

Fast spherical drawing of planar triangulations: an experimental study of graph drawing tools

june 27-29 2018, SEA2018, L'Aquila

Luca Castelli Aleardi

Gaspard Denis

Éric Fusy

(planar) graph drawing

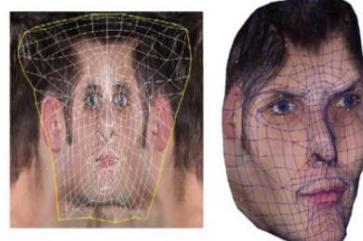
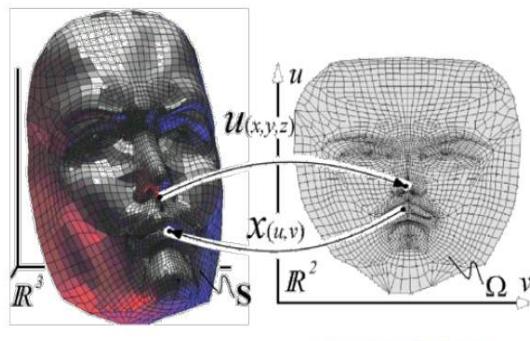
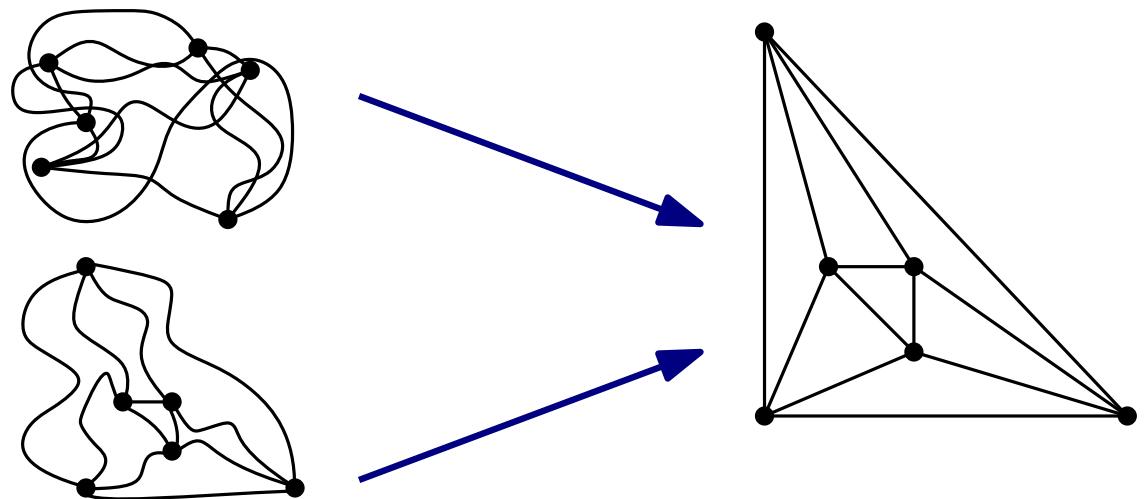
(a very short survey)

Straight-line planar drawings of planar graphs

Problem definition

Input: a planar graph (or planar map)

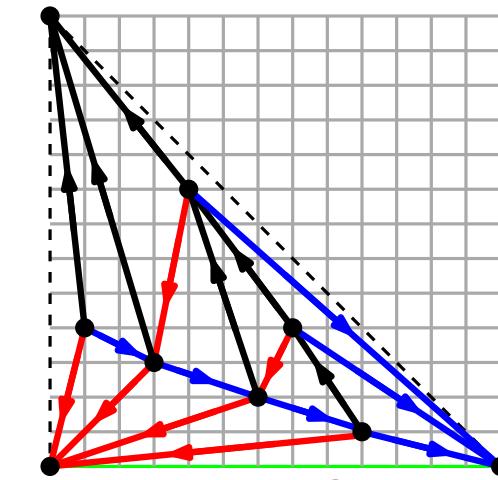
Output: a straight-line planar drawing
(crossing-free)



Bennis et al., 1991
Maillot et al., 1993

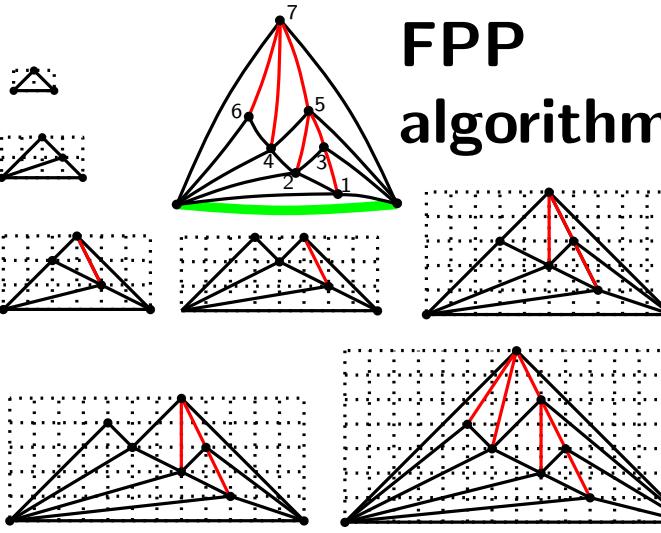
Straight-line planar drawings of planar graphs

Thm (Schnyder 1990)



face counting via Schnyder woods

Thm (De Fraysseix, Pack Pollack 1989)



**FPP
algorithm**

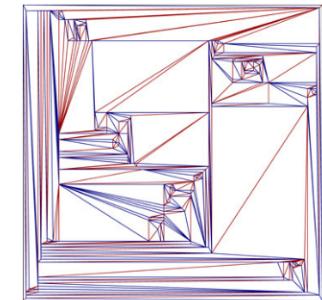
shift algorithm via Canonical orderings

linear time algorithms

$O(n) \times O(n)$ grid drawings

not trivial to implement

extremely fast: they can process millions of vertices per second



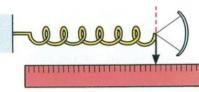
(images by E. Fusy)

**Spring embedder (Eades, 1984)
(Fruchterman and Reingold, 1991)**

force-directed paradigm

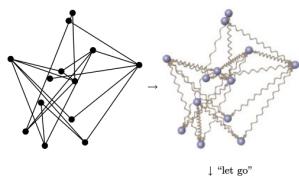
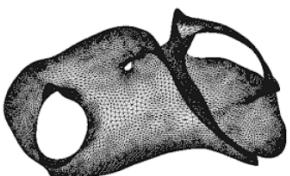
easy to implement

pretty slow: $O(n^2)$ or $O(n \log n)$ time per iteration



$$\mathbf{F}_a(v) = c_1 \cdot \sum_{(u,v) \in E} \log(\text{dist}(u,v)/c_2)$$

$$\mathbf{F}_r(v) = c_3 \cdot \sum_{u \in V} \frac{1}{\sqrt{\text{dist}(u,v)}}$$



images from Kaufman Wagner (Springer, 2001)

[Tutte'63] Tutte barycentric embedding

minimize the spring energy

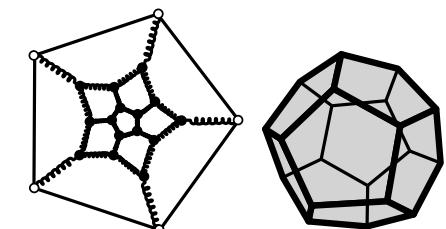
$$E(\rho) := \sum_{(i,j) \in E} |\mathbf{x}(v_i) - \mathbf{x}(v_j)|^2 = \sum_{(i,j) \in E} (x_i - x_j)^2 + (y_i - y_j)^2$$

solve large sparse linear systems

$$\mathbf{x}(v_i) = \sum_{j \in \mathcal{N}(i)} \frac{1}{\deg(v_i)} \mathbf{x}(v_j)$$

easy to implement

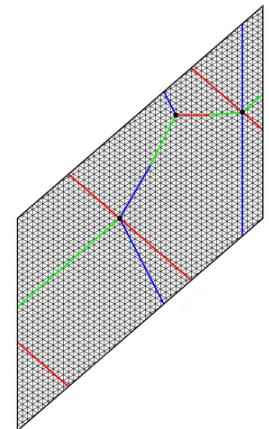
not very fast: they can process $\approx 10^4$ vertices per second



Straight-line drawings of toroidal graphs ($g = 1$)

(Goncalves Leveque 2014)

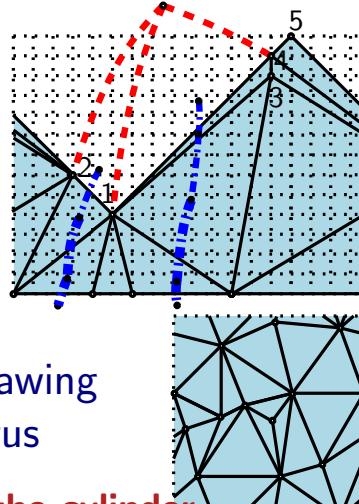
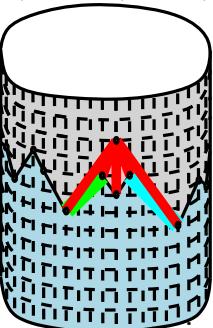
$O(n^2) \times O(n^2)$ grid drawing



face counting in the universal cover

(Castelli Aleardi, Devillers, Fusy 2012)

$O(n) \times O(n\sqrt{n})$ grid drawing

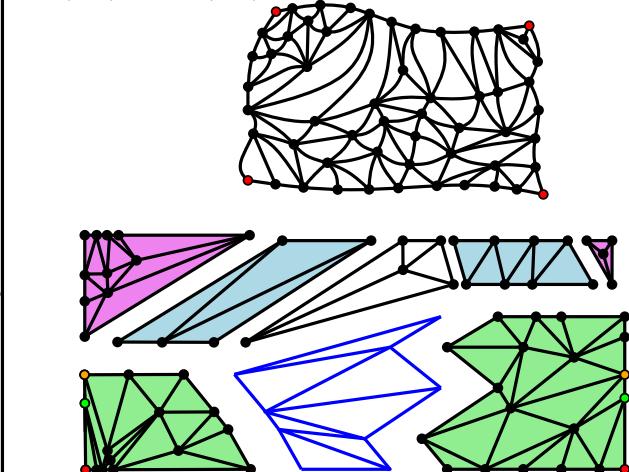


periodic drawing
flat torus

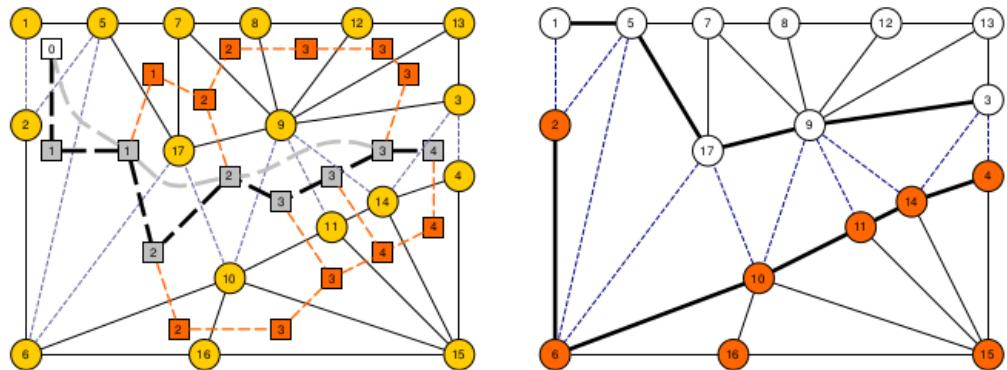
shift algorithm on the cylinder

(Castelli Aleardi, Fusy, Kostrygin 2014)

