		0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A			0	0	0	0	2	2	2	2	4	4	5	7	7	8	10	
U				0	0	0	0	0	0	2	2	4	5	7	7	8	10	
A					0	0	0	0	0	2	2	2	5	5	5	8	8	
C						0	0	0	0	0	0	2	5	5	5	8	8	
U							0	0	0	0	0	2	3	5	5	6	7	
U								0	0	0	0	2	3	5	5	5	7	
C									0	0	0	0	3	3	3	5	5	
U										0	0	0	0	2	2	2	3	
U											0	0	0	0	0	1	2	
A												0	0	0	0	0	0	
G													0	0	0	0	0	
A														0	0	0	0	
C															0	0	0	
G															0	0	0	
A																0		



Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	2	2	2	2	4	4	5	7	7	7	8	10		
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10		
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8		
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8		
U	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5			
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3			
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			



Nussinov/Jacobson

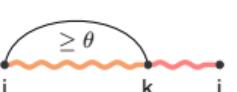
	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	2	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5			
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3			
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	2	2	2	2	4	4	5	7	7	8	10	
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10		
A	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8		
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7		
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0



Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	2	2	2	2	4	4	5	7	7	7	8	10		
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10		
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8		
C	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8		
U	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
G	i		j	=	i	i+1	j		+ 							0		
A																		

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	3	3	3	5	5	5	5	5	6	8	10	10	10
A	0	0	0	0	2	2	2	2	4	4	5	7	7	7	8	10		
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10		
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8		
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8		
U	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	3	3	3	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	2	2	2	2	4	4	5	7	7	7	8	10		
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10		
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8		
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8		
U	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5			
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
G	i					j	=	i	i+1			j	+	i		k	j	0
A																		

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	2	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7	7	7	7
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7	7	7	7
C	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	5	5	5
U	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	3	3	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	i				j	=	i		i+1		j	+	i		k	j	0	
A																		0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	2	2	2	2	2	4	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	8
U	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7	7	7	7
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7	7	7	7
C	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	5	5	5
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	3	3	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	i				j	=	i		i+1		j	+	i	≥ θ	k	j		0
A																		0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(.)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	2	2	2	2	4	4	5	7	7	7	8	10		
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10		
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8		
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8		
U	0	0	0	0	0	0	0	2	3	5	5	5	6	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5			
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3			
U	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2			
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			



Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	((.).)).	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	2	2	2	2	4	4	4	5	7	7	7	8	10	
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10		
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8		
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8		
U	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7			
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7			
C	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5			
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3			
U	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2			
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
G	i				j	=	i		i+1		j	+	i		k	j	$\geq \theta$	0
A																		0

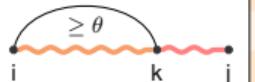
Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	((.))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	2	2	2	2	4	4	4	5	7	7	7	8	10	
U	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10		
A	0	0	0	0	0	0	0	2	2	2	5	5	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	
U	0	0	0	0	0	0	0	2	3	5	5	5	5	6	6	7		
U	0	0	0	0	0	0	0	2	3	5	5	5	5	5	5	7		
C	0	0	0	0	0	0	0	3	3	3	3	3	3	3	3	5	5	
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	i		j	=	i	i+1	j			i	k	j		$\geq \theta$				0
A																		

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	((.))	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	
A	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7		
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

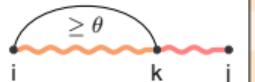
Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
C	(((.	.	.)))	.
G	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	
A	0	0	0	0	0	0	0	0	2	2	2	5	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	7			
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	i						j	=	i				j	+	i		j	0
C															0	0	0	0
G															0	0	0	0
A															0	0	0	0

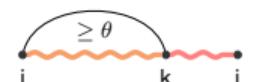
Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(((.	.	.)))	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A	0	0	0	0	0	0	0	2	2	2	2	2	4	4	5	7	7	8	10
U	0	0	0	0	0	0	0	0	0	2	2	2	4	5	7	7	8	10	
A	0	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	5	7			
C	0	0	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	0	0	0	2	2	2	2	3			
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2			
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
G	i					j	=	i		i+1		j	+	i	≥ θ	k	j		
A																		0	

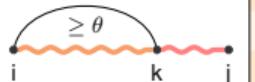
Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(((.	.	.)))	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A	0	0	0	0	0	0	0	2	2	2	2	2	4	4	5	7	7	8	10
U	0	0	0	0	0	0	0	0	0	2	2	2	4	5	7	7	8	10	
A	0	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	6	7	7	
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	5	7	7	7	
C	0	0	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5	5	
U	0	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	2	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	i						j	=	i				j	+	i		j	k	j
G																			
A																			0

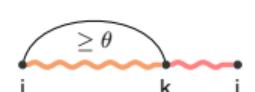
Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	2	2	4	4	5	7	7	10
U	0	0	0	0	0	0	0	0	0	2	2	4	5	7	7	8	10	
A	0	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	8	8
C	0	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7		
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	
U	0	0	0	0	0	0	0	0	0	0	0	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	i						j	=	i				j	+	i		j	0
A																		0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(((.	.	.)))	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A	0	0	0	0	0	0	0	2	2	2	2	2	4	4	5	7	7	8	10
U	0	0	0	0	0	0	0	0	0	2	2	2	4	5	7	7	8	10	
A	0	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	6	7	7	
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7	7	
C	0	0	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5	5	
U	0	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	i					j	=	i		i+1		j	+	i		k	j	0	

