

TD1 – Cartesian categories

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1 Categories and functors

1. Recall the definition of *category* and provide some examples (e.g. **Set**, **Top**, **Vect**, **Grp**).
2. Recall the definition of a *functor* and provide some examples.
3. Define the category **Cat** of categories and functors.

2 Cartesian categories

Suppose fixed a category \mathcal{C} . A *cartesian product* of two objects A and B is given by an object $A \times B$ together with two morphisms

$$\pi_1 : A \times B \rightarrow A \quad \text{and} \quad \pi_2 : A \times B \rightarrow B$$

such that for every object C and morphisms $f : C \rightarrow A$ and $g : C \rightarrow B$, there exists a unique morphism $h : C \rightarrow A \times B$ making the diagram

$$\begin{array}{ccc} & C & \\ & \downarrow h & \\ & A \times B & \\ \swarrow f & & \searrow g \\ A & & B \\ \swarrow \pi_1 & & \searrow \pi_2 \end{array}$$

commute. We also recall that a *terminal object* in a category is an object 1 such that for every object A there exists a unique morphism $f_A : A \rightarrow 1$. A category is *cartesian* when it has finite products, i.e. has a terminal object and every pair of objects admits a product.

1. Suppose that (E, \leq) is a poset. We associate to it category whose objects are elements of E and such that there exists a unique morphism between object a and b iff $a \leq b$. What is a terminal object and a product in this category?

Solution. A terminal object is an element $b \in E$ such that, for every other element $a \in E$, there is a unique morphism $a \rightarrow b$, i.e. $a \leq b$. A terminal object is thus precisely a maximal element of the set.

Given two elements a and b of E , a product is an element $a \times b$ equipped with two morphisms $\pi_1 : a \times b \rightarrow a$ and $\pi_2 : a \times b \rightarrow b$ such that for every element c equipped with two morphisms $f : c \rightarrow a$ and $f : c \rightarrow b$ there exists a unique morphism $h : c \rightarrow a \times b$ making some diagrams commute. This is thus precisely an element $a \times b$ such that $a \times b \leq a$ and $a \times b \leq b$ such that for every element c with $c \leq a$ and $c \leq b$, we have $c \leq a \times b$. Otherwise said, $a \times b$ is an infimum of a and b .

2. Show that the category **Set** of sets and functions is cartesian.

Solution. Given two sets A and B , we define

$$A \times B = \{(a, b) \mid a \in A \text{ and } b \in B\}$$

and

$$\begin{array}{ll} \pi_1 : A \times B \rightarrow A & \pi_2 : A \times B \rightarrow B \\ (a, b) \mapsto a & (a, b) \mapsto b \end{array}$$

Given a set C and functions $f : C \rightarrow A$ and $g : C \rightarrow B$, suppose that there exists a function $h : C \rightarrow A \times B$ making the following diagram commute:

$$\begin{array}{ccc} & C & \\ & \downarrow h & \\ & A \times B & \\ \swarrow f & & \searrow g \\ A & & B \\ \swarrow \pi_1 & & \searrow \pi_2 \end{array}$$

Given $c \in C$, we have $h(c) \in A \times B$, i.e. $h(c)$ is of the form $h(c) = (a, b)$. The commutation of the left triangle imposes

$$a = \pi_1(a, b) = \pi_1 \circ h(c) = f(c)$$

and the one of the right that $b = g(c)$. Therefore, necessarily, we have $h(c) = (f(c), g(c))$ (if h exists it is unique). Conversely, the function h thus defined makes the two triangle commutes (h actually exists).

3. Show that two terminal objects in a category are necessarily isomorphic.

Solution. Suppose given two terminal objects A and B . Since B is terminal, we have a unique morphism $f : A \rightarrow B$ and, since A is terminal, we have a unique morphism $g : B \rightarrow A$.

$$\text{id}_A \hookrightarrow A \begin{array}{c} \xrightarrow{f} \\ \xleftarrow{g} \end{array} B \hookleftarrow \text{id}_B$$

The composite $g \circ f : A \rightarrow A$ and $\text{id}_A : A \rightarrow A$ are both morphisms with the same source and A as target: since A is terminal, we therefore have $g \circ f = \text{id}_A$. Similarly, we have $f \circ g = \text{id}_B$ and we deduce that A and B are isomorphic.