$O(n^4) \times O(n^4)$ grid drawing



(Duncan, Goodrich, Kobourov 2011)



cut along a polygonal scheme + shift algorithm



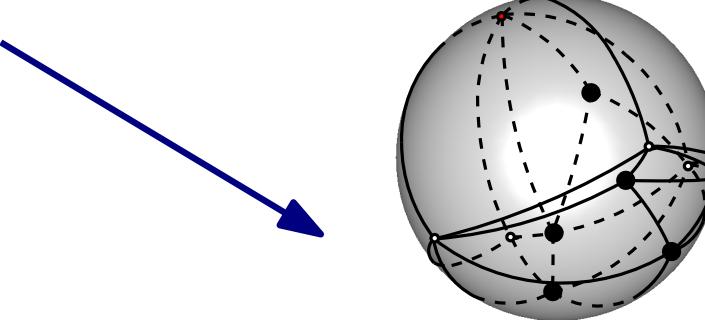
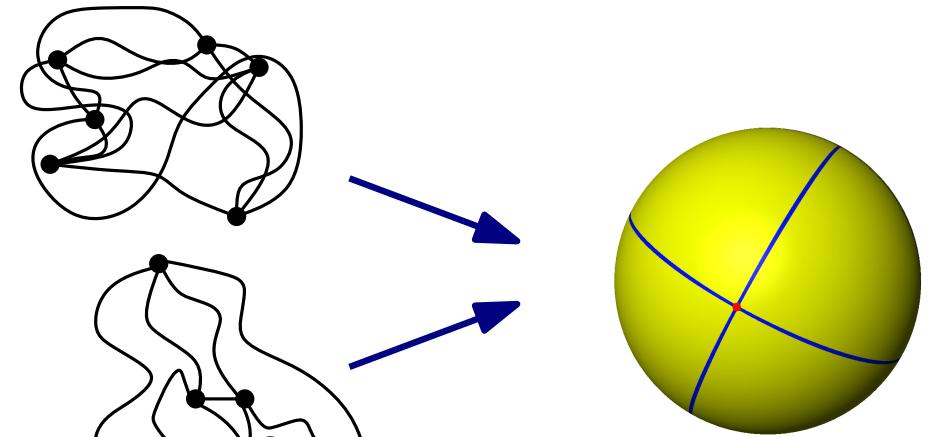
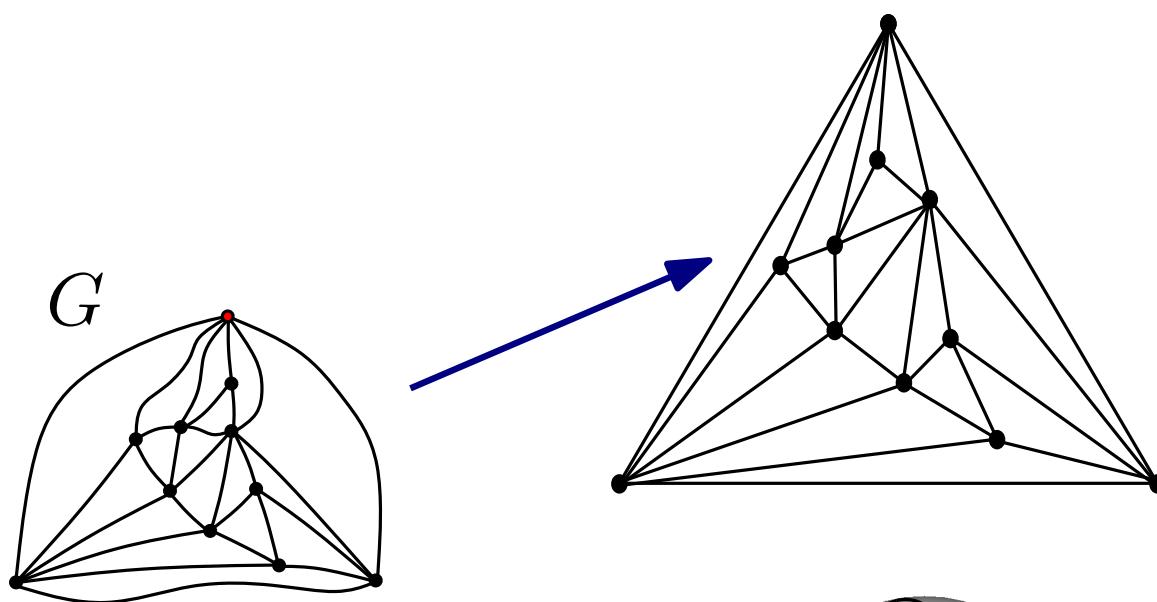
Spherical drawing and parameterizations

Spherical drawings/parameterizations of planar graphs

Problem definition

Input: a planar graph (or planar map)

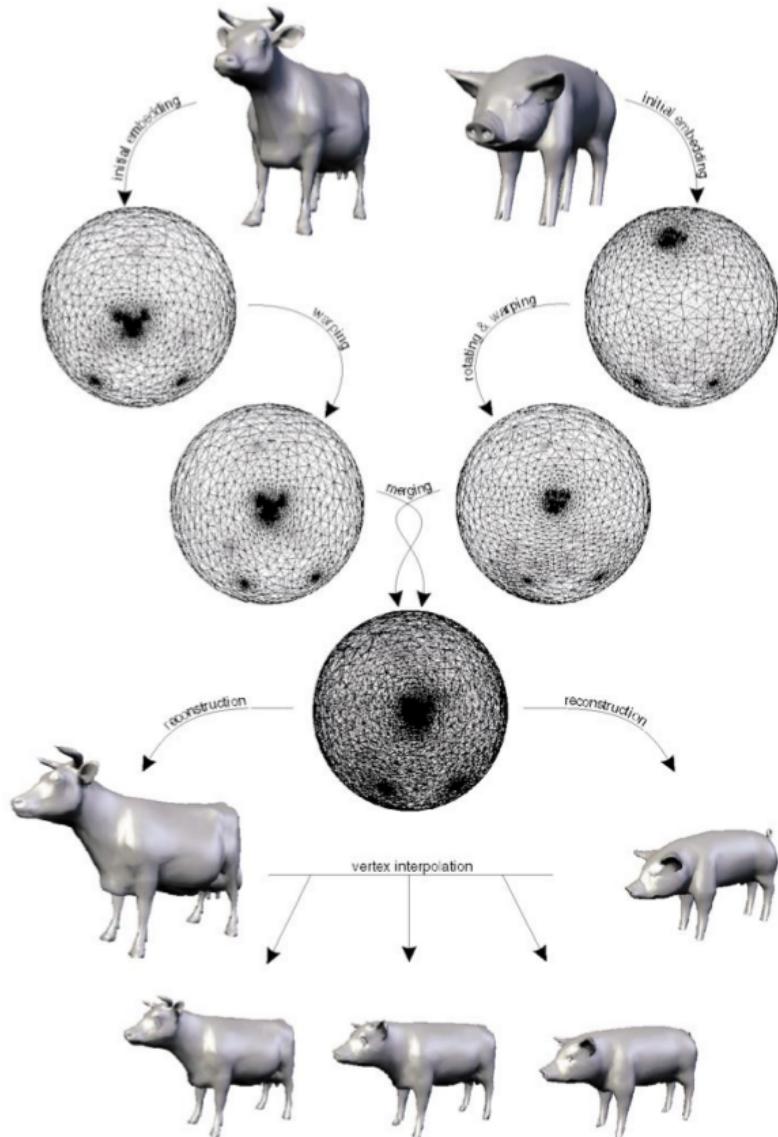
Output: a geodesic spherical drawing
(crossing-free)



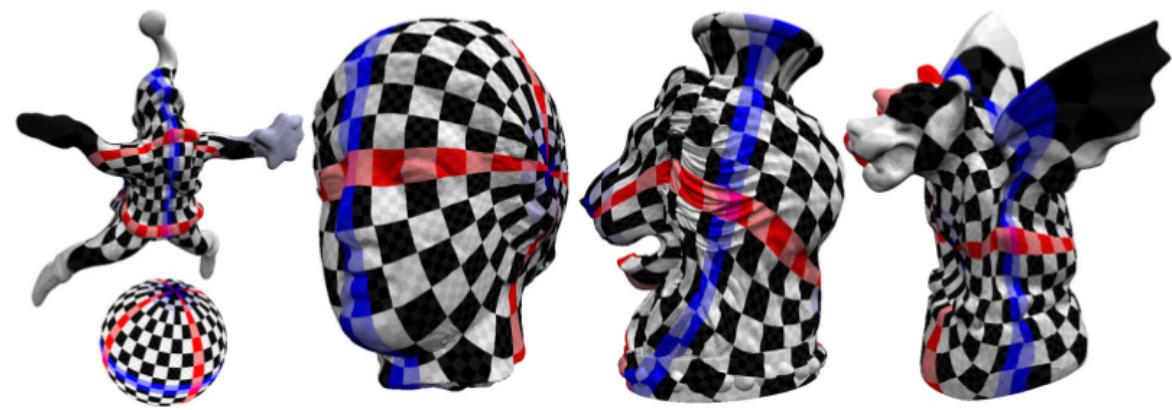
Spherical parameterization: applications

Morphing

(Alexa, 2000)



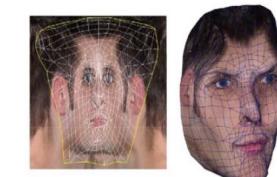
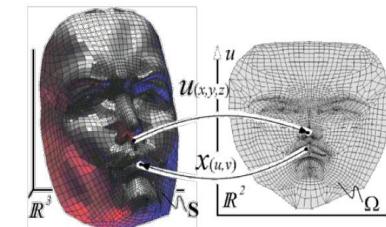
Texture mapping



(image by Friedel, Schroder, Desbrun, 2007)

Spherical vs. planar parameterization

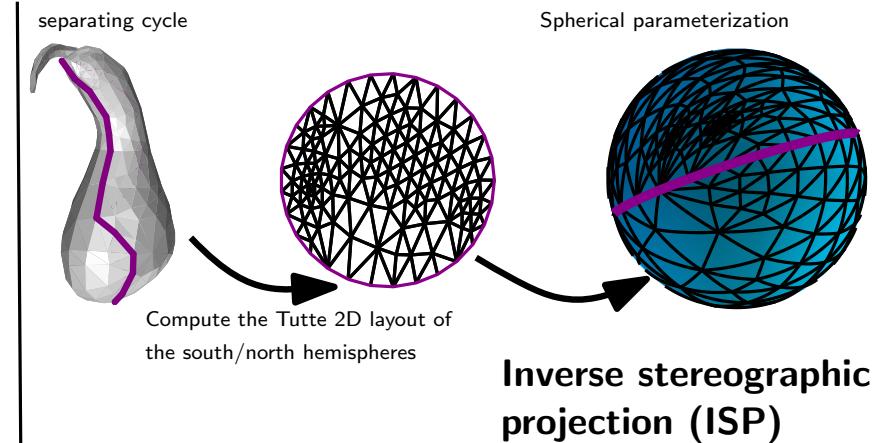
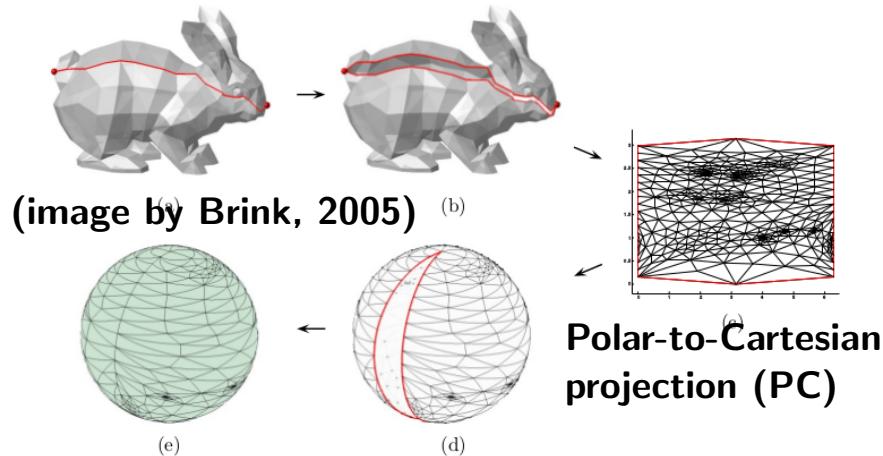
Main advantage: the spherical domain allows for seamless and continuous parameterization in the case of genus 0 meshes (without boundaries)



