

# On polynomials attached to curves

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# Plan

I. Introduction.

II. Warming up: the genus 0 case.

III. The genus 1 case: elliptic curves.

IV. A glimpse at genus 2.

V. Conclusions.

# I. Introduction and motivations

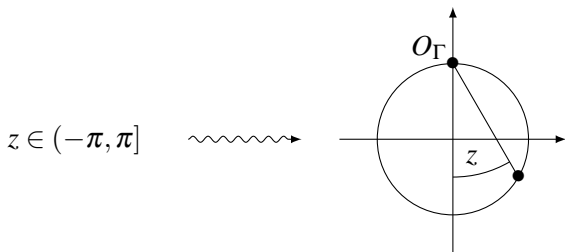
- ▶ Construct mathematical objects: cyclotomic fields, abelian extensions.
- ▶ Point counting on algebraic curves over finite fields. (Schoof-Elkies-Atkin algorithm, Pila, etc.).

In this talk: division polynomials, modular polynomials in genus  $\leq 2$ .

## **II. Warming up: the genus 0 case**

# Circles

$$\Gamma : y^2 = 1 - x^2$$



We have a map:

$$\begin{aligned} \wp : \mathbb{C}/(2\pi\mathbb{Z}) &\rightarrow \Gamma \\ z &\mapsto (\sin z, \cos z). \end{aligned}$$

**Group law:** inherited from  $z_1 + z_2$  on the left and the neutral element is  $\wp(0) = (0, 1)$ . Formulas are polynomials.

## Division polynomials

$$\begin{aligned}\wp([n](\sin z, \cos z)) &= \wp(\sin(nz), \cos(nz)) \\ &= (\mathcal{S}_n(\sin z, \cos z), \mathcal{C}_n(\sin z, \cos z)).\end{aligned}$$

**Computing  $\mathcal{S}_n, \mathcal{C}_n$ :**  $O(\log n)$  polynomial operations using

$$\mathcal{S}_n(z) + i\mathcal{C}_n(z) = \exp(inz) = (\mathcal{S}_1(z) + i\mathcal{C}_1(z))^n.$$

An  **$n$ -torsion** point is  $\sin(nz) = 0, \cos(nz) = 1$ , that is  $z = \exp(2ik\pi/n)$  for  $k \in \{0, \dots, n-1\}$ . Equivalently,  $z$  is a root of  $X^n - 1$ .

$C_d(X)$  = minimal polynomial of  $z$  of order **exactly**  $d$ :

$$X^n - 1 = \prod_{d|n} C_d(X).$$

## Some of the magical properties of $C_n(X)$

**Fact:**  $C_d(X) \in \mathbb{Z}[X]$  is reciprocal.

$$C_\ell(X) = (X^\ell - 1)/(X - 1) = X^{\ell-1} + X^{\ell-2} + \dots + X + 1;$$

$$C_{\ell^e}(X) = C_\ell(X^{\ell^{e-1}}).$$

For  $\gcd(n_1, n_2) = 1$ :

$$C_{n_1 n_2}(X) = \frac{C_{n_1}(X^{n_2})}{C_{n_2}(X)}.$$

**Möbius:**

$$C_n(X) = \prod_{1 \leq k \leq n, k|n} (X^k - 1)^{\mu(k/n)}.$$

**Algorithms:**

- ▶ euclidean division (exact, but possible overflow);
- ▶ Brent (power sums + Newton, *Math. Comp.*);
- ▶ Monagan/Arnold (recent): Möbius + division of sparse series.