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	
A	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	6	7		
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	0	3	3	3	3	3	5	5	
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3		
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	i						j	=	i				j	+	i		j	0
C															0	0	0	
G															0	0	0	
A															0	0	0	

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(((.	.	.)))	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	
A	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	6	7		
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	3	5	5	
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	i					j	=	i		i+1		j	+	i	≥ θ	k	j	
A																		0

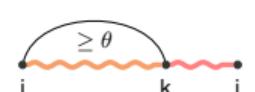
Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
C	(((.	.	.)	.	(.)))	.
G	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
A	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	2	2	4	4	5	7	7	10
U	0	0	0	0	0	0	0	0	0	2	2	2	4	5	7	7	8	10
A	0	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	8	8
C	0	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	6	7	7
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7	7
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	5
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	3	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	i				j	=	i		i+1		j	+	i	≥ θ	k	j	0	0
C																	0	0
G																	0	0
A																	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(((.	.	.)	.	(.)))	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A	0	0	0	0	0	0	0	2	2	2	2	2	4	4	5	7	7	8	10
U	0	0	0	0	0	0	0	0	0	2	2	2	4	5	7	7	8	10	
A	0	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	8	8	
C	0	0	0	0	0	0	0	0	0	0	2	5	5	5	5	8	8	8	
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7			
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	5	7			
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5		
U	0	0	0	0	0	0	0	0	0	0	0	2	2	2	2	3			
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2			
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
G	i					j	=	i		i+1		j	+	i	≥ θ	k	j	0	
A																		0	

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
C	(((.	.	.)	.	.	(.	.	.	.)))	.
G	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	0	2	2	2	2	2	4	4	5	7	7	10
U	0	0	0	0	0	0	0	0	2	2	2	4	5	7	7	8	10	
A	0	0	0	0	0	0	0	0	2	2	2	2	5	5	5	5	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	6	7		
U	0	0	0	0	0	0	0	0	0	2	3	5	5	5	5	7		
C	0	0	0	0	0	0	0	0	0	3	3	3	3	3	3	5	5	
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	3	
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
A	i				j	=	i		i+1		j	+	i		k	j	0	

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
C	(((.	.	.)	.	((.	.	.))))	.
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	2	2	2	2	4	4	5	7	7	8	10	10
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	10
A	0	0	0	0	0	0	0	2	2	2	2	5	5	5	5	8	8	8
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	8
U	0	0	0	0	0	0	0	0	2	3	5	5	5	6	6	7	7	7
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7	7	7
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	3	5	5	5
U	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	3	3
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	2	2
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	i				j	=	i		i+1		j	+	i	≥ θ	k	j	0	0
C																0	0	0
G																0	0	0
A																0	0	0

Nussinov/Jacobson

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
C	(((.	.	.)	.	((.	.	.))))	.
G	0	0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G	0	0	0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A	0	0	0	0	0	0	2	2	2	2	4	4	5	7	7	8	10	.
U	0	0	0	0	0	0	0	0	2	2	4	5	7	7	7	8	10	.
A	0	0	0	0	0	0	0	2	2	2	2	5	5	5	5	8	8	.
C	0	0	0	0	0	0	0	0	0	2	5	5	5	5	5	8	8	.
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	6	7	.	.
U	0	0	0	0	0	0	0	0	2	3	5	5	5	5	5	7	.	.
C	0	0	0	0	0	0	0	0	3	3	3	3	3	3	3	5	5	.
U	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	3	.	.
U	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	.	.
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	.
G	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	.
A	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	.
C	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	.
G	i					j	=	i		i+1		j	+	i		k	j	0
A																		0

References I

-  A. Condon, B. Davy, B. Rastegari, S. Zhao, and F. Tarrant.
Classifying RNA pseudoknotted structures.
Theoretical Computer Science, 320(1):35–50, 2004.
-  R. B. Lyngsø and C. N. S. Pedersen.
RNA pseudoknot prediction in energy-based models.
Journal of Computational Biology, 7(3-4):409–427, 2000.
-  N. Leontis and E. Westhof.
Geometric nomenclature and classification of RNA base pairs.
RNA, 7:499–512, 2001.
-  Jan Manuch, Chris Thachuk, Ladislav Stacho, and Anne Condon.
Np-completeness of the direct energy barrier problem without pseudoknots.
In Russell Deaton and Akira Suyama, editors, *DNA Computing and Molecular Programming*, volume 5877 of *Lecture Notes in Computer Science*, pages 106–115. Springer Berlin Heidelberg, 2009.
-  Lioudmila V Sharova, Alexei A Sharov, Timur Nedorezov, Yulan Piao, Nabeebi Shaik, and Minoru S H Ko.
Database for mRNA half-life of 19 977 genes obtained by DNA microarray analysis of pluripotent and differentiating mouse embryonic stem cells.
DNA Res, 16(1):45–58, Feb 2009.