Bennis et al., 1991
Maillot et al., 1993

Spherical drawings/parameterizations: related works

Spherical parameterizations: use planar parameterization + spherical projections (**not crossing-free**)



compute a **Steinitz representation** and project on the sphere
(Shapiro, Tal, 1998)

Non linear optimization: (try to) solve Tutte equations on the sphere

(Alexa, 2000)

(Gotsman Gu Sheffer, 2003)

(Sheffer Gotsman Dyn, 2004)

(Saba, Yavneh, Gotsman, Sheffer, 2005)

(Zayer, Rössl, Seidel 2006)

(Friedel, Schroder, Desbrun 2007)

(Aigerman, Lipman, 2015)

(Aigerman, Kowalsky, Lipman, 2017)

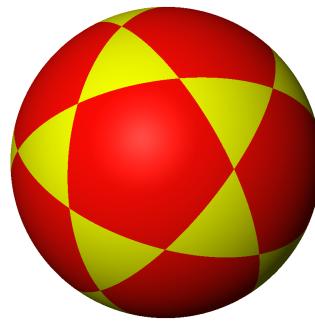
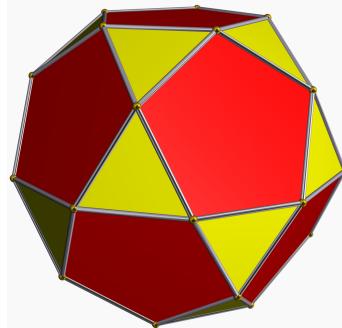


Extends graph drawing tools on the sphere: use non-euclidean spring embedders
(Kobourov Wampler, 2005)

Spherical drawings/parameterizations: related works

Spherical parameterizations: use planar parameterization + spherical projections (**not crossing-free**)

(layouts from Wikipedia)



(Pak, Wilson, 2017)

best upper bound on the grid size:
 $O(n^3) \times O(n^5) \times \exp(O(\sqrt{n} \log n))$

compute a **Steinitz polyhedral representation** and project on the sphere
(Shapiro, Tal, 1998) it takes $O(n^2)$ in the worst case

Non linear optimization: (try to) solve Tutte equations on the sphere

(Alexa, 2000)

(Gotsman Gu Sheffer, 2003)

(Sheffer Gotsman Dyn, 2004)

(Saba, Yavneh, Gotsman, Sheffer, 2005)

(Zayer, Rössl, Seidel 2006)

(Friedel, Schroder, Desbrun 2007)

(Aigerman, Lipman, 2015)

(Aigerman, Kowalsky, Lipman, 2017)

Extends graph drawing tools on the sphere: use non-euclidean spring embedders

(Kobourov Wampler, 2005)

Spherical drawings/parameterizations: related works

Spherical parameterizations: use planar parameterization + spherical projections (**not crossing-free**)

Non linear optimization: (try to) solve Tutte equations on the sphere

(Alexa, 2000)

(Gotsman Gu Sheffer, 2003)

(Sheffer Gotsman Dyn, 2004)

(Saba, Yavneh, Gotsman, Sheffer, 2005)

(Zayer, Rössl, Seidel 2006)

(Friedel, Schroder, Desbrun 2007)

(Aigerman, Lipman, 2015)

(Aigerman, Kowalsky, Lipman, 2017)

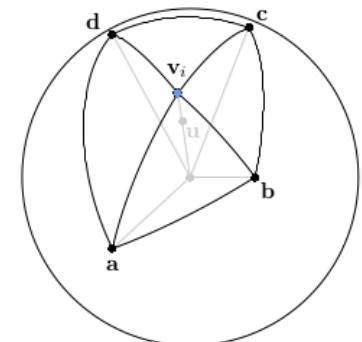
$$x_i^2 + y_i^2 + z_i^2 = 1 \quad i=1,..,n$$

$$\alpha_i x_i - L_w[i]x = 0 \quad i=1,..,n$$

$$\alpha_i y_i - L_w[i]y = 0 \quad i=1,..,n$$

$$\alpha_i z_i - L_w[i]z = 0 \quad i=1,..,n$$

$$\mathbf{v}_i = \frac{\mathbf{u}}{\|\mathbf{u}\|}.$$



perform iterative relaxation

Projected Gauss-Seidel($\mathbf{x}^0, \lambda, \varepsilon$)

$r = 0$; // **iteration counter**

do {

 for($i = 1; i \leq n; i++$) {

$$\mathbf{s} = (1 - \lambda)\mathbf{x}^r(v_i) + \lambda \sum_j w_{ij}\mathbf{x}^r(v_j)$$

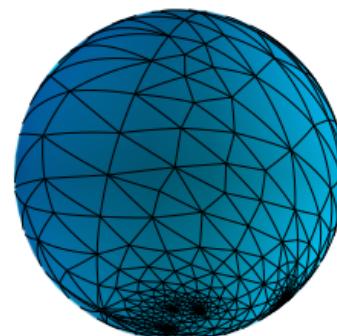
$$\mathbf{x}^{r+1}(v_i) = \frac{\mathbf{s}}{\|\mathbf{s}\|}$$

}

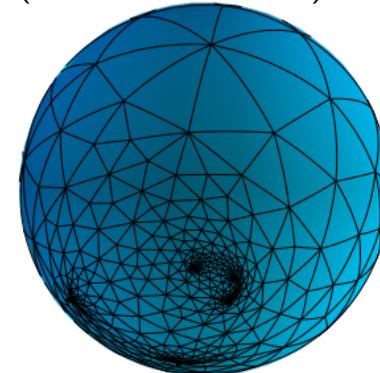
$r++$;

} while ($\|\mathbf{x}^r - \mathbf{x}^{r-1}\| > \varepsilon$)

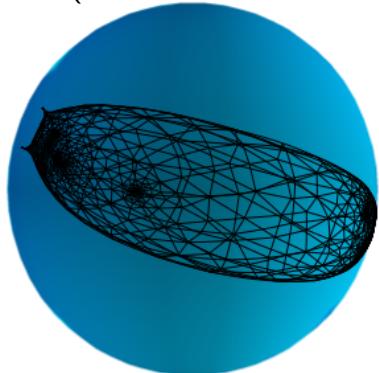
x^0 (initial layout)



x^{50} (after 50 iterations)



x^{1600} (after 1600 iterations)



collapse to degenerate solution

Extends graph drawing tools on the sphere: use non-euclidean spring embedders

Spherical drawings/parameterizations: related works

Spherical parameterizations: use planar parameterization + spherical projections (**not crossing-free**)

Non linear optimization: (try to) solve Tutte equations on the sphere

(Alexa, 2000)

(Gotsman Gu Sheffer, 2003)

(Sheffer Gotsman Dyn, 2004)

(Saba, Yavneh, Gotsman, Sheffer, 2005)

(Zayer, Rössl, Seidel 2006)

(Friedel, Schroder, Desbrun 2007)

(Aigerman, Lipman, 2015)

(Aigerman, Kowalsky, Lipman, 2017)

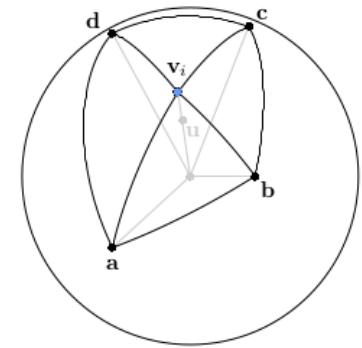
$$x_i^2 + y_i^2 + z_i^2 = 1 \quad i=1,..,n$$

$$\alpha_i x_i - L_w[i]x = 0 \quad i=1,..,n$$

$$\alpha_i y_i - L_w[i]y = 0 \quad i=1,..,n$$

$$\alpha_i z_i - L_w[i]z = 0 \quad i=1,..,n$$

$$\mathbf{v}_i = \frac{\mathbf{u}}{\|\mathbf{u}\|}.$$



perform iterative relaxation

Alexa($\mathbf{x}^0, \varepsilon$)

Idea: penalize long edges

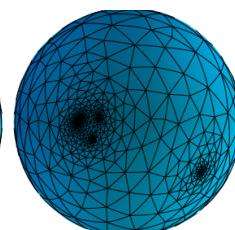
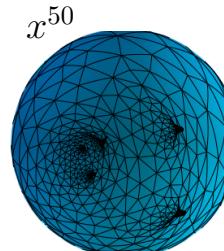
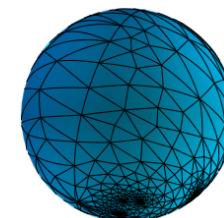
(Alexa, 2000)

$$c_i^r = [\max_{j \in N(i)} \{\|\mathbf{x}^r(v_i) - \mathbf{x}^r(v_j)\|\}]^{-1}$$

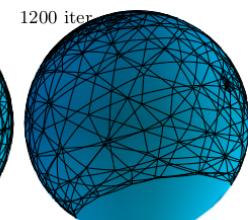
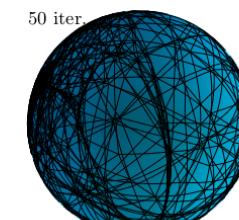
$$\mathbf{q}^r = \frac{c_i^r}{\deg(v_i)} \sum_{j \in N(i)} (\mathbf{x}^r(v_i) - \mathbf{x}^r(v_j)) \|\mathbf{x}^r(v_i) - \mathbf{x}^r(v_j)\|$$

$$\mathbf{x}^{r+1}(v_i) = \frac{\mathbf{x}^r(v_i) - \mathbf{q}_i^r}{\|\mathbf{x}^r(v_i) - \mathbf{q}_i^r\|}$$

\mathbf{x}^0 (initial layout)



\mathbf{x}^0 (initial random layout)



Extends graph drawing tools on the sphere: use non-euclidean spring embedders

Spherical drawings/parameterizations: related works

Spherical parameterizations: use planar parameterization + spherical projections (**not crossing-free**)

compute a **Steinitz representation** and project on the sphere
(Shapiro, Tal, 1998)

Non linear optimization: (try to) solve Tutte equations on the sphere

(Alexa, 2000)

(Gotsman Gu Sheffer, 2003)

(Sheffer Gotsman Dyn, 2004)

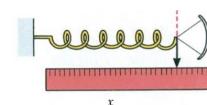
(Saba, Yavneh, Gotsman, Sheffer, 2005)

(Zayer, Rössl, Seidel 2006)

