

Non-Crossing Tree Realizations of Ordered Degree Sequences

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Abstract. We investigate the enumeration of non-crossing tree realizations of integer sequences, and we consider a special case in four parameters, that can be seen as a four-dimensional tetrahedron that generalizes Pascal's triangle and the Catalan numbers. This work is motivated by the study of ambiguities in categorial grammars.

Keywords: Proof nets · Non-crossing trees · Integer sequences · Catalan's triangle · Pascal-Catalan-tetrahedron

1 Introduction

A *non-crossing* tree t is a labeled tree on a *sequence* of vertices $\langle v_0, v_1, \dots, v_n \rangle$ drawn in counterclockwise order on a circle, and whose edges are straight line segments that do not cross. For any index $0 \leq i \leq n$, let d_i stand for the number of edges incident with v_i (that is the degree of v_i). Then as any other tree on $n+1$ vertices, t satisfies $\sum_{i=0}^n d_i = 2n$. Thus, the *sequence* $\langle d_0, d_1, \dots, d_n \rangle$ defines a *composition* of $2n$ into $n+1$ positive summands (two sequences of integers that differ only in the order of their elements define *distinct* compositions of the same integer). Stated otherwise, t is a non-crossing tree *realization* of the composition $\langle d_0, d_1, \dots, d_n \rangle$.

For any composition $c = \langle d_0, d_1, \dots, d_n \rangle$ of $2n$ into $n+1$ positive summands, let $nct(c)$ stand for the number of non-crossing tree realizations of c , that is the number of non-crossing trees on $n+1$ vertices $\langle v_0, v_1, \dots, v_n \rangle$ such that vertex v_i has degree d_i for any index $0 \leq i \leq n$ (there always exists at least one, see Proposition 2.3). We aim at computing nct . Note that here the input is more specific than the degree partition, as for example in [5].

From Proof Nets to Non-Crossing Trees

Our interest for these non-crossing tree realizations comes from linguistics and proof theory. The starting point for this work was the following linguistic problem: How many different readings can an ambiguous sentence at most have? Particularly, which sentence of a given length has the most different readings? When using categorial grammars based on the Lambek calculus [9] or related systems, a parse tree is a formal proof in a deductive system. Thus, our questions become: How many different formal proofs can a formula have? Particularly, which formula of a given length has the most different formal proofs? In

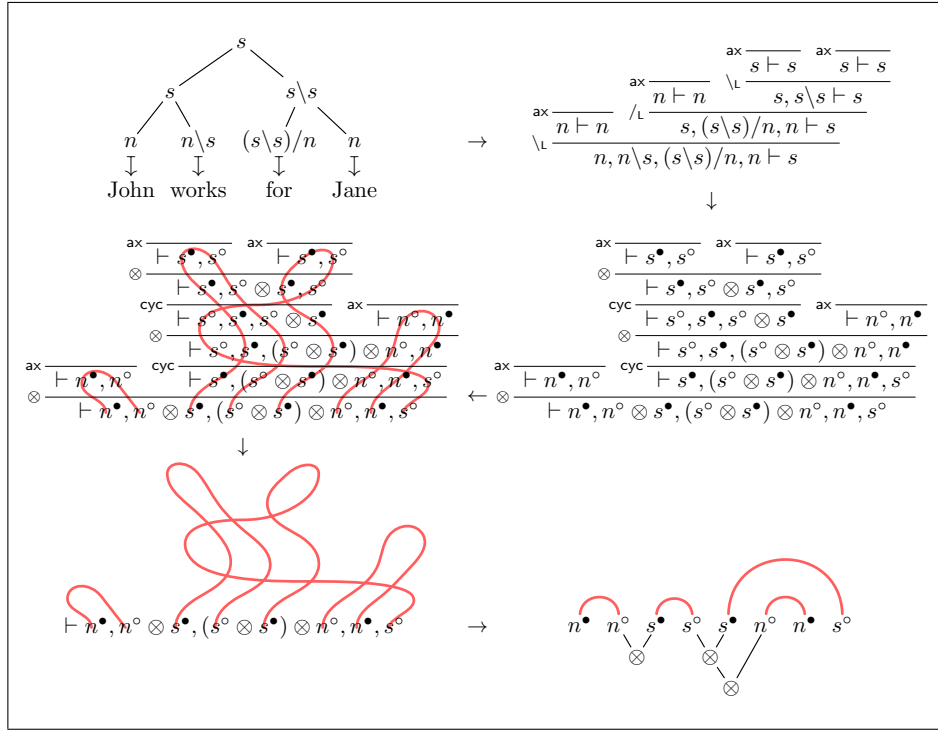


Fig. 1: From parse trees via proof trees to proof nets

category theoretical terms these questions come down to the cardinality of the Hom-sets in a free non-commutative star-autonomous category [2]. The corresponding logic is a variant of non-commutative intuitionistic linear logic [17,8] for which formal proofs can be represented as planar *proof nets*. It would go too far beyond the scope of this paper to go into the details of this correspondence. However, to give the reader an idea, we have shown in Figure 1 the transformation of a parse tree into a proof net. The first step transforms the parse tree into a formal proof according to Lambek’s work [9]. In the second step, this proof is embedded into a one-sided multiple conclusion system using the binary connectives \wp and \otimes [17]. In order not to lose the information on positive and negative positions in the formulas we use polarities (see, e.g., [8] for details). The final three steps show how this one-sided sequent proof is translated into a proof net by simply drawing the flow graph on the atoms appearing in the proof (for more details, see [8,3,15]). It is a well-known fact of linear logic that such a graph G does indeed correspond to a sequent proof if and only if every switching (that is, every graph obtained from G by removing for each \wp -node one of the two edges that it to its children) is a connected and acyclic graph [4]. If G does not contain any \wp -nodes, as in our example, then G itself has to be connected and acyclic. Furthermore, we have that G is planar if and only if the sequent proof does not contain the exchange rule, as it is the case for the Lambek calculus [17].

