Trees, maps and Hurwitz numbers

GILLES SCHAEFFER CNRS & École Polytechnique ERC Research Starting Grant 208471 "ExploreMaps"

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General summary of the 3 lectures:

Factorizations, maps and ramified coverings

Orientations and decompositions of maps into trees

Applications to Hurwitz numbers

First lecture

Factorizations, maps and ramified coverings

Permutations, factorizations and increasing maps Hurwitz original motivation, ramified coverings Ramified coverings provide bijections "for free" Permutations, factorizations, increasing maps

Permutations in cycle notation: $\sigma = (1, 2, 5)(3, 6)(4)(7) = (1, 2, 5)(3, 6)$ Cycle type = distribution of cycle lengths: $\lambda(\sigma) = 1^2 2 3$ Transpositions = permutations with type $\lambda = 2 1^{n-1}$: $\tau = (2, 5)$.

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The graph of a factorization $\tau_1 \dots \tau_m = \sigma \in S_n$: - vertices represent the permuted elements: $\{1, \dots, n\}$ - edges represent transpositions: an edge (i, j) with index k if $\tau_k = (i, j)$ $1 \int_{0}^{2} \int_{0}^$

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Cayley trees n^{n-2} with n nodes n^{n-2} (non-embedded) edge indexing (n-1)!



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Minimal factorizations

Proposition: Let $\lambda = 1^{\ell_1} \dots n^{\ell_n}$ with $\sum_i \ell_i = \ell$. A minimal factorization of a permutation of cycle type λ has $m = n - \ell$ factors.

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Their number is $\frac{n!}{\prod_i \ell_i! i^{\ell_i}} \prod_i (i^{i-2})^{\ell_i} \frac{m!}{\prod_i (i-1)!^{\ell_i}} = m! n! \prod_i \frac{1}{\ell_i!} \left(\frac{i^{i-2}}{i!}\right)^{\ell_i}.$ Corollaries:

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Hurwitz formula for the number of minimal transitive factorizations in transpositions

Theorem. Let $\lambda = 1^{\ell_1}, \ldots, n^{\ell_n}$ be a partition n, and $\ell = \sum_i \ell_i$. The number of *m*-uples of transpositions (τ_1, \ldots, τ_m) such that

- (product cycle type) $\tau_1 \cdots \tau_m = \sigma$ has cycle type λ
- (transitivity) the associated graph is connected
- (minimality) the number of factors is $m = n + \ell 2$

$$n^{\ell-3} \cdot m! \cdot n! \cdot \prod_{i \ge 1} \frac{1}{\ell_i!} \left(\frac{i^i}{i!}\right)^{\ell_i}$$

Proofs:

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(Hurwitz 1891, Strehl 1996) (Goulden–Jackson 1997) (Lando–Zvonkine 1999) (Bousquet-Mélou–Schaeffer 2000) (recurrences, Abel identities) (gfs and differential eqns) (geometry of LL mapping) (bijection + inclusion/exclusion)

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Combinatorial interpretation and proof?

How do we compute the product directly on the graph











The computation is performed along the boundary of the graph because I have consistently drawn edges increasingly in counterclockwise direction around each vertex $e_1 < e_2 < \ldots < e_k$





 $e_1 < e_2 < \ldots < e_k$



Any tree with indexed edges has a unique such increasing embedding





pointed indexed tree with n vertices (admit a unique embedding as increasing plane trees)









Proof: there is unique way to put labels so that the computation works!

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(4,6)(1,6)(1,5)(2,8)(3,4)(1,7)(5,8)(2,3)(4,8)(2,7)(3,8) = (1,6,7)(2,5)(3)(4)(8)



 $-\ell$ faces (face with k crosses = cycle of length k in the product)

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Origin of the problem...

and a common setup for all permutations/maps relations?

Ramified coverings of the sphere by itself

See book Lando-Zvonkin for more details.

Ramified coverings of the sphere by itself

Let $D = \{z \mid |z| < 1\} \subset \mathbb{C}$ be the unit open disc, and let \sim denote equivalence up to homeomorphisms (bijective, bicontinuous mappings).

