Representation for the radical of a finitely generated differential ideal

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Abstract
We give an algorithm which represents the radical \( \mathcal{J} \) of a finitely generated differential ideal as an intersection of radical differential ideals. The computed representation provides an algorithm for testing membership in \( \mathcal{J} \). This algorithm works over either an ordinary or a partial differential polynomial ring of characteristic zero. It has been programmed. We also give a method to obtain a characteristic set of \( \mathcal{J} \), if the ideal is prime.


1 Introduction
Let \( \Sigma \) be a finite subset of a differential polynomial ring \( K \{ y_1, \ldots, y_n \} \), where \( K \) denotes a differential field, ordinary or with partial derivatives, of characteristic zero. Let \( \mathcal{R} \) be a ranking of the set of derivatives of these \( y_i \).

We present an algorithm, called Rosenfeld–Gröbner, which represents the least radical differential ideal containing \( \Sigma \) as a finite intersection of radical differential ideals \( \mathcal{J}_i \):

\[ \{ \Sigma \} = \mathcal{J}_1 \cap \cdots \cap \mathcal{J}_r. \]

Each radical differential ideal \( \mathcal{J}_i \) is described by a differential system of polynomial equations and inequations \( \Omega_i \), and a (non-differential) Gröbner basis \( B_i \) satisfying:

1) \( \Omega_i \) and \( B_i \) provide an algorithm for testing membership in \( \mathcal{J}_i \), through simple reductions,

2) \( B_i \) depends only on the differential ideal \( \mathcal{J}_i \) and the ranking \( \mathcal{R} \).

Thus, the set of tuples \( (\Omega_i, B_i) \) allows to decide the membership in the differential ideal \( \{ \Sigma \} \) by simple reductions.

The intersection computed may not be minimal. Unfortunately, we do not know how to test redundancy, which is a problem close to the open problem related in [Ko], page 166. However, when we know that the differential ideal \( \{ \Sigma \} \) is prime, the formula mentioned above may be simplified to:

\[ \{ \Sigma \} = \mathcal{J}_1, \]

and we give a method for calculating, starting with the Gröbner basis \( B_1 \), a characteristic set of the differential ideal \( \{ \Sigma \} \), in the sense of Ritt, relative to the ranking \( \mathcal{R} \).

The Rosenfeld–Gröbner algorithm relies essentially on three theorems:

1) the theorem of zeros of Hilbert, which states that a polynomial \( p \) belongs to the radical of an ideal given by a finite family of generators \( \Sigma \) if and only if the system of equations and inequations \( \Sigma = 0 \), \( p \neq 0 \) has no solutions; we use this theorem, in the algebraic case, and in the differential case.

2) a lemma of Rosenfeld [Ro], which gives a sufficient condition so that a system of polynomial equations and inequation admits a differential solution if and only if it admits a purely algebraic solution; the systems \( \Omega_i \) described above satisfy the condition of Rosenfeld.

3) a lemma of D. Lazard, which establishes in particular that the ideals \( \mathcal{J}_i \) described above are radical.

The algorithm which we describe utilizes only the operations and equality test with zero in the base field \( K \): we refer to the reduction algorithm of Ritt, the computations of Gröbner bases, and splittings similar to those in the elimination methods of Seidenberg [Se1]. It does not need any factorization.

An implementation of Rosenfeld–Gröbner has been realized [Bo1], in the language C. It makes calls to the big number library of PARI and the software GB [FGM] for the calculation of Gröbner bases.

In order to place the interest of this algorithm,
let us describe in a few words the principals of existing methods.

Ritt gave [Ri] a method to decompose the radical of a differential ideal as an intersection of prime differential ideals, providing a characteristic set for each of these ideals. That algorithm is inconvenient because it is only partially effective: it proceeds by factorization over a tower of algebraic field extensions of the field of coefficients. To our knowledge, it has not been implemented.

Ollivier [O] and Carrà-Ferro [Ca] have independently tried to generalize to differential algebra the Gröbner bases invented by Buchberger [Bu] for the study of polynomial ideals in commutative algebra. These differential Gröbner bases are in general however infinite.

Another attempt to define differential Gröbner bases has been done by E. Mansfield [Ma]. The algorithm DIFGBASIS, implemented in MAPLE, utilizes Ritt’s algorithm of reduction and then always terminates. In general however, it cannot guarantee its output to be a differential Gröbner basis.

We may remark that the membership problem in an arbitrary differential ideal is undecidable [GMO], and the membership problem of a finitely generated differential ideal is still open.

The elimination algorithms of Seidenberg [Se] are more general. Rosenfeld–Gröbner borrows from them the idea to combine Hilbert’s theorem of zeros and Ritt’s algorithm of reduction. They decide the membership problem in the radical $\Sigma$ of a finitely generated differential ideal by successively eliminating all the unknowns appearing in the polynomials of $\Sigma$. They use only the operations of the base field $K$, but present two inconveniences: first, the description of the differential ideal $\Sigma$ they give is not usable to test the membership in the ideal afterwards; second, their behavior is a lot more explosive in practice than that of Rosenfeld–Gröbner, because they are restricted to the elimination rankings. This phenomenon is particularly striking in the case of systems with partial derivatives.

### 2 Preliminaries

**Differential algebra.** In this paper, $K$ denotes a differential field of characteristic zero endowed with a certain number of derivations denoted $\delta_1, \ldots, \delta_m$. Let $u$ be an element of $K$. We denote by $\theta$ the derivation operators $\frac{\partial u}{\partial \delta_i}$ and by $\theta u$ the element of $K$ obtained by differentiating $u$ at times by $\delta_1, \ldots, \delta_m$ times by $\delta_\theta$. The sum of the exponents $a_i$ is called the order of the operator $\theta$. The identity operator is of order 0. The other operators are said to be proper.

Let $S$ be a subset of a differential ring $R$ which contains $K$. We denote respectively by $K[S]$ and $K\{S\}$ the smallest subring and the smallest differential subring of $R$ containing $K$ and $S$ (denoting by $\Theta S$ the smallest subset of $R$ containing $S$ and stable under differentiation, we have $K[\Theta S] = K\{S\}$).

Let $S$ be a subset of a differential ring $R$. We denote by $(S)$ and $[S]$ the smallest ideal and the smallest differential ideal of $R$ which contains $S$ (we have $(\Theta S) = [S]$). The smallest radical differential ideal containing $S$, denoted by $[S]$, coincides with the radical of $[S]$.

Let $I$ be an ideal and $T$ be a multiplicatively stable family of $R$. We denote $I:T$ the ideal of all the elements $p$ of $R$ such that, for some $t \in T$, the element $tp$ belongs to $I$. If the ideal $I$ is differential or radical, then so is $I:T$. If $T \subseteq R$ is any set, then $T^{-1}$ denotes the smallest multiplicatively stable family of $R$ which contains $T$.