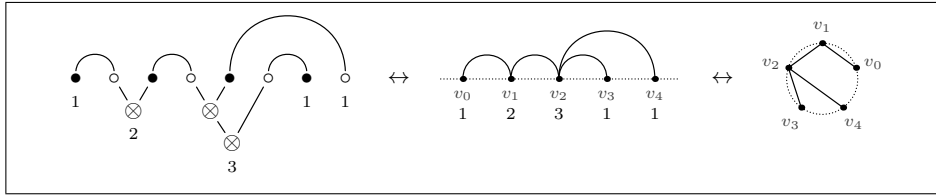


Fig. 2: Planar \emptyset -free proof nets as non-crossing trees

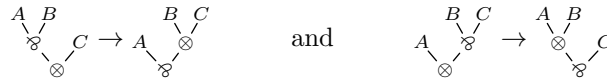
To summarize, a sentence corresponds to a sequent, and a parsing of the sentence to a planar proof net for the sequent. Thus, our question of how many readings does a sentence have becomes:

- (i) How many different planar proof nets can at most be defined over a given sequent?

Particularly, if we are interested in sentences that are as ambiguous as possible we have to ask:

- (ii) Over which sequent of a given length can the most different planar proof nets be defined?

In that respect, we can ignore the names of the atoms, and only \emptyset -free sequents are of interest: on the one hand, occurrences of \emptyset lying above an occurrence of \otimes can be moved down by the transformations



which preserve correctness without affecting linkings (see [6,7]); on the other hand, root occurrences of \emptyset are irrelevant and can be removed. Hence, for every sequent Γ there is a \emptyset -free sequent Γ' , such that for Γ' exist *at least* as many different planar proof nets as for Γ .

Finally, up to the associativity of \otimes , planar \emptyset -free proof nets are in bijection with non-crossing trees as shown in Figure 2.

Unfortunately, we were not able to find a closed formula for $nct\langle d_0, d_1, \dots, d_n \rangle$ depending only on the input composition $\langle d_0, d_1, \dots, d_n \rangle$, which would be needed to give a clear answer to Question (i) above. When investigating Question (ii), we found through experiments that at least up to $n = 25$, nct is maximized for compositions of the shape

$$1, \underbrace{2, 2, \dots, 2}_p \text{ summands}, \underbrace{1, 3, 1, 3, \dots, 1, 3}_{2q} \text{ summands}, 1, \underbrace{2, 2, \dots, 2}_r \text{ summands}, \underbrace{1, 3, 1, 3, \dots, 1, 3}_{2s} \text{ summands} \quad (1)$$

which we write as $12^p(13)^q12^r(13)^s$ and are the first step of our study (compositions maximizing nct for higher values of n may be of a different shape or involve summands higher than 3). The input is now reduced to four parameters p, q, r and s such that $n + 1 = 1 + p + 2q + 1 + r + 2s$. We write $nct(12^p(13)^q12^r(13)^s)$ as $N_{p,q,r,s}$, and we are interested in computing $N_{p,q,r,s}$.

A Four-Dimensional Generalization of Pascal's and Catalan's Triangles

Recall that Pascal's triangle $P_{p,r} = \binom{p+r}{p,r}$ can be generated recursively by:

$$P_{p,r} = \begin{cases} 1 & \text{if } p = 0 \text{ or } r = 0 \\ P_{p,r-1} + P_{p-1,r} & \text{if } p > 0 \text{ and } r > 0 \end{cases} \quad (2)$$

The first few values are shown below (see [14, A007318]):

$$\begin{array}{cccccc} & & p & & r & \\ & & \swarrow & & \searrow & \\ & & & 1 & & \\ & & & 1 & 1 & \\ & & 1 & 2 & 1 & \\ & 1 & 3 & 3 & 1 & \\ & 1 & 4 & 6 & 4 & 1 \\ 1 & 5 & 10 & 10 & 5 & 1 \\ 1 & 6 & 15 & 20 & 15 & 6 & 1 \end{array} \quad (3)$$

The Catalan numbers $C_q = \frac{1}{q+1} \binom{2q}{q,q}$ are generated recursively by $C_0 = 1$ and $C_{q+1} = \sum_{j=0}^q C_j \cdot C_{q-j}$. A combination of Pascal's triangle and the Catalan numbers is known as Catalan's triangle $Q_{p,q} = \frac{p+1}{p+q+1} \binom{p+2q}{p+q,q}$ which can be generated recursively by:

$$Q_{p,q} = \begin{cases} 1 & \text{if } q = 0 \\ C_q & \text{if } p = 0 \\ Q_{p+1,q-1} + Q_{p-1,q} & \text{if } p, q > 0 \end{cases} \quad (4)$$

The first few values are shown below (see [14, A009766]):

$$\begin{array}{cccccc} & & p & & q & \\ & & \swarrow & & \searrow & \\ & & & 1 & & \\ & & & 1 & 1 & \\ & & 1 & 2 & 2 & \\ & 1 & 3 & 5 & 5 & \\ & 1 & 4 & 9 & 14 & 14 \\ 1 & 5 & 14 & 28 & 42 & 42 \\ 1 & 6 & 20 & 48 & 90 & 132 & 132 \end{array} \quad (5)$$