A mapping $\phi : \mathcal{D} \to \mathcal{I}$ is a covering if, for all xin \mathcal{I} there exists $n \geq 1$ and a neighborhood Vof x such that $\phi^{-1}(V) \sim D \times \{1, \ldots, n\}$, and the restriction of ϕ to each sheet D_i (connected component of the preimage) is an homeomorphism $\phi_{|D_i} : D_i \xrightarrow{\sim} D$.

Example:

Let A_r be the annulus $\{z \mid r < |z| < 1\} \subset \mathbb{C}$. Consider $\phi_k : A_r \to A_{r^k}$ with $\phi_k(z) = z^k$.



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What is we try to extend from A_r to D?

Recall $\phi_k : A_r \to A_{r^k}$ with $\phi_k(z) = z^k$. Extend from A_r to D? The mapping $\phi_k : D^* \to D^*$ is a covering. but not $\phi_k : D \to D$.



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A mapping ϕ is ramified at x = 0 if

- there is a neighborhood V of the origin such that $\phi^{-1}(V) \sim D \times [1, \dots, p]$ and,
- the restriction of ϕ to each component of $\phi^{-1}(V)$ is homeomorphic to ϕ_k for some k.



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A mapping ϕ is a ramified covering of S by S if there exists a finite subset $X = \{x_1, \ldots, x_p\}$ such that:

- $\phi_{\mathbb{S}\setminus\phi^{-1}(X)}$ is a covering, and
- ϕ is ramified over each x_i



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the passport $\Lambda = (\lambda^{ig(1)}, \ldots, \lambda^{ig(p)})$ of a ramified covering



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To understand the "shape" of the covering, draw paths on \mathcal{I} and study its preimages.

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To understand the "shape" of the covering, \mathcal{T}

draw paths on \mathcal{I} and study its preimages.

- n independant preimages as long as we stay away from critical points
- a contractible loop on *I* yields *n* contractible loops on *D* but if we wind around critical points
 some sheets may get permuted
 - visiting critical points create multiple values or "vertices"

 $\Rightarrow \mbox{The partitions } \lambda^{(i)} \\ \mbox{are partitions of } n, \\ \mbox{degree of the covering.}$

Let us label $\{1, \ldots, n\}$ the preimages of a regular point.



 $\mathsf{Loop} \Rightarrow \mathsf{permutation}$ of sheet labels

Example: (1, 2)(3, 4)(5) in cyclic notation

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Monodromy, and permutations

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3 critical values $\lambda^{\bullet} = 2^3 1^2 \quad \lambda^{\circ} = 3^2 2 \quad \lambda^{\Box} = 62$







On \mathcal{I} , draw an edge between \bullet and \circ via the basepoint

We get a planar map:

that is, a graph embedded on the sphere with simply connected faces









1 regular value with labeled preimages

















1 regular value with labeled preimages

m+1 critical values, *m*-constellations, permutations



m+1 critical values, *m*-constellations, permutations



The preimage of the m-star is called a star-constellation.

Thm. Planar star-constellations with: - n labelled m-stars, - λ_j^{\Box} faces of degree j, - $\lambda_j^{(i)}$ color i vertices of degree jare in bijection with minimal transitive factorizations $\sigma_1 \cdots \sigma_m = \sigma_{\Box}$

with σ_i of cyclic type $\lambda^{(i)}$.

Monodromy, permutations, constellations summary

Theorem. There is a bijection between

- Labelled ramified covering of \mathbb{S} of type $\Lambda = (\lambda_0, \dots, \lambda_m)$
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Specializations.

- m = 2: bipartite maps with n edges
- m = 2 and $\lambda_{\bullet} = 2^{\frac{n}{2}}$: all \bullet have deg 2 \Leftrightarrow nonbipartite maps ($\frac{n}{2}$ edges)
- for all $i \ge 1$, $\lambda^{(i)} = 21^{n-2}$: factorizations in transpositions. coverings with almost only simple branch points; increasing maps

Ramified coverings and "trivial" bijections: combinatorial data structures

























Today

Factorizations, maps and ramified coverings Permutations, factorizations and increasing maps Hurwitz original motivation, ramified coverings Ramified coverings provide bijections "for free"

Later...