We work with differential polynomials in the ring $K[y_1, \ldots, y_n]$. We call the $y_i$’s letters and the $\theta y_i$’s derivatives.

An order $\mathcal{R}$ over the set of the derivatives ($\theta y_i$) is said to be a ranking ((Ko], page 75) if it is total and if it is compatible with the differentiations over the alphabet:

1) $\delta_i \theta y_j > \theta y_j$, for all derivation $\delta_i$, all operator $\theta$ and all letter $y_j$)
2) $\theta y_j > \theta y_i \Rightarrow \delta_i \theta y_j > \delta_i \theta y_i$ (for all derivations $\delta_i$, all operators $\theta_1, \theta_2$ and all letters $y_j, y_i$).

Let $p$ be a polynomial in $K[y_1, \ldots, y_n]$ and $\mathcal{R}$ a ranking on the $\theta y_i$. The leader $u$ of $p$ is the largest derivative with respect to the ranking $\mathcal{R}$ which appears in $p$. The two conditions mentioned above imply that the leader of $\theta y_i$ is $\theta u$ for all derivation operators $\theta$. Let $d$ be the degree of $u$ in $p$. The initial $I_u$ of $p$ is the coefficient of $u^d$ in $p$. The separatant $S_u$ of $p$ is the initial of all the proper derivatives of $p$ ($S_u = \delta y_i \theta u$). The rank of a polynomial $p = I_u \cdot u^d + R_u$ is the rank of $u$. The rank of a set $E$ is the rank of the elements of $E$.

Let $p$ and $q$ be two polynomials and let $u^d$ be the rank of $p$. The polynomial $q$ is said to be partially reduced with respect to $p$ if no proper derivative of $u$ appears in $q$. The polynomial $q$ is said to be reduced with respect to $p$ if $q$ is partially reduced with respect to $p$ and its degree in $u$ is less than $d$.

A set of polynomials $A$ is said to be triangular if its elements have different leaders. A set of polynomials $A$ is said to be autoreduced if each element of $A$ is reduced with respect to every other element of the set. Every autoreduced set is triangular. Every autoreduced set is finite ([Ko], page 77). Let $A$ be an autoreduced set. We denote $H_A$ the set of all the initials and the separatans of $A$. Hence $H^d_A$ denotes the set of all the products of powers of the initials and separatans of the elements of $A$.

Let $p$ and a polynomial and $A = p_1, \ldots, p_n$ be an autoreduced set. There exists ([Ko], page 77) an algorithm, called Ritt’s algorithm of reduction, which rewrites $p$ as a polynomial $r - p \text{ rem } A$, reduced with respect to $A$ (i.e. with respect to all the elements of $A$), satisfying the relation: $r \equiv t_1 p_1 + \ldots + t_n p_n + S \cdot p \pmod {\text{ideal } [A]}$ for some integers $a_i$ and $b_i$, where $t_i$ and $S_i$ denote respectively the initial and the separatant of $p_i$.

The algorithm begins by producing a partial remainder $q = p \text{ partial-rem } A$. The polynomial $q$ is partially reduced with respect to $A$ and satisfies for some integers $b_1, \ldots, b_t$, the relation: $q \equiv \sum b_i S_i + S \cdot p \pmod {\text{ideal } [A]}$. The algorithm then calculates

*The definitions which we give are only valid for polynomials $p \notin K$. In this paper, we don’t need to bother with the exception $p \in K$. 
lates \( r \equiv p \mod A \) by applying to \( q \) a simple algebraic reduction.

If \( p \in [A]; \ H_r^\infty \), then \( (p \mod A) \in [A]; \ H_r^\infty \).

Many such algorithms exist. We fix one of them.

An autoreduced subset \( C \) of a set \( E \) of polynomials is called a characteristic set\(^3\) of \( E \) if \( E \) does not contain any non-zero element reduced with respect to \( C \). All the characteristic sets of \( E \) have the same rank. A characteristic set \( C \) of an ideal \( J \) reduces to zero all elements of \( J \). If the ideal is prime, \( C \) reduces to zero only the elements of \( J \) and we have \( J = [C]; \ H_C^\infty \) ([Ko], lemma 2, page 167).

Let \( p_1 \) and \( p_2 \) be polynomials in an autoreduced subset \( A \) of \( K \{ y_1, \ldots, y_n \} \), whose leaders \( y_1, y_2 \) are \( \text{derivatives of some same letter } y_1 \) (this can only happen for partial differential systems). We denote \( \theta \) the operator of minimal order and \( \phi_1 \) and \( \phi_2 \) the two derivation operators such that \( \phi_1 \theta = \phi_2 \theta = \theta \). We define the \( \Delta \)-polynomial between \( p_1 \) and \( p_2 \) as the polynomial \( \Delta_{p_1, p_2} = S_{\phi_1, \phi_2} - S_{\phi_2, \phi_1} \). Its leader is strictly less than \( \theta y_1 \).

Denote \( \Delta \)-polynomial \( (A) \) the set of all possible \( \Delta \)-polynomials which can be formed between any two elements of \( A \). The set \( A \) is said to be coherent\(^4\) if it reduces to zero all its \( \Delta \)-polynomials: \( \Delta \)-polynomial \( (A) \mod A = \{ 0 \} \) (or \( = \theta \)).

**Gröbner Bases.** We shall have to calculate (non-differential) Gröbner bases of (non-differential) ideals of \( K \{ y_1, \ldots, y_n \} \). Let \( A \) be a finite subset of \( K \{ y_1, \ldots, y_n \} \) and let \( \mathcal{R}_1 \) be a ranking of the derivatives \( \theta y_1 \). We order following \( \mathcal{R}_1 \) the derivatives \( y_1 < \cdots < y_n \) which appear in the elements of \( A \). The order \( \mathcal{R}_2 \) induces an order of elimination \( \mathcal{R}_2 \) on the monomials of the ring \( K \{ w_1, \ldots, w_i \} \). Let \( m_1 = w_1^{a_1} \cdots w_i^{a_i} \) and \( m_2 = w_1^{b_1} \cdots w_i^{b_i} \). The order \( \mathcal{R}_2 \) is defined by: \( m_1 < m_2 \) if for the largest index \( i \) such that \( a_i < b_i \) and \( a_i \) are different, we have \( a_i < b_i \).

The largest monomial for the order \( \mathcal{R}_2 \) which appears in a polynomial \( p \) is called the **head monomial** of \( p \). Also, if \( u^P \) is the rank of a polynomial \( p \) for the order \( \mathcal{R}_1 \), then \( u^P \) appears as a factor in the head monomial of \( p \).

If \( A \) is a subset of \( K \{ w_1, \ldots, w_i \} \subset K \{ y_1, \ldots, y_n \} \), then the non-differential ideal generated by \( A \) in the ring \( K \{ w_1, \ldots, w_i \} \) coincides with the intersection between the non-differential ideal generated by \( A \) in \( K \{ y_1, \ldots, y_n \} \) and the polynomial ring \( K \{ w_1, \ldots, w_i \} \). Thus for the non-differential ideal \( (A) \), the property to be prime of radical is independent of the polynomial ring.