It is also possible to generalize the recursive formula of the Catalan numbers into a triangle $R_{q,s}$ generated by (assuming that $C_{-1} = 0$):

$$R_{q,s} = \begin{cases} 1 & \text{if } q = s = 0 \\ \sum_{j=0}^q \sum_{l=0}^s C_{j+l-1} \cdot R_{q-j,s-l} & \text{if } q + s > 0 \end{cases} \quad (6)$$

The first few values are shown below:

$$\begin{array}{ccccccc}
 & & q & & s & & \\
 & & \swarrow & & \searrow & & \\
 & & 1 & & 1 & & \\
 & & 1 & & 1 & & \\
 & & 2 & & 3 & & 2 \\
 & & 5 & & 9 & & 9 & & 5 \\
 & & 14 & & 28 & & 34 & & 28 & & 14 \\
 & & 42 & & 90 & & 123 & & 123 & & 90 & & 42 \\
 & & 132 & & 297 & & 440 & & 497 & & 440 & & 297 & & 132
 \end{array} \tag{7}$$

As an example,

$$\begin{aligned}
 R_{3,2} &= 123 \\
 &= C_0 \cdot (28 + 34) + C_1 \cdot (5 + 9 + 9) + C_2 \cdot (2 + 3 + 2) + C_3 \cdot (1 + 1) + C_4 \cdot 1
 \end{aligned}$$

We shall establish that $N_{p,q,r,s}$ is a four-dimensional “tetrahedron” that generalizes the three triangles P , Q and R above, insofar as:

$$N_{p,0,r,0} = N_{r,0,p,0} = P_{p,r}, \tag{8}$$

$$N_{p,q,0,0} = N_{0,0,p,q} = Q_{p,q}, \tag{9}$$

$$N_{0,q,0,s} = N_{0,s,0,q} = R_{q,s}. \tag{10}$$

Outline

The organization of this paper is as follows: First, in [Section 2](#), we study the general case of enumerating non-crossing tree realizations of integer compositions. Then, in [Sections 3–7](#), we concentrate on the four-parameter case. In particular, we will prove identities (8)–(10) in [Sections 4](#) and [5](#). Finally, we will provide the generating function for $N_{p,q,r,s}$ in [Section 7](#).

Missing proofs can be found in the technical report [\[10\]](#).

2 General Case

Any labeled tree on a *sequence* of vertices can be drawn in such a way that its vertices lie in counterclockwise order on a circle and its edges are straight line segments lying inside that circle. In that case, of course, some of its edges may cross each other. Let us call such a labeled tree a *crossing tree*. The order of summands in a composition does not matter regarding the number of its labeled tree realizations (there are six for any composition of $2 \cdot 4$ into $4 + 1$ summands in the multiset $\{1, 1, 2, 2, 2\}$). But it does as soon as we distinguish between non-crossing and crossing realizations. As an example, there are one non-crossing and five crossing realizations of $\langle 1, 1, 2, 2, 2 \rangle$, while there are three non-crossing and three crossing realizations of $\langle 1, 2, 1, 2, 2 \rangle$ (these are shown on [Figure 3](#)).

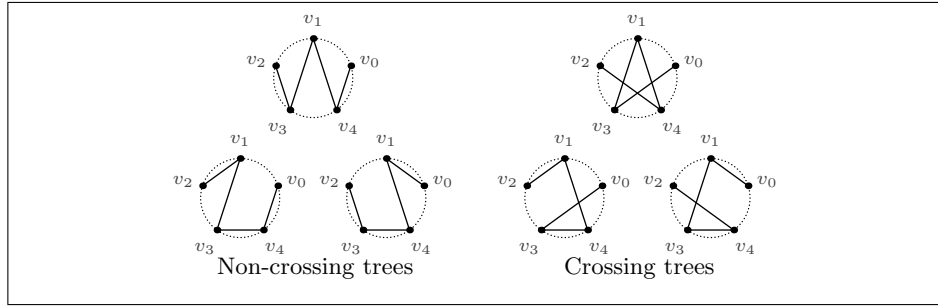


Fig. 3: The six labeled tree realizations of $\langle 1, 2, 1, 2, 2 \rangle$

Remark 2.1 A proof of Cayley’s formula (see e.g., [1]), which asserts that the number of labeled trees on $n + 1$ vertices is $(n + 1)^{n-1}$ (see [14, A000272]) relies on:

$$\sum_{\langle d_0, d_1, \dots, d_n \rangle} \binom{n-1}{d_0-1, d_1-1, \dots, d_n-1} = (n+1)^{n-1} \quad (11)$$

where the sum ranges over the $\binom{2n-1}{n, n-1}$ compositions of $2n$ into $n + 1$ positive summands. Noy established in [11, Corollary 1.2] that the number of non-crossing trees on $n + 1$ vertices is $\frac{1}{2n+1} \binom{3n}{2n, n}$ (see [14, A001764]). Recall that $nct\langle d_0, d_1, \dots, d_n \rangle$ stands for the number of non-crossing tree realizations of the composition $nct\langle d_0, d_1, \dots, d_n \rangle$. Then

$$\sum_{\langle d_0, d_1, \dots, d_n \rangle} nct\langle d_0, d_1, \dots, d_n \rangle = \frac{1}{2n+1} \binom{3n}{2n, n} \quad (12)$$

where the sum ranges over the $\binom{2n-1}{n, n-1}$ compositions of $2n$ into $n + 1$ positive summands.

