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# THREE GEOMETRIC FAMILIES



triangulation = maximal crossing-free set of edges

pseudotriangulation = maximal crossing-free pointed set of edges

*k*-triangulation = maximal (k + 1)-crossing-free set of edges

# **RELEVANT EDGES**



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**pseudotriangulation** = maximal crossing-free pointed set of edges

*k*-triangulation = maximal (k + 1)-crossing-free set of edges



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= decomposition into triangles

pseudotriangulation = maximal crossing-free pointed set of edges = decomposition into pseudotriangles

k-triangulation = maximal (k + 1)-crossing-free set of edges = decomposition into k-stars



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## DECOMPOSITIONS OF SURFACES















flip = exchange an internal edge with the common bisector of the two adjacent cells



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#### **FLIPS ON SURFACES**



fundamental group of the flip graph  $G_{n,k} \mapsto mapping$  class group of the surface  $S_{n,k}$ 

# THREE GEOMETRIC STRUCTURES



- associahedron  $\longleftrightarrow$
- pseudotriangulations polytope  $\longleftrightarrow$ 
  - multiassociahedron  $\leftrightarrow$



VP & M. Pocchiola, Multitriangulations, pseudotriangulations and primitive sorting networks, 2012.

# LINE SPACE OF THE PLANE = MÖBIUS STRIP
















































VP & M. Pocchiola, Multitriangulations, pseudotriangulations and primitive sorting networks, 2012.

**MULTIPSEUDOTRIANGULATIONS** 



VP & M. Pocchiola, Multitriangulations, pseudotriangulations and primitive sorting networks, 2012.



VP & M. Pocchiola, Multitriangulations, pseudotriangulations and primitive sorting networks, 2012.

# SORTING NETWORKS





# SUBWORD COMPLEX

(W,S) a finite Coxeter system,  $Q = q_1 q_2 \cdots q_m$  a word on S,  $\rho$  an element of W.

Subword complex  $S(Q, \rho) =$  simplicial complex of subsets of positions of Q whose complement contains a reduced expression of  $\rho$ .

A. Knutson & E. Miller, Subword complexes in Coxeter groups, 2004.



Classical situation of type A:

- Coxeter group  $W = \mathfrak{S}_{n+1}$
- simple system  $S = \{\tau_i \mid i \in [n]\}$ , where  $\tau_i = (i \ i+1)$
- word  $Q = q_1 q_2 \cdots q_m$  on S
- $\bullet~\rho$  element of W

The subword complex can be interpreted with a primitive sorting network:

- $\mathcal{N}_{\mathbf{Q}}$  formed by n+1 levels and m commutators
- facets of  $\mathcal{S}(Q, \rho) \longleftrightarrow$ pseudoline arrangements on  $\mathcal{N}_Q$



### FLIPS

flip = exchange a contact with the corresponding crossing





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C. Ceballos, JP. Labbé & C. Stump, Subword complexes, cluster complexes, & gener. multiassoc., 2011. VP & C. Stump, Brick polytopes of spherical subword complexes, 2015.

## **ROOT FUNCTION**



For a facet I of  $S(Q, \rho)$  and a position  $k \in [m]$ , define the root  $r(I, k) = Q_{[k-1]\setminus I}(\alpha_{q_k})$ , where  $Q_{[k-1]\setminus I}$  is the product of all reflections  $q_j$  for j from 1 to k-1 but not in I.

The root function of the facet I is  $\mathsf{r}(I,\cdot):[m]\longrightarrow \Phi$ 

The root configuration of I is  $R(I) = {r(I, i) | i \in I}$ 

# **ROOT FUNCTION & FLIPS**



**PROPOSITION**. The root function encodes flips in subword complexes:

- 1. The map  $r(I, \cdot)$  is a bijection from the complement of I to  $inv(\rho)$ .
- 2. If I and J are two adjacent facets of  $S(Q, \rho)$  with  $I \smallsetminus i = J \smallsetminus j$ , then j is the unique position in the complement of I such that  $r(I, i) = \pm r(I, j)$ .
- 3. In the situation of 2, the root function of J is obtained from that of I by

$$\mathsf{r}(J,k) = \begin{cases} s_{\mathsf{r}(I,i)}(\mathsf{r}(I,k)) & \text{if } \min(i,j) < k \le \max(i,j), \\ \mathsf{r}(I,k) & \text{otherwise.} \end{cases}$$

C. Ceballos, JP. Labbé & C. Stump, Subword complexes, cluster complexes, & gener. multiassoc., 2011.



S. Fomin & A. Zelevinsky, Cluster Algebras I, II, III, IV, 2002 – 2007. C. Hohlweg, C. Lange & H. Thomas, Permutahedra and generalized associahedra, 2011. C. Ceballos, JP. Labbé & C. Stump, Subword complexes, cluster complexes, & gener. multiassoc., 2011. VP & C. Stump, Brick polytopes of spherical subword complexes, 2015. C. Hohlweg, Permutahedra and associahedra, 2013.