### Theorems Used

**3.1 The theorem of zeros**

Let \( \Sigma \) be a polynomial system of equations and inequalities. A model of \( \Sigma \) is a solution of \( \Sigma \) in a

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**3.2 Regular systems**

A rapid computation shows that \( x^2 \in \langle x \rangle \) but that \( x^3 \notin \langle x \rangle \). More generally, if \( A \) defines a finite set of polynomials, the set of the elements of \( A \) partially reduced w.r.t. \( A \) may also contain polynomials which are not in \( A \). This phenomenon demonstrates well the importance of the following lemma.

**Lemma 1 (Rozenfeld).** If \( A \) is an autoreduced and coherent subset of the ring \( K \{ y_1, \ldots, y_n \} \) then every differential polynomial which belongs to \([A]; H_{\infty}^\infty \) and which is partially reduced with respect to \( A \) belongs also to \((A); H_{\infty}^\infty \).
Proof. See [Ro], page 397 or [Ko], lemma 5, page 135.

The regular systems are differential polynomial systems of equations and inequalities for which Rosenfeld’s lemma applies.

Definition 2. A system of differential equations and inequalities is said to be regular with respect to a ranking $R_1$, if the set of its equations is autoreduced and coherent, the initial and separant of each equation appear among the inequalities and if its other inequalities are partially reduced with respect to the equations:

$$
\begin{align*}
\Omega &= \left\{ \begin{array}{l}
p_1 = 0 \\
\vdots \\
p_s = 0 \\
I_1 \neq 0 \\
\vdots \\
S_s \neq 0 \\
q \neq 0 & q \text{ is partially reduced with respect to } A
\end{array} \right.
\end{align*}
$$

Notation. We use the letter $\Omega$ to denote regular systems (for instance: $\Omega$, $\Omega_1$, etc.). We use the letter $A$ to denote the set of the equations of $\Omega$ (for instance: $A_1$ and $A_2$, stand for the set of the equations of $\Omega_1$ and $\Omega_2$). We use $H_{\Omega_1}$ to denote the set of all the power products of the inequalities of $\Omega$ (for instance, $H_{\Omega_1}^{(3)}$ and $H_{\Omega_2}^{(3)}$ correspond to $\Omega_1$ and $\Omega_2$). We have $H_{\Omega}^{(3)} \subseteq H_{\Omega}^{(3)}$.

Theorem 2 (Rosenfeld). A regular system $\Omega$ admits a differential model if and only if it admits an algebraic model.

Proof. See [Ro], page 398. Suppose that $\Omega$ does not admit a differential model and we show that it then does not admit an algebraic model. By the theorem of zeros, $1 \in [A]; H_{\Omega}^{(3)}$. By Rosenfeld’s lemma, $1 \in (A); H_{\Omega}^{(3)}$. By Rosenfeld’s lemma, $1 \in (A); H_{\Omega}^{(3)}$ and $\Omega$ does not admit an algebraic model.

The other implication is immediate.

The lemma 1 and the theorem 2 are extensions of two results of Seidenberg ([Se1], theorems 6 and 7 pages 51 and 52) which provide his elimination algorithm for partial differential systems.

3.3 Regular ideals

We establish in this section some important properties of the ideals $[A]; H_{\Omega}^{(3)}$ and $(A); H_{\Omega}^{(3)}$. In particular, we show that they are always radical and that there exists an algorithm which decides if a given polynomial belongs to them.

The following lemma is interesting by itself. In particular, it generalizes a result of Kolchin (see [Ko], lemma 13, page 36).

The total ring of fractions of a ring $R$ is obtained by making invertible all the elements of $R$ which do not divide zero. We denote it $Q(R)$.

Lemma 2 (Lazard). Let $A = p_1, \ldots, p_s$ be a triangular set of a polynomial ring $K[w_1, \ldots, w_t]$, for the ranking $w_1 < \cdots < w_t$. Let $w_1 < \cdots < w_t$, be the leaders of the elements of $A$ and $S_{\mathcal{A}}$ denote the set of the separants of the elements of $A$. If the ideal $(A); S_{\mathcal{A}}$ is non-trivial, then the total ring of fractions $Q$ of the ring $K[w_1, \ldots, w_t](A); S_{\mathcal{A}}$ verifies the following properties:

(1) $P_1$: it is isomorphic to a product of fields.

(2) $P_2$: denoting $\mathcal{V}_i$ the image of $w_i$, we have: $\mathcal{V}_i$ satisfies an algebraic relation over $Q(K[w_{i+1}, \ldots, w_{t+1}])$ if and only if $w_i$ is some $w_j$.

Moreover, the properties above remain true if $S_{\mathcal{A}}$ is replaced by any multiplicative family $S$ which contains it, provided that the ideal $(A); S$ is non-trivial.

The following small lemmas are used in the proof.

Let $R$ be a ring. Let $I$ be an ideal and $S$ be a multiplicative family of $R$. Let $X$ be an indecomposable ring homomorphism $R \rightarrow S^{-1}R$, $R \rightarrow R/I$ and $R \rightarrow R[X]$ commute together. Moreover, if $S$ and $S'$ are two multiplicative families, the morphisms $i_v$ and $i_{v'}$ commute also. If $I \subset J$, then $R/J \simeq (R/I)/(p/J)$.

Let $R = R_1 \times \cdots \times R_\lambda$ be a product of rings and if $a = (a_1, \ldots, a_\lambda)$ is one of its elements, we have: $R/(a) = R_1/(a_1) \times \cdots \times R_\lambda/(a_\lambda)$.

$T_\lambda = T_{\lambda-1} \times \cdots \times T_1$.

Let $p \in R[X]$ be a polynomial and $S_{\mathcal{A}}$ its separant.

Then $f(S_{\mathcal{A}}) = S_{\mathcal{A}(p)}$.

If $p \in R[X]$ is a polynomial over a field. Let $p_{\mathcal{A}}^n : p_{\mathcal{A}}^n$ be the decomposition of $p$ into irreducible factors. Since the separant $S_{\mathcal{A}}$ of $p$ contains as factors the multiple factors $a_i \geq 1$ of $p$, the ideal $(p); S_{\mathcal{A}}^n$ is generated by the product of the simple factors of $p$. The ring $K[X]/(p); S_{\mathcal{A}}^n$ is hence either the zero ring (by (c), if $p$ has no simple factors), either a product of fields, according to the Chinese Remainders theorem.

Proof. We define a sequence of rings as follows:

$R_0 = K$,

$R_{i+1} = Q(R_i[w_{i+1}])$ if $w_{i+1} \neq w_i$ for each $w_i$,

$R_{i+1} = Q(R_i(w_i))$ if $w_{i+1} = w_i$,

where $\bar{w}_i$ and $\bar{S}_i$ denote the images of $w_i$ and $S_i$ in the ring $R_i[w_i]$.