The image under rotation of a non-crossing tree t on vertices $\langle v_0, v_1, \dots, v_n \rangle$ is a non-crossing tree on vertices $\langle v_{k+1}, \dots, v_n, v_0, \dots, v_k \rangle$ for some $k \leq n$. Moreover, t realizes a composition $\langle d_0, d_1, \dots, d_n \rangle$ iff its image under rotation realizes the composition $\langle d_{k+1}, \dots, d_n, d_0, \dots, d_k \rangle$. Thus, for any composition $\langle d_0, d_1, \dots, d_n \rangle$ and any $k \leq n$,

$$nct\langle d_0, d_1, \dots, d_n \rangle = nct\langle d_{k+1}, \dots, d_n, d_0, \dots, d_k \rangle \quad . \quad (13)$$

We shall refer to this property as *stability under rotation*.³

In the same way, the mirror image of a non-crossing tree t on vertices $\langle v_0, v_1, \dots, v_n \rangle$ is a non-crossing tree on vertices $\langle v_n, v_{n-1}, \dots, v_0 \rangle$, and t realizes a composition $\langle d_0, d_1, \dots, d_n \rangle$ iff its mirror image realizes the composition $\langle d_n, d_{n-1}, \dots, d_0 \rangle$. Thus, for any composition $\langle d_0, d_1, \dots, d_n \rangle$,

$$nct\langle d_0, d_1, \dots, d_n \rangle = nct\langle d_n, d_{n-1}, \dots, d_0 \rangle \quad . \quad (14)$$

³ In that respect, we may focus on *necklace*-compositions, i.e., compositions that are lexicographically minimal under rotation [13].

We shall refer to this property as *stability under mirror image*.⁴

We will now establish that for any positive integer n and any composition c of $2n$ into $n + 1$ positive summands, there exists a non-crossing tree realization of c (Proposition 2.3).

Lemma 2.2 *For any positive integer n and any sequence $\langle 1, d_1, \dots, d_n, d_{n+1} \rangle$ of $n + 2$ positive integers such that $1 + \sum_{i=1}^{n+1} d_i < 2(n + 1)$, there is an index $1 \leq k < n + 1$ such that $1 + \sum_{i=1}^k d_i = 2k$.*

Proof For any index $1 \leq l \leq n + 1$, let S_l stand for $1 + \sum_{i=1}^l d_i$. We prove the following implication by induction on l : if there is no index $1 \leq k < l$ such that $S_k = 2k$, then $S_l \geq 2l$. Since by hypothesis $S_{n+1} < 2(n + 1)$, there must exist an index $1 \leq k < n + 1$ such that $S_k = 2k$.

Base. Since $d_1 \geq 1$, $S_1 = 1 + d_1 \geq 2 \cdot 1$ and the stated implication holds trivially.

Induction. Assume that the stated implication holds for l (IH), and that there exists no index $1 \leq k < l + 1$ such that $S_k = 2k$. We reformulate the latter hypothesis as: (i) there exists no index $1 \leq k < l$ such that $S_k = 2k$, and (ii) $S_l \neq 2l$. By (IH) we get from (i), that $S_l \geq 2l$, and from (ii), that $S_l > 2l$, i.e., $S_l \geq 2l + 1$. Since $d_{l+1} \geq 1$, $S_{l+1} = S_l + d_{l+1} \geq 2(l + 1)$. \square

Proposition 2.3 *For any positive integer n and any composition c of $2n$ into $n + 1$ positive summands, there exists a non-crossing tree realization of c .*

Proof We proceed by induction on n .

Base. The unique composition $\langle 1, 1 \rangle$ of $2 \cdot 1$ into $1 + 1$ positive summands is realized by the unique (trivially non-crossing) tree on $1 + 1$ vertices.

Induction. Assume that the stated property holds for any positive integer up to n (IH), and let $\langle d_0, d_1, \dots, d_{n+1} \rangle$ be a composition of $2(n + 1)$ into $n + 2$ positive summands. Since n is a positive integer, $2(n + 1) > n + 2$ and there must exist at least one summand $d_k > 1$. By stability under rotation, we can assume without loss of generality that d_0 is such a summand, i.e. that $d_0 > 1$. Then $1 + \sum_{i=1}^{n+1} d_i < 2(n + 1)$ and by Lemma 2.2, there exists an index $1 \leq k < n + 1$ such that $1 + \sum_{i=1}^k d_i = 2k$. By difference, $(d_0 - 1) + \sum_{i=k+1}^{n+1} d_i = 2(n - k + 1)$. Then by (IH):

- there exists a non-crossing tree on vertices $\langle t_0, t_1, \dots, t_k \rangle$ realizing the composition $\langle 1, d_1, \dots, d_k \rangle$ of $2k$ into $k + 1$ positive summands,
- there exists a non-crossing tree on vertices $\langle u_0, u_1, \dots, u_{n-k+1} \rangle$ realizing the composition $\langle d_0 - 1, d_{k+1}, \dots, d_{n+1} \rangle$ of $2(n - k + 1)$ into $n - k + 2$ positive summands.

Let T and U stand for the respective edge sets of these non-crossing trees (where edges are defined as couples of vertices). We “merge” t_0 and u_0 into a single

⁴ In that respect, we may focus on *bracelet*-compositions, i.e., necklace-compositions that are lexicographically minimal under mirror image [12].

vertex v_0 to get a tree on vertices $\langle v_0, v_1, \dots, v_{n+1} \rangle$ which edge set is defined as

$$\begin{aligned} & \{ \{v_i, v_j\} \mid \{t_i, t_j\} \in T \} \\ & \cup \{ \{v_0, v_{j+k}\} \mid \{u_0, u_j\} \in U \} \\ & \cup \{ \{v_{i+k}, v_{j+k}\} \mid \{u_i, u_j\} \in U, i > 0, j > 0 \} \end{aligned} \quad (15)$$

