

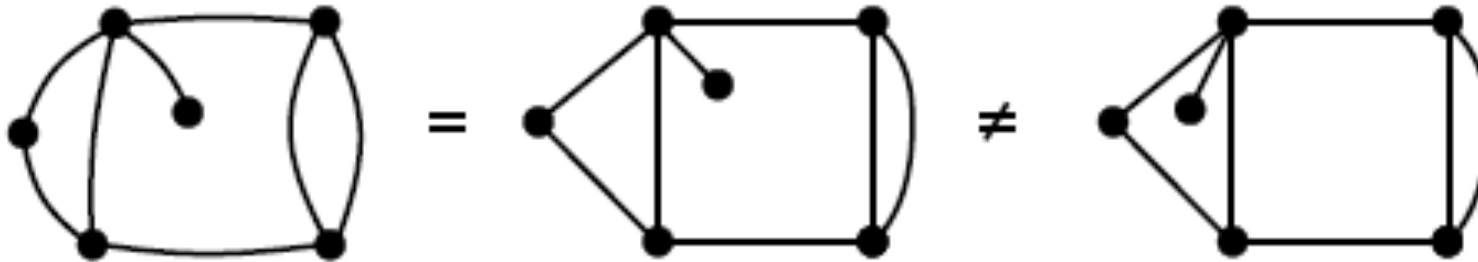
On the diameter of random planar graphs

Eric Fusy (LIX, Ecole Polytechnique)

joint work with Guillaume Chapuy, Omer Gimenez, Marc Noy

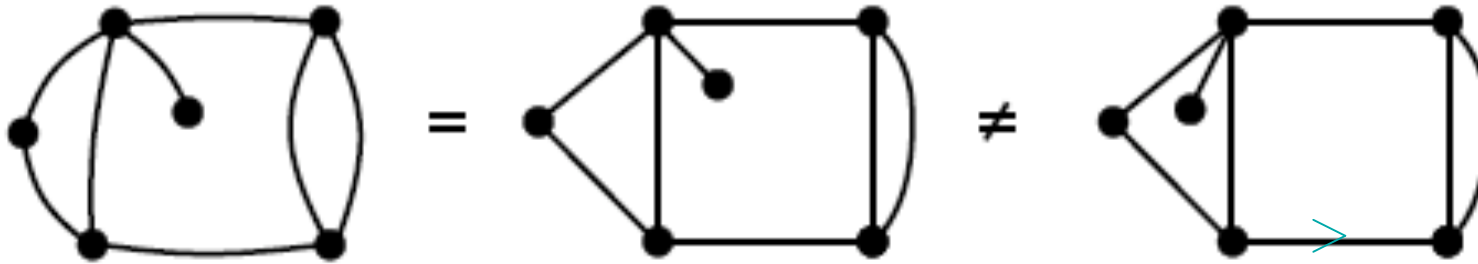
Planar maps / planar graphs

- Planar map = planar embedding of a graph



Planar maps / planar graphs

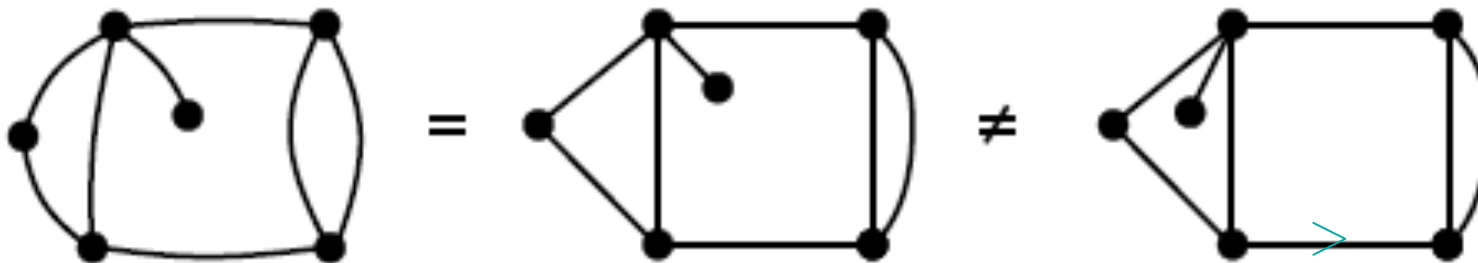
- Planar map = planar embedding of a graph



(rooting is enough to avoid symmetries)

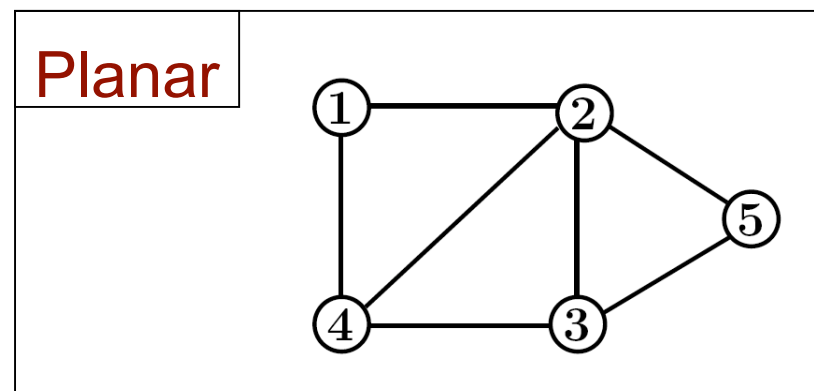
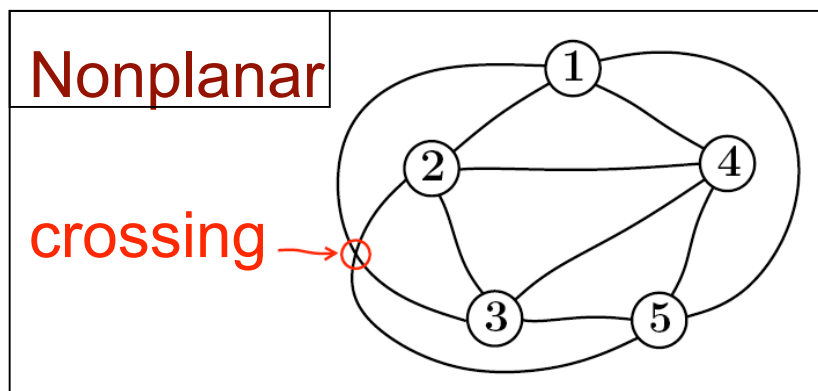
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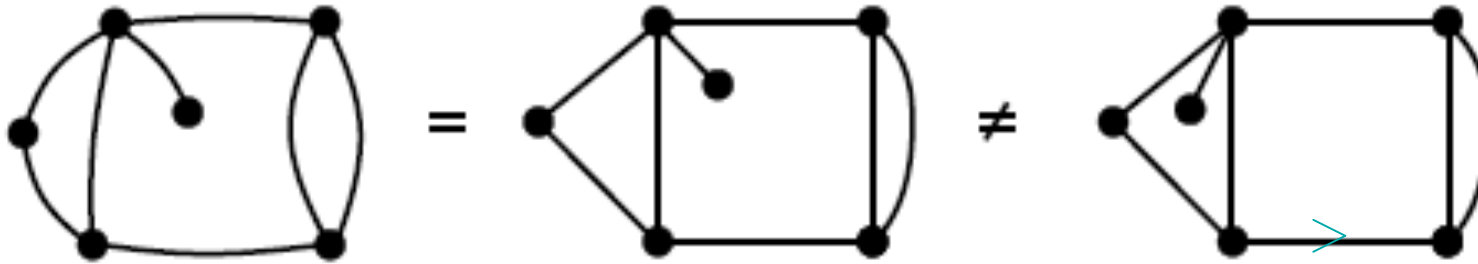
- Planar graph = graph with at least one planar embedding