To prove the lemma 2 we are going to establish, first that $R_0$ verifies (P1) and (P2), second that $R_0$ is isomorphic to $Q$. Last, we consider the case of multiplicative families which contain $S_{\mathcal{A}}$.

1. We show by induction on $i$ that $R_i$ verifies (P1) and (P2). Clearly, $R_0$ satisfies them. Assume that $R_i \simeq R_1 \times \cdots \times R_{i-1}$ verifies these two properties and let us show that $R_{i+1}$ verifies (P1) and (P2) also.

If $w_i$ is not a leader $w_j$, using (d), $R_{i+1}$ is isomorphic to $\prod_{k=1}^{i} K_{\bar{w}_k}$. It verifies (P1) and (P2).

Now, let us consider the case $w_i = w_j$. Let $\bar{w}_i = \bar{w}_j = \bar{w}_w$ and $\bar{S}_i = (\bar{S}_i, \ldots, \bar{S}_j)$. By (d) we have: $R_i(w_i)/(\bar{w}_w) : \bar{S}_i \simeq \prod_{k=1}^{i} K_{\bar{w}_k}(w_i)/(\bar{w}_w) : \bar{S}_i$.

Let $k < m$. Then
If \( f_{i,j} \in K\alpha \), then \( \tilde{S}_{i,j} = 0 \) and the \( k^{th} \) factor of the product above is the zero ring, by (e) and (c).  
If \( f_{i,j} \notin K\alpha \), then by (e), \( K\alpha[u_j]/(f_{i,j}) : \tilde{S}_{i,j} \) is either the zero ring, either isomorphic to some product of algebraic field extensions of \( K\alpha \).  
Thus \( R_{i+1} \) verifies (P1) and (P2).

2. We show by induction on \( i \) that \( R_i \cong \mathbb{Q} \).  
The main point to check is that the inversion of the non zero divisors commute with the other ring homomorphisms.  
Let us denote \( T_i = K[w_1, \ldots, w_i]/(p_1, \ldots, p_{i-1} ; (S_1, \ldots, S_{i-1})^{\infty}) \), where \( w_{i-1} \leq w_i < u_i \).  
We have \( \mathbb{Q} = Q(T_0) \).  
Clearly, \( R_0 \cong Q(T_0) \).  
Assume that \( R_i \cong Q(T_i) \) and let us prove that \( R_{i+1} \cong Q(T_{i+1}) \).

If \( w_{i+1} = u_i \), then \( R_{i+1} = Q(R_i[w_{i+1}]) \).  
Every non zero divisor in \( T_i \) is still a non zero divisor in \( T_i[w_{i+1}] \), so \( R_{i+1} \cong Q(T_i[w_{i+1}]) \) and by (a), \( R_{i+1} \cong Q(T_i[w_{i+1}]) \).

If \( w_{i+1} \neq u_i \), then \( R_{i+1} = Q(R_i[u_i]) \).  
Since \( R_{i+1} \) verifies (P2), every non zero divisor in \( T_i \) is still a non zero divisor in \( T_i[u_i] \).  
By (a) and (b) \( R_{i+1} \cong Q(T_{i+1}) \).

3. By (c) and (d), the inversion of an element \( p \) of a product of fields only suppresses the fields of the product for which \( p \) has a zero component.

**Definition 3**  
A differential ideal \( J \) is said to be regular if there exists a regular system \( \Omega \) such that \( J = [A] : H_1^\infty \).  
An algebraic ideal \( J \) is said to be regular if there exists a regular system \( \Omega \) such that \( J = (A) : H_1^\infty \).

**Theorem 3**  
Every regular ideal is radical.

**Proof.** Let \( \Omega \) be a regular system.  
Let \( p \) be a polynomial for which a power \( p^m \) belongs to \( (A) : H_1^\infty \).  
The image of \( p^m \) in \( K[w_1, \ldots, w_i] \ldots (A) : H_n^\infty \) is zero.  
That ring has nilpotent element, since its total ring of fractions is a product of fields, according to the lemma 2.  
Hence the image of \( p \) is zero, \( p \) belongs to \( (A) : H_1^\infty \) and that ideal is radical.

Let us show that the regular differential ideal \( [A] : H_1^\infty \) is also radical.  
Let \( p \) be a polynomial for which a power \( p^m \) belongs to \( (A) : H_n^\infty \).  
The polynomial \( f \) is equivalent to some \( S_1 \ldots S_n \) modulo \( [A] : H_1^\infty \).  
By Rosenfeld’s lemma and the fact that \( H_1^\infty \) is the ideal of the ring \( \tilde{S} \), we have \( f \in [A] : H_1^\infty \).  
This ideal is thus radical.

The following lemma is a consequence of lemma 2 (property (P2)), the proof of which is left to the reader.  
It is used in the proof of the lemma 5 and shows that we may read the transcendance degree of a system without calculating the Gröbner basis of \( (A) : H_1^\infty \) except the condition to ascertain that the ideal is non trivial.

**Lemma 3**  
Let \( \Omega \) be a regular system for a ranking \( R_\lambda \).  
Let \( u_1 \ldots u_i \) be the leaders of the equations of the system.  
Let \( B \) be a Gröbner basis of \( (A) : H_1^\infty \) for the order \( R_\lambda \) induced by \( R_\lambda \).  
If \( (A) : H_1^\infty \) is not the unit ideal, then the leaders of the polynomials of \( B \) are the derivatives \( u_1, \ldots, u_i \).

Let \( \Omega \) be a regular system and \( A \) be the set of its equations.  
We give in section 4 a method to calculate a Gröbner basis \( B \) of \( (A) : H_1^\infty \), and in section 6 an example of a regular system without models.

The following lemma shows how to decide the membership problem in a regular differential ideal.  
Its proof is an easy consequence of Rosenfeld’s lemma.

**Lemma 4**  
Let \( \Omega \) be a regular system, \( A \) be the set of its equations, and \( B \) be a Gröbner basis of \( (A) : H_1^\infty \).  
For each polynomial \( p \) of \( K(y_1, \ldots, y_m) \), we have:

\[
p \in [A] : H_1^\infty \iff (p \text{ partial-rem } A) \in (B).
\]

We would like to clarify the correspondence between systems of regular algebraic ideals and regular differential ideals.  
An example suffices to show that two different regular systems may define the same regular ideals:

\[
x + 1 = 0, \quad (x + 1)(x + 2) = 0, \quad (x + 2)(3x + 4) = 0.
\]

**Question:** Is the correspondence between regular algebraic ideals and regular differential ideals bijective?  
In other words, do two regular differential systems define the same regular algebraic ideal if and only if they define the same regular differential ideal?  
The following lemma, which shows the implication from right to left, is a step in the proof of theorem 6.  
The converse implication, which we have not established, seems to be in keeping with the open problem: to decide the inclusion of two prime differential ideals each given by a characteristic set (see [Köl], page 166).