# Height

$$C_n(X) = \sum_{k=0}^{\varphi(n)} a_n(k)x^k.$$

$$A(n) = \max_k |a_n(k)|, \quad \beta(n) = \sum_k |a_n(k)|$$

**Thm.** (Bateman, 1949)  $\beta(n) \leq n^{d(n)/2}$ .

**Thm.** (J.-L. Nicolas, G. Terjanian, 1999) for  $n \geq 7$ ,  $n \neq 10$ ,

$$\beta(n) < (\sqrt{2})^{\varphi(n)}.$$

**Thm.** (S. Konyagin, H. Maier, E. Wirsing, 2004) For fixed  $C > 2/\log 2$  and  $\mathcal{E}_C = \{n, \omega(n) \geq C \log \log n\}$ , then for all  $\varepsilon > 0$ , there exists  $n_\varepsilon$  s.t.  $n \geq n_\varepsilon$ ,

$$A(n) \geq \exp\left((\log n)^{C \log 2/2 - \varepsilon}\right).$$

**Monagan:**

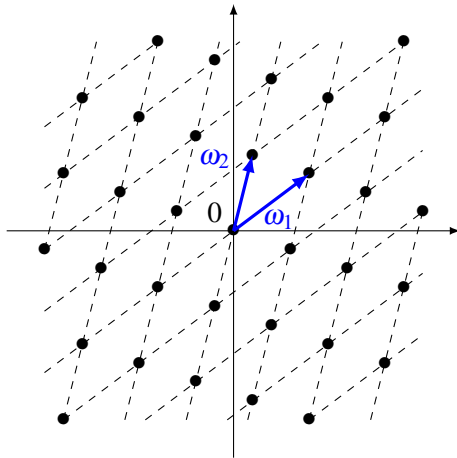
$$A(1880394945) = 64540997036010911566826446181523888971563.$$



### **III. The genus 1 case: elliptic curves**

## A) Lattices and the classical theory over $\mathbb{C}$

$$\mathcal{L} = \mathbb{Z}\omega_1 + \mathbb{Z}\omega_2, \quad \tau = \omega_2/\omega_1 \in \mathcal{H} \{z \in \mathbb{C}, \Im(z) > 0\}$$



$\Rightarrow$  What are the periodic functions over  $\mathcal{L}$ ?

# Weierstrass's function

**Def.**  $f$  is an **elliptic function** iff

- ▶  $f$  is doubly periodic:  $f(z + \omega_i) = f(z)$ ;
- ▶  $f$  is analytic (except at poles), with no finite singularities (except at poles)

$$\wp(z) = \frac{1}{z^2} + \sum_{\omega \in \mathcal{L}, \omega \neq 0} \left( \frac{1}{(z - \omega)^2} - \frac{1}{\omega^2} \right)$$

**Thm.**  $\wp$  is differentiable and:

$$\wp'(z) = -2 \sum_{\omega \in \mathcal{L}} \frac{1}{(z - \omega)^3}.$$

**Prop.**  $\wp'$  and  $\wp$  are periodic on  $\mathcal{L}$ .

## Expansion of $\wp$ near the origin

$$\frac{1}{(z-\omega)^2} = \frac{1}{\omega^2(1-\frac{z}{\omega})^2} = \frac{1}{\omega^2} + \frac{2z}{\omega^3} + \cdots + \frac{kz^{k-1}}{\omega^{k+1}} + \cdots$$

Eisenstein series ( $k \geq 2$ ):

$$G_k(\mathcal{L}) = \sum_{\omega \in \mathcal{L}, \omega \neq 0} \frac{1}{\omega^k}$$

$$\wp(z) = \frac{1}{z^2} + 3z^2 G_4 + 5z^4 G_6 + \cdots$$

**Rem.** Fast expansion of  $\wp$  in BoMoSaSc08.

## Link with elliptic curves

$$g_2 = 60G_4, \quad g_3 = 140G_6$$

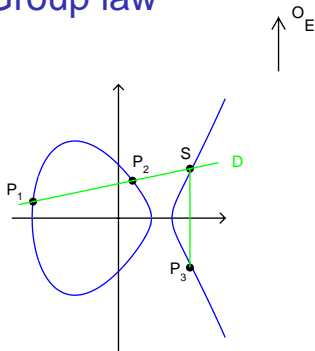
$$\forall z \in \mathbb{C} - \mathcal{L}, \wp^2(z) = 4\wp(z)^3 - g_2\wp(z) - g_3$$

We get a parametrization

$$\begin{array}{ccc} \mathbb{C} - \mathcal{L} & \rightarrow & E \\ z & \mapsto & (\wp(z), \wp'(z)) \end{array}$$

(and we send  $\mathcal{L}$  to  $O_E$ .)

