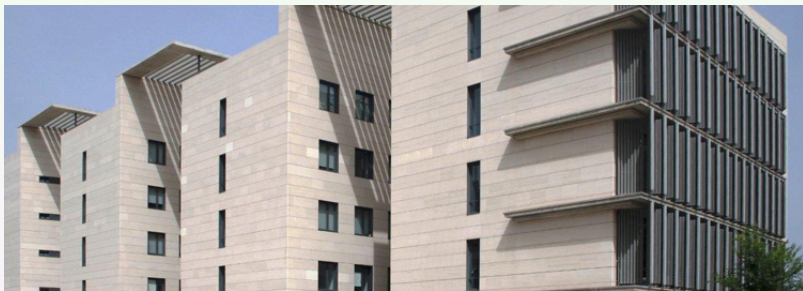


## Connectivity degeneracy measures

Dimitrios M. Thilikos



AIGCo project team, CNRS, LIRMM

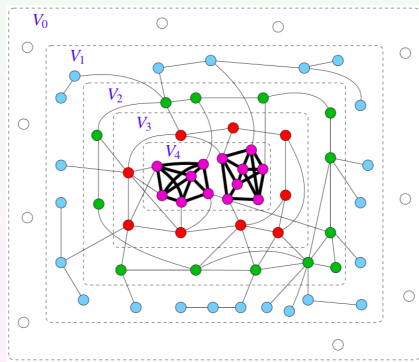
ESIGMA Project Kick-off meeting, May–June 2018

LAMSADE, Paris, May 31, 2018

$k$ -core( $G$ ) = maximum subgraph  $H$  where  $\delta(H) \geq k$ .

Core decomposition:  $V_0, V_1, \dots, V_k$ ,

$V_i = i$ -core( $G$ )  $\setminus$   $(i - 1)$ -core( $G$ )



Graph degeneracy:  $\delta^*(G) = \max\{k \mid k$ -core( $G$ ) is non-trivial}

Core decomposition: graph hierarchization

Graph degeneracy:  $\delta^*(G) = \max\{k \mid k\text{-core}(G) \text{ is non-trivial}\}$

[P.Erdős. On the structure of linear graphs. Israel J. Math., 1:156–160, 1963.]

[S.Seidman. Network structure and minimum degree. Social Networks, 5(3):269–287, 1983.]

[R. Anderson & E. Mayr. Parallelism and greedy algorithms, Adv. Comput. Res., 4 (1987) pp. 17-38;]

Good news:  $O(kn)$  algorithm to compute the core decomposition:

▶ Remove vertices of minimum degree... the degrees determine the core levels.

Graph degeneracy as a measure of **collaboration**:

just being a *hub* is not enough in order to be "collaborative"

Other variants of (degree) degeneracy:

Directed graphs:

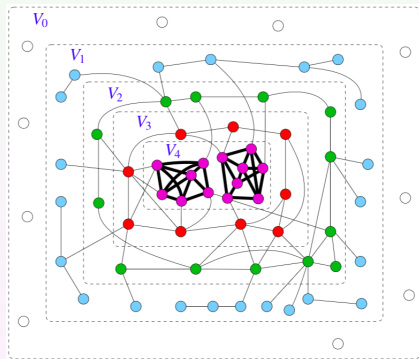
[C. Giatsidis, D.M. Thilikos, M. Vazirgiannis: D-cores: measuring collaboration of directed graphs based on degeneracy. Knowl. Inf. Syst. 35(2): 311-343 (2013)]

Signed graphs:

[C. Giatsidis, B. Cautis, S. Maniu, D.M. Thilikos, M.Vazirgiannis: Quantifying trust dynamics in signed graphs, the S-Cores approach. SDM 2014: 668-676]

## Clustering:

[C. Giatsidis, F. D. Malliaros, D. M. Thilikos, M. Vazirgiannis: CoreCluster: A Degeneracy Based Graph Clustering Framework. AAAI 2014: 44-50]



Back to the definition:

$k$ -core( $G$ ) = maximum subgraph  $H$  where  $\delta(H) \geq k$ .