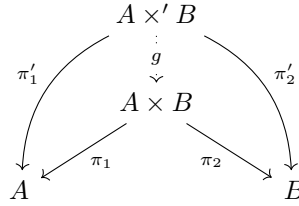
4. Similarly, show that the cartesian product of two objects is defined up to isomorphism.

Solution. Fix two objects A and B and suppose that they admit two products $(A \times B, \pi_1, \pi_2)$ and $(A \times' B, \pi'_1, \pi'_2)$. From the following diagram,

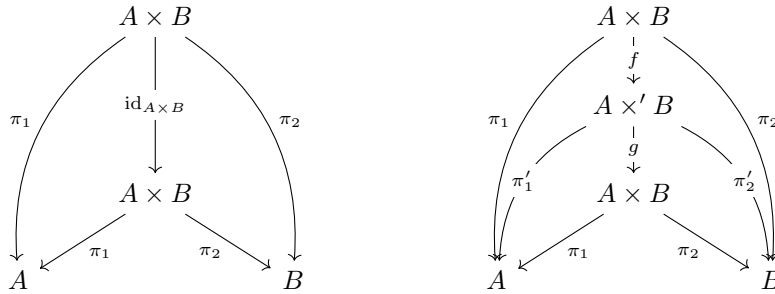
$$\begin{array}{ccc} & A \times B & \\ & \downarrow f & \\ & A \times' B & \\ \swarrow \pi_1 & & \searrow \pi_2 \\ A & & B \\ \swarrow \pi'_1 & & \searrow \pi'_2 \end{array}$$

by definition of the product $A \times' B$, we deduce the existence of a unique morphism $f : A \times B \rightarrow A \times' B$ such that $\pi'_1 \circ f = \pi_1$ and $\pi'_2 \circ f = \pi_2$. Similarly, there exists a unique morphism

$g : A \times' B \rightarrow A \times B$ such that $\pi_1 \circ g = \pi'_1$ and $\pi_2 \circ g = \pi'_2$:



Now, we have two morphisms from $A \times B$ to $A \times B$, namely $\text{id}_{A \times B}$ and $g \circ f$:



and those make the two triangles commute: we have

$$\pi_1 \circ \text{id}_{A \times B} = \pi_1$$

$$\pi_2 \circ \text{id}_{A \times B} = \pi_2$$

and

$$\pi_1 \circ g \circ f = \pi'_1 \circ f = \pi_1$$

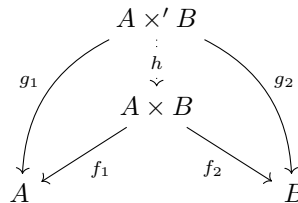
$$\pi_2 \circ g \circ f = \pi'_2 \circ f = \pi_2$$

Therefore by universal property of the product, we have $g \circ f = \text{id}_{A \times B}$. Similarly, we have $f \circ g = \text{id}_{A \times' B}$ and thus $A \times B$ and $A \times' B$ are isomorphic.

5. How could you show previous question using question 3.?

Solution. Suppose fixed two objects A and B of our ambient category \mathcal{C} . The idea is to construct another category \mathcal{D} in which a terminal object is precisely a product of A and B . We therefore define the category

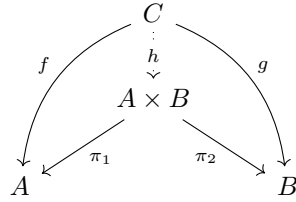
- whose objects are triple (C, f_1, f_2) where C is an object of \mathcal{C} and $f_1 : C \rightarrow A$ and $f_2 : C \rightarrow B$ are morphisms of \mathcal{C} ,
- a morphism in $\mathcal{D}((C, f_1, f_2), (D, g_1, g_2))$ is a morphism $h : C \rightarrow D$ of \mathcal{C} such that $g_1 \circ h = f_1$ and $g_2 \circ h = f_2$:



- identities are identities of \mathcal{C} ,
- composition is the same as in \mathcal{C} .