# Group law



On  $E : y^2 = x^3 + ax + b$ , to find

$$(x_3, y_3) = (x_1, y_1) \oplus (x_2, y_2),$$

relate  $\wp(z_1 + z_2)$  to  $\wp(z_1)$  and  $\wp(z_2)$ .

$$P_1 P_2 : y = \lambda x + \mu, \quad x_1 + x_2 + x_3 = -\lambda^2$$

with

$$\lambda = \begin{cases} \frac{y_2 - y_1}{x_2 - x_1} & \text{if } P_2 \neq P_1 \\ \frac{3x_1^2 + a}{2y_1} & \text{if } P_2 = P_1 \end{cases}$$

$\Rightarrow$  algebraic formulas; ditto for  $[k]P = \underbrace{P \oplus \dots \oplus P}_{k \text{ times}}$ .

# Division polynomials

for  $E : y^2 = x^3 + Ax + B$

$$[n](X, Y) = \left( \frac{\phi_n(X, Y)}{\psi_n(X, Y)^2}, \frac{\omega_n(X, Y)}{\psi_n(X, Y)^3} \right)$$

$$\phi_n = X\psi_n^2 - \psi_{n+1}\psi_{n-1}, \quad 4Y\omega_n = \psi_{n+2}\psi_{n-1}^2 - \psi_{n-2}\psi_{n+1}^2$$

$$\phi_n, \psi_{2n+1}, \psi_{2n}/(2Y), \omega_{2n+1}/Y, \omega_{2n} \in \mathbb{Z}[A, B, X]$$

$$f_n(X) = \begin{cases} \psi_n(X, Y) & \text{for } n \text{ odd} \\ \psi_n(X, Y)/(2Y) & \text{for } n \text{ even} \end{cases}$$

$$f_{-1} = -1, \quad f_0 = 0, \quad f_1 = 1, \quad f_2 = 1$$

$$f_3(X, Y) = 3X^4 + 6AX^2 + 12BX - A^2$$

$$f_4(X, Y) = X^6 + 5AX^4 + 20BX^3 - 5A^2X^2 - 4ABX - 8B^2 - A^3$$

## Division polynomials (cont'd)

$$f_{2n} = f_n(f_{n+2}f_{n-1}^2 - f_{n-2}f_{n+1}^2)$$

$$f_{2n+1} = \begin{cases} f_{n+2}f_n^3 - f_{n+1}^3f_{n-1}(16Y^4) & \text{if } n \text{ is odd} \\ (16Y^4)f_{n+2}f_n^3 - f_{n+1}^3f_{n-1} & \text{otherwise.} \end{cases}$$

$$\deg(f_n(X)) = \begin{cases} (n^2 - 1)/2 & \text{if } n \text{ is odd} \\ (n^2 - 4)/2 & \text{otherwise.} \end{cases}$$

**Main use:** Schoof-Elkies-Atkin algorithm.

**Qi Cheng:** direct use of these formulas require keeping only nine values of  $f_{n+k}$  at the same time.

**McKee:** exact formula for  $[X^i]f_n$  (it is homogeneous in  $A$  and  $B$  with respective weights 4 and 6).

**Dewaghe:** polynomials of exact division (cyclotomic polynomials).



## The invariant $j$

$$\Delta(\tau) = g_2^3(\tau) - 27g_3^2(\tau) = q \prod_{n \geq 1} (1 - q^n)^{24}$$

$$j(\tau) = 1728 \frac{g_2^3(\tau)}{\Delta(\tau)}$$

$$j(q) = \frac{1}{q} + 744 + \sum_{n \geq 1} c_n q^n, c_n \in \mathbb{N}$$

where  $q = \exp(2i\pi\tau)$ .

