

Algebraic Biochemistry: a Framework for On-line Analog Computation in Cells

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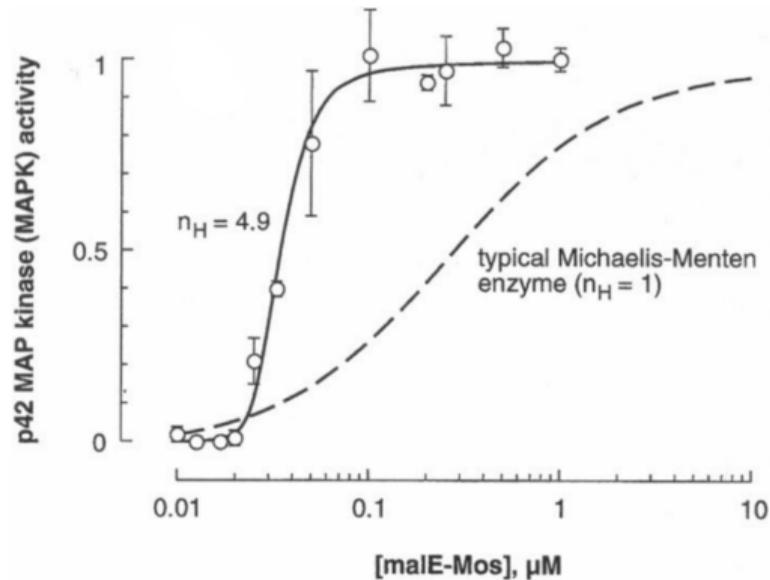
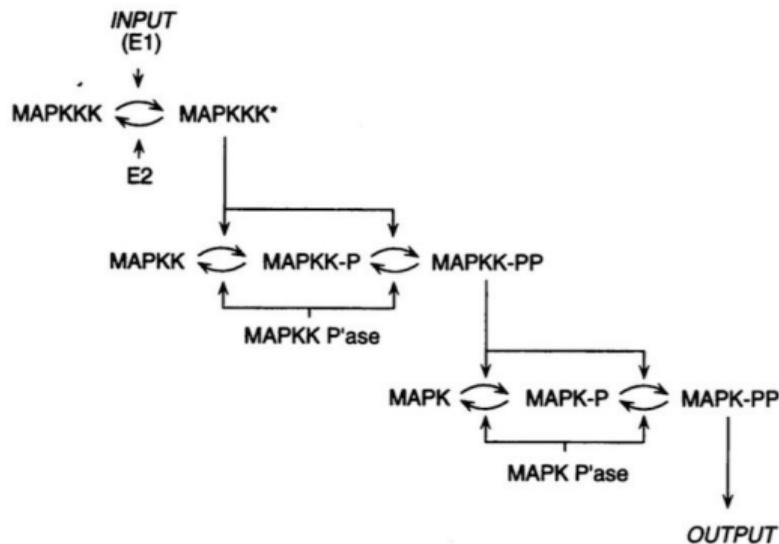
The logo for Inria, featuring the word "Inria" in a stylized, red, cursive script font.

EPI Lifeware, Centre Inria Saclay Île de France

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Mitogen-Activated Protein Kinase (MAPK)

An “on-line” analog digital converter: $f(x) \approx \frac{x^{4.9}}{1+x^{4.9}}$!



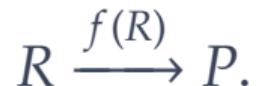
Huang, C.-Y. et Ferrell J. Ultrasensitivity in the MAPK cascade. PNAS, 1996, 93,19.

Chemical Reactions Networks 101

A Chemical Reaction Network (CRN) is:

- a finite set of species X_i
- a set of reactions:
 - a multiset of reactants: R
 - a multiset of products: P
 - a kinetic rate function: $f(R)$

and denoted:



Chemical Reactions Networks 101

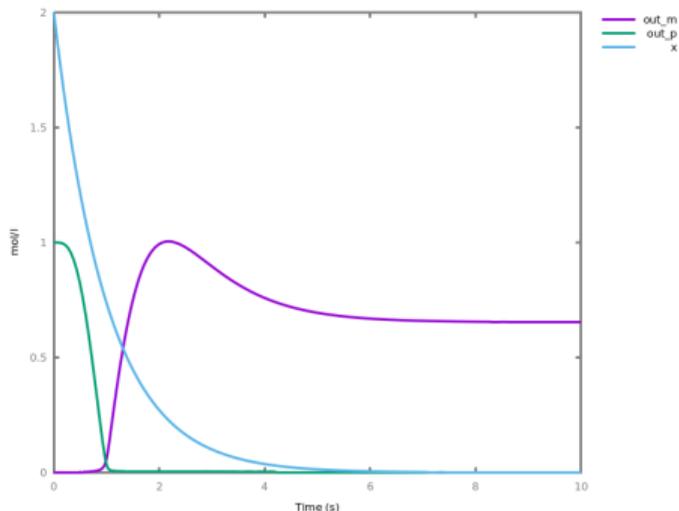
In the *differential semantics*, a CRN defines an Ordinary Differential Equations (ODE), $A \xrightarrow{f(A)} 2 \cdot B$:

$$\begin{aligned}\frac{dA}{dt} &= -f(A), \\ \frac{dB}{dt} &= 2 \cdot f(A).\end{aligned}\tag{1}$$

With *Mass Action Law*, f is a monomial: $f(R) = k \prod_{X \in R} X$, and we have Polynomial ODE (PODE).

Chemical Reactions Networks are Turing Complete

$$\lim_{t \rightarrow \infty} \text{out} = \cos x^2$$



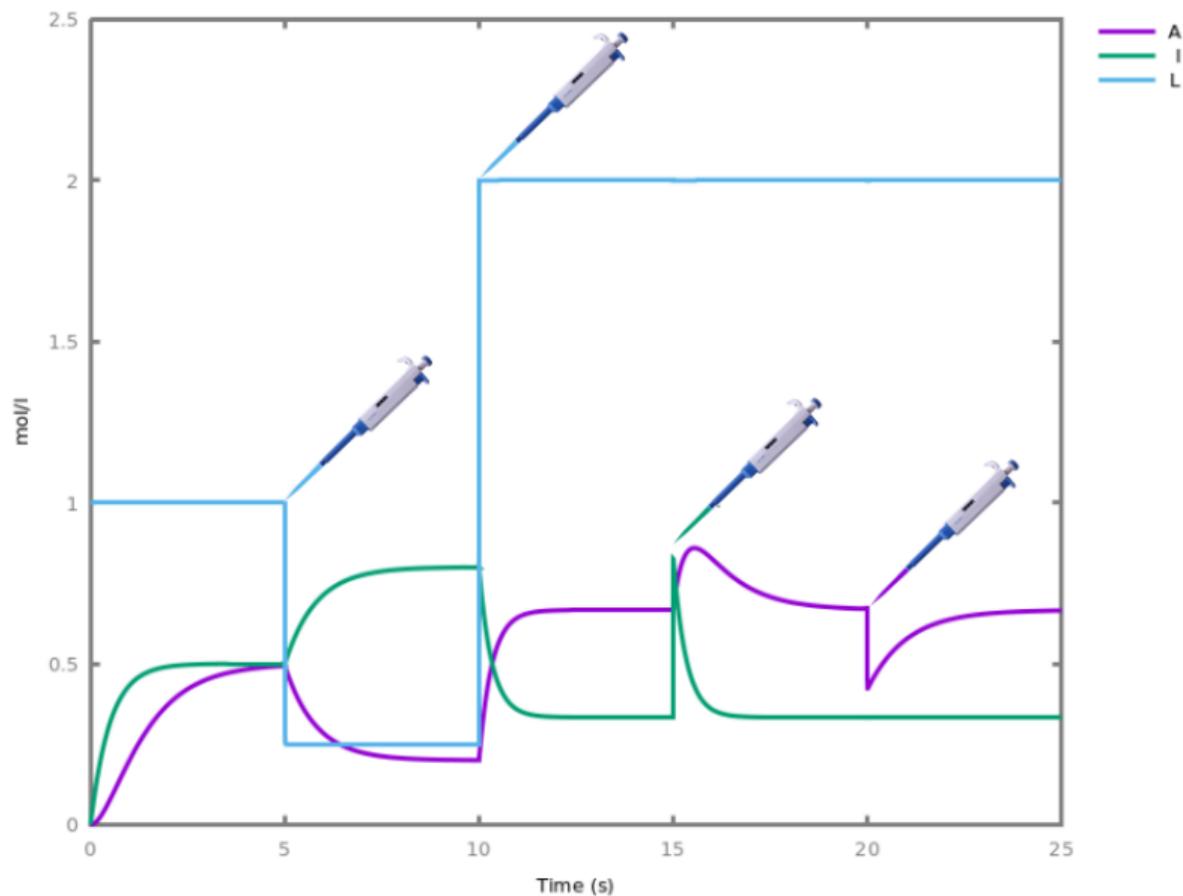
$f : I \subset \mathbb{R}_+ \rightarrow \mathbb{R}_+$ is CRN-computable if there exist a CRN over species \bar{y} , with ODE semantics \bar{p} , and a polynomial $q \in \mathbb{R}_+^n[\mathbb{R}_+]$, such that for all $x \in I$:

- $\bar{y}(0) = \bar{q}(x)$
- $\bar{y}'(t) = \bar{p}(\bar{y}(t))$
- $\forall t > 1, \quad |y_1(t) - f(x)| \leq y_2(t),$
- $y_2(t) \geq 0$ and $\lim_{t \rightarrow \infty} y_2(t) = 0.$

Fages F, Le Guludec G, et al. Strong turing completeness of continuous CRN..., CMSB 2017.

Computing “on-line” with perturbations

A toy model:
the “1-level MAPK”



Battle plan

What is the class of functions that can be computed
On-line by a CRN?

- Definition
- Proof & Theorem
- Compilation pipeline

Def – F_S : functions stabilized by a CRN

A CRN over m inputs X , 1 output y and n auxiliary Z , stabilizes $f : I \subset \mathbb{R}_+^m \mapsto \mathbb{R}_+$, over the domain $D \subset \mathbb{R}_+^{m+1+n}$ if:

- 1 $\forall X^0 \in I, \{(X, y, Z) \in D | X = X^0\}$ is of plain dimension $n + 1$,
- 2 In the differential semantic with pinned input species X and initial conditions $X^0, y^0, Z^0 \in D$:

$$\boxed{\lim_{t \rightarrow \infty} y(t) = f(X^0)}.$$

It can be extended to functions $f : \mathbb{R}^m \rightarrow \mathbb{R}$ by dual-rail encoding:

$$\lim_{t \rightarrow \infty} (y^+ - y^-)(t) = f(X^+ - X^-).$$

Example: 1-level MAPK

We use the differential semantics with Mass Action Law.

At steady state, we have:

$$\frac{dI}{dt} = 1 - I - LI = 0, \quad \frac{dA}{dt} = LI - A = 0, \quad \frac{dL}{dt} = 0. \quad (3)$$

Thus by elimination:

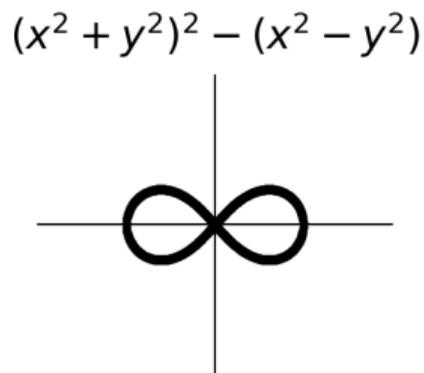
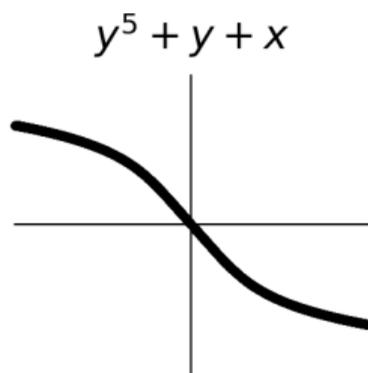
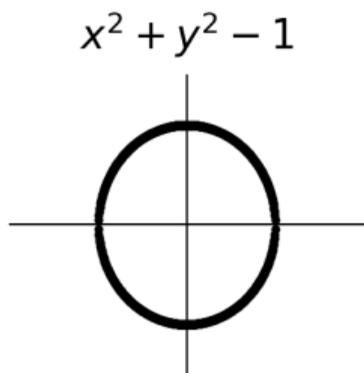
$$P(L, A) = \boxed{L - A - LA = 0}. \quad (4)$$

$$\text{That is: } A(L) = \frac{L}{1+L}.$$

Def – F_A : real algebraic functions

A function $f : I \subset \mathbb{R}^m \mapsto \mathbb{R}$ is algebraic iff there exists a polynomial P_f of $m + 1$ variables such that:

$$\forall X \in I, \quad \boxed{P_f(X, f(X)) = 0}. \quad (5)$$



$F_S \subset F_A$ (sketch of the proof)

As $y(t)$ is projectively polynomial [1]:

$$P(X, y, y^{(1)}, \dots, y^{(n)}) = 0 \quad (6)$$

with X^0, y^0 as initial condition we have: $\forall k, y^{(k)}(t) = 0$, thus:

$$P^\star(X^0, y^0) = 0. \quad (7)$$

Can we stabilize any algebraic functions?