# **CLUSTER ALGEBRAS**

cluster algebra = commutative ring generated by distinguished cluster variables grouped into overlapping clusters

clusters computed by a mutation process :

cluster seed = algebraic data  $\{x_1, \ldots, x_n\}$ , combinatorical data B (matrix or quiver) cluster mutation =  $(\{x_1, \ldots, x_k, \ldots, x_n\}, B) \xleftarrow{\mu_k} (\{x_1, \ldots, x'_k, \ldots, x_n, \mu_k(B))$ 

$$x_{k} \cdot x_{k}' = \prod_{i, b_{ik} > 0} x_{i}^{b_{ik}} + \prod_{i, b_{ik} < 0} x_{i}^{-b_{ik}}$$
$$\left(\mu_{k}(B)\right)_{ij} = \begin{cases} -b_{ij} & \text{if } k \in \{i, j\} \\ b_{ij} + |b_{ik}| \cdot b_{kj} & \text{if } k \notin \{i, j\} \text{ and } b_{ik} \cdot b_{kj} > 0 \\ b_{ij} & \text{otherwise} \end{cases}$$

cluster complex = simplicial complex w/ vertices = cluster variables & facets = clusters

S. Fomin & A. Zelevinsky, Cluster Algebras I, II, III, IV, 2002 – 2007.

# **CLUSTER ALGEBRAS**

THEOREM. (Laurent phenomenon)

All cluster variables are Laurent polynomials in the variables of the initial cluster seed.

S. Fomin & A. Zelevinsky, Cluster algebras I: Fundations, 2002.

**THEOREM**. (Classification)

Finite type cluster algebras are classified by the Cartan-Killing classification for crystallographic root systems.

S. Fomin & A. Zelevinsky, Cluster algebras II: Finite type classification, 2003.

In fact, for a root system  $\Phi,$  there is a bijection

 $\begin{array}{lll} \text{cluster variables} & \longleftrightarrow & \Phi_{\geq -1} = \Phi^+ \cup -\Delta \\ y = \frac{F(x_1, \dots, x_n)}{x_1^{d_1} \cdots x_n^{d_n}} & \longleftrightarrow & \beta = d_1 \alpha_1 + \cdots + d_n \alpha_n \\ & \text{cluster} & \longleftrightarrow & \text{c-cluster} \\ & \text{cluster complex} & \longleftrightarrow & \text{c-cluster complex} \end{array}$ 

# FINITE CLUSTER COMPLEXES ARE SUBWORD COMPLEXES

New approach to the combinatorics and geometry of the cluster complex:

THEOREM. The subword complex  $\mathcal{S}(cw_{\circ}(c))$  is isomorphic to the cluster complex.

C. Ceballos, JP. Labbé & C. Stump, Subword complexes, cluster complexes, & gener. multiassoc., 2011.

$$\begin{array}{cccc} \mathsf{cluster variables} & \longleftrightarrow & \Phi_{\geq -1} = \Phi^+ \cup -\Delta & \longleftrightarrow & \mathsf{position in } \operatorname{cw}_\circ(\operatorname{c}) \\ y = \frac{F(x_1, \ldots, x_n)}{x_1^{d_1} \cdots x_n^{d_n}} & \longleftrightarrow & \beta = d_1 \alpha_1 + \cdots + d_n \alpha_n & \longleftrightarrow & \begin{cases} i & \mathsf{if } \beta = -\alpha_{c_i} \\ j & \mathsf{if } \beta = \mathsf{r}([n], j) \\ \mathsf{cluster} & \longleftrightarrow & \mathsf{c-cluster} & \longleftrightarrow & \mathsf{facet of } \mathcal{S}(\operatorname{cw}_\circ(\operatorname{c})) \\ \mathsf{cluster complex} & \longleftrightarrow & \mathsf{c-cluster complex} & \longleftrightarrow & \mathsf{subword complex } \mathcal{S}(\operatorname{cw}_\circ(\operatorname{c})) \end{array}$$

## TYPE $D_n$ AS PSEUDOTRIANGULATIONS



C. Ceballos & VP, Cluster algebras of type D: pseudotriangulations approach, 2015<sup>+</sup>.

## TYPE $D_n$



C. Ceballos & VP, Cluster algebras of type D: pseudotriangulations approach, 2015<sup>+</sup>.



VP & F. Santos, The brick polytope of a sorting network, 2012. VP & C. Stump, Brick polytopes of spherical subword complexes, 2015.