Where  $\delta(H) = \min\{\mathbf{deg}_H(v) \mid v \in V(H)\}$

Other definitions of  $\mathbf{deg}_H(v)$ ?

For instance:  $\mathbf{deg}_H(v) = \#$  triangles containing  $v$ :  
(i.e.,  $|E(G[N_H(v)])|$ )

[F.D. Malliaros, M.G. Rossi, and Michalis Vazirgiannis. Locating influential nodes in complex networks. Scientific reports, Nature Publishing Group, 2016.]

Another example:  $\mathbf{deg}_H(v) = \#$  edge-disjoint triangles on  $v$ :  
(i.e., max matching in  $G[N_H(v)]$ ). Better? ... still polynomial!

Which “degree” function  $\mathbf{deg}_H(v)$  define good “core decompositions”?

- ▶ Combinatorially Good:  $i\text{-core} \subseteq (i - 1)\text{-core}$ . (hierarchization)
- ▶ Algorithmically Good: **polynomially** computable.  
(*polynomially?*)

A major advance:

[V. Batagelj, M. Zaversnik: Fast algorithms for determining (generalized) core groups in social networks. *Adv. Data Analysis and Classification* 5(2): 129-145 (2011)]

Let  $H \subseteq G$

$\deg_H(v)$  is *monotone* if:

$H_1 \subseteq H_2 \Rightarrow \forall v \in V : \deg_{H_1}(v) \leq \deg_{H_2}(v)$

$\deg_H(v)$  is *local* if:  $\deg_H(v) = \deg_{N_H[v]}(v)$

Theorem

If  $\deg_H(v)$  is *monotone* then  $\delta^*$  is hierarchical

Theorem

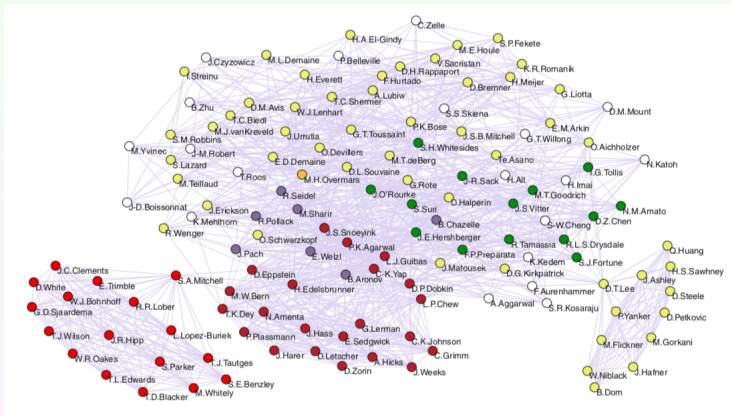
If  $\deg_H(v)$  is *local* then  $\delta^*$  can be computed in

$O(m \cdot \max(\Delta, \log n))$  steps

not just polynomial!



## Some reflection on (degree) cores...



The 10-core of the “Computational geometry” network

► It might be quite *disconnected*! Degree Locality is not enough!

## Connectivity degeneracy

$$S \subseteq H \subseteq G$$

edge connectivity of a vertex set  $S$  of a subgraph  $H$ :

$$\lambda_H(S) = \#\text{cut-edges in } H \text{ between } S \text{ \& } H \setminus S$$

$$H \subseteq G$$

edge connectivity of a subgraph  $H$ :

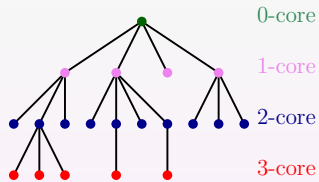
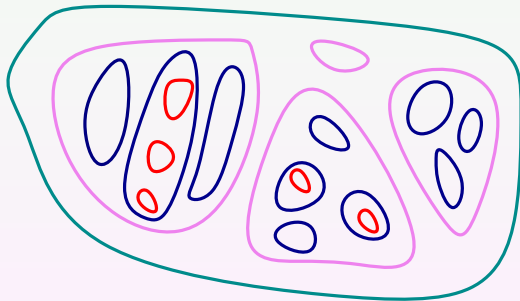
$$\lambda(H) = \min\{\lambda_H(S) \mid \emptyset \subsetneq S \subsetneq H\}$$

(edge) connectivity degeneracy of  $G$ :

$$\lambda^*(G) = \max\{\lambda(H) \mid H \subseteq G\}$$

What is a  $k$ - $\lambda$ -core?