(needs to label all vertices to avoid symmetries)

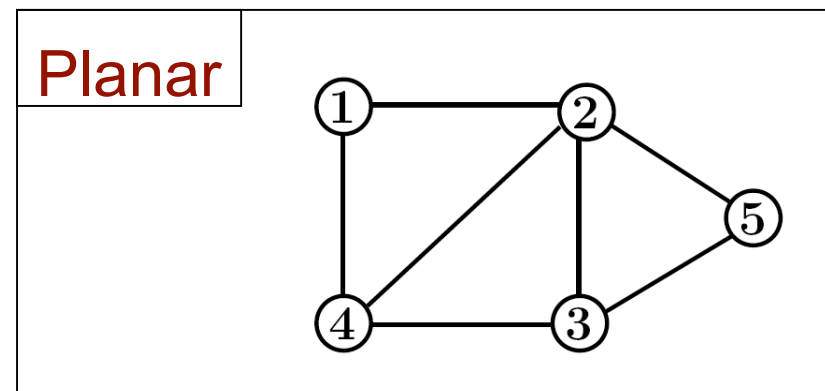
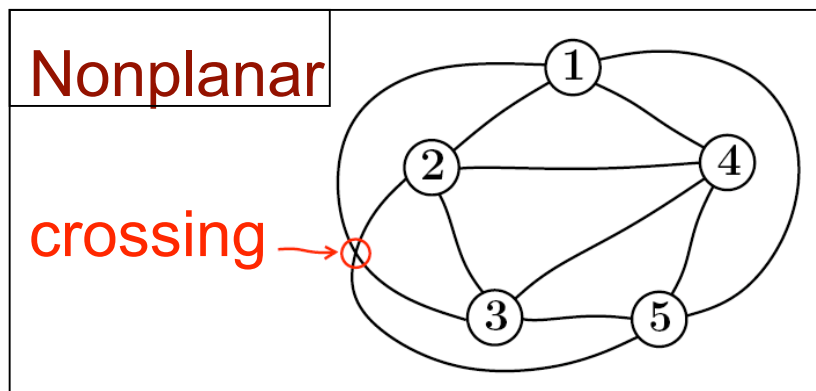
Planar maps / planar graphs

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- Random planar graph $G_n \neq$ random planar map M_n

Main result

Theorem: Let G_n be the random (unembedded) planar graph with n vertices. Then

$$\mathbb{P}(\text{Diam}(G_n) \notin [n^{1/4-\epsilon}, n^{1/4+\epsilon}]) = O(\exp(-n^{\Theta(\epsilon)}))$$

Conjecture:

$$\frac{1}{n^{1/4}} \text{Diam}(G_n) \text{ converges in law}$$

Related results

Random planar structures

Random map M_n :

- $\text{Radius}(M_n)/n^{1/4}$ converges in law

[Chassaing-Schaeffer'04], [Marckert-Miermont'06],...

Random tree T_n :

- $h(T_n)/n^{1/2}$ converges in law
- $\mathbb{P}(h(T_n) \notin [n^{1/2-\epsilon}, n^{1/2+\epsilon}]) = O(\exp(-n^{\Theta(\epsilon)}))$

[Flajolet et al'93]

Random graph (Erdős-Rényi)

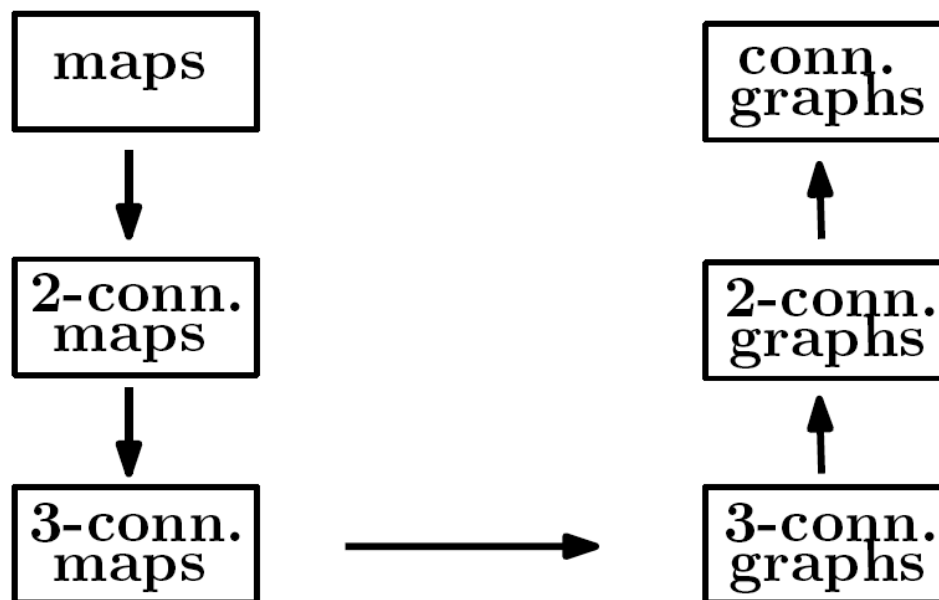
For each fixed $\lambda > 1$

$\text{Diam}(G(n, p = \lambda/n)) \sim c_\lambda \log(n)$
with high probability

[Riordan-Wormald'09]