 ${\cal N}$  a sorting network with n+1 levels



 $\mathcal{N}$  a sorting network with n+1 levels  $\Lambda$  pseudoline arrangement supported by  $\mathcal{N} \longmapsto \text{brick vector } B(\Lambda) \in \mathbb{R}^{n+1}$ 



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 ${\cal N}$  a sorting network with n+1 levels

Λ pseudoline arrangement supported by  $\mathcal{N}$  → brick vector B(Λ) ∈  $\mathbb{R}^{n+1}$ B(Λ)<sub>j</sub> = number of bricks of  $\mathcal{N}$  below the *j*th pseudoline of Λ

Brick polytope  $\mathcal{B}(\mathcal{N}) = \text{conv} \{ \mathsf{B}(\Lambda) \mid \Lambda \text{ pseudoline arrangement supported by } \mathcal{N} \}$ 



# WEIGHT FUNCTION, BRICK VECTOR & BRICK POLYTOPE

(W, S) a finite Coxeter system,  $Q = q_1 q_2 \cdots q_m$  a word on S,  $w_\circ$  longest element of W.  $S(Q) = S(Q, w_\circ)$  spherical subword complex.

To a facet I of S(Q) and a position  $k \in [m]$ , associate a weight  $w(I, k) = Q_{[k-1]\setminus I}(\omega_{q_k})$ , where  $Q_{[k-1]\setminus I}$  is the product of all reflections  $q_j$  for j from 1 to k-1 but not in I. The brick vector of I is the vector  $B(I) = \sum_{k \in [m]} w(I, k)$ .

The brick polytope is the convex polytope  $\mathcal{B}(Q) = \text{conv} \{ \mathsf{B}(I) \mid I \text{ facet of } \mathcal{S}(Q) \}.$ 



In type A,  $w(I, k) = \text{characteristic vector of the pseudolines passing above the <math>k$ th brick. B(I) = (number of bricks below the jth pseudoline of  $I)_{j \in [n+1]}$ 

## BRICK VECTORS AND FLIPS



If  $\Lambda$  and  $\Lambda'$  are two pseudoline arrangements supported by  $\mathcal{N}$  and related by a flip between their *i*th and *j*th pseudolines, then  $B(\Lambda) - B(\Lambda') \in \mathbb{N}_{>0} (e_j - e_i)$ .

THEOREM. The cone of the brick polytope  $\mathcal{B}(Q)$  at the brick vector B(I) is generated by -R(I), for any facet I of  $\mathcal{S}(Q)$ .
## BRICK POLYTOPE

The brick polytope is the convex polytope  $\mathcal{B}(Q) = \text{conv} \{B(I) \mid I \text{ facet of } \mathcal{S}(Q)\}.$ 

THEOREM. The polar of the brick polytope  $\mathcal{B}(Q)$  realizes the subword complex  $\mathcal{S}(Q)$  $\iff Q$  is such that R(I) is linearly independent, for I facet of  $\mathcal{S}(Q)$ .

THEOREM. If Q is root-independent, the cone of the brick polytope  $\mathcal{B}(Q)$  at the brick vector B(I) is generated by -R(I), for any facet I of  $\mathcal{S}(Q)$ .

THEOREM. If  ${\rm Q}$  is root-independent, the Coxeter fan refines the normal fan of the brick polytope. More precisely,

normal cone of B(I) in  $\mathcal{B}(Q) = \bigcup_{\substack{w \in W \\ R(I) \subset w(\Phi^+)}} w($ fundamental cone ).

## NORMAL FAN

THEOREM. The Coxeter fan refines the normal fan of the brick polytope.



THEOREM. The brick polytope  $\mathcal{B}(cw_{\circ}(c))$  realizes the subword complex  $\mathcal{S}(cw_{\circ}(c))$ .

THEOREM. The brick polytope  $\mathcal{B}(cw_{\circ}(c))$  is a translate of the known realizations of the generalized associahedron.

F. Chapoton, S. Fomin & A. Zelevinsky, Polytopal realizations of generalized associahedra, 2002.
C. Hohlweg, C. Lange & H. Thomas, Permutahedra and generalized associahedra, 2011.
S. Stella, Polyhedral models for generalized associahedra via Coxeter elements, 2013.
C. Hohlweg, Permutahedra and associahedra, 2013.

## GENERALIZED ASSOCIAHEDRA



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