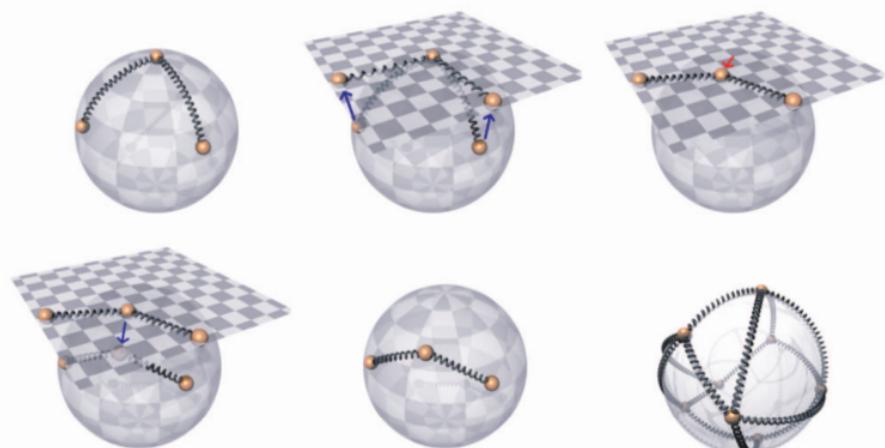
(Friedel, Schroder, Desbrun 2007)

(Aigerman, Lipman, 2015)

(Aigerman, Kowalsky, Lipman, 2017)



non-euclidean spring embedders



Extends graph drawing tools on the sphere:

use non-euclidean spring embedders

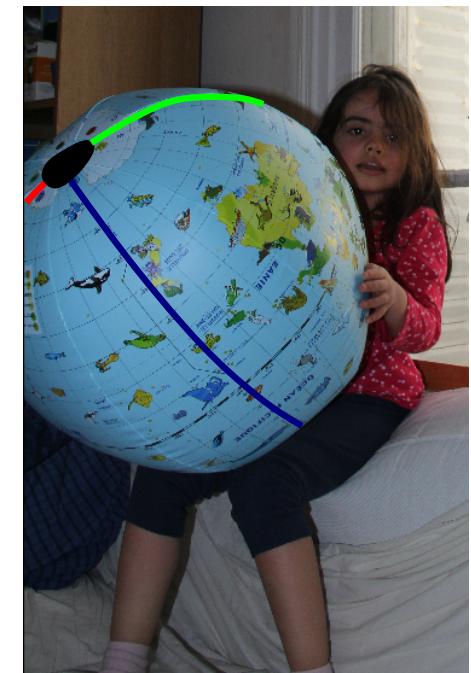
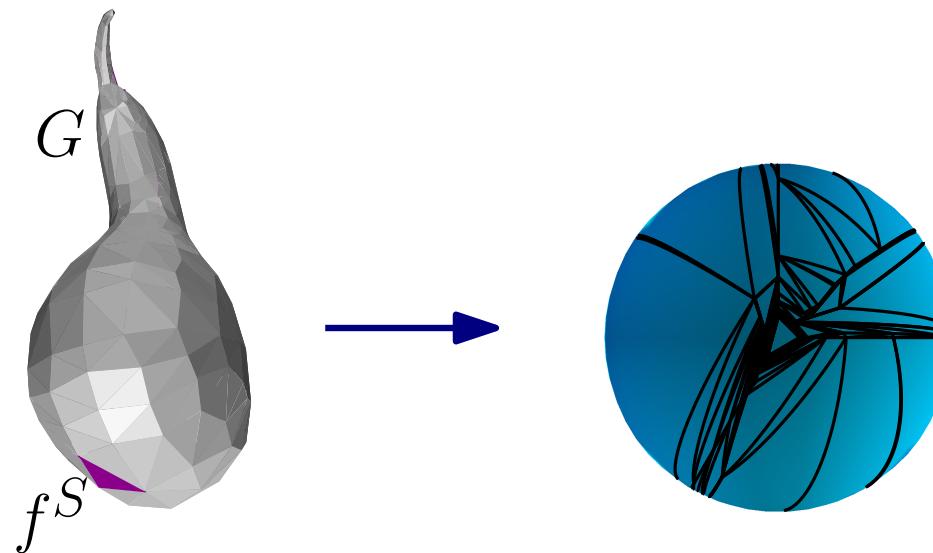
(Kobourov Wampler, 2005)

Fast and robust spherical drawing: our solution

Theorem

Let G be a planar triangulation with n vertices, having two non-adjacent faces f^S and f^N .

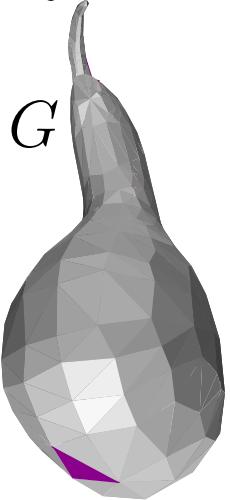
Then one can compute in $O(n)$ time a spherical drawing of G , where edges are drawn as (non-crossing) geodesic arcs of length at least $\Omega(\frac{1}{n})$.



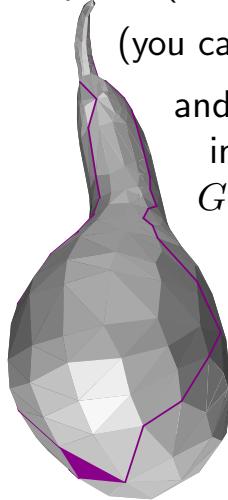
our SFPP algorithm: spherical FPP

Our combinatorial solution

take 2 vertex
disjoint faces



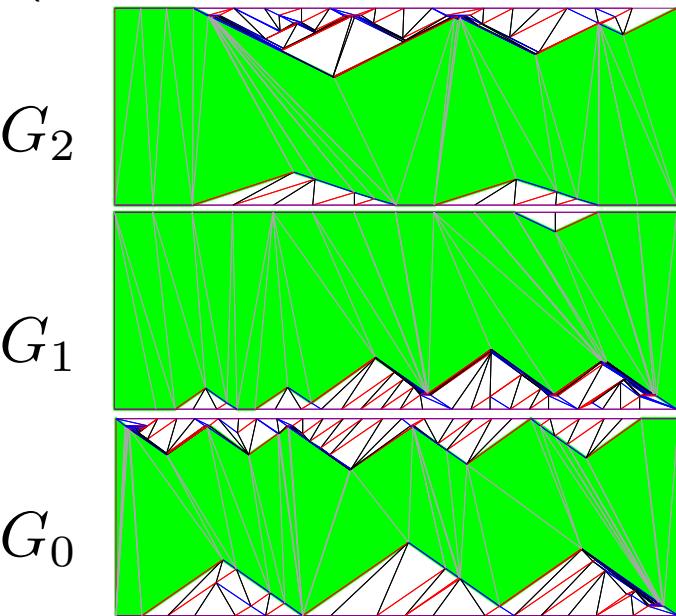
compute 3 vertex-disjoint
paths (Menger theorem)
(you can use **Schnyder woods**)



and decompose G
into 3 regions
 G_0, G_1 and G_2

G

compute and align the 6 rectangular layouts
(Castelli Aleardi, Fusy, Kostrygin 2014)
(Castelli Aleardi, Devillers, Fusy, 2012)

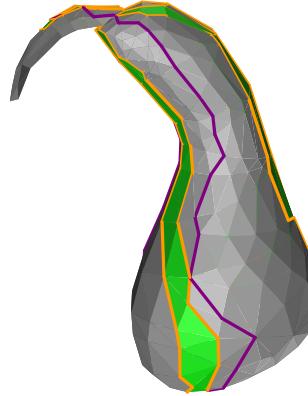


G_2

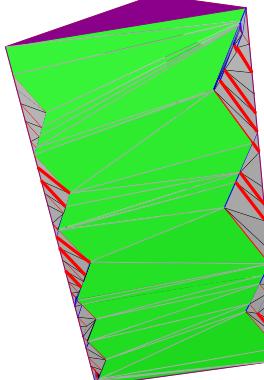
G_1

G_0

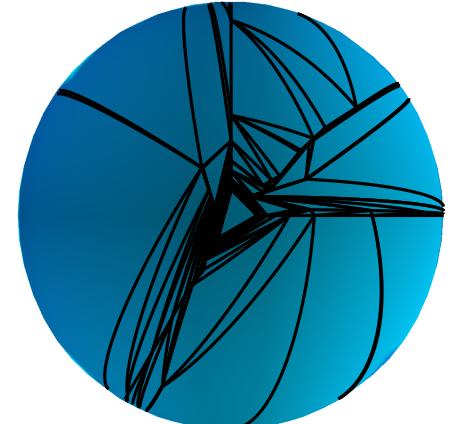
compute 3 **rivers** + **labeling** for each sub-region
(Duncan, Goodrich, Kobourov 2011)



each region G_i is decomposed by
the its river into 2 bottom and top
sub-regions



glue together the
rectangular drawings
to obtain a prism



perform central
projection

incremental shift algorithm (original FPP)

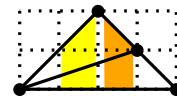
[de Fraysseix, Pollack, Pach'89]

use the canonical
ordering

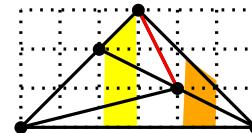
1.



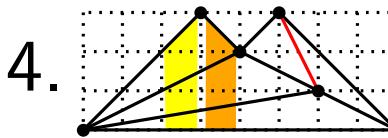
2.



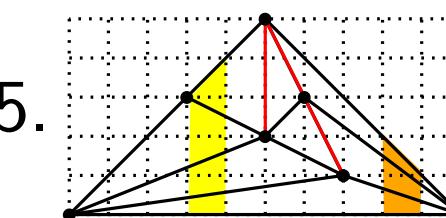
3.



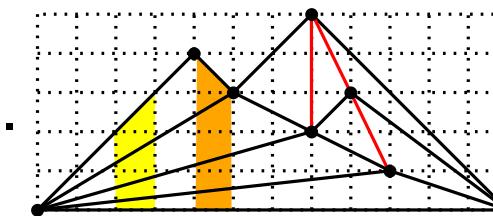
4.



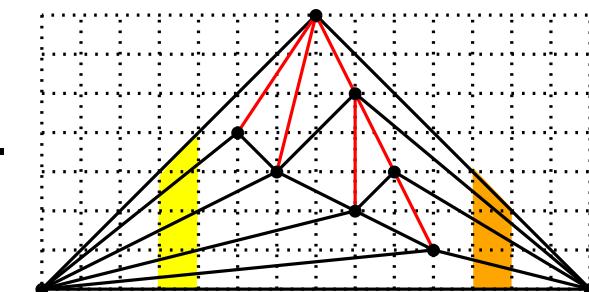
5.



6.



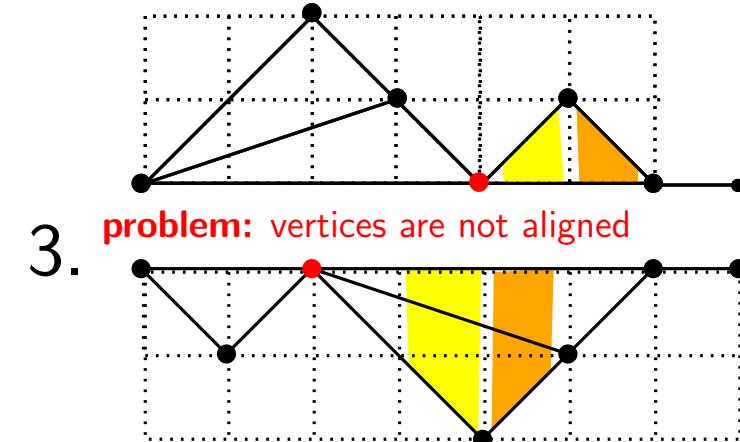
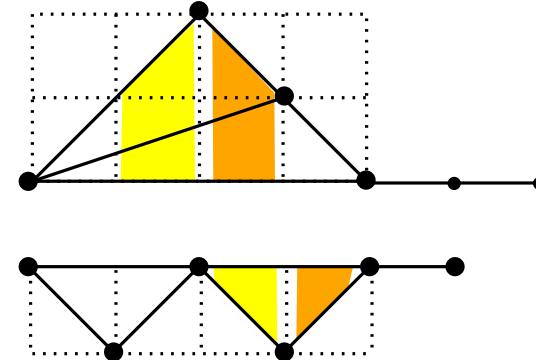
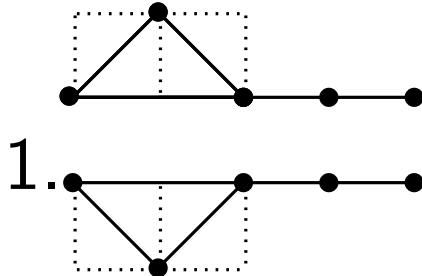
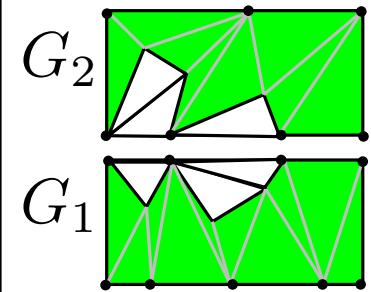
7.