**Lemma 5**  
Two regular systems which define the same regular differential ideal define also the same regular algebraic ideal.

**Proof.** Let \( \Omega \) and \( \Omega' \) be two regular systems defining the same regular differential ideal \( [A] : H_1^\infty = [A'] : H_1^\infty \).  
Let \( B \) and \( B' \) be the Gröbner bases respectively of the ideals \( (A) : H_1^\infty \) and \( (A') : H_1^\infty \) for the order \( R_\lambda \) induced by \( R_\lambda \).  
We suppose \( B \) is different from \( B' \) and we seek a contradiction.

We order the polynomials of \( B \) by \( b_0, b_1, \ldots, b_m \) and of \( B' \) by \( b'_0, b'_1, \ldots, b'_n \), by increasing order.  
Let \( i \) be the least index such that the head monomials of the polynomials \( b_i \) and \( b'_i \) are different and suppose \( b_i < b'_i \).  
Since \( b'_i \) belongs to the differential ideal \( [A] : H_1^\infty \), by lemma 4 \( (b'_i \text{ partial-rem } A) \in (B) \).

Let \( u_1 \) and \( u'_1 \) be the leaders of the polynomials in the basis \( B \) and \( B' \) and let \( u_i \) and \( u'_i \) be the leader of \( b_i \).

We have \( u_1 - u'_1, \ldots, u_{i-1} - u'_{i-1} \).

By the lemma 3, each polynomial of the basis \( B' \) (respectively \( B \)) is partially reduced w.r.t. each other.  
Since \( u_1 - u'_1, \ldots, u_{i-1} - u'_{i-1} \), and since \( b'_i < b_i \), the partial reduction of \( b'_i \) by \( A \) does not modify \( b'_i \) and we have \( b'_i \in (B) \).  
In view of the hypothesis made on \( i \), the head monomial of \( b'_i \) can not be reduced by any rule from \( B \).

This contradiction proves the lemma.
4 The Rosenfeld–Gröbner Algorithm

The program Rosenfeld–Gröbner gathers at entry a differential system of equations and inequations \( \Sigma \) and a ranking \( \mathcal{R}_1 \). It produces by splittings a finite family \( (\Omega_i) \) of consistent (with models) regular systems whose differential models form a partition of the differential models of \( \Sigma \).

The Greek letters \( \Lambda, \Omega, \Gamma_i \), denote systems of equations and inequations. \( \Lambda_{eq} \) and \( \Lambda_{in} \) stand respectively for the set of the equations and for the set of the inequations of the system \( \Lambda \).

The function \texttt{obviouslyInconsistent} returns \texttt{true} if a non-zero element of \( K \) appears among the equations, or if \( \omega \) appears among the inequations of the system.

```plaintext
program Rosenfeld-Gröbner (\Lambda, \mathcal{R}_1) begin
  if not obviously inconsistent (\Lambda) then
    A := a characteristic set of the finite set \( \Lambda_{eq} \)
    Let \( \{h_1, \ldots, h_n\} \) denote the set of the initials
    and of the separants of the elements of \( A \)
    such that \( h_i \notin K \).
    \( R := (\Lambda_{eq} \setminus A \cup \Delta\text{-polys}(A)) \rem A \)
    if \( R = 0 \) or \( R = \emptyset \) then
      \( \Omega_{eq} := A \)
      \( \Omega_{in} := (\Lambda_{in} \setminus \Delta\text{-polys}(A)) \}
    \( B := \text{a Gröbner basis of } (A) : H_{\Omega_{eq}}^\infty \)
    if \( B \neq \{1\} \) then
      produce on output \( \Omega \) and \( B \)
    endif
  else
    \( \Gamma_{i+1,eq} := A \cup R \)
    \( \Gamma_{i+1,in} := \Lambda_{in} \cup \{h_1 \neq 0, \ldots, h_n \neq 0\} \)
    Rosenfeld-Gröbner (\( \Gamma_{i+1,eq}, \mathcal{R}_1 \))
  endif
for i := r down to 1 do
  \( \Gamma_{i,eq} := \Lambda_{eq} \cup \{h_i = 0\} \)
  \( \Gamma_{i,in} := \Lambda_{in} \cup \{h_{i-1} \neq 0, \ldots, h_1 \neq 0\} \)
  Rosenfeld-Gröbner (\( \Gamma_{i,eq}, \mathcal{R}_1 \))
end
end
```

Some other ways exist to do the splitting of \( \Lambda \) into the \( \Gamma_i \) (see [Bo]). This one was used by Seidenberg in [Sc].

The Gröbner basis \( B \) of the ideal \( (A) : H_{\Omega_{eq}}^\infty \) is computed by the method below. It is classical [Tr]; It detects regular systems without models: those with basis \{1\}.

1. The system \( \Omega \) is transformed into a system of equations: the algorithm introduces a new unknown \( z_i \) for each inequation \( h_i \neq 0 \) of the system and rewrites \( h_i \neq 0 \) as \( z_i = 1 \).
2. A basis \( B_k \) is computed following any elimination order \( \mathcal{R}_2 \) satisfying \( \theta_i < z_i \) (for all derivatives \( \theta_i \), and all unknowns \( z_i \)).
3. The desired basis \( B \) is obtained by truncating \( B_k \). Only those polynomials of \( B_k \) which do not involve \( z_i \) are retained.

4.1 Proofs

Lemma 6 The Rosenfeld–Gröbner algorithm stops.

Proof. The set of the equations of each system \( \Gamma_i \), not obviously inconsistent, produced from \( \Lambda \), contains \( A \) and at least one polynomial \( p \notin K \) reduced w.r.t. \( A \).

Thus, the characteristic sets of the sets of the equations of the systems \( \Gamma_i \), not obviously inconsistent are lower than \( A \), for the usual ranking over autoreduced sets ([Ko], page 81).

This ranking is a well ordering ([Ko], proposition 3, page 81). Since the algorithm discards obviously inconsistent systems, Rosenfeld-Gröbner stops.

The two lemmas below deal with the correction of the algorithm. Since \( \Omega \) corresponds to a particular case of \( \Gamma_{i+1} \), we do not distinguish it from \( \Gamma_{i+1} \), in order to simplify the statements.

Lemma 7 \( \phi \) is a differential model of \( \Lambda \) if and only if \( \phi \) is a differential model of some \( \Gamma_i \) \( (1 \leq i \leq r+1) \).

Moreover, the differential models of the systems \( \Gamma_i \) are disjoint.

We only give the main argument of the proof.

Let \( \phi \) be a differential model of some system \( A - 0, h_1 \neq 0, \ldots, h_n \neq 0 \). Let \( p \) be any polynomial and let \( \overline{p} = p \rem A \). According to the definition of the models and to the specifications of Ritt’s algorithm of reduction, we have \( \phi(p) = 0 \iff \phi(\overline{p}) = 0 \).