**Def.**  $\mathcal{L}'$  and  $\mathcal{L}$  are **isomorphic** iff there exists  $P$  in  $SL_2(\mathbb{Z})$  s.t.

$$\begin{pmatrix} \omega'_1 \\ \omega'_2 \end{pmatrix} = P \begin{pmatrix} \omega_1 \\ \omega_2 \end{pmatrix}.$$

**Thm.**  $\mathcal{L}$  and  $\mathcal{L}'$  are isomorphic iff  $j(\mathcal{L}) = j(\mathcal{L}')$ .

## Isogenous lattices

**Def.**  $\mathcal{L}$  and  $\mathcal{M}$  are **isogenous** iff  $\exists \alpha \in \mathbb{C}, \alpha \mathcal{L} \subset \mathcal{M}$ .

**Most interesting case:**  $\mathcal{M}$  is a sublattice of  $\mathcal{L}$  s.t.  $\mathcal{L}/\mathcal{M}$  is cyclic of finite index. In other words:

$$\mathcal{M} = (a\omega_1 + b\omega_2)\mathbb{Z} + (c\omega_1 + d\omega_2)\mathbb{Z}$$

and  $ad - bc = m$  with  $\gcd(a, b, c, d) = 1$ .

**Fundamental theorem (modular polynomial):**  $\exists \alpha \in \mathbb{C}$  s.t.  $\alpha \mathcal{M} \subset \mathcal{L}$  iff  $\exists m$  s.t.  $\Phi_m(j(\mathcal{M}), j(\mathcal{L})) = 0$  where

$$\Phi_m(X, \tau) = \prod_{A \in \mathcal{S}_m} (X - j(A\tau)) = \sum_{k=0}^{\mu_0(m)} C_k(\tau) X^k,$$

with

$$\mathcal{S}_m = \left\{ \begin{pmatrix} a & b \\ 0 & d \end{pmatrix}, ad = m, \gcd(a, b, d) = 1, a > 0, d > b \geq 0 \right\}$$

of cardinality  $\mu_0(m) = m \prod_{p|m} (1 + 1/p)$ .

## SEA again

If  $j'$  is any root of  $\Phi_\ell(X, j(E))$ , then there exists an isogeny between  $E$  and  $E' = E/C$ :

$$\begin{array}{ccc} E & \xrightarrow{[\ell]} & E \\ & \searrow I & \nearrow I^* \\ & E' = E/C & \end{array}$$

$$I(X, Y) = \left( \frac{k(X)}{h(X)^2}, \dots \right), \quad \deg(k) = \ell, \deg(h) = (\ell - 1)/2$$

$$I(P) = O_E \Leftrightarrow h(X) = 0$$

$$I(P) = O_E \Rightarrow I^*(I(P)) = [\ell]P = O_E$$

$\rightarrow f_\ell(X)$  has a factor of degree  $(\ell - 1)/2$

**Rem.** Computing  $E'$ ,  $I$  is done via Weierstrass functions, see Bostan/Morain/Salvy/Schost (*Math. Comp.* 2008).

# Properties of the modular polynomial

**Thm.**

- ▶  $\Phi_m(X, Y) \in \mathbb{Z}[X, Y]$ ;
- ▶  $\Phi_m(Y, X) = \Phi_m(X, Y)$ ;
- ▶ if  $m$  is squarefree, then the coefficient of highest degree of  $\Phi_m(X, X)$  is  $\pm 1$ .