(We leave the reader check the composite is well-defined, i.e. that the composite of two morphisms is still a morphism, and that identities are actually morphisms). A terminal object in \mathcal{D} is an object $(A \times B, \pi_1, \pi_2)$ such that for every other object (C, f, g) there is a

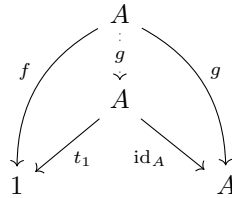
unique morphism $h : (C, f, g) \rightarrow (A \times B, \pi_1, \pi_2)$ in \mathcal{D} , i.e. there exists a unique morphism $h : C \rightarrow A \times B$ in \mathcal{C} making the following diagram commute:



i.e. $A \times B$ is a product of A and B . By question 3, two such objects are isomorphic in \mathcal{D} , and they will thus be isomorphic in \mathcal{C} since composition and identities in \mathcal{D} are induced by those of \mathcal{C} .

6. Show that for every object A of a cartesian category, the objects $1 \times A$, A and $A \times 1$ are isomorphic.

Solution. Of course the same reasoning “by hand” as above can be performed here. Another way to proceed in order to show that $1 \times A$ and A are isomorphic is to show that (A, t_A, id_A) is a product of 1 and A (where $t_A : 1 \rightarrow A$ is the terminal map) and conclude by question 4. Namely, given two morphisms $f : C \rightarrow 1$ and $g : C \rightarrow A$, we have the morphism g which make the following diagram commute:



The left triangle commutes because 1 is terminal and therefore $\pi_1 \circ g = f$, and the right triangle commutes by definition of identities: $\text{id}_A \circ g = g$. Conversely, g is the only such morphism by commutation of the right triangle.

7. Show that for every objects A and B , $A \times B$ and $B \times A$ are isomorphic.
 8. Show that for every objects A , B and C , $(A \times B) \times C$ and $A \times (B \times C)$ are isomorphic.

3 Examples of cartesian categories

1. Show that the category **Rel** of sets and relations is cartesian.

Solution. Let us first properly define the category **Rel**. An object is a set, a morphism in $\mathbf{Rel}(A, B)$ is a relation between A and B , i.e. a subset of $A \times B$:

$$\mathbf{Rel}(A, B) = \mathcal{P}(A, B)$$

Composition of $R : A \rightarrow B$ and $S : B \rightarrow C$, i.e. $R \subseteq A \times B$ and $S \subseteq B \times C$, is the following subset of $A \times C$:

$$S \circ R = \{(a, c) \mid \exists b \in B. (a, b) \in R \wedge (b, c) \in S\}$$

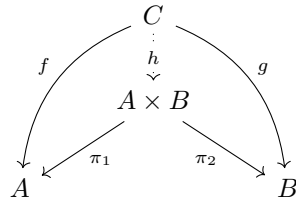
the identity on A is the diagonal subset of $A \times A$:

$$\text{id}_A = \{(a, a) \mid a \in A\}$$

It turns out that that product of A and B is the disjoint union $A \sqcup B$ of the sets A and B (we write $\iota_1(a)$, resp. $\iota_2(b)$, for the canonical injections of an element $a \in A$, resp. $b \in B$, in $A \sqcup B$). The projections are

$$\pi_1 = \{(\iota_1(a), a) \mid a \in A\} \subseteq (A \sqcup B) \times A \quad \pi_2 = \{(\iota_2(b), b) \mid b \in B\} \subseteq (A \sqcup B) \times B$$

Given a pair of morphisms $f : C \rightarrow A$ and $g : C \rightarrow B$,



one can check that the unique mediating morphism $h : C \rightarrow A \times B$ is

$$h = \{(c, \iota_1(a)) \mid (c, a) \in f\} \cup \{(c, \iota_1(b)) \mid (c, b) \in g\} \subseteq C \times (A \sqcup B)$$

2. We write **Vect** for the category of \mathbb{k} -vector spaces (where \mathbb{k} is a fixed field) and linear functions. Show that this category is cartesian. Given a basis for A and B , describe a basis for $A \times B$.
3. Show that the category **Cat** is cartesian.

4 Cartesian product as a functor

1. Given a cartesian category \mathcal{C} , show that the cartesian product induces a functor $\mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$.