M2 AMI2B - Lecture 2 Boltzmann ensemble

Yann Ponty

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December 2nd, 2016

Outline

Foreword

- Boltzmann ensemble
 - Nussinov: Minimisation ⇒ Counting
 - Computing the partition function
 - Statistical sampling
 - Inside/outside

Grammars and DP schemes over sequences

$$= \underbrace{\qquad \qquad \qquad }_{i} = \underbrace{\qquad \qquad }_{j} + \underbrace{\qquad \qquad \geq \theta}_{k}$$

$$\begin{array}{lcl} \textit{N}_{i,t} & = & 0, & \forall t \in [i,i+\theta] \\ \textit{N}_{i,j} & = & \min \left\{ \begin{array}{ll} \textit{N}_{i+1,j} & \textit{i} \text{ unpaired} \\ \displaystyle \min_{k=i+\theta+1} \Delta \textit{G}_{i,k} + \textit{N}_{i+1,k-1} + \textit{N}_{k+1,j} & \textit{i} \text{ paired with } k \end{array} \right. \end{array}$$

Ambiguity? Consider *i*: Either **unpaired**, or **paired** to k. Sets of structures generated in these two cases are clearly disjoint. (also holds for various values of k) \Rightarrow **Unambiguous** decomposition

Completeness? True, since scheme explores every possible outcome for i. + Induction on interval length \Rightarrow **Complete** decomposition

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Foreword

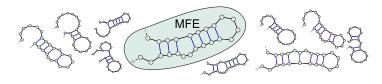
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The canonical Boltzmann Ensemble

RNA *breathes* ⇒ There is no more than a single conformation.

New paradigm

The conformations of an RNA coexist in the **Boltzmann distribution**.



Consequence: The MFE probability can be arbitrarily small.

 \Rightarrow To understand how RNA acts, one must account for the set of alternative structures.

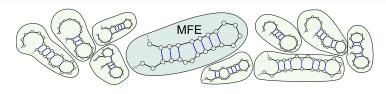
In particular, structurally close structures may *ally*, and become the most realistic candidate in the search for a functional conformation.

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Boltzmann Distribution: Definition

For each structure S compatible with an RNA ω , the Boltzmann distribution associates a **Boltzmann factor** $\mathcal{B}_{S,\omega}=e^{-\mathcal{E}_{S,\omega}\over RT}$, where:

- ▶ $E_{S,\omega}$ is the free-energy S (kCal.mol⁻¹)
- ► *T* is the temperature (K)
- ► R is the perfect gaz constant (1.986.10⁻³ kCal.K⁻¹.mol⁻¹)

To obtain a distribution, one simply renormalizes by the partition function

$$\mathcal{Z}_{\omega} = \sum_{S \in \mathcal{S}_{\omega}} e^{rac{-\mathcal{E}_{S,\omega}}{RT}}$$

where S_{ω} is the set of conformations that are compatibles with ω .

The **Boltzmann probability** of a structure S is simply given by

$$P_{\mathcal{S},\omega} = rac{e^{rac{-\mathcal{E}_{\mathcal{S},\omega}}{RT}}}{\mathcal{Z}_{\omega}}.$$

Nussinov/Jacobson DP scheme

$$\begin{array}{lcl} \textit{N}_{i,t} & = & 0, & \forall t \in [i,i+\theta] \\ \textit{N}_{i,j} & = & \min \left\{ \begin{array}{ll} \textit{N}_{i+1,j} & \textit{i} \text{ unpaired} \\ \displaystyle \min_{k=i+\theta+1} \Delta \textit{G}_{i,k} + \textit{N}_{i+1,k-1} + \textit{N}_{k+1,j} & \textit{i} \text{ paired with } k \end{array} \right. \end{array}$$

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Nussinov/Jacobson DP scheme

Recurrence for minimal free-energy of a fold :

$$\begin{array}{lcl} \textit{N}_{i,t} & = & 0, \quad \forall t \in [i,i+\theta] \\ \textit{N}_{i,j} & = & \min \left\{ \begin{array}{ll} \textit{N}_{i+1,j} & \textit{(i unpaired)} \\ \min_{k=i+\theta+1}^{j} \textit{E}_{i,k} + \textit{N}_{i+1,k-1} + \textit{N}_{k+1,j} & \textit{(i comp. with k)} \end{array} \right. \end{array}$$

Recurrence for counting compatible structures:

$$\begin{array}{lcl} C_{i,t} & = & 1, & \forall t \in [i,i+\theta] \\ \\ C_{i,j} & = & \sum \left\{ \begin{array}{ll} C_{i+1,j} & \text{(i unpaired)} \\ \sum_{k=i+\theta+1}^{j} 1 \times C_{i+1,k-1} \times C_{k+1,j} & \text{(i comp. with k)} \end{array} \right. \end{array}$$

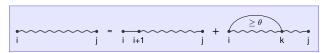
Decomposition matters, and the rest (MFE, count...) follows!