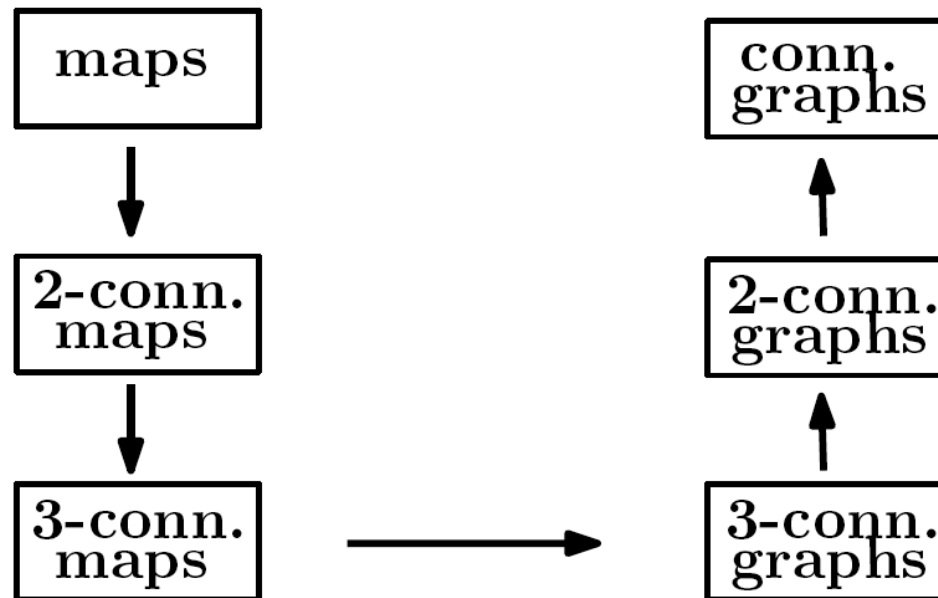
Overview

- We recall the classical scheme:



Overview

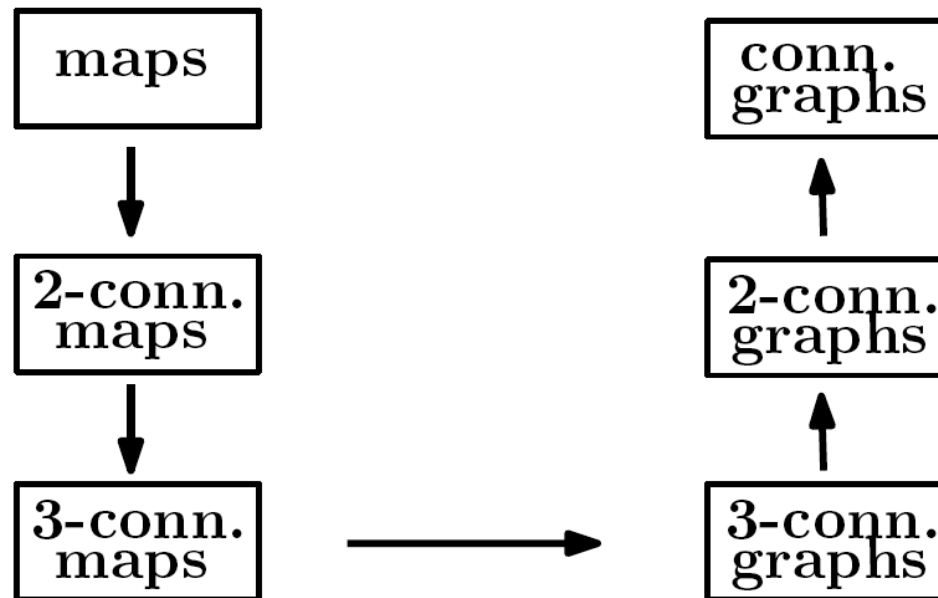
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- Can carry out **counting** (exact/asymptotic) along the scheme:
[Bender, Gao, Wormald'02], [Giménez, Noy'05]

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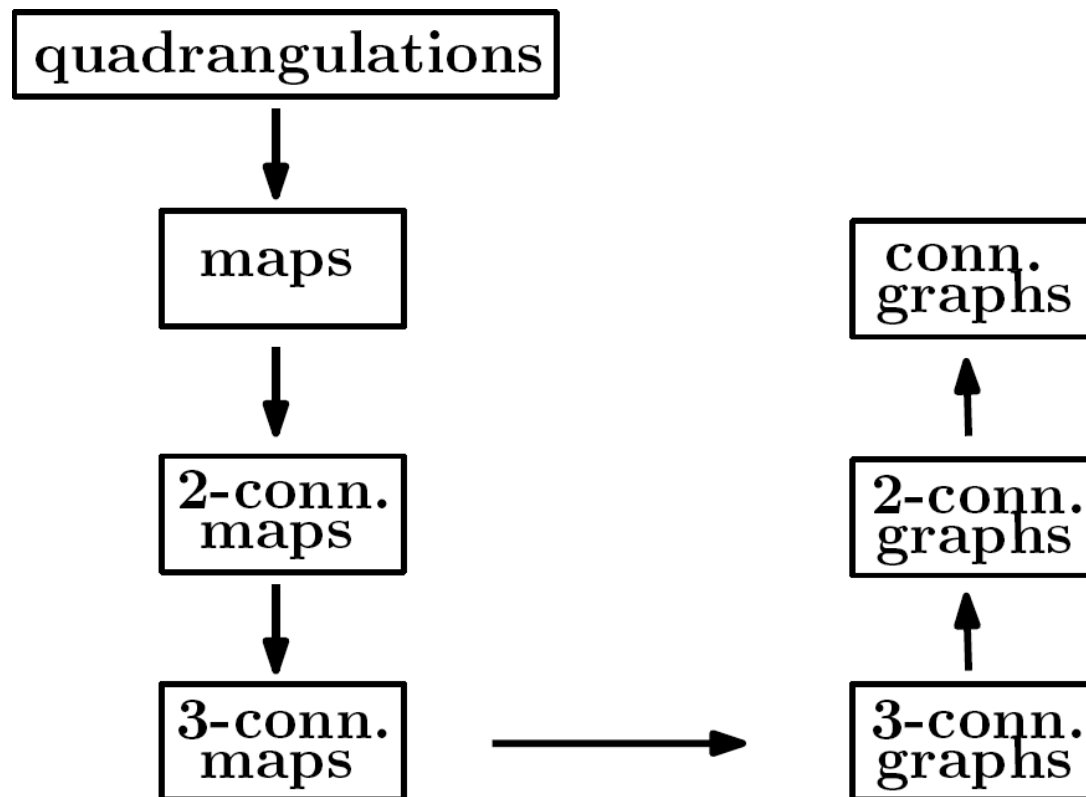