This tree is non-crossing and it realizes the composition $\langle d_0, d_1, \dots, d_n, d_{n+1} \rangle$ of $2(n+1)$ into $n+2$ positive summands (see Figure 4 for an example). \square

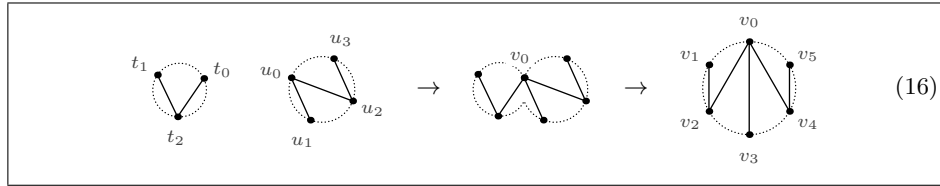


Fig. 4: Merging two non-crossing trees into a single one

The previous proof suggests a recursive definition of nct :

- The unique composition $\langle 1, 1 \rangle$ of $2 \cdot 1$ into $1 + 1$ positive summands is realized by the unique non-crossing tree on $1 + 1$ vertices. Thus,

$$nct\langle 1, 1 \rangle = 1. \quad (17)$$

- Let n be strictly greater than 1 and $\langle d_0, d_1, \dots, d_n \rangle$ be a composition of $2n$ into $n+1$ positive summands. Let k be the smallest index such that $d_k > 1$ (there exists at least one). By stability under rotation,

$$nct\langle 1, \dots, 1, d_k, \dots, d_n \rangle = nct\langle d_k, \dots, d_n, 1, \dots, 1 \rangle. \quad (18)$$

Thus we can assume that $d_0 > 1$. In that case,

$$nct\langle d_0, d_1, \dots, d_n \rangle = \sum_k (nct\langle 1, d_1, \dots, d_k \rangle \cdot nct\langle d_0 - 1, d_{k+1}, \dots, d_n \rangle) \quad (19)$$

where the sum ranges over the set of indices $1 \leq k < n$ such that $1 + \sum_{i=1}^k d_i = 2k$ (there exists at least one).

Remark 2.4 Applying the recursive formula in a row to $d_0, d_0 - 1, \dots, 1$, we get

$$nct\langle d_0, d_1, \dots, d_n \rangle = \sum_{\langle k_0, \dots, k_{d_0} \rangle} \prod_{j=1}^{d_0} nct\langle 1, d_{k_{j-1}+1}, \dots, d_{k_j} \rangle \quad (20)$$

where the sum ranges over the set of sequences $\langle k_0, \dots, k_{d_0} \rangle$ of $d_0 + 1$ indices such that $0 = k_0 < \dots < k_{d_0} = n$ and such that for all $0 < j \leq d_0$,

$$1 + \sum_{i=k_{j-1}+1}^{k_j} d_i = 2(k_j - k_{j-1}). \quad (21)$$

Remark 2.5 The construction in the proof of Proposition 2.3 is similar to the proof of the sequentialization theorem of linear logic using the *splitting tensor*. It is then easy to see that every non-crossing tree can be obtained this way as a merging of smaller ones, as indicated in Figure 4. However, it is not clear how this insight can be used for counting non-crossing trees since there is no unique decomposition for a given tree.

3 The Four Parameters Case

We focus now on the special case where compositions $\langle d_0, d_1, \dots, d_n \rangle$ are of the shape

$$1, \underbrace{2, 2, \dots, 2}_p, \underbrace{1, 3, 1, 3, \dots, 1, 3}_{2q}, \underbrace{1, 2, 2, \dots, 2}_r, \underbrace{1, 3, 1, 3, \dots, 1, 3}_{2s} \quad (22)$$

which we write as $12^p(13)^q12^r(13)^s$. For the following, recall that $N_{p,q,r,s}$ stands for $nct\langle 12^p(13)^q12^r(13)^s \rangle$.

Lemma 3.1 *For any p, q, r and s , $N_{p,q,r,s} = N_{r,q,p,s} = N_{p,s,r,q}$.*

Proof This follows from stability under rotation and mirror image. We give the formal calculations here in full, because we use similar arguments later on without showing them explicit.

$$\begin{aligned} N_{p,q,r,s} &= nct\langle 12^p(13)^q12^r(13)^s \rangle && \text{by definition of } N \\ &= nct\langle 2^p(13)^q12^r(13)^s1 \rangle && \text{by stability under rotation} \\ &= nct\langle 2^p1(31)^q2^r1(31)^s \rangle && \text{by reparentesizing} \\ &= nct\langle (13)^s12^r(13)^q12^p \rangle && \text{by stability under mirror image} \\ &= nct\langle 12^r(13)^q12^p(13)^s \rangle && \text{by stability under rotation} \\ &= N_{r,q,p,s} && \text{by definition of } N. \end{aligned}$$

The same way, $N_{p,q,r,s} = N_{p,s,r,q}$. □

According to Lemma 3.1, so as to get a recursive definition of $N_{p,q,r,s}$, we need to consider only $N_{0,0,0,0}$ on the one hand, $N_{p+1,q,r,s}$ and $N_{p,q,r,s+1}$ on the other hand.