We need the notations below for the lemma 8, which is used for the calculus of characteristic sets in section 5.

Let \( \Sigma \) be a system of equations and inequations, \( \{\Omega_{eq} \} \) be the radical differential ideal generated by the equations of the system, and \( H_{\Sigma}^\infty \) be the multiplicative family generated by its inequations. We denote \( J(\Sigma) \) the radical differential ideal \( \{\Omega_{eq} \} : H_{\Sigma}^\infty \).

Lemma 8 If \( J(\Lambda) \) is prime and if \( i \) is the greatest index such that \( \Gamma_i \) has a differential model, then

\[ J(\Lambda) = J(\Gamma_i) \]

Proof. According to the lemma 7 above and to the theorem of zeros,

\[ J(\Lambda) = J(\Gamma_{i+1}) \cap \cdots \cap J(\Gamma_1), \]

so the index \( \ell \) exists. We consider thus two cases.

1. No polynomial \( h_i \) \((1 \leq i \leq r) \) belongs to \( J(\Lambda) \). We prove that \( J(\Gamma_{i+1}) \subset J(\Lambda) \) and the equality follows from the formula above \((\ell - r + 1)\).

We have \( \Gamma_{i+1,eq} \subset [\Lambda_{eq} \subset J(\Lambda) \). Since the ideal is prime and since no polynomial \( h_i \) belongs to it, we have \( H_{\Gamma_{i+1,eq}}^\infty \cap J(\Lambda) = \emptyset \). Now, assume that \( p \in J(\Gamma_{i+1}) \) i.e. that for some \( h \in H_{\Gamma_{i+1,eq}}^\infty \) we have \( hp \in \{\Gamma_{i+1,eq} \} \subset J(\Lambda) \). Since \( h \notin J(\Lambda) \) and that ideal is prime, we have \( p \notin J(\Lambda) \).

2. Let \( t \leq r \) be the smallest index such that \( h_t \in J(\Lambda) \). We prove that all the ideals \( J(\Gamma_t) \) \((t \leq i \leq r + 1) \) are trivial and that \( J(\Gamma_t) \subset J(\Lambda) \). The equality follows from the formula above \((\ell - t)\).

By the formula above, for \( t \leq r \leq r + 1 \) we have \( h_t \in J(\Lambda) \subset J(\Gamma_t) \) but, according to the way the \( \Gamma_i \) are computed, we have also \( h_t \in H_{\Gamma_{i+1,eq}}^\infty \). These ideals \( J(\Gamma_t) \) are hence trivial and the corresponding systems \( \Omega \), are not produced on the output of the program.
According to the hypothesis $\Gamma_{\text{eq}} \subseteq \mathcal{J}(\Lambda)$. Since this ideal is prime and no polynomial $h_i (1 \leq i < t)$ belongs to it, we have $H_{\Pi_0} \cap \mathcal{J}(\Lambda) = \emptyset$ whence as in 1. above, $\mathcal{J}(\Gamma_1) = \mathcal{J}(\Lambda)$.

4.2 Properties of the computed representation

A basis in the sense of Ritt and Raudenbush of a radical differential ideal $\mathcal{J}$ is any finite family $\Sigma$ such that $\mathcal{J} = \{ \Sigma \}$, Ritt and Raudenbush established [R3], page 10 that every radical differential ideal admitted a basis.

The Rosenfeld-Gröbner algorithm decomposes a differential ideal $\{ \Sigma \}$ given by a finite basis as an intersection of regular differential ideals each described by a regular system. This decomposition is also an algorithm for membership testing in $\{ \Sigma \}$.

Consider a system $\Sigma : p_1 = 0, \ldots, p_m = 0$ of differential polynomial equations of $K[y_1, \ldots, y_n]$. Let $\Omega_1, \ldots, \Omega_\ell$ be the successive regular systems produced by the Rosenfeld-Gröbner algorithm applied to $\Sigma$ for some ranking $R_\ell$.

For each system $\Omega_i$, we denote $A_i$ the set of its equations and $H_{\Omega_i}^\infty$ the multiplicative family generated by its inequations.

**Theorem 4** With notations as above, we have:

1) $\phi$ is a differential model of $\Sigma$ if and only if $\phi$ is a differential model of some $\Omega_i (1 \leq i \leq \ell)$. Moreover, the differential models of the regular systems $\Omega_i$ are disjoint.

2) the radical differential ideal $\{ \Sigma \}$ is the intersection of the regular differential ideals $[A_i] : H_{\Omega_i}^\infty$.

**Proof.**

1) It is an easy consequence of the lemma 7.

2) By 1) and the theorem of zeros, a polynomial $p$ belongs to $\{ \Sigma \}$ if and only if, for each $i \in [1, \ell]$, the system obtained by adjoining the equation $p \neq 0$ to $\Omega_i$ has no differential models. By the theorem of zeros and the theorem 3, these systems have no differential models if and only if $p \in [A_i] : H_{\Omega_i}^\infty$. 

The description of the ideal $\{ \Sigma \}$ computed by Rosenfeld-Gröbner allows us to decide the membership problem in $\{ \Sigma \}$, using a few reductions. This is expressed in the following theorem, whose proof is an immediate consequence of the theorem 4 and the lemma 4.

**Theorem 5** With notations as above, we have:

$p \in \{ \Sigma \} \iff \forall i \in [1, \ell], (p \text{ partial-rem } A_i) \in (B_i)$.

5 Computation of characteristic sets

We give a method to compute the characteristic set of a prime differential ideal given by a basis in the sense of Ritt and Raudenbush. We generalize here the result [O], page 89, of Ollier.

We retain the notations of the preceding section.

**Lemma 9** If the differential ideal $\{ \Sigma \}$ is prime then

$\{ \Sigma \} = [A_i] : H_{\Omega_i}^\infty$.

**Proof.** The inclusion from left to right comes from the theorem 4. The other one is a consequence of the lemma 9. 

To our knowledge, there does not exist any algorithm which decides if a differential ideal given by a basis (either in the classical sense or in the sense of Ritt and Raudenbush) is prime.

The coherent and autoreduced set $A_1$ satisfies a property of characteristic sets of the ideal: if $C$ is a characteristic set of a prime differential ideal $\{ \Sigma \}$, then we have $\{ \Sigma \} = C : H_{\Omega_1}^\infty$. However, $A_1$ is not necessarily a characteristic set of the ideal. Consider the (algebraic) example below:

$A_1 : (x + 1)(x + 2) - 0, (x + 1)y + 2 - 0$.

$A_1$ is autoreduced with respect to the order $x < y$, the ideal $(A_1) : H_{\Omega_1}^\infty$ is prime but its characteristic set is $C : x + 2 = 0, y - 2 = 0$.