**Prop.** If  $(m_1, m_2) = 1$ , then

$$\Phi_{m_1 m_2}(X, J) = \text{Resultant}_Z(\Phi_{m_1}(X, Z), \Phi_{m_2}(Z, J)).$$

**Prop.** If  $m = \ell^e$  with  $e > 1$ , then

$$\Phi_{\ell^e}(X, J) = \text{Resultant}_Z(\Phi_{\ell}(X, Z), \Phi_{\ell^{e-1}}(Z, J)) / \Phi_{\ell^{e-2}}(Z, J)^{\ell}.$$

**Thm.**[Kronecker] If  $\ell$  is prime, then

$$\Phi_{\ell}(X, Y) \equiv (X^{\ell} - Y)(Y^{\ell} - X) \pmod{\ell}.$$

## Numerical examples

$$\begin{aligned}\Phi_2(X, Y) &= X^3 + X^2 (-Y^2 + 1488 Y - 162000) \\ &\quad + X (1488 Y^2 + 40773375 Y + 8748000000) \\ &\quad + Y^3 - 162000 Y^2 + 8748000000 Y - 157464000000000.\end{aligned}$$

$$\begin{aligned}\Phi_3(X, Y) &= X^4 + Y^4 + 1855425871872000000000X + 452984832000000X^2 + 36864000X^3 \\ &\quad + 1855425871872000000000Y - 770845966336000000XY + 8900222976000X^2Y \\ &\quad - 1069956X^3Y + 452984832000000Y^2 + 8900222976000XY^2 + 2587918086X^2Y^2 \\ &\quad + 2232X^3Y^2 + 36864000Y^3 - 1069956XY^3 + 2232X^2Y^3 - X^3Y^3.\end{aligned}$$

# Height

**Thm.** (P. Cohen)

$$H(\Phi_m) = 6\mu_0(m)(\log m - 2\sum_{p|m}(\log p)/p + O(1)).$$

$\ell$	101	211	503	1009	2003
$H(\Phi_\ell)$	3985	9256	24736	53820	115125
PCohen	2768	6743	18736	41832	91320

**Trick:**  $H \geq \log|\Phi_\ell(0,0)|$  and

$$\Phi_\ell(X, X) = \prod_{x^2 - Df^2y^2 = 4\ell} H_{-Df^2}(X)^*,$$

where  $H_{-Df^2}(X)$  is the Hilbert class polynomial (complex multiplication).

# Computing the modular polynomial

Remember that

$$j(q) = \frac{1}{q} + 744 + \sum_{n \geq 1} c_n q^n.$$

Then  $\Phi_\ell(X, Y)$  is such that  $\Phi_\ell(j(q), j(q^\ell))$  vanishes identically.

**Naive method:** indeterminate coefficients (over  $\mathbb{Q}$  or small  $p$ 's); at least  $\tilde{O}((\ell^2)^\omega)$  operations over  $\mathbb{Q}$ .

**Atkin** (analysis by Elkies):

1. use  $q$ -expansion of  $j$  and  $j$  with  $O(\ell^2)$  terms;
2. compute power sums of roots of  $\Phi_\ell$ , write them as polynomials in  $J$ ;
3. go back to coefficients of  $\Phi_\ell(X, J)$  via Newton's formulas;
4. use CRT on small primes.  $\tilde{O}(\ell^3 M(p))$ ; used for  $\ell \leq 1000$  fifteen years ago.

## Computing modular polynomials (cont'd)

**Charles+Lauter (2005):** compute  $\Phi_\ell$  modulo  $p$  using supersingular invariants mod  $p$ , Mestre *méthode des graphes*,  $\ell$  torsion points defined over  $\mathbb{F}_{p^{O(\ell)}}$  and interpolation.  $\tilde{O}(\ell^4 M(p))$

**Enge (2004); Dupont (2004):** use complex floating point evaluation and interpolation.  $\tilde{O}(\ell^3)$



Write

$$\Phi_\ell(X, J) = X^{\ell+1} + \sum_{u=0}^{\ell} c_u(J) X^u$$

where  $c_u(J) \in \mathbb{Z}[J]$ ,  $\deg(c_u(J)) \leq \ell + 1$ . All computations are done using precision  $H = O(\ell \log \ell)$ .