Grid size of G_k : $2k \times k$

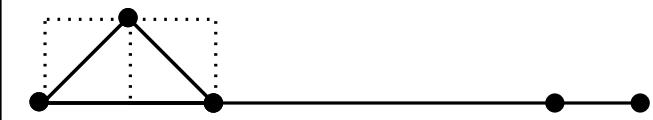
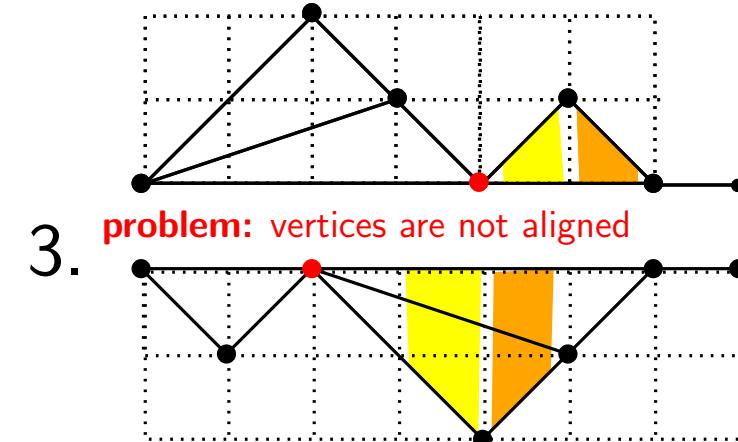
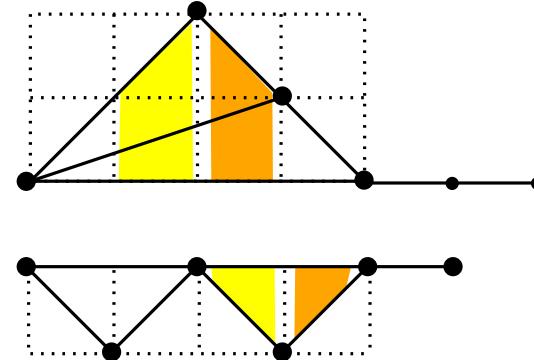
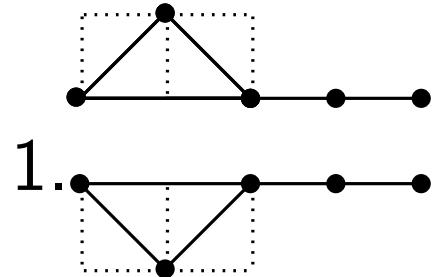
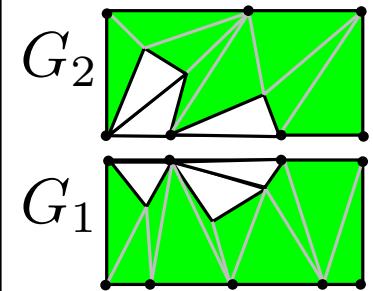
Drawing rectangular drawings with our adaptation of the FPP algorithm

(Castelli Aleardi, Fusy, Kostrygin 2014)
(Castelli Aleardi, Devillers, Fusy, 2012)

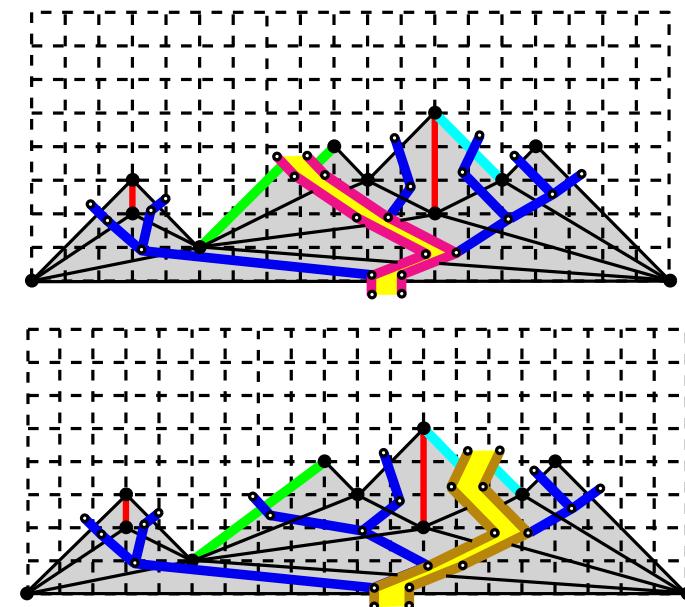
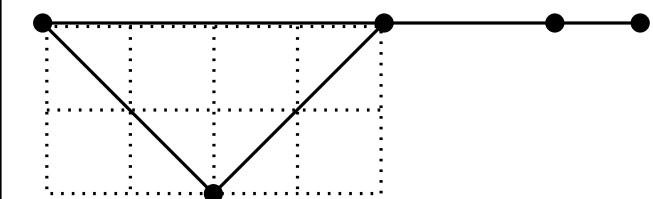


Drawing rectangular drawings with our adaptation of the FPP algorithm

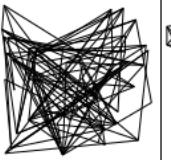
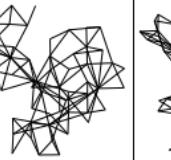
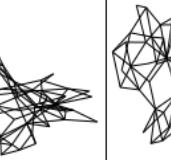
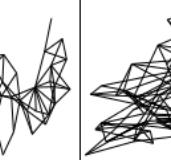
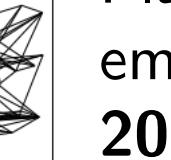
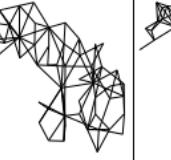
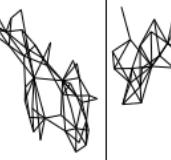
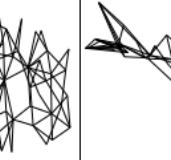
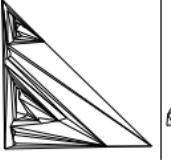
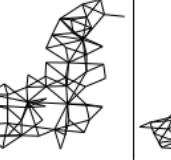
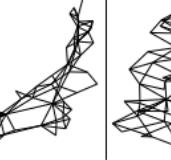
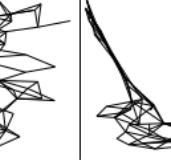
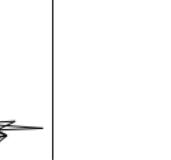
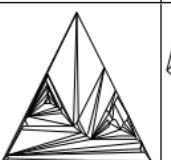
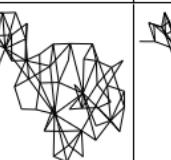
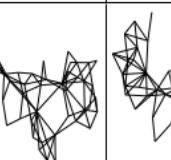
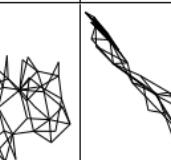
(Castelli Aleardi, Fusy, Kostrygin 2014)
(Castelli Aleardi, Devillers, Fusy, 2012)



1'. solution: perform a second pass, adding x -stretches



Experimental evaluation: use our spherical drawing as initial placer for iterative schemes

Initial Placer	Spring Embedder				
	None	FR	FRE	KK	SM
RANDOM					
	c=1349, ma=0.22, el=0.67, ns=0.75	c=113, ma=0.41, el=0.75, ns=0.80	c=124, ma=0.32, el=0.67, ns=0.78	c=52, ma=0.40, el=0.71, ns=0.81	c=314, ma=0.30, el=0.64, ns=0.66
FPP					
	c=0, ma=0.35, el=0.50, ns=0.55	c=71, ma=0.46, el=0.70, ns=0.86	c=50, ma=0.40, el=0.69, ns=0.80	c=58, ma=0.37, el=0.72, ns=0.82	c=156, ma=0.29, el=0.61, ns=0.61
SCHNYDER					
	c=0, ma=0.38, el=0.44, ns=0.57	c=53, ma=0.45, el=0.74, ns=0.83	c=59, ma=0.41, el=0.75, ns=0.79	c=69, ma=0.40, el=0.67, ns=0.77	c=68, ma=0.32, el=0.71, ns=0.71
KANT					
	c=0, ma=0.44, el=0.48, ns=0.59	c=58, ma=0.42, el=0.70, ns=0.79	c=57, ma=0.38, el=0.71, ns=0.77	c=52, ma=0.43, el=0.70, ns=0.76	c=168, ma=0.26, el=0.66, ns=0.61

Planar preprocessing for spring
embedders (**Fowler and Kobourov, 2012**)

Table 5: Layouts of a representative 55-node graph from the FUSY library.

Edge length statistic (Fowler and Kobourov, 2012)
(average percent deviation of edge length)
high values indicates more uniform edge length

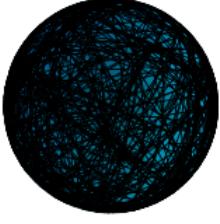
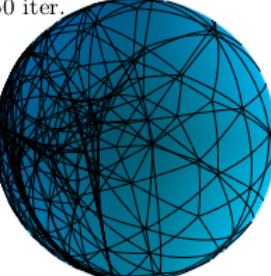
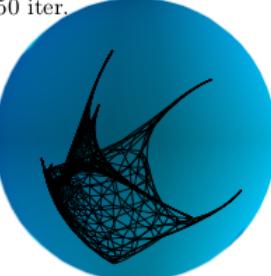
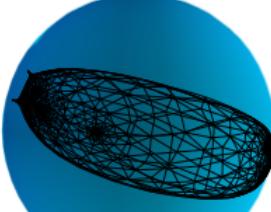
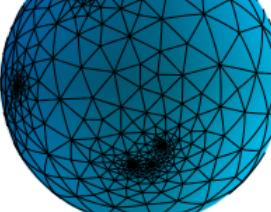
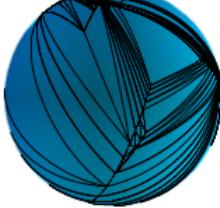
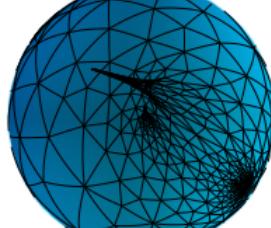
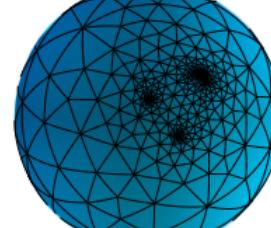
$$l(e) := \text{edge length of } e$$

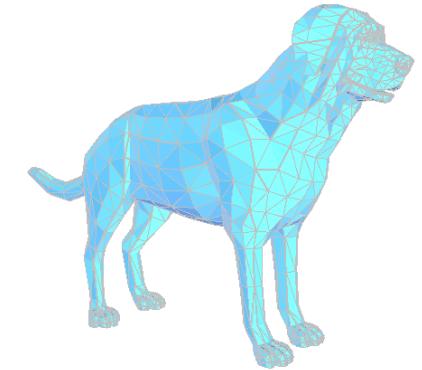
$$\epsilon_l := 1 - \left(\frac{1}{|E|} \sum_{e \in E} \frac{|l(e) - l_{avg}|}{\max(l_{avg}, l_{max} - l_{avg})} \right)$$