[1] Carothers D., Parker G., et al. Some properties of solutions to PODE. EJDE 2005, vol. 2005, No. 40, 17.

$F_A \subset F_S$ (sketch of the proof)

Take $f : I \mapsto \mathbb{R}$ a real algebraic function

$\exists P_f$ in **reduced form**, $\forall X \in I, P_f(X, f(X)) = 0$,

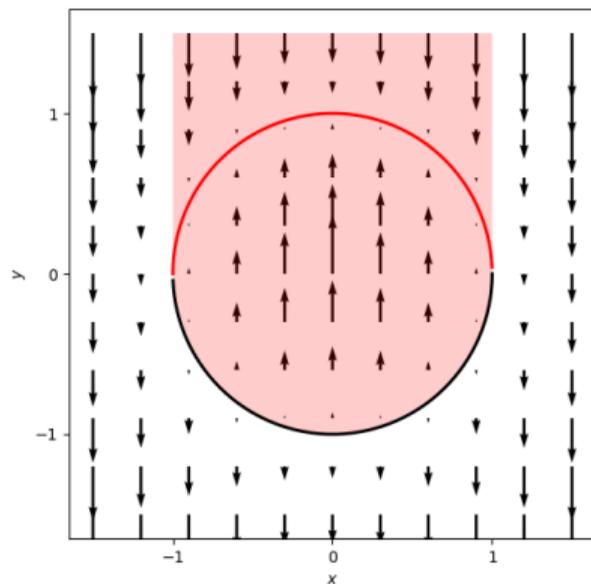
$$\begin{aligned} \frac{dy}{dt} &= \pm P_f(X, y), \\ \frac{dX}{dt} &= 0, \end{aligned} \tag{8}$$

$Y = f(X)$ is a fixed point.

Choose + or - to stabilize the desired branch.

Example: Unit circle

$$P(x, y) = 1 - x^2 - y^2 \quad (9)$$

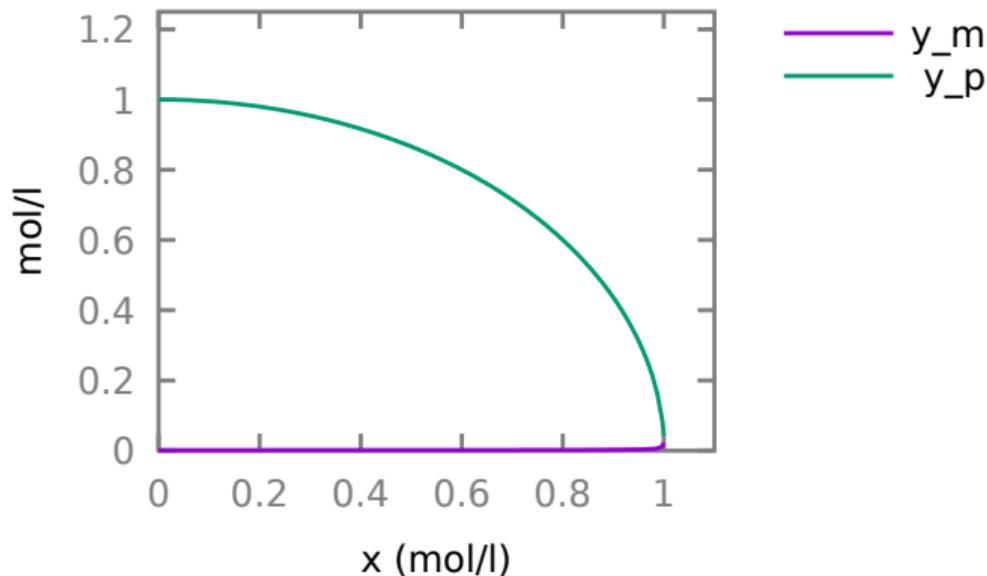
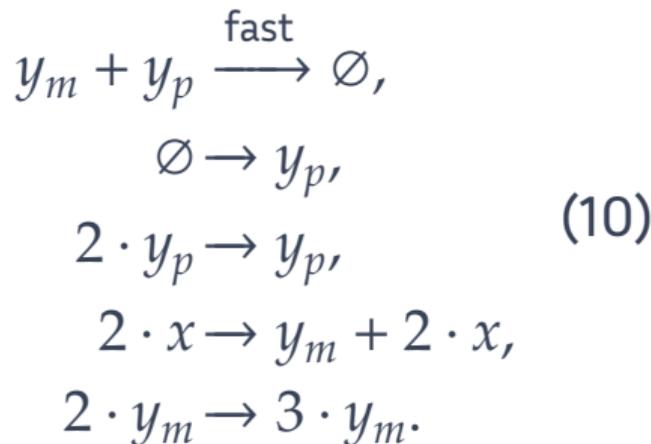