Orientations and decompositions of maps into trees Applications to Hurwitz numbers

A formula for general factorizations [BMS00]

Theorem. Let $\lambda = 1^{\ell_1}, \ldots, n^{\ell_n}$ be a partition of n, and $\ell = \sum_i \ell_i$. The number of *m*-uple of permutations $(\sigma_1, \ldots, \sigma_m)$ such that

- (factorization) $\sigma_1 \cdots \sigma_m = \sigma$ with cycle type λ
- (transitivity) $\langle \sigma_1, \ldots, \sigma_m \rangle$ acts transitively on $\{1, \ldots, n\}$
- (minimality) the total rank of factors is $\sum_i r(\sigma_i) = n + \ell 2$

$$m \frac{((m-1)n-1)!}{(mn-(n+\ell-2))!} \cdot n! \cdot \prod_{i} \frac{1}{\ell_{i}!} {mi-1 \choose i}^{\ell_{i}}$$

Proofs:

is

(Bousquet-Mélou–Schaeffer 2000) (Goulden–?? 2009) (bijection + inclusion/exclusion)(gfs and differential eqns)

$$\lambda = n$$
, factorizations of *n*-cycles: $\frac{1}{(mn+1)} \binom{mn+1}{n} \cdot (n-1)!$

 $\lambda = 1^n$, identity factorizations: $\frac{m}{(m-2)n+2} \frac{(m-1)^{n-1}}{(m-2)n+1} \binom{(m-1)n}{n} \cdot (n-1)!$
Our aim in the rest of the lectures

Prove the following two results using two bijective methods:

Factorization in transpositions: $\lambda = 1^n$, factorizations of the identity: $n^{n-3} \cdot (2n-2)!$

Need to count fully increasing quadrangulations

Factorizations in arbitrary factors:

 $\lambda = 1^n$, factorizations of the identity: $m \frac{(mn-n-1)!}{(mn-2n+2)!} (m-1)^n$

Need to count (m + 1)-constellations.

The two methods extend to general λ .

The second method extends to non minimal factorizations (higher genus)

Second lecture

Orientation and the decomposition of maps into trees

- A quick reminder about trees
- General idea: decompose a map into two trees
- 2 strategies explain (almost) all known bijections
 - minimal orientations and direct opening
 - left accessible orientations and the master bijection

A quick reminder about trees

Dyck paths and plane trees

Dyck path of length 2n = contour of a plane tree with n edges



The Dyck code of a tree is obtained during the walk around it upon:

- writing u the first time a vertex is visited (up steps)
- writing d the last time a vertex is visited (down steps)

Encodings of trees by words

We shall need two other classical codes:

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Encodings of trees by words

We shall need two other classical codes:

- the height code: write the height of each vertex during its first visit



- degree code: write the degree of each vertex during its first visit



after 45^o rotation: let P(2n) denote paths from (0,0) to (n,n+1) ending by an horizontal step









take a fonction f of $[n] \rightarrow [n+1]$

1	2	3	4	5	6	7
2	6	8	8	2	8	6

































Put labels on edges:



From maps to trees (I): tree-rooted maps

first strategy: Mullin primal dual decomposition

Recall a planar map is a proper embedding of a connected graph on the sphere (considered up to homeomorphisms).

From now on, map means rooted planar map.



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Planar maps, spanning trees and duality

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Starting at a root corner turn around the tree



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Starting at a root corner turn around the tree

non visited edges = balanced parenthesis word





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Code of the tree-rooted map = tree decorated by a balanced parenthesis word



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Using the code of the tree by its contour word:

uuuududududddddddddddd

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Using the code of the tree by its contour word:

Code of the tree-rooted map = tree decorated by a balanced parenthesis word = shuffle of two Dyck words

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Using the code of the tree by its contour word:

Code of the tree-rooted map = tree decorated by a balanced parenthesis word = shuffle of two Dyck words

The number of tree rooted planar maps with n edges is $\sum_{i=0}^{n} {2n \choose i} C_i C_{n-i}$ where C_n denotes Catalan numbers.

From maps to trees (I): tree-rooted maps

first strategy: Mullin primal dual decomposition

intermede: minimal orientations

second strategy: unfolding

Bernardi's master bi-theorem

Orient the tree edges away from the root



Orient the tree edges away from the root

Orient the other edges couterclockwise around the tree



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The tree is recovered by reconstructing its contour (or equivalently by leftmost depth first search).