The basis $B$ of $(A_1) : H_{\Omega_1}^\infty$, computed with respect to the order $R_2$ induced by $R_1$ is almost a characteristic set of $\{ \Sigma \}$, but not quite. We give in the following section an example which shows that this is not necessarily the case.

Theorem 6 Let $\{ \Sigma \}$ be a prime differential ideal and $\Omega$ be a regular system with respect to a ranking $R$, such that $\{ \Sigma \} = [A] : H_{\Omega_1}^\infty$. Let $B$ be a Gröbner basis of $(A) : H_{\Omega_1}^\infty$ computed with respect to the order $R_2$ induced by $R_1$.

The following algorithm calculates a characteristic set $C$ of the ideal $\{ \Sigma \}$, with respect to the ranking $R_1$, from the basis $B$.

begin Assume that the elements of $B = b_1 < \cdots < b_m$ are arranged in increasing order.

$C := \{ b_1 \}$

for $i := 2, \ldots, m$ do

let $u_i$ and $u_{i - 1}$ be the leaders of $b_i$ and $b_{i - 1}$

if $u_i \neq u_{i - 1}$ then

$C := C \cup \{ b_i \text{ rem } C \}$

end

end

**Proof.** We are going to successively establish the following points:

1. To determine $C$ amounts to determining a characteristic set of the prime ideal $(C) : H_{\Omega_1}^\infty$ with respect to the order $R_1$.

2. $B$ is a Gröbner basis of $(C) : H_{\Omega_1}^\infty$.

3. Let $p - I_1 : u_1 b_1 + R_0$ be a polynomial of $(B) - (C) : H_{\Omega_1}^\infty$, whose initial $I_1$ is not in the ideal. There exists then in $B$ a polynomial $b - I_1 : u_1 b_1 + R_0$ with $d_0 \leq d_p$ and there exists in $C$ a polynomial $c - I_1 : u_1 c + R_0$ with $d_c \leq d_p$.

Since neither $I_1$, nor $I_0$, appear in $(B) - (C) : H_{\Omega_1}^\infty$, and since $B$ and $C$ are two subsets of the ideal, the algorithm described in the theorem extracts from $B$ a set of polynomials of the same rank as $C$, but which is not necessarily autoreduced in the sense of Ritt. The proof of the theorem is completed by:
4. The reductions carried out by the algorithm may not reduce the rank of the polynomials extracted from the basis. □

Proof of 1. See [R]. The characteristic set of a differential ideal, antireduced by definition, is coherent since it reduces to zero every polynomial (in particular the $\Delta$-polynomials) of the ideal. We apply the lemma of Rosenfeld. \((C): H_{\infty}^{C}\) is the intersection of the prime differential ideal \([C]: H_{\infty}^{C}\) and the ring of partially reduced polynomials with respect to \(C\). The ideal \((C): H_{\infty}^{C}\) is then prime.

To say that a coherent and antireduced set \(C\) is not a characteristic set of \([C]: H_{\infty}^{C}\), is to say that there exists in that ideal a non-zero polynomial \(p\), reduced with respect to \(C\). By the lemma of Rosenfeld, this is to say that \(p\) belongs to \((C): H_{\infty}^{C}\) and hence \(C\) is not a characteristic set of \([C]: H_{\infty}^{C}\). □

Proof of 2. Since \([\Sigma] = [C]: H_{\infty}^{C}\) (see [Ko], lemma 2, page 167) and \([\Sigma] = [\Lambda]: H_{\infty}^{\Lambda}\), by lemma 5, we have \((B) = (C): H_{\infty}^{C}\). □

Proof of 3. Let \(p = I_{p} \cdot u^{d_{p}} + R_{p}\) be a polynomial of \((B) = (C): H_{\infty}^{C}\), whose initial \(I_{p}\) is not in the ideal. Suppose the head monomial \(m_{i}\) of \(I_{p}\) under normal form modulo \(B\). Since \(p\) is reduced to zero by \(\Lambda\), there exists a polynomial \(b\) of \(B\) whose head monomial divides the head monomial \(m_{i} - m_{i}^{d_{p}}\) of \(p\), but does not divide \(m_{i}\). The rank of \(b\) is then \(u_{b}^{k}\) with \(0 < d_{b} \leq d_{p}\).

Since \((B) = (C): H_{\infty}^{C}\) is prime and \(I_{p}\) does not belong to the ideal, \(C\) does not reduce \(I_{p}\). There exists then in the characteristic set a polynomial \(c = I_{c} \cdot u^{d_{c}} + R_{c}\) with \(0 < d_{c} \leq d_{p}\). □

Proof of 4. This is immediate since the initials of the polynomials of \(B\) do not belong to \((B) = (C): H_{\infty}^{C}\), the ideal is prime and the characteristic set of a prime ideal reduces to zero only the elements of the ideal. □

6 Examples

The algorithms described in the preceding sections have been programmed (see [Bo], VI) in the language C. The manipulations of big numbers are effected by the library of PARI. The Gröbner bases computations are by the software GB (see [FGLM]).

The computations are done on the IBM RS/6000 station columbia.polytechnique.fr. The timing of computation are given by the UNIX command time.

6.1 Membership testing

We first give a very simple example to illustrate splittings and to show how to test membership in radical differential ideals. We deal with

$$\Sigma \left\{ \begin{array}{l} (2x + 1)y + y = 0 \\ x^{2} + x - 0. \end{array} \right.$$ 

For the ranking \(\theta x < \phi y\) (for all derivation operators \(\theta\) and \(\phi\)), the first equation may be reduced by the second one. Its remainder is \(xy\). \(\Sigma\) is thus split into two systems:

$$\Gamma_{2} \left\{ \begin{array}{l} xy = 0 \\ x^{2} + x = 0 \end{array} \right. \quad \text{and} \quad \Gamma_{1} \left\{ \begin{array}{l} (2x + 1)y + y = 0 \\ x \neq 0 \end{array} \right. \quad \text{and} \quad \Gamma_{1} \left\{ \begin{array}{l} x^{2} + x = 0 \\ x = 0. \end{array} \right.$$

The system \(\Gamma_{2}\) gives immediately a regular system \(\Omega_{2} = \Gamma_{2}\). The Gröbner basis computation only simplifies the factor \(x\) in the first equation: \(B_{1} = \{y, x^{2} + x\}\). The system \(\Gamma_{1}\) leads with no splittings to the regular system below. The Gröbner basis computation is useless.

$$\Omega_{1} \left\{ \begin{array}{l} y + y = 0 \\ x = 0 \end{array} \right.$$

Now, we may verify that \((2x + 1) \notin \{\Sigma\}\). We apply the theorem 5: \((2x + 1)\) is reduced to zero by \((x^{2} + x)\) in \(\Omega_{1}\), but reduced to \(1\) in \(\Omega_{2}\).

We may also verify that \(x(2x + 1)\), which is the first derivative of \((x^{2} + x)\), belongs to \(\{\Sigma\}\). This polynomial is reduced to zero by both systems \(\Omega_{1}\) and \(\Omega_{1}\).