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$$\mathcal{Z}_{i,t} = 1, \quad \forall t \in [i, i+\theta]$$

$$\mathcal{Z}_{i,j} = \sum \left\{ \sum_{k=i+\theta+1}^{j} 1 \times \mathcal{Z}_{i+1,k-1} \times \mathcal{Z}_{k+1,j} \right.$$



$$\mathcal{Z}_{i,t} = 1, \quad \forall t \in [i, i+\theta]$$

$$\mathcal{Z}_{i,j} = \sum \left\{ \sum_{k=i+\theta+1}^{j} \frac{e^{-\mathcal{E}_{\text{DB}}(i,k)}}{\theta^{T}} \times \mathcal{Z}_{i+1,k-1} \times \mathcal{Z}_{k+1,j} \right.$$

$$\mathcal{M}'_{i,j} = \operatorname{Min} \left\{ \begin{aligned} & \mathcal{E}_{\mathcal{H}}(i,j) \\ & \mathcal{E}_{\mathcal{S}}(i,j) + \mathcal{M}'_{i+1,j-1} \\ & \operatorname{Min}(\mathcal{E}_{\mathcal{B}}(i,i',j',j') + \mathcal{M}'_{\mathcal{V},j'}) \\ & a + c + \operatorname{Min}\left(\mathcal{M}_{i+1,k-1} + \mathcal{M}^{1}_{k,j-1}\right) \end{aligned} \right. \\ & \mathcal{M}_{i,j} = \operatorname{Min} \left\{ \operatorname{Min}\left(\mathcal{M}_{i,k-1},b(k-1)\right) + \mathcal{M}^{1}_{k,j} \right\} \\ & \mathcal{M}^{1}_{i,j} = \operatorname{Min} \left\{ b + \mathcal{M}^{1}_{i,j-1},c + \mathcal{M}'_{i,j} \right\}$$

$$\mathcal{M}'_{i,j} = \operatorname{Min} \left\{ \begin{array}{c} e^{\frac{-E_{\mu}(i,j)}{RT}} \\ e^{\frac{-E_{\mu}(i,j)}{RT}} + \mathcal{M}'_{i+1,j-1} \\ \operatorname{Min} \left(e^{\frac{-E_{\mu}(i,i',j',j)}{RT}} + \mathcal{M}'_{i',j'} \right) \\ e^{\frac{-(a+c)}{RT}} + \operatorname{Min} \left(\mathcal{M}_{i+1,k-1} + \mathcal{M}^1_{k,j-1} \right) \end{array} \right.$$

$$\mathcal{M}_{i,j} = \operatorname{Min} \left\{ \operatorname{Min} \left(\mathcal{M}_{i,k-1}, e^{\frac{-C(k-1)}{RT}} \right) + \mathcal{M}^1_{k,j} \right\}$$

$$\mathcal{M}^1_{i,j} = \operatorname{Min} \left\{ e^{\frac{b}{RT}} + \mathcal{M}^1_{i,j-1}, e^{\frac{-c}{RT}} + \mathcal{M}'_{i,j} \right\}$$

Partition function = Weighted count over compatible structures

$$\mathcal{M}'_{i,j} = \operatorname{Min} \left\{ \begin{array}{l} e^{\frac{-E_{H}(i,j)}{RT}} \\ e^{\frac{-E_{S}(i,j)}{RT}} \mathcal{M}'_{i+1,j-1} \\ \operatorname{Min} \left(e^{\frac{-E_{B}(i,j',j',j)}{RT}} \mathcal{M}'_{i',j'} \right) \\ e^{\frac{-(a+c)}{RT}} \operatorname{Min} \left(\mathcal{M}_{i+1,k-1} \mathcal{M}^{1}_{k,j-1} \right) \\ \mathcal{M}_{i,j} = \operatorname{Min} \left\{ \operatorname{Min} \left(\mathcal{M}_{i,k-1}, e^{\frac{-E_{k}(a-1)}{RT}} \right) \mathcal{M}^{1}_{k,j} \right\} \\ \mathcal{M}^{1}_{i,j} = \operatorname{Min} \left\{ e^{\frac{-b}{RT}} \mathcal{M}^{1}_{i,j-1}, e^{\frac{-c}{RT}} \mathcal{M}'_{i,j} \right\} \end{array} \right.$$