A partition of  $G$  so that each part induces a maximal  $k$ -edge-connected subgraph.



Connectivity degeneracy vs degree degeneracy:

$$\delta^*(G) \leq \lambda^*(G) \leq 2 \cdot \delta^*(G)$$

[W. Mader, Existenz  $n$ -fach zusammenhängender Teilgraphen in Graphen genügend großer Kantendichte, Abh. Math. Sem Univ. Hamburg 37, 86–97 (1972).]

A “**polynomial**” algorithm (also parallel approximation):

[L. Kirousis, M. Serna, P. Spirakis, The parallel complexity of the subgraph connectivity problem, SIAM J. Comput., 22 (1993), pp. 573–586.]

More on connectivity degeneracy in the next talk...

Which “cut” functions  $\lambda_H(S)$  define good “core decompositions”?

▶ Combinatorially Good:

$i$ -core is a **refinement** of  $(i - 1)$ -core. (hierarchization)

▶ Algorithmically Good: **polynomially** computable.

(*polynomially?*)

Consider a universe  $U$  and a function  $\lambda : U \rightarrow \mathbb{N}$ .

We say that  $\lambda$  is a *connectivity function on  $U$*  if the following hold:

- ▶  $\lambda(\emptyset) = 0$ .
- ▶  $\forall X \subseteq U, \lambda(X) = \lambda(U \setminus X)$  (symmetry).
- ▶  $\forall X, Y \subseteq U, \lambda(X \cup Y) + \lambda(X \cap Y) \leq \lambda(X) + \lambda(Y)$   
(submodularity).

We say a family  $\{\lambda_S : S \subseteq U\}$  of functions  $\lambda_S : 2^S \rightarrow \mathbb{N}$  is *subset monotone* if  $\lambda_{S'}(X) \leq \lambda_S(X)$  whenever  $S' \subseteq S$ .

## An abstract notion of connectivity degeneracy

Let  $U$  be a finite set (a “universe”).

Given a family  $\mathcal{F} = \{\lambda_S : S \subseteq U\}$  of functions  $\lambda_S : 2^S \rightarrow \mathbb{N}$ , we define

$$\lambda_{\mathcal{F}}^*(U) = \max\{\min\{\lambda_S(X) \mid \emptyset \subsetneq X \subsetneq S\} \mid S \subseteq U\}$$

### Our results:

#### Theorem

If  $\mathcal{F}$  is a subset-monotone family, then  $\lambda_{\mathcal{F}}^*$  is *hierarchical*.

#### Theorem

If  $\mathcal{F}$  is a family of connectivity functions, then  $\lambda_{\mathcal{F}}^*$  can be computed in *polynomial time*.

## Conclusions and research directions

- ▶ We set up a theoretical framework for connectivity degeneracy.
- ▶ The theorems of [Batagelj & Zaversnik] follow as the degree analogues of our results.
- ▶ We consider several connectivity functions/measures fitting in the above framework.
- ▶ We study  $\lambda_H(S) = \max \#$  edge-disjoint paths between  $S$  and  $H \setminus S$ .
- ▶ We work on an analogous framework on vertex-connectivity (clusters may overlap).
- ▶ We work on optimizing the running times of the polynomial algorithms for certain interesting instantiations of  $\lambda$
- ▶ We work on clustering problems under alternative connectivity measures.



Ongoing project Involving:

- ▶ EunJung Kim
- ▶ Christophe Paul
- ▶ Mark Jones

and

- ▶ Joannie Perret
- ▶ Stratis Limnios

and **hopefully more!**

Merci!