Comparison of initial placers

Comparison of convergence rates and aesthetic statistics

Dataset: 3D meshes from [aim@shape](#) repository

Placer	Initial layout	Projected Gauss-Seidel		Alexa method	
Random			50 iter.		250 iter.
PC			50 iter.		1600 iter.
ISP	$\rho = 0.42$ 		50 iter.		1190 iter.
SFPP			50 iter.		1058 iter.
		$\mathcal{E} = 47.92$	$\text{cl} = 0.89$	$\mathcal{E} = 12.26$	$\text{cl} = 0.91$
		$\mathcal{E} = 49.2$	$\text{cl} = 0.89$	$\mathcal{E} = 45.21$	$\text{cl} = 0.926$
		$\mathcal{E} = 51.37$	$\text{cl} = 0.864$	$\mathcal{E} = 48.02$	$\text{cl} = 0.790$
		1.73sec		987 iter.	1.06sec
				636 iter.	0.88sec
				1096 iter.	1.10sec
				$\mathcal{E} = 50.60$	$\text{cl} = 0.868$
				$\mathcal{E} = 48.02$	$\text{cl} = 0.801$

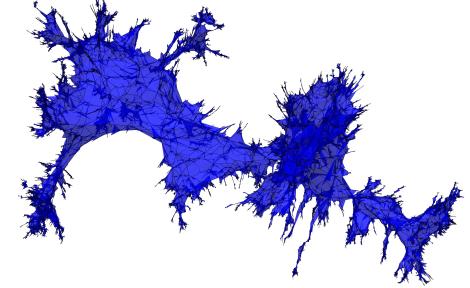
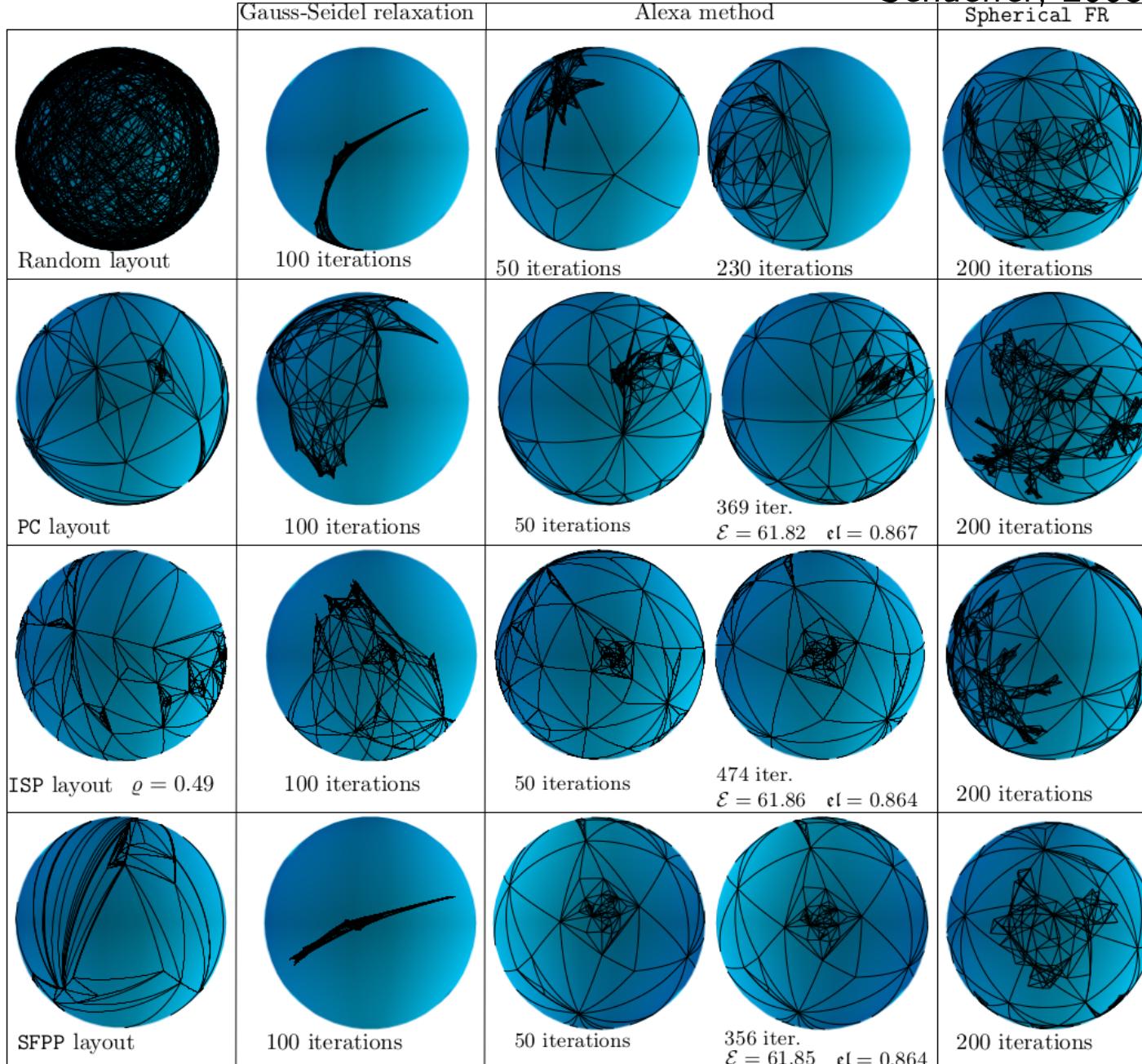


Comparison of initial placers

Comparison of convergence rates and aesthetic statistics

Dataset: random planar triangulations

(uniform random sampler by Poulalhon and Schaeffer, 2003)



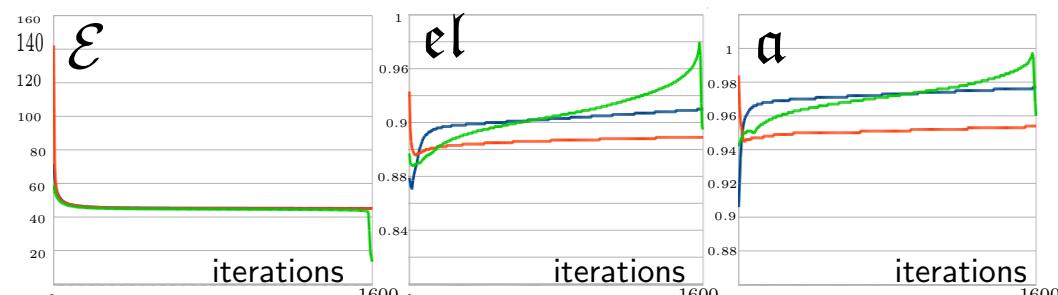


Comparison of initial placers

Comparison of convergence rates and aesthetic statistics

Placer	Initial layout	Projected Gauss-Seidel		Alexa method	
Random					
PC		50 iter. $\mathcal{E} = 47.92$ $\text{cl} = 0.89$	1600 iter. $\mathcal{E} = 12.26$ $\text{cl} = 0.91$	1.73sec 50 iter. $\mathcal{E} = 51.37$ $\text{cl} = 0.864$	987 iter. 1.06sec $\mathcal{E} = 48.02$ $\text{cl} = 0.790$
ISP	$\rho = 0.42$ 	50 iter. $\mathcal{E} = 49.2$ $\text{cl} = 0.89$	1190 iter. $\mathcal{E} = 45.21$ $\text{cl} = 0.926$	50 iter. $\mathcal{E} = 51.32$ $\text{cl} = 0.859$	636 iter. $\mathcal{E} = 47.82$ $\text{cl} = 0.849$
SFPP		50 iter. $\mathcal{E} = 48.8$ $\text{cl} = 0.89$	1058 iter. $\mathcal{E} = 45.02$ $\text{cl} = 0.908$	50 iter. $\mathcal{E} = 50.60$ $\text{cl} = 0.868$	1096 iter. $\mathcal{E} = 48.02$ $\text{cl} = 0.801$

projected Gauss-Seidel



(geodesic) edge length statistic

$$\text{el} := 1 - \left(\frac{1}{|E|} \sum_{e \in E} \frac{|l_g(e) - l_{avg}|}{\max(l_{avg}, l_{max} - l_{avg})} \right)$$

Area statistic

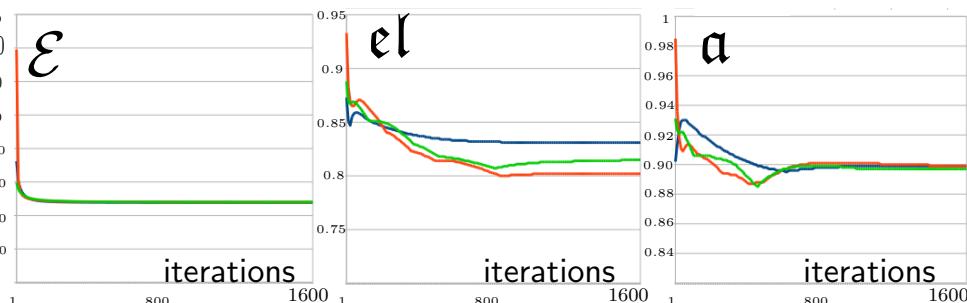
$$\text{al} := 1 - \left(\frac{1}{|F|} \sum_{t \in F} \frac{|a(f) - a_{avg}|}{\max(a_{avg}, a_{max} - a_{avg})} \right)$$

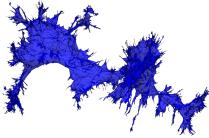
Spring energy

$$\mathcal{E} = \frac{1}{2} \sum_{i=1}^n \sum_{j \in N(i)} w_{ij} \|\mathbf{x}(v_i) - \mathbf{x}(v_j)\|^2$$