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$$Z'(i,j) = \sum \begin{cases} e^{\frac{-E_{R}(i,j)}{Rl}} \\ e^{\frac{-E_{S}(i,j)}{Rl}} Z'(i+1,j-1) \\ + \sum_{j} \left(e^{\frac{-E_{B}(i,i',j',j)}{Rl}} Z'(i',j') \right) \\ + e^{\frac{-(a+c)}{Rl}} \sum_{j} \left(Z(i+1,k-1) Z^{1}(k,j-1) \right) \end{cases}$$

$$Z(i,j) = \sum_{j} \left(Z(i,k-1) + e^{\frac{-c(k-1)}{Rl}} \right) Z^{1}(k,j)$$

$$Z^{1}(i,j) = e^{\frac{-b}{Rl}} Z^{1}(i,j-1) + e^{\frac{-c}{Rl}} Z'(i,j)$$

Partition function = Weighted count over compatible structures

$$\begin{split} \mathcal{Z}_{i,t} &= 1, \quad \forall t \in [i, i+\theta] \\ \mathcal{Z}_{i,j} &= \sum \left\{ \sum_{k=i+\theta+1}^{\mathcal{Z}_{i+1,j}} e^{\frac{-\mathcal{E}_{bp}(i,k)}{RT}} \times \mathcal{Z}_{i+1,k-1} \times \mathcal{Z}_{k+1,j} \right. \end{split}$$

Validity of a partition function computation:

- Completeness/Unambiguity of decomposition scheme

$$\begin{split} e^{-E_{bp}(i,k)/RT} \times \mathcal{Z}_{i+1,k-1} \times \mathcal{Z}_{k+1,j} &= \cdot \sum_{x} e^{-E(x)/RT} \cdot \sum_{y} e^{-E(y)/RT} \\ &= \sum_{x,y} e^{-a/RT} \cdot e^{-E(x)/RT} \cdot e^{-E(y)/RT} \\ &= \sum_{x,y} e^{-(E_{bp}(i,k)+E(x)+E(y))/RT} \end{split}$$

Partition function = Weighted count over compatible structures

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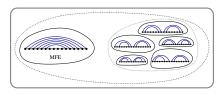
- Completeness/Unambiguity of decomposition scheme
- ▶ Correctness of Boltzmann factor
 Weight induced by backtrack = Product of derivations weights $e^{-E/RT}$ → Weight products \Leftrightarrow Summing energy terms

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Statistical sampling of RNA 2^{ary} structures

MFE (\Leftrightarrow Max probability) may be **heavily dominated** by a set \mathcal{B} of **structurally similar** suboptimal structures.

 \Rightarrow Functional conformation probably closer to \mathcal{B} than to MFE.



Proof-of-concept: [DCL05]

- Sample structures within Boltzmann probability
- Cluster structures
- ▶ Build and return consensus structure of the heaviest cluster
- \Rightarrow Relative improvement for specificity (+17.6%) and sensitivity (+21.74%, except group II introns)

Problem

How to sample from the Boltzmann ensemble?

Goal [DL03]: From sequence ω , draw S with prob. $e^{-E_S/RT}/\mathcal{Z}$

- ① Draw uniform random number $r \in [0, \mathcal{Z}'(i, j))$
- ② Subtract from r the contributions of $\mathbb{Z}'(i,j)$ until r < 0
- 3 Recurse over associated regions/matrices

$$\mathbf{Z}'(i,j) \in \left\{ \begin{array}{l} -\rightarrow e^{\frac{-E_{H}(i,j)}{RT}} + e^{\frac{-E_{S}(i,j)}{RT}} \mathbf{Z}'(i+1,j-1) \\ \\ \bullet e^{\frac{-(a+c)}{RT}} \sum \left(\mathbf{Z}(i+1,k-1) \mathbf{Z}^{1}(k,j-1) \right) \end{array} \right.$$

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\downarrow & \downarrow \\
\mathbf{A}_{1} | \mathbf{A}_{2} | \mathbf{B}_{i} | \mathbf{B}_{i+1} | \dots | \mathbf{B}_{j-1} | \mathbf{B}_{j} | \mathbf{C}_{i} | \mathbf{C}_{i+1} | \dots | \mathbf{C}_{j-1} | \mathbf{C}_{j}
\end{cases}$$

Goal [DL03]: From sequence ω , draw S with prob. $e^{-E_S/RT}/\mathcal{Z}$

- **1** Draw uniform random number $r \in [0, \mathbb{Z}'(i, j))$
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\end{cases}$$

Goal [DL03]: From sequence ω , draw S with prob. $e^{-E_S/RT}/\mathcal{Z}$

Principle: Choose derivation with prob. prop. to its contribution to part. fun. **Precomputation:** Compute part. fun. versions of matrices $(\mathcal{Z}, \mathcal{Z}', \mathcal{Z}^1)$. Stochastic backtrack:

- **1** Draw uniform random number $r \in [0, \mathbb{Z}'(i, j))$
- ② Subtract from r the contributions of $\mathcal{Z}'(i,j)$ until r < 0

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Correctness: Each $S \in S_{\omega}$ uniquely generated (DP scheme unambiguity) Therefore the probability of generated S is

$$p_S = \frac{\mathcal{B}(E_1)}{\mathcal{B}(S_w)} \cdot \frac{\mathcal{B}(E_2)}{\mathcal{B}(E_1)} \cdot \frac{\mathcal{B}(E_3)}{\mathcal{B}(E_2)} \cdots \frac{\mathcal{B}(\{S\})}{\mathcal{B}(E_m)}$$

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$$p_{S} = \frac{1}{\mathcal{B}(S_{w})} \cdot \frac{1}{1} \cdot \frac{1}{1} \cdot \dots \frac{\mathcal{B}(\{S\})}{1}$$

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$$p_{S} = \frac{\mathcal{B}(\{S\})}{\mathcal{B}(S_{W})} = \frac{e^{-E_{S}/RT}}{\mathcal{Z}} = P_{S,\omega}$$

Complexity

Goal [DL03]: From sequence ω , draw S with prob. $e^{-E_S/RT}/\mathcal{Z}$

Stochastic backtrack:

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$$\otimes e^{\frac{-(a+c)}{RT}} \sum \left(\mathcal{Z}(i+1,k-1) \mathcal{Z}^{1}(k,j-1) \right)$$

Average-case complexity in $\Theta(k \times n\sqrt{n})$ (homopolymer model) [Pon08]. Boustrophedon search $\Rightarrow \mathcal{O}(k \times n \log n)$ worst-case [Pon08].

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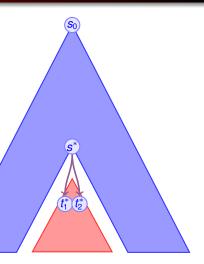
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\end{cases}$$

After $\Theta(n)$ operations, recurse over region of length n-1 \Rightarrow Worst-case complexity in $\mathcal{O}(k \times n^2)$ for k samples

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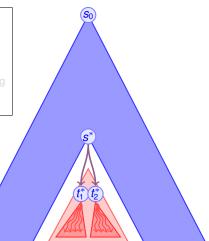
Structure including base pair (i, k):

- ▶ Inside: Structures over [i+1, k-1]
- ightharpoonup Outside: **Contexts** of interval (i, k)
 - ▶ \forall interval $[i,j], i < j \le k$
 - ► Complete structure by generating brother intervals ([k + 1, j]) and recurse over the father of [i, k].



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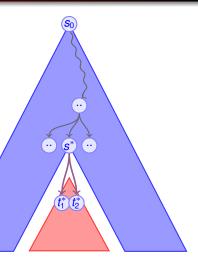
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 - ► Complete structure by generating brother intervals ([k + 1, j]) and recurse over the father of [i, k].



Structure including base pair (i, k): ▶ Inside: Structures over [i+1, k-1]► Outside: Contexts of interval (i, k) ▶ \forall interval $[i, j], i < j \le k$ ► Complete structure by generating brother intervals ([k+1, j]) and recurse over the **father** of [i, k].

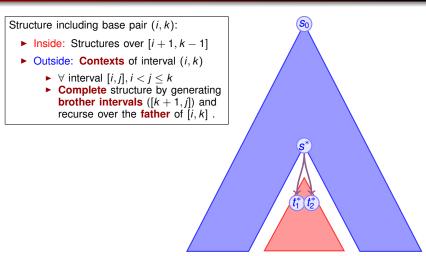
Structure including base pair (i, k):

- ▶ Inside: Structures over [i+1, k-1]
- ► Outside: Contexts of interval (i, k)
 - ▶ \forall interval $[i,j], i < j \le k$
 - Complete structure by generating brother intervals ([k + 1,j]) and recurse over the father of [i, k].



Structure including base pair (i, k): ▶ Inside: Structures over [i+1, k-1]► Outside: Contexts of interval (i, k) ▶ \forall interval $[i, j], i < j \le k$ ► Complete structure by generating brother intervals ([k+1, j]) and recurse over the **father** of [i, k].

Structure including base pair (i, k): ▶ Inside: Structures over [i+1, k-1]► Outside: Contexts of interval (i, k) ▶ \forall interval $[i, j], i < j \le k$ ► Complete structure by generating brother intervals ([k+1, j]) and recurse over the **father** of [i, k].



Whenever some further **technical conditions** are satisfied, this decomposition is **complete** and **unambiguous**, and implies a **simple recurrence** for computing the base pair probability matrix in $\Theta(n^3)$.

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