

# Distributive laws between monads

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A *distributive law* between two monads  $(S, \eta^S, \mu^S)$  and  $(T, \eta^T, \mu^T)$  on category  $\mathcal{C}$  is a natural transformation

$$\lambda : S \circ T \Rightarrow T \circ S$$

such that the following diagrams commute

$$\begin{array}{ccc}
 SST \xrightarrow{S\lambda} STS \xrightarrow{\lambda S} TSS & & \begin{array}{ccc} & T & \\ \eta^S T \swarrow & & \searrow T \eta^S \\ ST & \xrightarrow{\lambda} & TS \end{array} \\
 \mu^S T \Downarrow & & \\
 ST \xrightarrow{\lambda} TS & & 
 \end{array}$$
  

$$\begin{array}{ccc}
 STT \xrightarrow{\lambda T} TST \xrightarrow{T\lambda} TTS & & \begin{array}{ccc} & S & \\ S \eta^T \swarrow & & \searrow \eta^T S \\ ST & \xrightarrow{\lambda} & TS \end{array} \\
 S \mu^T \Downarrow & & \\
 ST \xrightarrow{\lambda} TS & & 
 \end{array}$$

1. Can we always compose monads?
2. Draw those diagrams as string diagrams in the 2-category **Cat**.
3. Draw the laws for monads as string diagrams.
4. Show that the distributive law  $\lambda$  induces a structure of monad on the functor  $T \circ S$ .
5. Consider on **Set** the monads  $S$  of free monoid and  $T$  of free abelian group. Construct a distributive law  $\lambda : ST \Rightarrow TS$  so that the composite monad is the monad of free ring.
6. How can we compose three monads with distributive laws?

Suppose given two monads  $S$  and  $T$  as above. We write  $U^S : \mathcal{C}^S \rightarrow \mathcal{C}$  for the forgetful functor from the category of algebras. A *lift* of the monad  $T$  to  $\mathcal{C}^S$  is a monad  $(\tilde{T}, \tilde{\eta}^T, \tilde{\mu}^T)$  such that

$$U^S \tilde{T} = T U^S \qquad U^S \tilde{\eta}^T = \eta^T U^S \qquad U^S \tilde{\mu}^T = \mu^T U^S$$

7. Show that distributive laws between  $S$  and  $T$  correspond to lifts of  $T$  to  $\mathcal{C}^S$  (hint: for lift-to-distributive-law direction use the fact that  $U^S$  has a left adjoint such that the induced monad is  $S$  and notice that the first equality above can be seen as an invertible natural transformation).
8. [Optional] Show that the distributive laws in a 2-category  $\mathcal{C}$  correspond to monads in the 2-category of monads in  $\mathcal{C}$ :  $\text{DLaw}(\mathcal{C}) = \text{Mnd}(\text{Mnd}(\mathcal{C}))$ .

A *strict factorization system* on a category  $\mathcal{C}$  consists of a pair of subcategories  $\mathcal{L}$  and  $\mathcal{R}$  with the same objects as  $\mathcal{C}$  such that every morphism  $f$  of  $\mathcal{C}$  factors uniquely as  $f = r \circ l$  with  $l \in \mathcal{L}$  and  $r \in \mathcal{R}$ .

9. Show that a distributive law between monads in the 2-category **Span** corresponds to a strict factorization system.