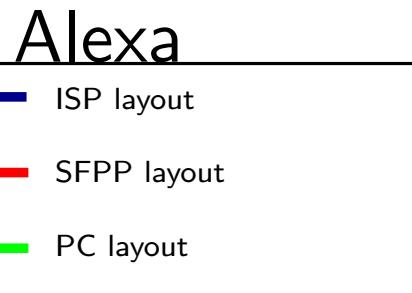
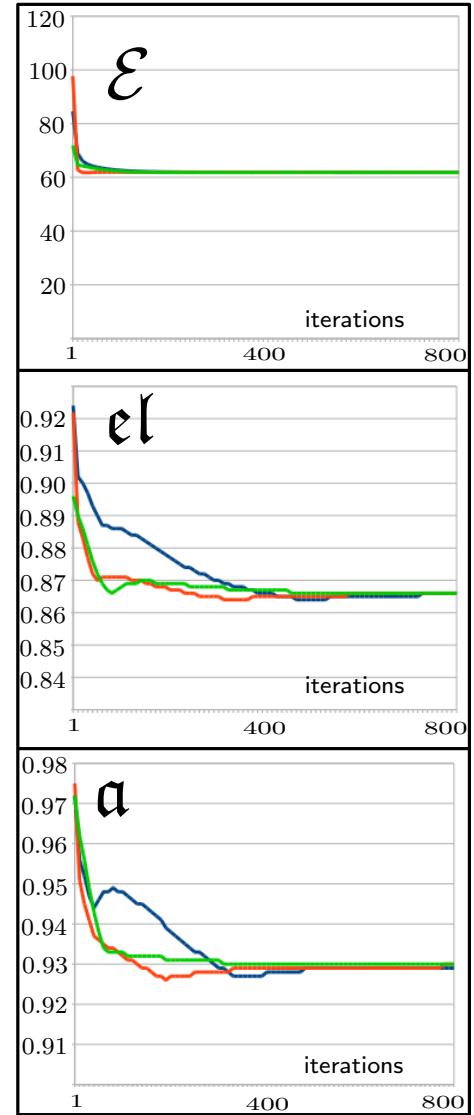
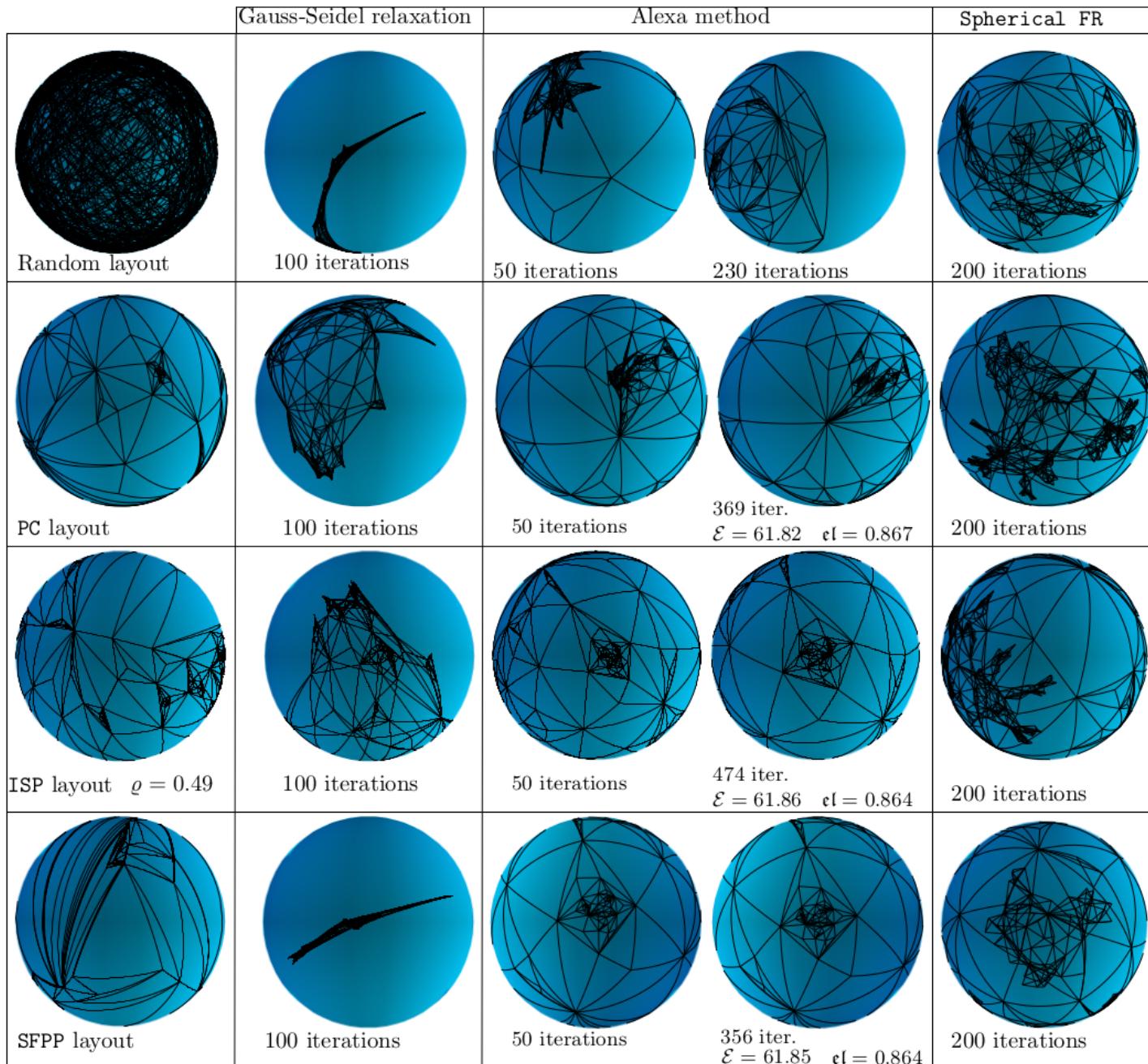
- ISP layout
- SFPP layout
- PC layout

Alexa





Comparison of convergence rates and aesthetic statistics

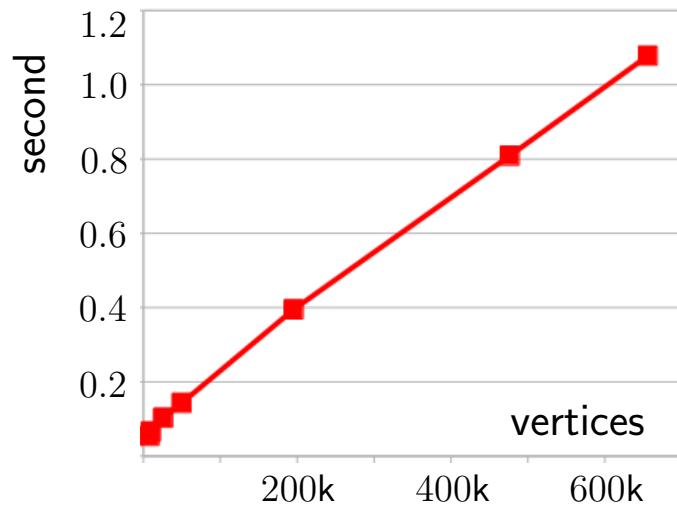


Comparison of initial placers

Comparison of timing performances

			our SFPP algorithm				Planar parameterizations	
mesh	vertices	faces	preprocessing		Layout computation		PC	ISP
			rivers comput.	canonical labeling	shift algorithm	prism projection		
Egea	8268	16K	0.015	0.017	0.005	0.017	0.24	0.16
Gargoyle	10002	20K	0.016	0.018	0.007	0.025	0.26	0.22
Bunny	26002	52K	0.017	0.031	0.019	0.036	1.14	0.75
Iphigenia	49922	99K	0.023	0.049	0.025	0.046	2.38	1.44
Camille's hand	195557	391K	0.076	0.121	0.073	0.125	17.02	7.92
Eros	476596	950K	0.162	0.260	0.132	0.255	50.54	29.99
Chinese dragon	655980	1.3M	0.174	0.314	0.157	0.433	89.64	53.12

solve sparse linear systems with the conjugate gradient solver of MTJ (Java) library (numeric precision 10^{-6})



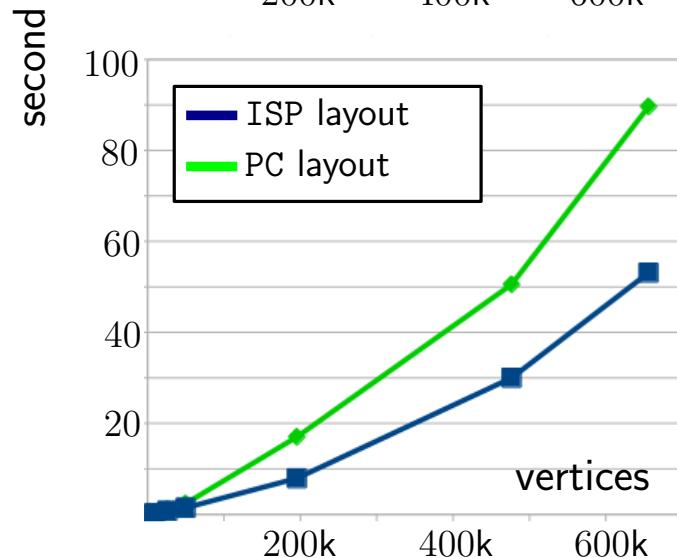
Dataset: 3D meshes from [aim@shape](#) repository

Spherical drawing of the
Chinese dragon (655K vert.)



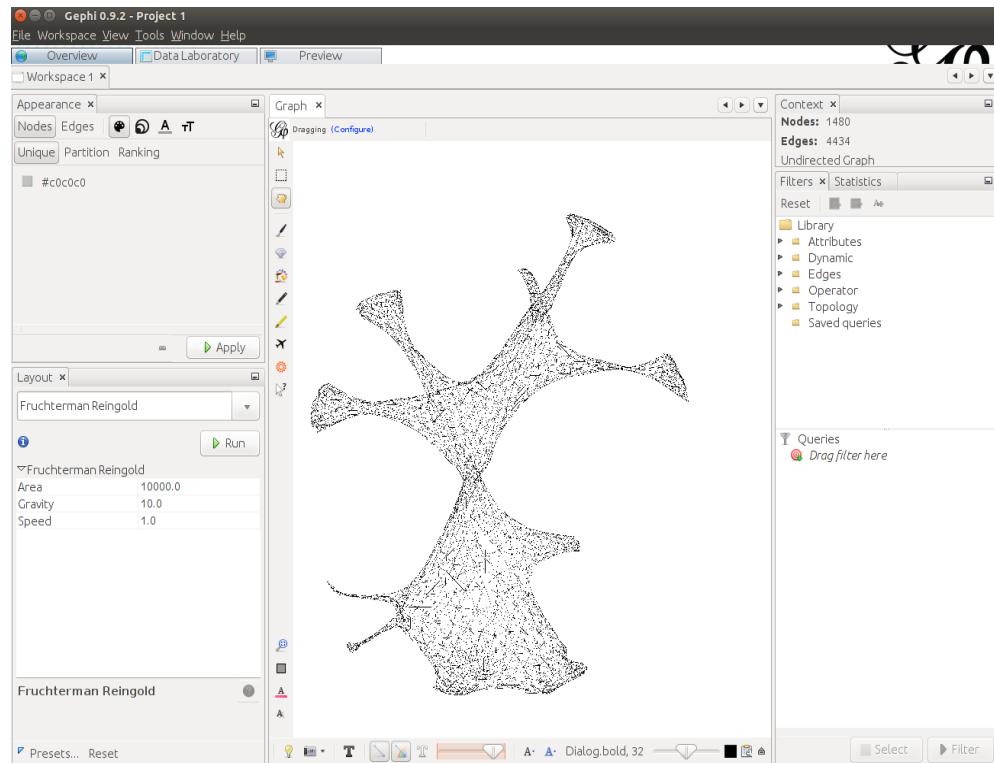
Aigerman Lipman (2015): 19 seconds (solving linear systems)
(Matlab, 3.50GHz Intel i7 CPU)

our SFPP algorithm: 1.07 seconds (total cost)
(Java, 2.66GHz Intel i7 CPU)