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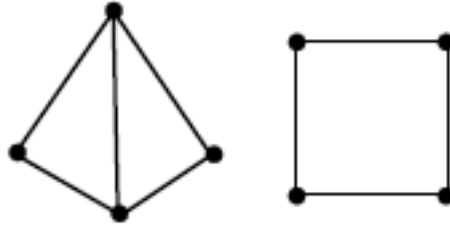
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Part I: counting scheme for connected planar graphs

Connectivity in graphs

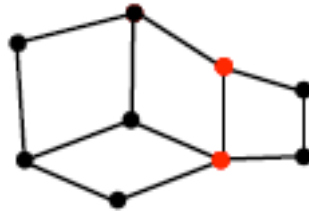
General



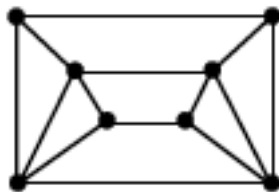
Connected



2-connected

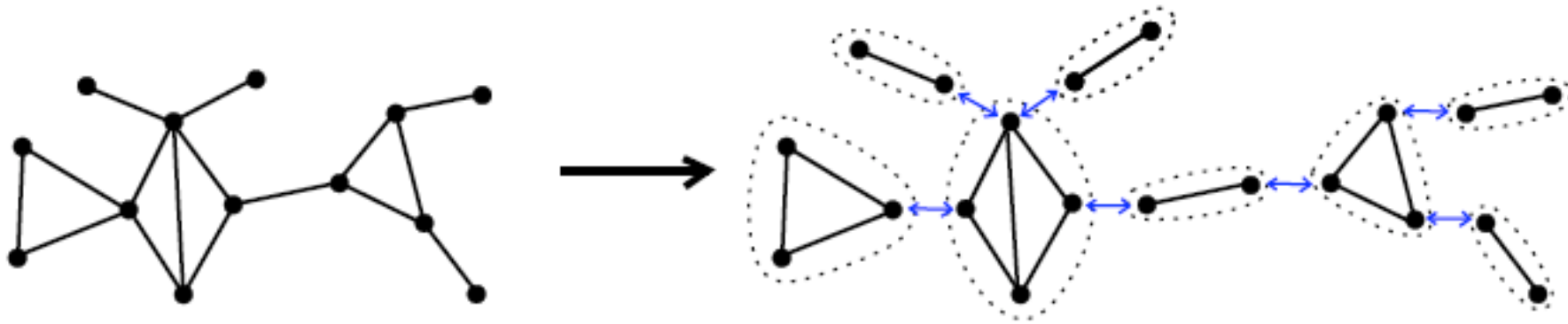


3-connected



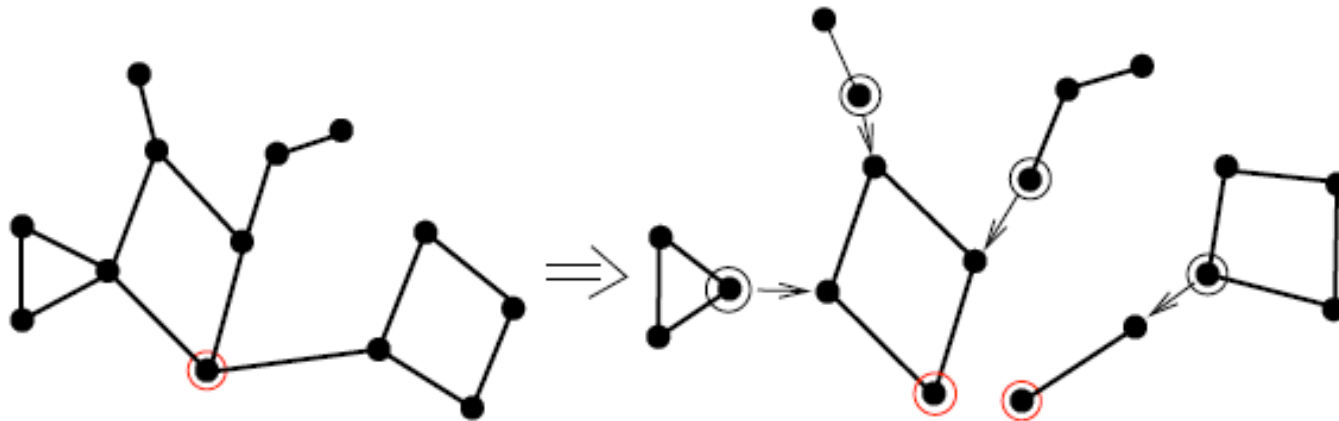
- [Tutte'66]:
 - a connected graph decomposes into 2-connected components
 - a 2-connected graph decomposes into 3-connected components

