

EXERCISES N° 2, DUALITY

Exercise 1. Let $C \subset \mathbb{F}_q^n$ be a code. Let $\mathcal{I} \subseteq \{1, \dots, n\}$. We define the following codes constructed from C :

- The punctured code on \mathcal{I} is defined as:

$$\mathcal{P}_{\mathcal{I}}(C) := \{(c_i)_{i \in \mathcal{I}} \mid c \in C, \} \subseteq \mathbb{F}_q^{|\mathcal{I}|}.$$

Roughly speaking, it is the set of codewords of C where the positions out of \mathcal{I} are removed.

- The shortened code on \mathcal{I} is defined as:

$$\mathcal{S}_{\mathcal{I}}(C) := \{(c_i)_{i \in \mathcal{I}} \mid c \in C, \forall i \notin \mathcal{I}, c_i = 0\} \subseteq \mathbb{F}_q^{|\mathcal{I}|}.$$

It is the set of codewords supported by \mathcal{I} which is punctured at \mathcal{I}

Prove that $(\mathcal{P}_{\mathcal{I}}(C))^{\perp} = \mathcal{S}_{\mathcal{I}}(C^{\perp})$ and $(\mathcal{S}_{\mathcal{I}}(C))^{\perp} = \mathcal{P}_{\mathcal{I}}(C^{\perp})$

Exercise 2. Let $\mathbb{F}_{q^m}/\mathbb{F}_q$ be an extension of finite fields. Recall that the *trace* of $\mathbb{F}_{q^m}/\mathbb{F}_q$ is defined as:

$$\mathrm{Tr}_{\mathbb{F}_{q^m}/\mathbb{F}_q} : \begin{cases} \mathbb{F}_{q^m} & \longrightarrow & \mathbb{F}_q \\ x & \longmapsto & x + x^q + x^{q^2} + \dots + x^{q^{m-1}} \end{cases} .$$

- (1) Prove that this map is an \mathbb{F}_q -linear form over \mathbb{F}_{q^m} .
- (2) Prove that this map is surjective.

Indication: use the fact that the polynomial $X + X^q + \dots + X^{q^{m-1}}$ cannot have q^m roots.

- (3) Prove that the map

$$\begin{cases} \mathbb{F}_{q^m} \times \mathbb{F}_{q^m} & \longrightarrow & \mathbb{F}_q \\ (x, y) & \longmapsto & \mathrm{Tr}_{\mathbb{F}_{q^m}/\mathbb{F}_q}(xy) \end{cases}$$

is \mathbb{F}_q -bilinear, symmetric and non degenerated.

- (4) Deduce from the previous question that for all linear form $\varphi : \mathbb{F}_{q^m} \rightarrow \mathbb{F}_q$, there exists a unique $a_{\varphi} \in \mathbb{F}_{q^m}$ such that

$$\forall x \in \mathbb{F}_{q^m}, \varphi(x) = \mathrm{Tr}_{\mathbb{F}_{q^m}/\mathbb{F}_q}(a_{\varphi}x).$$

- (5) Let $C \subseteq \mathbb{F}_{q^m}^n$, we recall the definitions of subfield subcodes and trace codes:

$$\begin{aligned} C_{|\mathbb{F}_q} &:= C \cap \mathbb{F}_q^n \\ \mathrm{Tr}(C) &:= \{(\mathrm{Tr}_{\mathbb{F}_{q^m}/\mathbb{F}_q}(c_1), \dots, \mathrm{Tr}_{\mathbb{F}_{q^m}/\mathbb{F}_q}(c_n)) \mid c \in C\}. \end{aligned}$$

Prove that we always have $C_{|\mathbb{F}_q} \subseteq \mathrm{Tr}(C)$.

Indication: Because of the surjectivity of $\mathrm{Tr}_{\mathbb{F}_{q^m}/\mathbb{F}_q}$, there exists $\gamma \in \mathbb{F}_{q^m}$ such that $\mathrm{Tr}_{\mathbb{F}_{q^m}/\mathbb{F}_q}(\gamma) = 1$.

Exercise 3. ★

Prove additive Hilbert's 90 Theorem for finite fields:

$$\forall x \in \mathbb{F}_{q^m}, \mathrm{Tr}_{\mathbb{F}_{q^m}/\mathbb{F}_q}(x) = 0 \iff \exists a \in \mathbb{F}_{q^m}, x = a^q - a.$$

Exercise 4. ★

The goal of this exercise is to prove Delsarte's Theorem: For all code $C \subseteq \mathbb{F}_{q^m}^n$,

$$(C_{|\mathbb{F}_q})^\perp = \text{Tr}(C^\perp).$$

- (1) Prove inclusion " \supseteq ".
- (2) To prove the converse inclusion, we will prove the equivalent one:

$$(\text{Tr}(C^\perp))^\perp \subseteq C_{|\mathbb{F}_q}.$$

For that we assume this inclusion to be wrong and take $y \in (\text{Tr}(C^\perp))^\perp \setminus C_{|\mathbb{F}_q}$.

- (a) Regarding y as an element of $\mathbb{F}_{q^m}^n$ (instead of \mathbb{F}_q^n), prove the existence of $x \in C^\perp$ such that $\langle x, y \rangle_{\mathbb{F}_{q^m}} \neq 0$.
- (b) Prove the existence of $\gamma \in \mathbb{F}_{q^m}$, such that

$$\text{Tr}_{\mathbb{F}_{q^m}/\mathbb{F}_q}(\gamma \langle x, y \rangle_{\mathbb{F}_{q^m}}) \neq 0.$$

- (c) Prove that $\langle \text{Tr}_{\mathbb{F}_{q^m}/\mathbb{F}_q}(\gamma x), y \rangle_{\mathbb{F}_q} \neq 0$.
- (d) Conclude.

- (3) Prove that if C is $[n, k, d]_{q^m}$ then $C_{|\mathbb{F}_q}$ is $[n, \geq n - m(n - k), \geq d]_q$.

Exercise 5. Let C be the binary Hamming code with parity-check matrix

$$\begin{pmatrix} 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 \end{pmatrix}$$

- (1) Prove that C is $[7, 4, 3]_2$.
- (2) Prove that $(1\ 1\ 1\ 1\ 1\ 1\ 1) \in C$ and deduce that the weight enumerator $P_C^\sharp(x, y)$ is symmetric: $P_C^\sharp(x, y) = P_C^\sharp(y, x)$.
- (3) Using McWilliams' identity, compute the polynomials P_C^\sharp and $P_{C^\perp}^\sharp$ without enumerating the codes.