Proposition 3.2 *For any p, q, r and s ,*

$$N_{0,0,0,0} = 1 \tag{23}$$

$$\begin{aligned} N_{p+1,q,r,s} &= N_{0,0,0,0} \cdot N_{p,q,r,s} + \sum_{j=1}^q (N_{0,j,0,0} \cdot N_{p,q-j,r,s}) \\ &+ \sum_{k=1}^r (N_{0,q,k,0} \cdot N_{p,0,r-k,s}) + \sum_{l=1}^s (N_{0,q,r,l} \cdot N_{p,0,0,s-l}) \end{aligned} \tag{24}$$

$$\begin{aligned} N_{p,q,r,s+1} &= N_{0,0,0,0} \cdot N_{1+p,q,r,s} + \sum_{i=1}^p (N_{i,0,0,0} \cdot N_{1+p-i,q,r,s}) \\ &+ \sum_{j=1}^q (N_{p,j,0,0} \cdot N_{1,q-j,r,s}) \end{aligned} \tag{25}$$

Proof We have $N_{0,0,0,0} = nct\langle 12^0(13)^0 12^0(13)^0 \rangle = 1$ by definition of N and (17). Next we have $N_{p+1,q,r,s} = nct\langle 12^p 2(13)^q 12^r(13)^s \rangle = nct\langle 21(31)^q 2^r(13)^s 12^p \rangle$ by definition of N , reparenthesizing and stability under rotation. Applying (19) we get

$$\begin{aligned} N_{p+1,q,r,s} &= nct\langle 11 \rangle \cdot nct\langle 1(31)^q 2^r(13)^s 12^p \rangle \\ &+ \sum_{j=1}^q (nct\langle 11(31)^j \rangle \cdot nct\langle 1(31)^{q-j} 2^r(13)^s 12^p \rangle) \\ &+ \sum_{k=1}^r (nct\langle 11(31)^q 2^k \rangle \cdot nct\langle 12^{r-k}(13)^s 12^p \rangle) \\ &+ \sum_{l=1}^s (nct\langle 11(31)^q 2^r(13)^l \rangle \cdot nct\langle 1(13)^{s-l} 12^p \rangle). \end{aligned}$$

Notice that there is no other way to “split” $\langle 21(31)^q 2^r(13)^s 12^p \rangle$ into two compositions such that the first one is of the form $\langle 1, d_1, \dots, d_k \rangle$ and satisfies $1 + \sum_{i=1}^k d_i = 2k$. Then we get (24) by reparenthesizing and stability under rotation. A similar argument applies to the proof of (25). \square

We have the following immediate consequences.

Corollary 3.3 *For any p and s ,*

$$N_{p,0,0,0} = 1 \tag{26}$$

$$N_{0,0,0,s} = C_s \tag{27}$$

$$N_{0,0,r,s+1} = N_{1,0,r,s} \tag{28}$$

where C_s stands for the s -th Catalan number.

Proof Both (26) and (27) are easily proved by induction, and (28) follows from (25) and (26). \square

4 Pascal's & Catalan's Triangles

In this section we are going to establish identities (8) and (9) mentioned in the introduction. First, recall Pascal's second identity

$$\sum_{k=0}^b \binom{a+k}{a, k} = \binom{a+1+b}{a+1, b} \quad (29)$$

and as a consequence for all a, b , and c

$$\sum_{k=c+1}^b \binom{a+k}{a, k} = \binom{a+1+b}{a+1, b} - \binom{a+1+c}{a+1, c}. \quad (30)$$

Proposition 4.1 For all p, r and s ,

$$N_{p,0,r,0} = \binom{p+r}{p, r} = P_{p,r} \quad (31)$$

$$N_{p,0,0,s} = \binom{p+2s}{p+s, s} - \frac{s}{p+s+1} \binom{p+2s}{p+s, s} \quad (32)$$

$$= \frac{p+1}{p+s+1} \binom{p+2s}{p+s, s} = Q_{p,s} \quad (33)$$

$$N_{p,0,r,s} = \binom{p+r+2s}{p+s, r+s} - \frac{s}{p+r+s+1} \binom{p+r+2s}{p+r+s, s} \quad (34)$$

A proof can be found in [10]. The triangles defined by $N_{p,0,r,0}$ and $N_{p,0,0,s}$ are known as *Pascal's triangle* and *Catalan's triangle* respectively ([A007318](#), [A009766](#)).

Remark 4.2 We may have deduced (27) from (34). Indeed, by (34)

$$\begin{aligned} N_{0,0,0,s} &= \binom{2s}{s, s} - \frac{s}{s+1} \binom{2s}{s, s} \\ &= \frac{1}{s+1} \binom{2s}{s, s} \\ &= C_s \end{aligned}$$

5 A Triangular Catalan Recurrence

In this section we establish the identity (10) from the introduction. We need the following lemma:

Lemma 5.1 For any t and u ,

$$\sum_{i=0}^t \sum_{j=0}^u C_i \cdot C_j \cdot C_{(t-i)+(u-j)} = C_{t+u+1}. \quad (35)$$

Proof By induction on u (the base case is the usual recurrence for Catalan numbers). \square

Proposition 5.2 (A recurrence for $N_{0,q,0,s}$) For any q and s ,

$$N_{1,q,0,s} = \sum_{j=0}^q \sum_{l=0}^s C_{j+l} \cdot N_{0,q-j,0,s-l} \quad (36)$$

$$N_{2,q,0,s} = \sum_{j=0}^q \sum_{l=0}^s C_j \cdot C_l \cdot N_{1,q-j,0,s-l} \quad (37)$$

$$= \sum_{j=0}^q \sum_{l=0}^s C_{j+l+1} \cdot N_{0,q-j,0,s-l} \quad (38)$$

$$N_{0,q+1,0,s+1} = N_{1,q+1,0,s} + N_{1,q,0,s+1} - N_{2,q,0,s} \quad (39)$$

$$N_{0,q,0,s} = \begin{cases} 1 & \text{if } q = s = 0 \\ \sum_{j=0}^q \sum_{l=0}^s C_{j+l-1} \cdot N_{0,q-j,0,s-l} & \text{if } q + s > 0 \end{cases} \quad (40)$$

Hence $N_{0,q,0,s} = R_{q,s}$.