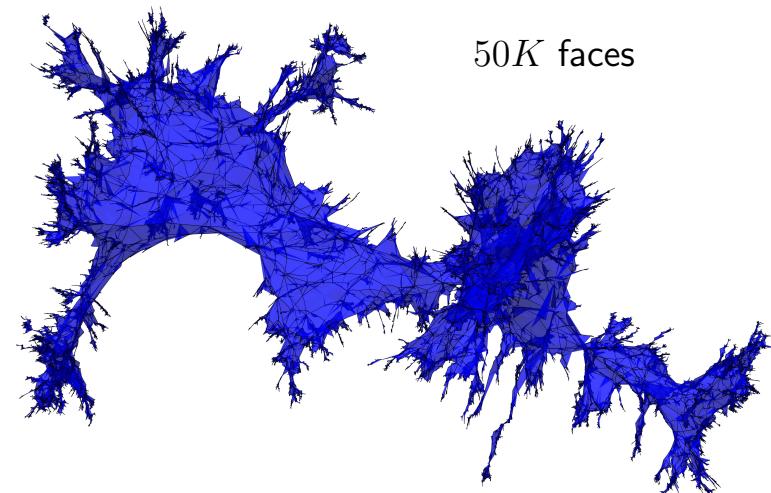


Application

spherical preprocessing for euclidean 3D spring embedders



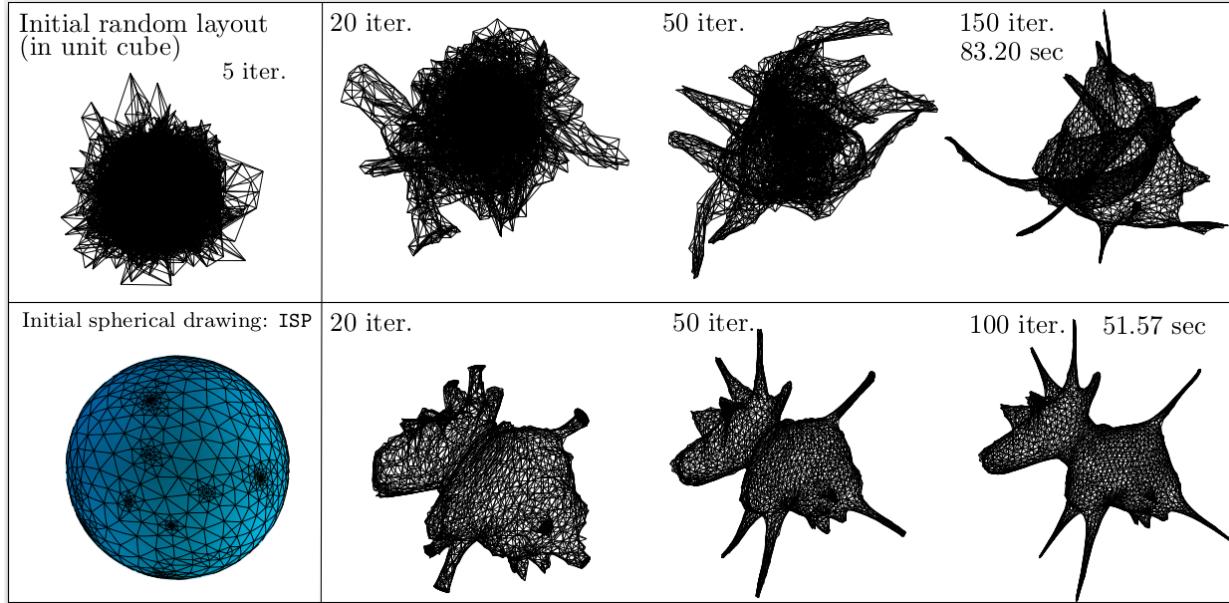
Yifan Hu layout of the dog graph (Gephi)



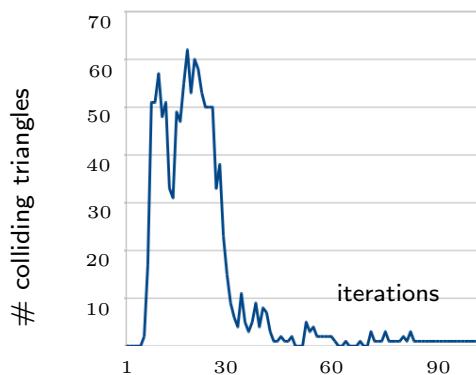
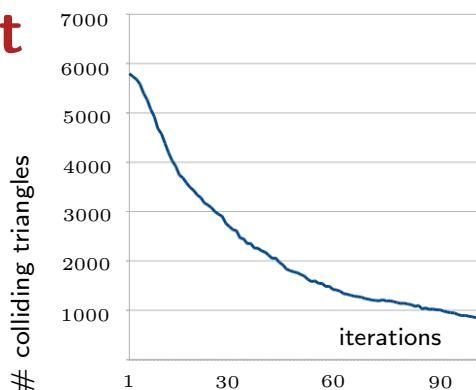
3D layout of a random planar map
by J. Bettinelli
(force-directed layout of Mathematica)

Spherical preprocessing for euclidean spring embedders

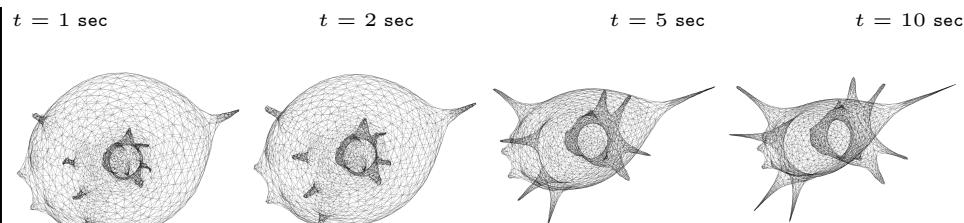
Use spherical drawings as initial layouts for 3D spring embedders: this allows us to better untangle the layout



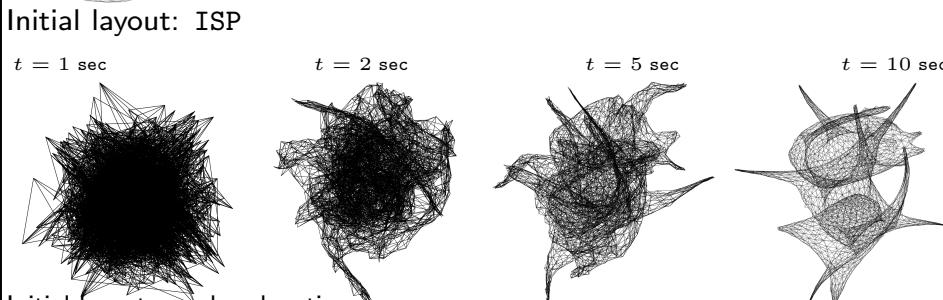
count triangle collisions



Our Java implementation of the FR91 spring embedder (exact computation of repulsive forces)



Gephi implementation of the Yifan Hu layout (fast approximate computation of repulsive forces)

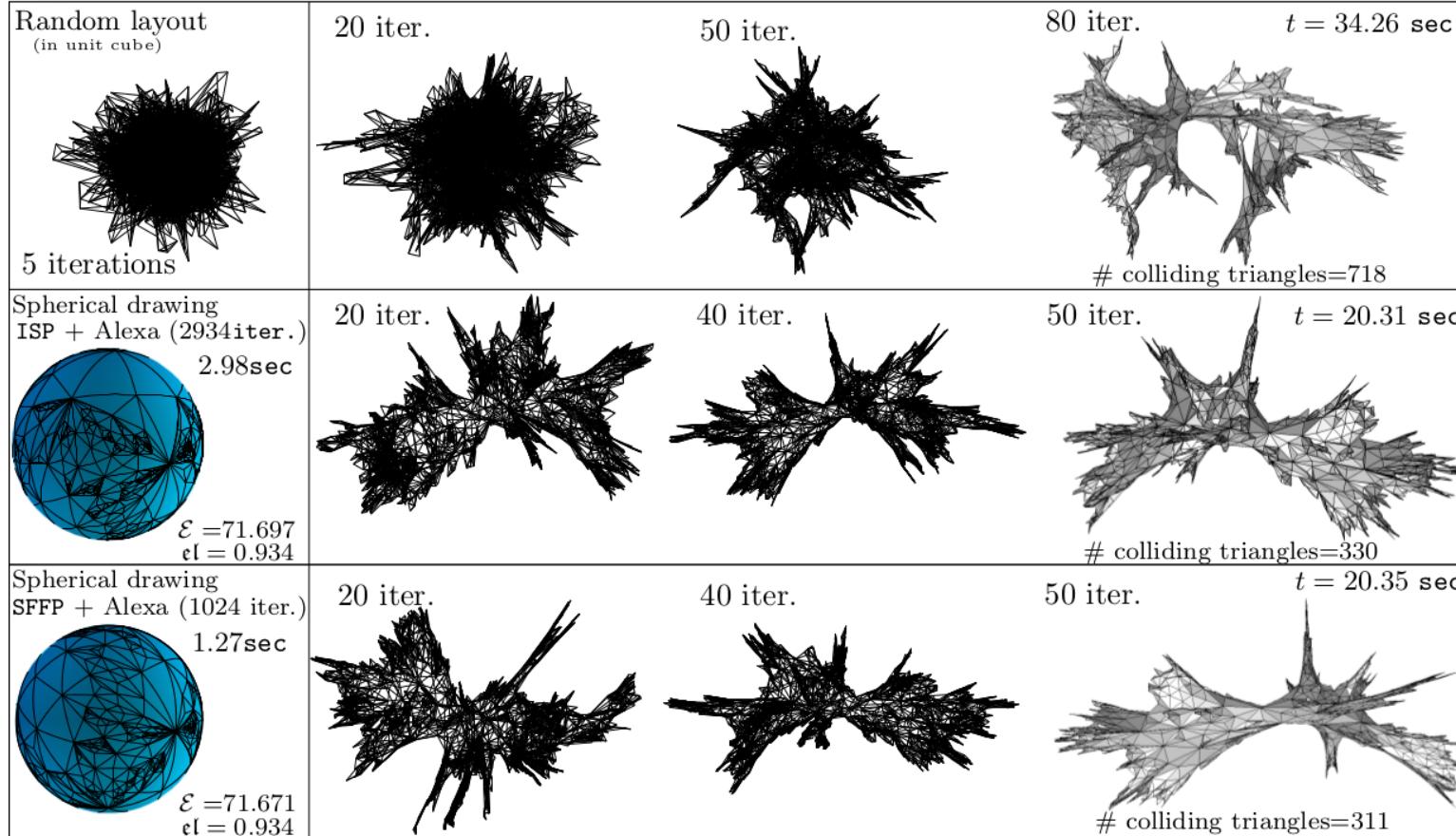


Spherical preprocessing for euclidean spring embedders

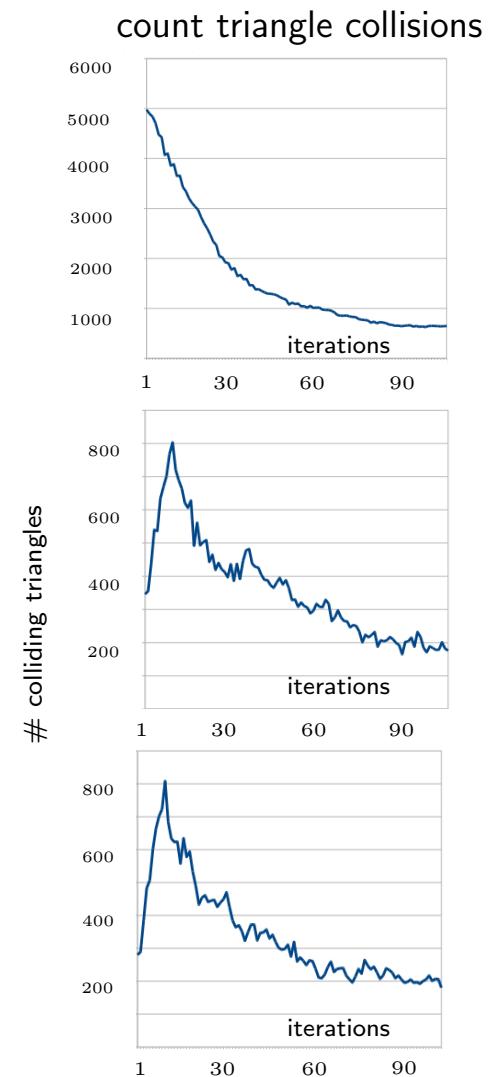
Use spherical drawings as initial layouts for 3D spring embedders: this allows us to better untangle the layout

random planar triangulation

with $5K$ triangles (generated with an uniform random sampler)



Layouts obtained with our Java implementation of the FR91 spring embedder
(exact computation of repulsive forces)



Thank you for your attention