Connected \rightarrow 2-connected



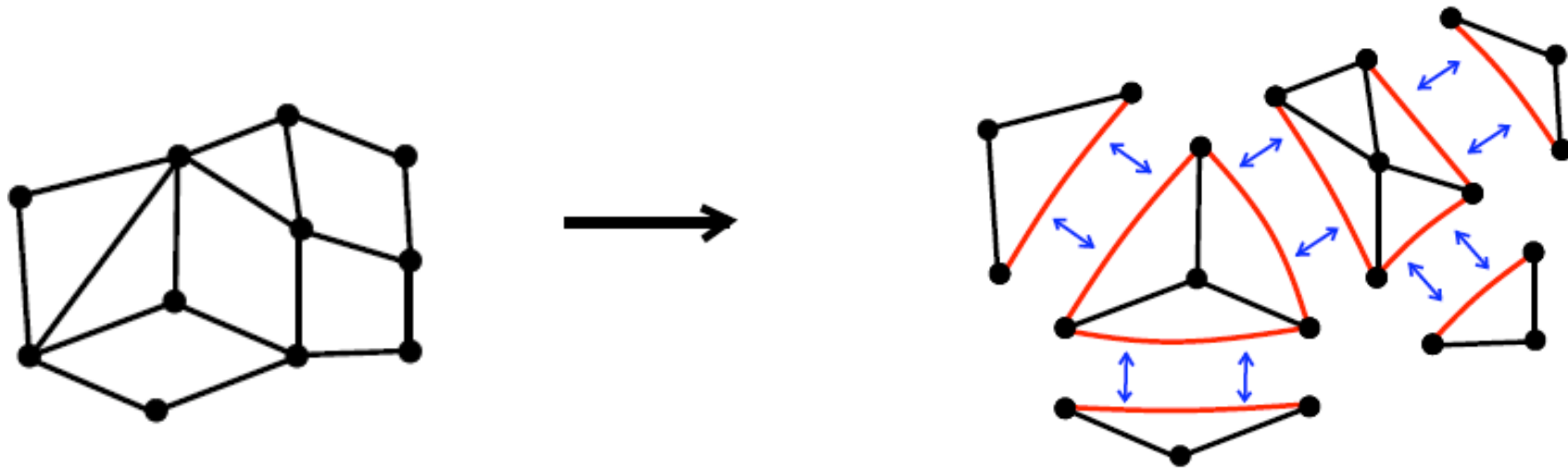
The incidences between cut-vertices and blocks forms a tree

- **Vertex-rooted** formulation:

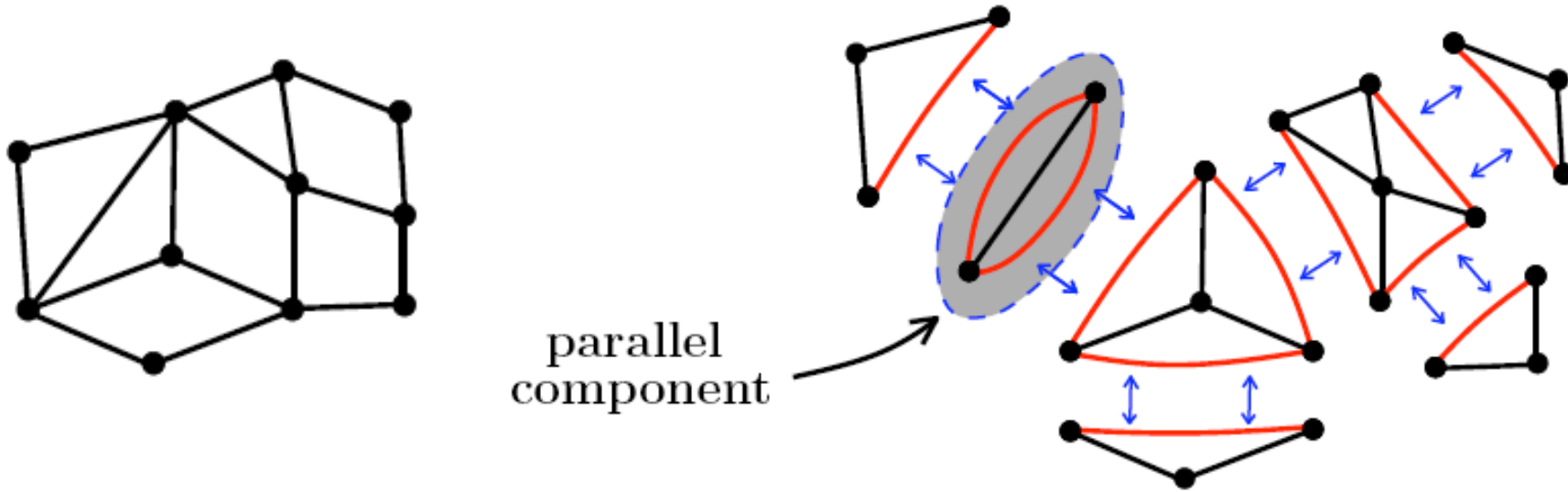


$$\Rightarrow C^\bullet(x, y) = x \exp(B'(C^\bullet(x, y), y))$$

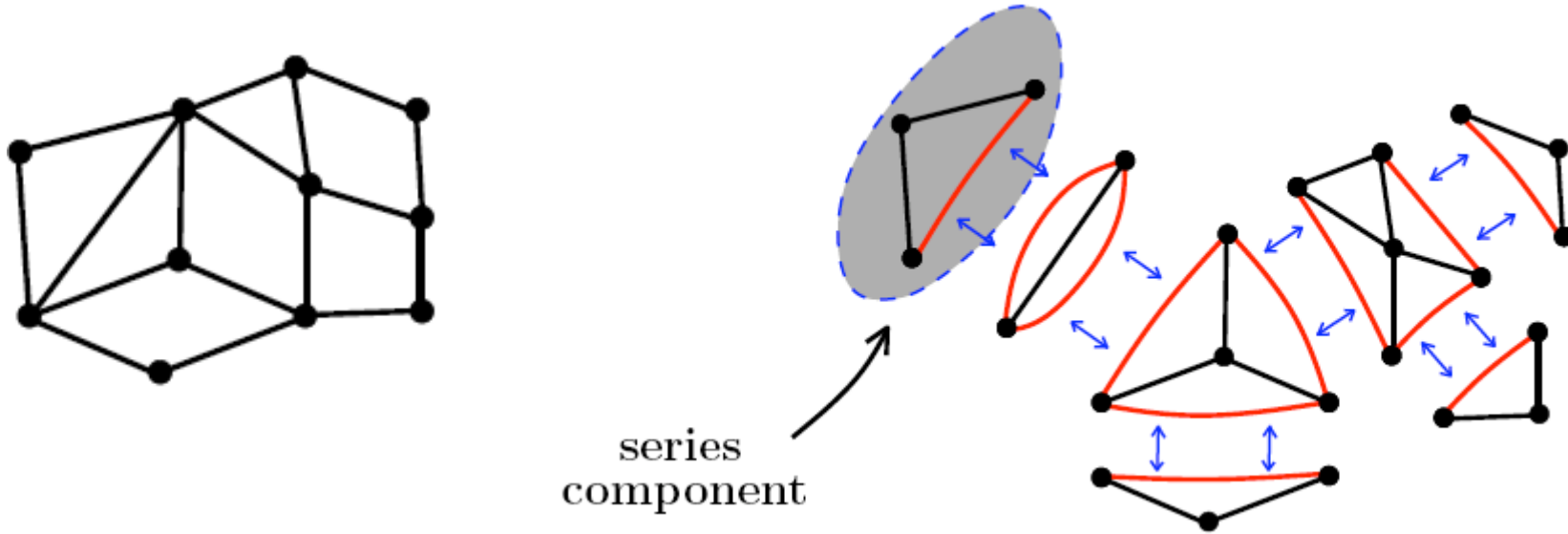
2-connected \rightarrow 3-connected



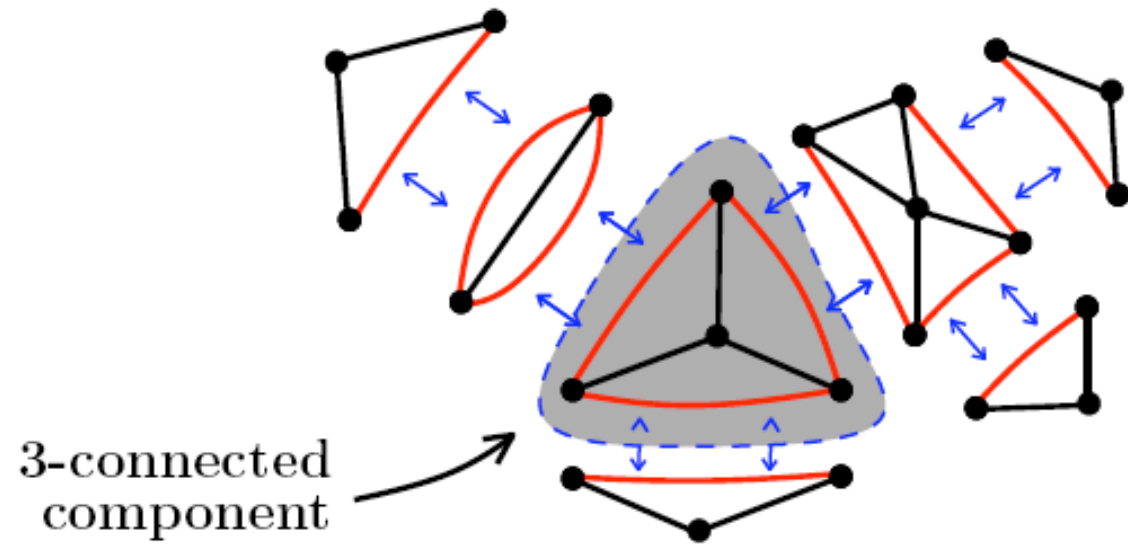
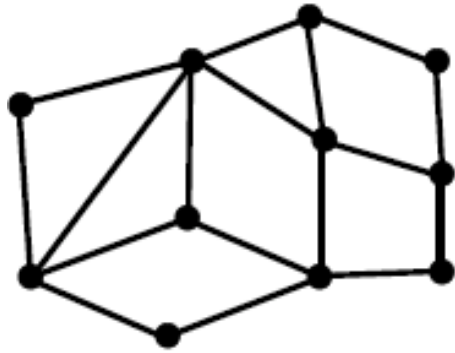
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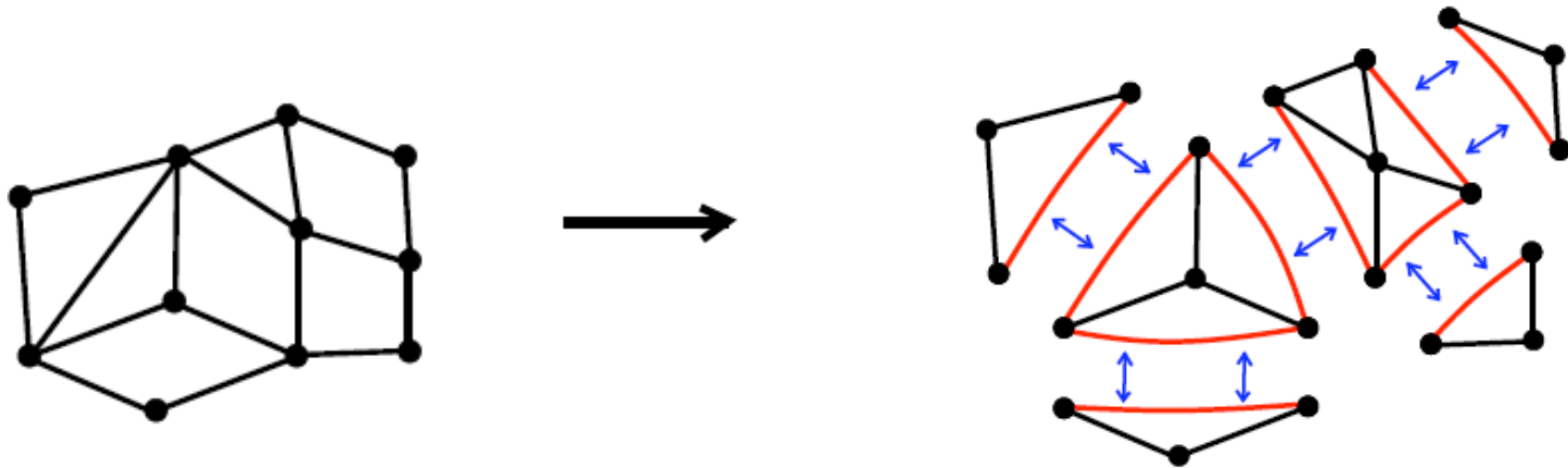
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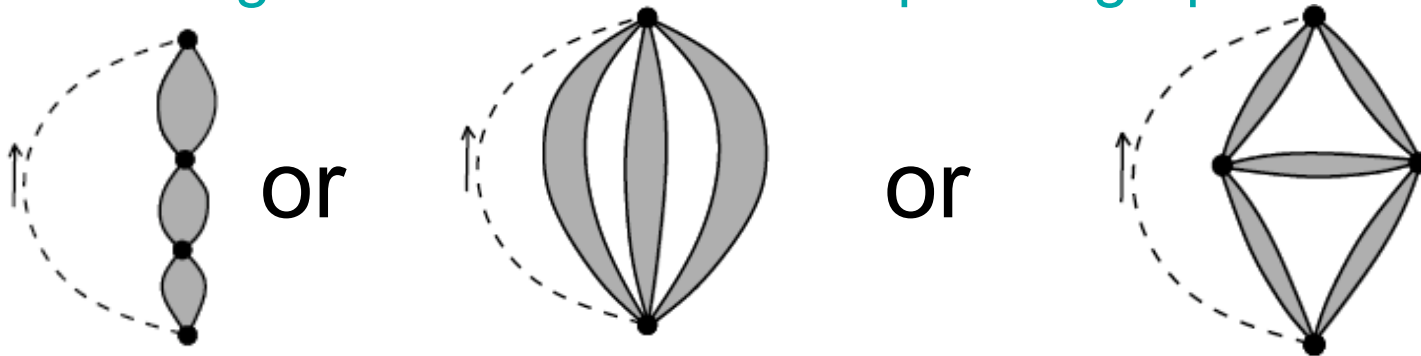
2-connected \rightarrow 3-connected