6 Triangles and Tetrahedra

The value of $N_{p,q,r,s}$ depends on four parameters p, q, r, s . If we fix two of them, we can obtain triangles. For example, for $q = s = 0$ we get *Pascal's triangle* (3), and for $r = s = 0$ we get the *Catalan triangle* (5) (which should more precisely be called *Pascal-Catalan triangle*). If we let $p = r = 0$, then we get the triangle (7), which could also be called *Catalan triangle*.

If we fix only one parameter, we obtain a tetrahedron. For example, let $s = 0$, and let us define $T_{p,q,r} = N_{p,q,r,0}$. Then we get from (34) and Lemma 3.1:

$$T_{p,q,r} = \binom{p+r+2q}{p+q, r+q} - \frac{q}{p+r+q+1} \binom{p+r+2q}{p+r+q, q} \quad (41)$$

This defines a tetrahedron where one side is Pascal's triangle and the other two sides are the Catalan triangle. Thus we can call it the *Pascal-Catalan tetrahedron*. We have the following recursive identities:

Proposition 6.1 For all p, q , and r , we have

$$T_{p+1,q+1,r+1} = T_{p+1,q+1,r} + T_{p,q+1,r+1} + T_{(p+1)+(r+1),q,0} \quad (42)$$

$$T_{p+1,q,r+1} = T_{p,q+1,r} + T_{(p+1)+(r+1),q,0} \quad (43)$$

$$T_{p+1,q+1,r} = T_{p+2,q,r} + T_{p,q+1,r} \quad (44)$$

Proof Easy calculation using (41). \square

For $r = 0$, we get the tetrahedron $V_{p,q,s} = N_{p,q,0,s}$, which we can call the *Catalan tetrahedron*, because two of its sides are the Catalan triangle (5) and the third side is the new Catalan triangle (7). Unfortunately, we could not find a closed formula for $V_{p,q,s}$. However, in Section 7 we will give the generating function. We also have the following:

Proposition 6.2 *For all p, q , and s , we have:*

$$\sum_{i=2}^{p+2} V_{i,q,s} = \sum_{j=0}^q C_j \cdot V_{p,q+1-j,s} \tag{45}$$

Observe that identities (42)–(44) establish close relationships among the triangles with $s = 0$ (or $q = 0$), i.e., the triangles that live inside the tetrahedron $T_{p,q,r}$. For example, from (43) we can get $N_{p+1,q,7,0} = N_{p,q+1,6,0} + N_{p+8,q,0,0}$.

Below, we exhibit some identities between triangles where $q \neq 0$ and $s \neq 0$.

$$N_{1,q,1,s} = N_{0,q,0,s+1} + N_{0,q+1,0,s} \tag{46}$$

$$N_{1,q,2,s} = N_{0,q+1,0,s+1} \tag{47}$$

$$N_{1,q+1,2,s+1} = N_{2,q,2,s} + N_{1,q,1,s+2} + N_{1,q+2,1,s} \tag{48}$$

$$N_{1,q+1,0,s+1} = N_{0,q+1,0,s+1} + N_{3,q,1,s} \tag{49}$$

$$N_{1,q+1,1,s+1} = N_{2,q+1,0,s+1} + N_{2,q,0,s} \tag{50}$$

$$N_{p,q,0,1} = \sum_{i=2}^{p+2} N_{i,q,0,0} \tag{51}$$

They can all be proved by using Lemma 3.1, (24) and (25) by easy but tedious calculations.

Now we can derive another recurrence for the triangle in (7), i.e., different from the one given in (6):

Proposition 6.3 *For all q and s , we have*

$$R_{q,s+2} = \sum_{j=0}^q C_{j+1} \cdot (R_{q-j,s+1} + R_{q+1-j,s}) \tag{52}$$

To see an example for (52) consider again the triangle (7):

$$\begin{array}{ccccccc}
 & & q & & s & & \\
 & & \swarrow & & \searrow & & \\
 & & & 1 & & & \\
 & & 1 & & 1 & & \\
 & & \swarrow & & \searrow & & \\
 & & 2 & 3 & 2 & & \\
 & & \swarrow & & \searrow & & \\
 & & 5 & 9 & 9 & 5 & \\
 & & \swarrow & & \searrow & & \\
 & & 14 & 28 & 34 & 28 & 14 \\
 & & \swarrow & & \searrow & & \\
 & & 42 & 90 & 123 & 123 & 90 & 42 \\
 & & \swarrow & & \searrow & & \\
 & & 132 & 297 & 440 & 497 & 440 & 297 & 132
 \end{array} \tag{53}$$

We have $123 = C_1 \cdot (28 + 34) + C_2 \cdot (9 + 9) + C_3 \cdot (3 + 2)$.

Remark 6.4 In the next section we will make use of the following identity:

$$N_{2,q,0,s} = \sum_{j=0}^q N_{0,j,0,0} \cdot N_{0,q+1-j,0,s} = \sum_{l=0}^s N_{0,0,0,l} \cdot N_{0,q,0,s+1-l} \quad (54)$$

It is a special case of (45), but it can also be shown directly: By (37), we have

$$N_{2,q,0,s} = \sum_{j=0}^q \sum_{l=0}^s C_j \cdot C_l \cdot N_{1,q-j,0,s-l}$$

Then (54) follows immediately by (25) and (27).