6.2 Hidden algebraic contradictions

The following example does not have any physical significance. It shows the necessity to assure that the regular systems have models.

Consider the following system in \(\mathbb{Q}[u, v]\) equipped with three derivations \(\delta_{x}, \delta_{y}\), and \(\delta_{v}\). The derivation operators are denoted by subscripts.

$$u_{x}^{2} - 2u_{x}u_{y} - 1, \quad u_{x}v_{x} - v_{x}, \quad v_{y} = u_{x}v_{x}, \quad v_{y} - u_{x}v_{x}, \quad u_{x}^{2} - u_{y}v_{x}.$$ 

The ranking \(\mathcal{R}_{1}\) used for the calculations is the following:

1) \(\theta_{u} > \phi v\) for all derivation operators \(\theta, \phi\).

2) \(\theta_{u} > \phi v\) if \(\theta > \phi\) for the lexicographic order given by \(\delta_{x} > \delta_{y} > \delta_{v}\) (same choice for \(v\)).

The Rosenfeld–Gröbner algorithm computes two regular systems in a little more than 6 seconds:

$$\left\{ \begin{array}{l} v_{x}v_{y} - v_{x}v_{y} = 0, \\ (v_{x}^{2}v_{x}v_{y} + (v_{x}^{2}v_{x}v_{y})v_{x} - v_{x}^{2}v_{y})v_{x} - v_{x}^{2}v_{y} \\ + (v_{x}^{2}v_{x}v_{y} - v_{x}^{2}v_{y})v_{x} = 0, \end{array} \right. \quad \left\{ \begin{array}{l} u_{x}^{2} - 1 = 0, \quad u_{x}v_{y} - v_{x} = 0, \\ (v_{x}^{2}v_{x}v_{y} + (v_{x}^{2}v_{x}v_{y})v_{x} - v_{x}^{2}v_{y})v_{x} - v_{x}^{2}v_{y} \\ + (v_{x}^{2}v_{x}v_{y} - v_{x}^{2}v_{y})v_{x} = 0, \end{array} \right. \quad \left\{ \begin{array}{l} v_{x} \neq 0, \quad v_{x}v_{y}v_{x} - v_{x}^{2}v_{y}v_{x} + v_{x}^{2}v_{y} = 0, \quad v_{y} \neq 0, \quad v_{x}^{2}v_{y} - v_{x}^{2}v_{y} = 0, \end{array} \right. \quad \left\{ \begin{array}{l} v_{x}^{2}v_{y}v_{x} + (v_{x}^{2}v_{x}v_{y})v_{x} - v_{x}^{2}v_{y} \\ + (v_{x}^{2}v_{x}v_{y} - v_{x}^{2}v_{y})v_{x} = 0. \end{array} \right.$$ 

The first is inconsistent. Remark that it is not detected before the Gröbner bases computation, although our implementation of the splitting process looks for simple contradictions: the final algebraic
treatment is necessary. Here is the Gröbner basis associated with the second system:

\[ v, \quad u_x^2 - 1, \quad u_y, \quad u_x u_y - 1. \]

6.3 Computing a characteristic set

The example below, which has no physical significance, shows the necessity to proceed with the described reductions in theorem 6 to obtain a characteristic set of a differential ideal that (we know) is prime.

Let \( \Sigma \) be the following system of ordinary differential equations:

\[
\begin{align*}
    \dot{x} &= -y^2 + y + 1 \\
    \dot{y} &= -2xy + 2y^2x + y + 2x \\
    \dot{z} &= y
\end{align*}
\]

The differential ideal \( \Sigma \) is prime, since the system is orthonomic (\( H^n_T \sim \{1\} \)) and autoreduced\(^5\) with the ranking:

1. \( \theta y > \phi y \) and \( \theta z > \psi x \) for all derivation operators \( \theta, \phi \) and \( \psi \);
2. \( \phi y > \psi x \) if the order of \( \phi \) is larger or equal to that of \( \psi \);
3. \( \psi x > \phi y \) if the order of \( \psi \) is strictly larger than that of \( \phi \).

We now apply the Rosenfeld-Gröbner algorithm to \( \Sigma \) with the elimination order:

1. \( \theta y > \phi y > \psi x \) for all operators \( \theta, \phi \) and \( \psi \).

We obtain (in a bit more than 3 seconds) a unique regular system \( \Omega \). The Gröbner basis \( B_1 \) that is associated to it is:

\[
\begin{align*}
p_1 &= (\dot{x} - x^2 + 1)x(3) - x^2 + 3x^2 - \dot{x} - 2x^4 + 12x^2 + 2x - \dot{x} - 6x^2 + 2x^2 + (-2x - 2x^4 + 2x^5 - 2x^4 - 2x^4 - 2x^4 - 2x^4) \\
p_2 &= (\dot{x} - x^2 + 1)y - x^{(3)} + (x + 2x)\dot{x} - x \\
p_3 &= (\dot{x} - x^2 + 1)y - x^{(3)} + (x + 2x)\dot{x} - x \\
p_4 &= (\dot{x} - x^2 + 1)y - x^{(3)} + (x + 2x)\dot{x} - x
\end{align*}
\]

If we extract a characteristic set from \( B_1 \), without effecting the described reduction in theorem 6, we obtain an autoreduced set \( A = p_1, p_2 \), which is not a characteristic set of the ideal \( \Sigma \) (the polynomial \( p_3 \) in fact does not leave since it is not reduced with respect to \( p_2 \)). Apply the theorem 6 by reducing \( p_3 \) by \( p_2 \). We obtain then a characteristic set \( C \) of \( \Sigma \):

\[
p_1, \quad p_2, \quad (\dot{x} - x^2 + 1)z - x^{(3)} + (x + 2x)\dot{x} - x.
\]

Remark. In practice, we may often do without Rosenfeld-Gröbner for generating the system \( \Omega \); the majority of the time, we know that an ideal \( \Sigma \) is prime by showing a ranking \( R \) for which \( \Sigma \) is orthonomic, autoreduced and coherent (cf. the example above). The ranking \( R \) furnishes a characteristic set, hence a membership testing algorithm of the ideal \( \Sigma \), which permits avoiding the splittings.

7 Conclusion

Although the models of the regular systems produced by the Rosenfeld-Gröbner algorithm are disjoint, the regular differential ideals which are defined by them may be redundant. In particular, the algorithm may produce many systems, even when the differential ideal \( \Sigma \) is prime.

We do not know how to decide the inclusion of two regular differential ideals. It is a problem very close to the problem of Ritt; to decide the inclusion of two prime differential ideals each given by a characteristic set, which “seems very far from solution” (Kr, page 166). Its solution would allow us to decide if a differential ideal given by a finite family of generators is prime.

References


\(^5\) It is necessarily coherent, since it only involves ordinary differential equations.