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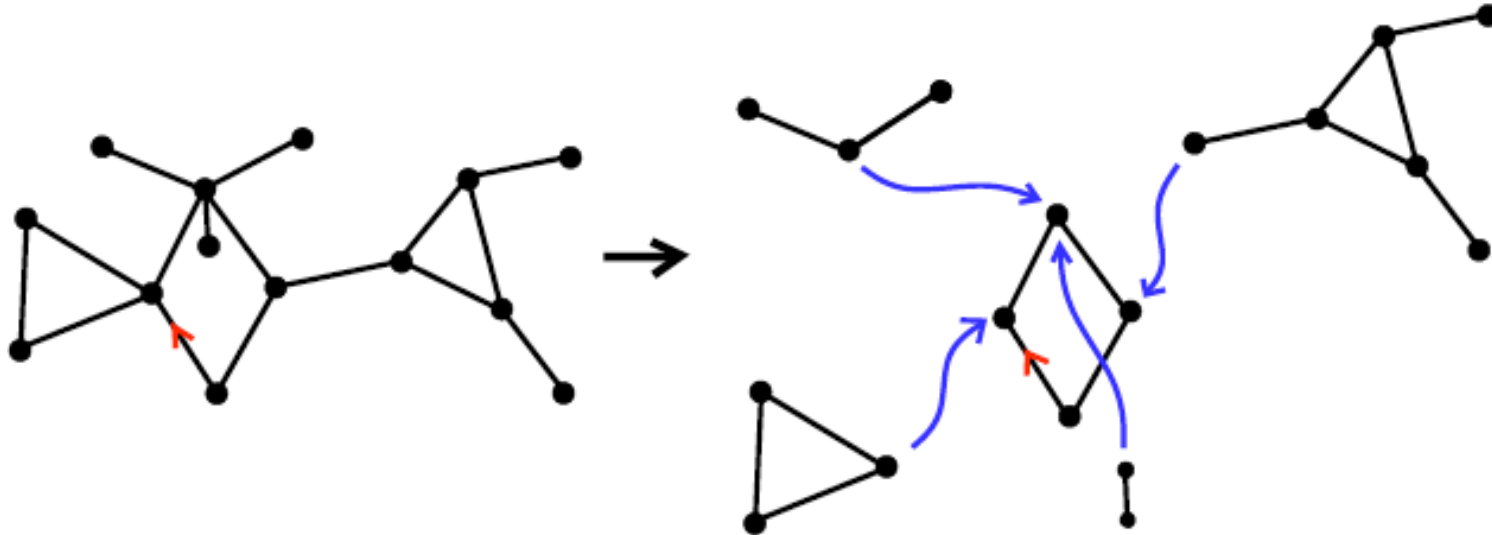
- **Edge-rooted** formulation [Trakhtenbrot'58]:
An **edge-rooted 2-connected planar graph** is either:



$$\overrightarrow{G}_2(x, y) \approx \overrightarrow{G}_3 \left(x, \overrightarrow{G}_2(x, y) \right) \quad (\text{modulo series-parallel operations})$$

3-connected graphs \rightarrow maps

- [Whitney]: 3-connected planar graphs \simeq 3-connected maps
- [Tutte'63]: Use the decomposition the reverse way and take the embedding into account