1. **for**  $\ell + 1$  values of  $z_i$  **do**:

1.1 Compute floating point approximations to the  $\ell + 1$  roots  $j_r(z_i)$  of  $\Phi_\ell(X, j(z_i))$  to precision  $H$ ;

1.2 Build  $\prod_{r=1}^{\ell+1} (X - f_r(z_i)) = X^{\ell+1} + \sum_{u=0}^{\ell} c_u(j(z_i)) X^u$ ;  
 $O(M(\ell) \log \ell)$  ops.

2. Perform  $\ell + 1$  interpolations for the  $c_u$ 's:  $O((\ell + 1)M(\ell) \log \ell)$  ops.

All 1.2 + 2 has cost  $O(\ell M(\ell)(\log \ell)M(H)) = \tilde{O}(\ell^3)$ .

## Examples

$\ell$	$r$	$H$	$\deg(J)$	eval( $s$ )	interp( $s$ )	tot (d)	Mb gz
3011	5	7560	200				368
3079	97	9018	254	7790	640	23	547
3527	13	9894	268	799	1440	3	746
3517	97	10746	290	12400	1110	42	850
4003	13	11408	308	1130	2320	4	1127
5009	5	13349	334	880	3110	3	1819
6029	5	16418	402	1550	6370	7	3251
7001	5	19473	466	2440	11700	13	5182
8009	5	22515	534	3500	20000	22	7905
9029	5	25507	602	5030	33100	35	11460
10079	5	28825	672	7690	56300	61	16152

## IV. A glimpse at genus 2

- ▶ Points  $\rightarrow$   $g$ -uplet of points on the Jacobian of  $y^2 = x^{2g+1} + \dots$ ;
- ▶ Lattice:  $\mathbb{C}/\mathbb{Z}^{2g}$ ;  $\wp(z_1, z_2, \dots, z_g)$ , theta-functions, etc.
- ▶  $j \rightarrow$  many more (3 for  $g = 2$ ).
- ▶ Division polynomials  $\rightarrow$  polynomial ideals.  $\Rightarrow$  Much harder, heavy computer algebra.
- ▶ Modular polynomials: very hard to compute; only  $\Phi_2$  known (R. Dupont).

$\Rightarrow$  ask Éric (Schost) and Pierrick (Gaudry) for more!

# Conclusions

- ▶ Interesting polynomials.
- ▶ Need computer algebra everywhere (fast techniques).
- ▶ More to be done as  $g$  increases.

**Complements (if needed!!!)**

## Oddities

**Binary** polynomials  $\Phi_{pq}(X)$  have  $\{-1, 0, 1\}$  coeffs; exact formula known for  $a_{pq}(k)$  (Kaplan, 2007).

**Ternary** polynomials  $\Phi_{pqr}$ :

- ▶ The smallest  $n$  for which  $C_n(X)$  has a coeff not  $\leq 1$  is  $n = pqr = 105$  and  $a_{105}(7) = a_{105}(41) = -2$ .
- ▶ Gallot/Moree recent preprint on arXiv: If  $p < q < r$  are primes s.t.

$$p > 3, \quad q \equiv 2 \pmod{p}, \quad r \equiv (pq - 1)/2 \pmod{pq}.$$

Put  $n = pqr$ ,  $k = (p - 1)(qr + 1)/2$ . Then

$$\{a_n(k - r) = -(p - 1)/2, \dots, a_n(k) = (p + 1)/2\} = [-(p - 1)/2, (p + 1)]$$

and for all  $i$ ,  $|a_n(i + 1) - a_n(i)| \leq 1$  (**jump property**).

**Ex.**  $p = 5$ ,  $q = 7$ ,  $r = 17$  yields

$$[-2, -1, 0, 1, 1, 2, 2, 1, 0, 0, -1, -1, -1, 0, 1, 2, 2, 3].$$