From maps to trees (I): tree-rooted maps

first strategy: Mullin primal dual decomposition

intermede: minimal orientations

second strategy: unfolding

Bernardi's master bi-theorem

Consider a minimal accessible map



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Define its vertex unfolding: In the unfolded map, the plain edges form a spanning tree. (clockwise cycles are ruled out by external edges) (a counterclockwise cycle would be non accessible from the outside)

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Bernardi's master bijection for tree-rooted maps

The primal and dual trees of the unfolded maps are glued canonically (no shuffling of the codes required)



Bernardi's master bijection for tree-rooted maps

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Conversely gluying an arbitrary tree with n edges with an arbitrary tree with s + f = n + 2 vertices yields a left-accessible map

Theorem(Bernardi) This is a bijection between such pairs of trees and minimal accessible maps with n edges, (and tree-rooted maps via previous Theorem).

Corollary: $\sum_{i=0}^{n} {\binom{2n}{i}} C_i C_{n-i} = C_{n+1} C_n$

Summary: two strategies for tree-rooted maps













From maps to trees (II): eulerian maps

first strategy: blossoming trees

Let us recycle the first idea used for tree-rooted maps

using a canonical spanning tree



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using a canonical spanning tree





Then write the code of the primal tree on the chosen canonical tree



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Then write the code of the primal tree on the chosen canonical tree The map is recovered from the code by *closure*. Our code of the map will be a canonical decorated tree Question is "How do we choose the canonical spanning tree ?" Minimal orientations and canonical spanning trees

Idea: Use Bernardi's first bijection the other way round: Choose a minimal accessible orientation to get a spanning tree

Our pb becomes:

How to choose a canonical accessible minimal orientation?
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How to choose a canonical accessible minimal orientation?

A function $\alpha: V \to \mathbb{N}$ is feasible on a plane map M if there exists an orientation of M such that for each vertex v the outdegree of v is f(v).

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Fact: For many subclasses \mathcal{F} of planar maps, there exists an $\alpha_{\mathcal{F}}$ s.t.:

A planar map is in \mathcal{F} if and only if it admits an $\alpha_{\mathcal{F}}$ -orientation.

A map is eulerian if it admits a cycle that visits every edge exactly once. Let $\frac{1}{2}$ deg denote the $\frac{1}{2}$ degree function.

Proposition. A map is eulerian if and only if its admits a $\frac{1}{2}$ deg-orientation.





a map is 2-connected \Leftrightarrow it admits a bipolar orientation $\Leftrightarrow \text{ its quadrangulation admits} \\ \text{ an orientation with } \alpha(v) = 2$

a map is a simple triangulation \Leftrightarrow it admits an orientation with $\alpha(v)=3$



endow with min orient











Corrolary. This is a bijection between eulerian map with d_i vertices of degree i and rooted^{*} plane trees with d_i vertices of total degree 2i s.t.

- a vertex of total degree 2i has i-1 incoming half-edges



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- the tree is balanced (half-edges must form balanced parentheses)

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However, 2 among n + 2 are balanced:

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(Tutte 1964, S. 1997)

Summary of the blossoming tree strategy

To enumerate maps admitting $\alpha_{\mathcal{F}}$ -orientations:

- endow them with their minimal $\alpha_{\mathcal{F}}$ -orientation (hope it is accessible)
- construct the associated canonical spanning trees (Bernardi)
- open the resulting tree-rooted maps (Mullin)
- count the encoding balanced trees

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In Bernardi original bijection, the basepoint must be in the outer face. But in some cases the orientation is not outerface accessible. This approach was further extended in (Albenque, Poulalhon 2012) to

cover essentially all known blossoming bijections, including Bernardi-Fusy's fractional orientations. Third lecture

Applications to factorization problems Which factorisations, which maps? *m*-eulerian maps Hurwitz problem

Conclusion

What do we want to enumerate?