In the remainder of this section we derive a closed formula for the tetrahedron $s = 1$. We need the following observation:

Lemma 6.5 For any p, r, s and t ,

$$N_{p,0,r,s+t} + N_{p+r,t,0,t,s} = N_{p+t,0,r+t,s} + N_{p+r,0,0,s+t} \quad (55)$$

Proof Easy calculation, using (34). \square

For any t and n , let $U_t(n)$ be defined as

$$U_t(n) = \frac{1}{t+2n} \binom{t+2n}{t+n, n}. \quad (56)$$

Notice that for any t ,

$$tU_t(0) = \frac{t}{t} \binom{t}{t, 0} = 1. \quad (57)$$

The following two identities are called *Rothe's identities* (see [16, identities 14–15, p. 329]).

$$\sum_{k=0}^n U_t(k) \cdot U_u(n-k) = \frac{t+u}{tu} U_{t+u}(n) \quad (58)$$

and

$$\sum_{k=0}^n kU_t(k) \cdot U_u(n-k) = \frac{n}{u} U_{t+u}(n). \quad (59)$$

As a consequence, we get

$$\sum_{k=0}^n (t+k)U_t(k) \cdot uU_u(n-k) = (t+u+n)U_{t+u}(n) \quad . \quad (60)$$

Lemma 6.6 For any p, r and s ,

$$N_{p,0,r,s} = (p + s + 1)U_{p-r+1}(r + s) - sU_{p+r+1}(s) \quad (61)$$

$$N_{p,0,0,s} = (p + 1)U_{p+1}(s) \quad (62)$$

Proof The identity (61) follows immediately from (34) and (56). The identity (62) is a special case of (61). \square

Lemma 6.7 For any p, r, s and t ,

$$\begin{aligned} \sum_{l=0}^s N_{p,0,r,l} \cdot N_{t,0,0,s-l} &= N_{p+t+1,0,r,s} \\ &- \sum_{l=0}^{r-1} (p - r + 1 + l)U_{p-r+1}(l) \cdot (t + 1)U_{t+1}(r + s - l) \end{aligned} \quad (63)$$

Now we can give a closed formula for the $N_{p,q,r,1}$ tetrahedron.

Proposition 6.8 For all p, q and r ,

$$N_{p,q,r,1} = T_{p,q+1,r} + T_{p+r+1,q+1,0} - T_{0,q+1,r} - T_{p,q+1,0} \quad (64)$$

$$= T_{p+1,q,r+1} + T_{p+r,q,1} - T_{1,q,r} - T_{p,q,1} \quad (65)$$

7 Generating Functions

We can use the identities (24) and (54) for calculating the generating function for $N_{p,q,r,s}$. Recall that we use the following abbreviations:

$$\begin{aligned} P_{p,r} &= N_{p,0,r,0} & C_q &= N_{0,q,0,0} \\ Q_{p,q} &= N_{p,q,0,0} & T_{p,q,r} &= N_{p,q,r,0} \\ R_{q,s} &= N_{0,q,0,s} & V_{p,q,s} &= N_{p,q,0,s} \end{aligned} \quad (66)$$

Theorem 7.1 We have

$$(i) \quad C(y) = \sum_q C_q y^q = \frac{1 - \sqrt{1 - 4y}}{2y}$$

$$(ii) \quad P(x, z) = \sum_{p,r} P_{p,r} x^p z^r = \frac{1}{1 - x - z}$$

$$(iii) \quad Q(x, y) = \sum_{p,q} Q_{p,q} x^p y^q = \frac{C(y)}{1 - x \cdot C(y)}$$

$$(iv) \quad R(y, w) = \sum_{q,s} R_{q,s} y^q w^s = \frac{C(y) \cdot C(w) \cdot (w - y)}{w \cdot C(y) - y \cdot C(w)}$$

$$\begin{aligned} (v) \quad T(x, y, z) &= \sum_{p,q,r} T_{p,q,r} x^p y^q z^r = \frac{(1 - x - z - x \cdot z \cdot C(y)) \cdot C(y)}{(1 - x \cdot C(y)) \cdot (1 - z \cdot C(y)) \cdot (1 - x - z)} \\ &= P(x, z) \cdot (1 + y \cdot Q(x, y) \cdot Q(z, y)) \end{aligned}$$

$$\begin{aligned}
\text{(vi)} \quad V(x, z, w) &= \sum_{p,q,s} V_{p,q,s} x^p y^q w^s \\
&= \frac{C(y) \cdot C(w) \cdot (w - y - x \cdot (w \cdot C(y) - y \cdot C(w)))}{(w \cdot C(y) - y \cdot C(w)) \cdot (1 - x \cdot C(y)) \cdot (1 - x \cdot C(w))} \\
&= Q(x, y) \cdot Q(x, w) \cdot \left(\frac{R(y, w)}{C(y) \cdot C(w)} - x \right) \\
\text{(vii)} \quad N(x, y, z, w) &= \sum_{p,q,r,s} N_{p,q,r,s} x^p y^q z^r w^s \\
&= \frac{C(y)C(w)((1-x-z)(w-y) + (wC(y) - yC(w))((1-x-z)(xz(C(y) + C(w)) - x - z) + x^2z^2C(y)C(w)))}{(1-xC(y))(1-xC(w))(1-zC(y))(1-zC(w))(1-x-z)(wC(y) - yC(w))} \\
&= Q(x, y)Q(x, y)Q(x, y)Q(x, y) \left(\frac{R(y, w)}{C(y)^2 \cdot C(w)^2} + x^2z^2P(x, z) - \frac{x}{C(w)Q(z, y)} - \frac{z}{C(y)Q(x, w)} \right)
\end{aligned}$$

Proof The formulas in (i) and (ii) are well-known. For the others, the calculation can be found in [10]. \square

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