Theorem

$$F_S = F_A$$

Example: Unit circle in Biocham

`stabilize_expression(x^2 + y^2 - 1, y, [x = 0, y=1]).`



Compilation pipeline in Biocham

```
stabilize_expression(polynomial, output_name, point)
```

Polynomialization preprocessing to ease the user interface

Stabilization choose the \pm to stabilize the selected branch

Quadratization reduce to a quadratic PODE to enforce elementary reactions

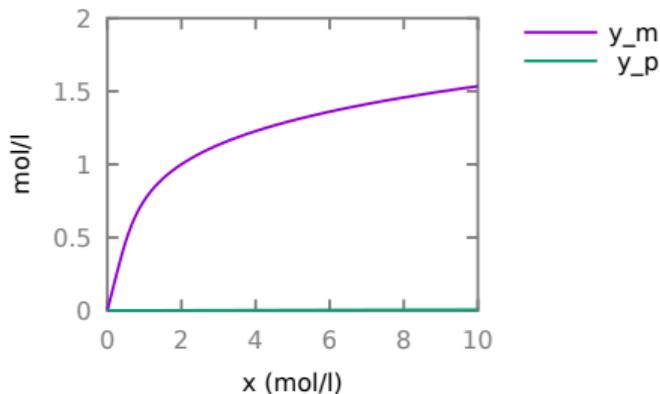
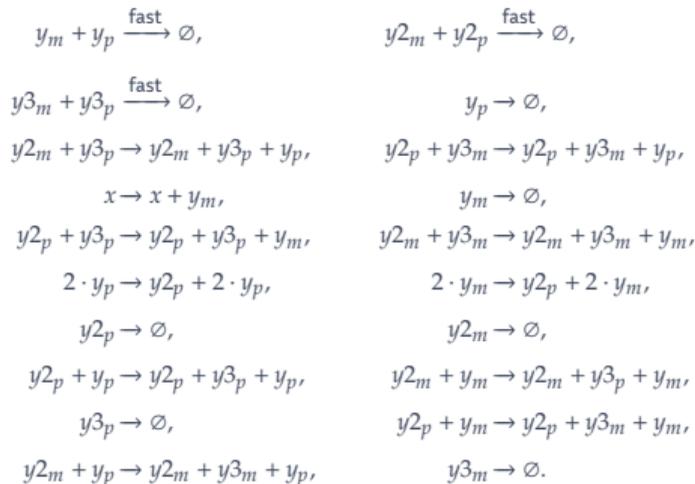
Dual-rail encoding split variable that may become negative into their positive and negative part

CRN generation insert a reaction for each monomial

Hemery M., Fages F. et Soliman S., On the complexity of quadratization for PODE. CMSB 2020.

Example: Bring radical

`stabilize_expression(y^5 + y + x, y, [x = 2, y=1]).`



Hermite, C. Sur la résolution de l'équation du cinquième degré. Comptes Rendus de l'Académie des Sciences, 1858.

A two inputs example: Computing the norm

$$\text{norm}(x, y) = \sqrt{x^2 + y^2} \quad \Rightarrow \quad P(x, y, \text{norm}) = \text{norm}^2 - x^2 - y^2$$

`stabilize_expression(norm^2-x^2-y^2, norm,
[norm = 1.4, x = 1, y=1]).`

$$2 \cdot x \rightarrow \text{norm} + 2 \cdot x,$$

$$2 \cdot y \rightarrow \text{norm} + 2 \cdot y,$$

$$2 \cdot \text{norm} \rightarrow \text{norm}.$$

Conclusions & Perspectives

- On-line analog computation is restricted to the set of algebraic functions,
- We can compile any algebraic function into an abstract stabilizing CRN.
- Towards a concrete compiler with real reactions from BRENDA,
- and its *in-vitro* implementation [1].

[1] Courbet A, Amar P, Fages F, et al. Computer-aided biochemical programming of synthetic microreactors as diagnostic devices. Mol. sys. bio., 2018.