Recall. There is a bijection between

- Labelled ramified covering of \mathbb{S} of type $\Lambda = (\lambda_0, \dots, \lambda_m)$
- Factorizations $(\sigma_1 \cdots \sigma_m = \sigma_0)$ of type Λ
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Today. Minimal transitive factorizations of $\sigma_0 = id$.

m arbitray factors $\Rightarrow \sum_{i=1}^{m} (n - \ell_i) = 2n - 2$

m-constellations



transpositions $\Rightarrow m = 2\ell - 2$ increasing maps

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increasing maps



increasing quadrangulations



From maps to trees: constellations

$\alpha \text{-orientations}$ for m-eulerian maps

Bipartite map with black and white vertices of degree m such that:

- faces with labels in $\{1,\ldots,m\}$
- around black vertices, face labels read $1, \ldots, m$ in cw order
- around white vertices, face labels read $1, \ldots, m$ in ccw order



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Proposition: A bipartite map is *m*-eulerian iff it admits an α_c -orientation. This orientation is accessible, in fact strongly connected. We can apply our strategy!





endow with min $\alpha_{\it C}\text{-orient}$

(return cycles)






Openning a *m*-eulerian map



Corrolary. This is a bijection between m-eulerian maps and rooted^{*} plane trees with black and white vertices of total degree m s.t.

- every non-root black vertex has indegree 1 and m-2 half-edges

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- every non-root black vertex has indegree 1 and m-2 half-edges
- half-edges are incoming at black, outgoing at white, the tree is balanced

The enumeration of constellations

Theorem:[Bousquet-Mélou–S. 2000] *m*-eulerian maps are in bijection* with trees such that:

- white vertices carry m-1 sibblings (black vertices or half-edges)
- black vertices carry m-2 half-edges and a white child.



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- white vertices carry m-1 sibblings (black vertices or half-edges)
- black vertices carry m-2 half-edges and a white child.



Counting the trees: this is a familly of simple tree

$$A_{-\Box}(t) = (1 + A_{-\bullet}(t))^{m-1}, \quad A_{-\bullet}(t) = (m-1) \cdot A_{-\Box}(t)$$

or observe directly that they are (m-1)-ary trees with (m-1) types of edges

The enumeration of constellations

Theorem:[Bousquet-Mélou–S. 2000] *m*-eulerian maps are in bijection* with trees such that:

- white vertices carry m-1 sibblings (black vertices or half-edges)
- black vertices carry m-2 half-edges and a white child.



- that are balanced

Counting the trees: this is a familly of simple tree

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or observe directly that they are (m-1)-ary trees with (m-1) types of edges

$$\Rightarrow \frac{m}{(m-2)n+2} \frac{1}{(m-2)n+1} \binom{(m-1)n}{n} \cdot (m-1)^{n-1}$$

A formula for general factorizations [BMS00]

Theorem. Let $\lambda = 1^{\ell_1}, \ldots, n^{\ell_n}$ be a partition of n, and $\ell = \sum_i \ell_i$. The number of *m*-uple of permutations $(\sigma_1, \ldots, \sigma_m)$ such that

- (factorization) $\sigma_1 \cdots \sigma_m = \sigma$ with cycle type λ
- (transitivity) $\langle \sigma_1, \ldots, \sigma_m \rangle$ acts transitively on $\{1, \ldots, n\}$
- (minimality) the total rank of factors is $\sum_i r(\sigma_i) = n + \ell 2$

$$m \frac{((m-1)n-1)!}{(mn-(n+\ell-2))!} \cdot n! \cdot \prod_{i} \frac{1}{\ell_{i}!} {mi-1 \choose i}^{\ell_{i}}$$

Proofs:

is

(Bousquet-Mélou–Schaeffer 2000) (Goulden–?? 2009) (bijection + inclusion/exclusion)(gfs and differential eqns)

$$\lambda = n$$
, factorizations of *n*-cycles: $\frac{1}{(mn+1)} \binom{mn+1}{n} \cdot (n-1)!$

 $\lambda = 1^n$, identity factorizations: $\frac{m}{(m-2)n+2} \frac{(m-1)^{n-1}}{(m-2)n+1} \binom{(m-1)n}{n} \cdot (n-1)!$

From maps to trees: Hurwitz formula

$\alpha\text{-orientations}$ for increasing quadrangulations

Planar quadrangulations (faces are 4-gons) such that:

- faces have labels in $\{1, \ldots, 2n-2\}$
- around labeled vertices, face labels increase in ccw order
- around white vertices, face labels increase in cw order



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Orient each edge so that the minimum incident label is on the left Then each black vertex has indegree $\alpha_h(black) = m - 1$, outdegree 1 each white vertex has indegree $\alpha_h(white) = 1$.

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As before, this orientation is accessible, in fact strongly connected.



























A local rule to create increasing half edges



Two half-edges with same label \Rightarrow edge and face of degree 4



Iterate the local rules as long as possible...











adding buds







adding buds

Parings and adding buds again





adding buds



Parings and adding buds again





adding buds





Parings and adding buds again



Lemma. When it stops, there are only white half-edges left.



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Face number i defines transposition au_i . Lemma: the product is the identity permutation.

(6,7)(4,5)(3,4)(3,6)(2,5)(1,2)(5,6)(1,4)(2,7)(1,7)(3,7)(2,6)=(1)(2)(3)(4)(5)(6)(7)



Theorem[Duchi-Poulalhon-S. 2012] Closure is the reverse bijection between - simple Hurwitz trees of size n, and

– minimal transitive factorizations of the identity in S_n .
From simple Hurwitz trees to factorizations



Theorem[Duchi-Poulalhon-S. 2012] Closure is the reverse bijection between - simple Hurwitz trees of size n, and

- minimal transitive factorizations of the identity in S_n . type λ

Hurwitz formula for the number of minimal transitive factorizations in transpositions

Theorem. Let $\lambda = 1^{\ell_1}, \ldots, n^{\ell_n}$ be a partition n, and $\ell = \sum_i \ell_i$. The number of *m*-uples of transpositions (τ_1, \ldots, τ_m) such that

- (product cycle type) $\tau_1 \cdots \tau_m = \sigma$ has cycle type λ
- (transitivity) the associated graph is connected
- (minimality) the number of factors is $m = n + \ell 2$

$$n^{\ell-3} \cdot m! \cdot n! \cdot \prod_{i \ge 1} \frac{1}{\ell_i!} \left(\frac{i^i}{i!}\right)^{\ell_i}$$

Proofs:

is

(Hurwitz 1891, Strehl 1996) (Goulden–Jackson 1997) (Lando–Zvonkine 1999) (Bousquet-Mélou–Schaeffer 2000) (recurrences, Abel identities) (gfs and differential eqns) (geometry of LL mapping) (bijection + inclusion/exclusion)

 $\lambda = n$, factorizations of *n*-cycles: $n^{n-2} \cdot (n-1)!$ $\lambda = 1^n$, factorizations of the identity: $n^{n-3} \cdot (2n-2)!$

Arbres de Hurwitz de type λ et formule d'Hurwitz

Pour traiter le cas général de la formule il faut définir des arbres de Hurwitz de type λ : ce sont des arbres plans avec

- n-1 sommets noirs de degré 2 ou 1, étiqueté avec $\{1,\ldots,n-1\}$
- ℓ sommets blancs dont ℓ_i portent i séparateurs et i-1 feuilles noires
- $m = n + \ell 2$ arêtes avec étiquettes distintes dans $\{1, \ldots, m\}$
- les arêtes sont croissantes en sens direct entre 2 séparateurs



$$H_n = n^{n-2}(n-1)! \quad H_1 n = n^{n-3}(2n-2)! \quad H_{\lambda} = n^{\ell-3}m!n! \prod_{i \ge 1} \frac{1}{\ell_i!} \left(\frac{i^i}{i!}\right)^{\ell_i}$$

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- les arêtes sont croissantes en sens direct entre 2 séparateurs

Lemme. Le nb d'arbres d'Hurwitz de type λ est $n^{\ell-3}m!n!\prod_{i\geq 1}\frac{1}{\ell_i!}\left(\frac{i^i}{i!}\right)^{\ell_i}$

Théorème La clôture s'étend en une bijection des arbres de Hurwitz de type λ avec les factorisations minimales transitives en transpositions de permutations de type cyclique λ .

Corollare: La formule d'Hurwitz.

$$H_n = n^{n-2}(n-1)! \quad H_1 n = n^{n-3}(2n-2)! \quad H_{\lambda} = n^{\ell-3}m!n! \prod_{i \ge 1} \frac{1}{\ell_i!} \left(\frac{i^i}{i!}\right)^{\ell_i}$$

Conclusion

- Cayley trees are plane trees and
 - Hurwitz formula counts variant of Cayley trees
- A second strategy (and proof) using Hurwitz mobiles also extends to higher genus
- Open problems:
 - double Hurwitz numbers
 - inequivalent factorisations in transpositions

Post Scriptum

Lately we realized that we should have found this much earlier...



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That's all!