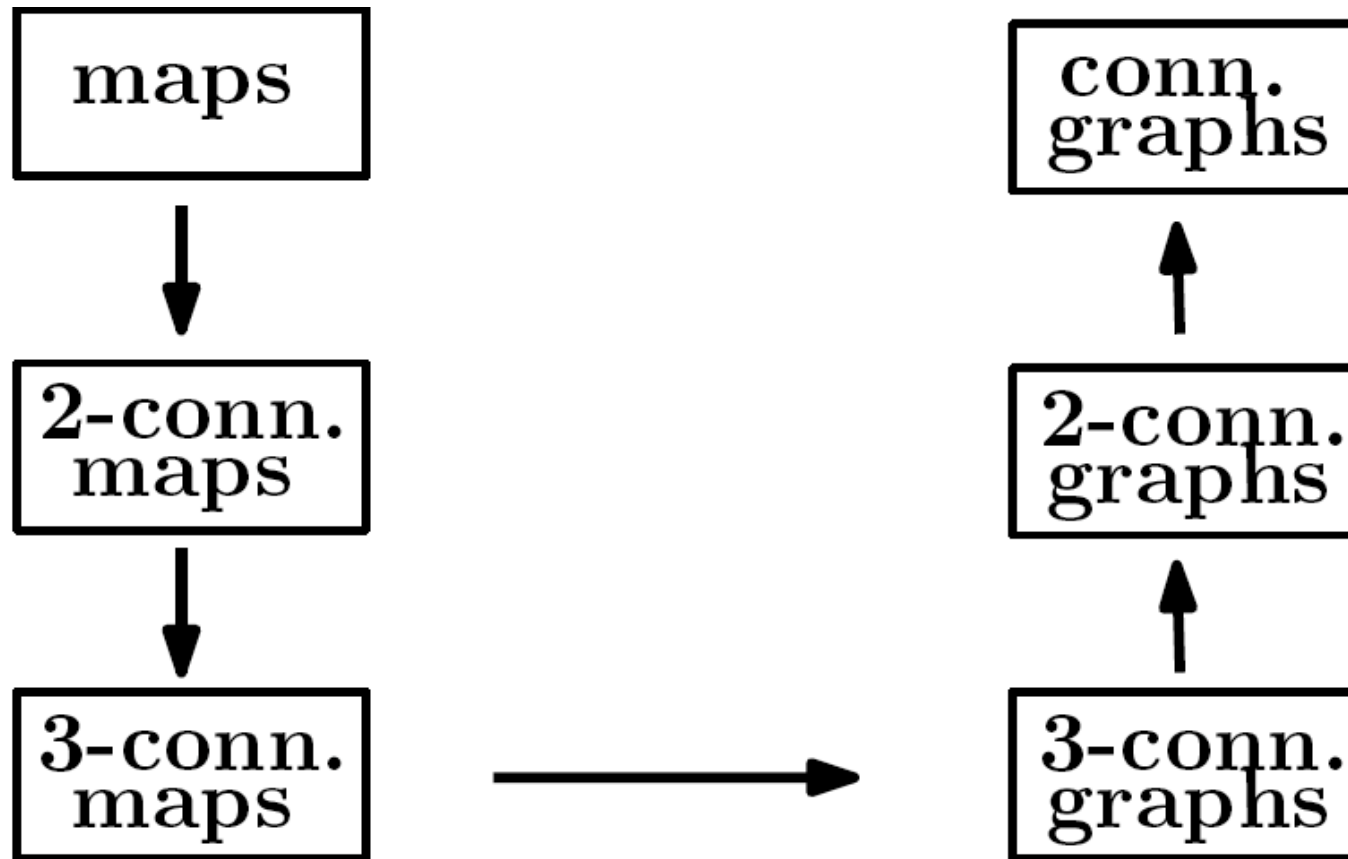


$$\vec{M}_1(y) = \vec{M}_2(y(1 + \vec{M}_1(y))^2)$$

\Rightarrow Can extract $\vec{M}_1(y)$ from $\vec{M}_2(y)$

(similarly can extract $\vec{M}_2(y)$ from $\vec{M}_3(y)$)

Count graphs reduces to count maps

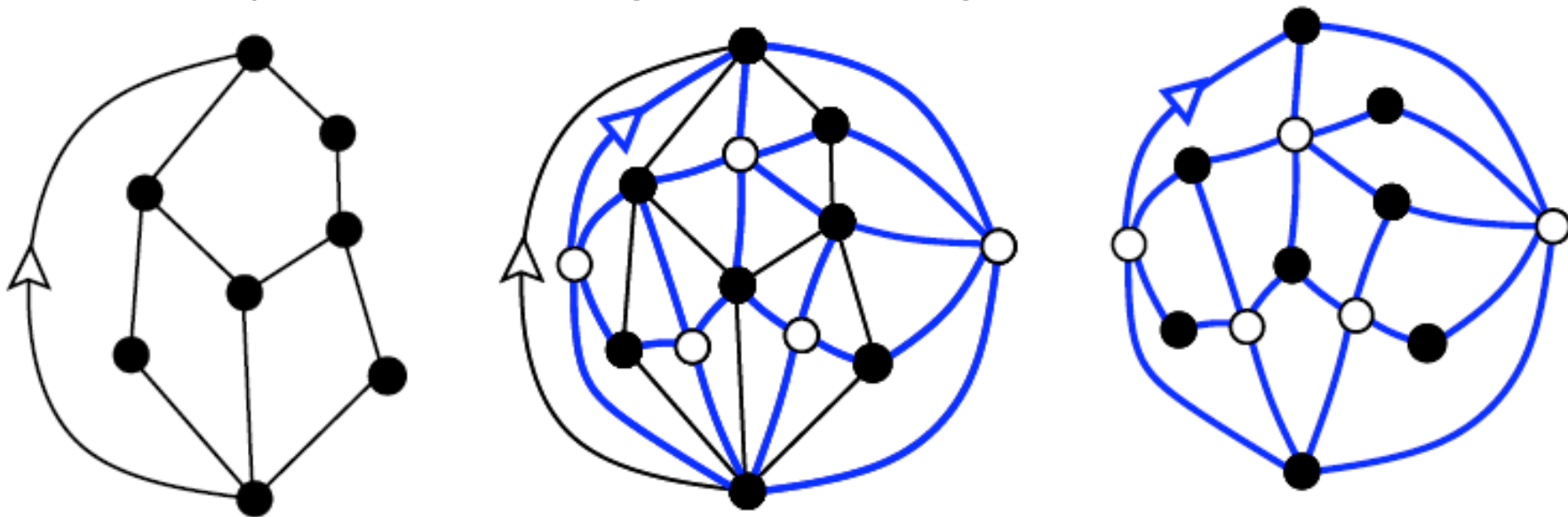


Count maps via quadrangulations

- There are several methods to count (rooted) maps:
 - **recursive** [Tutte'63]
 - **bijjective**: either directly or **via quadrangulations**

Count maps via quadrangulations

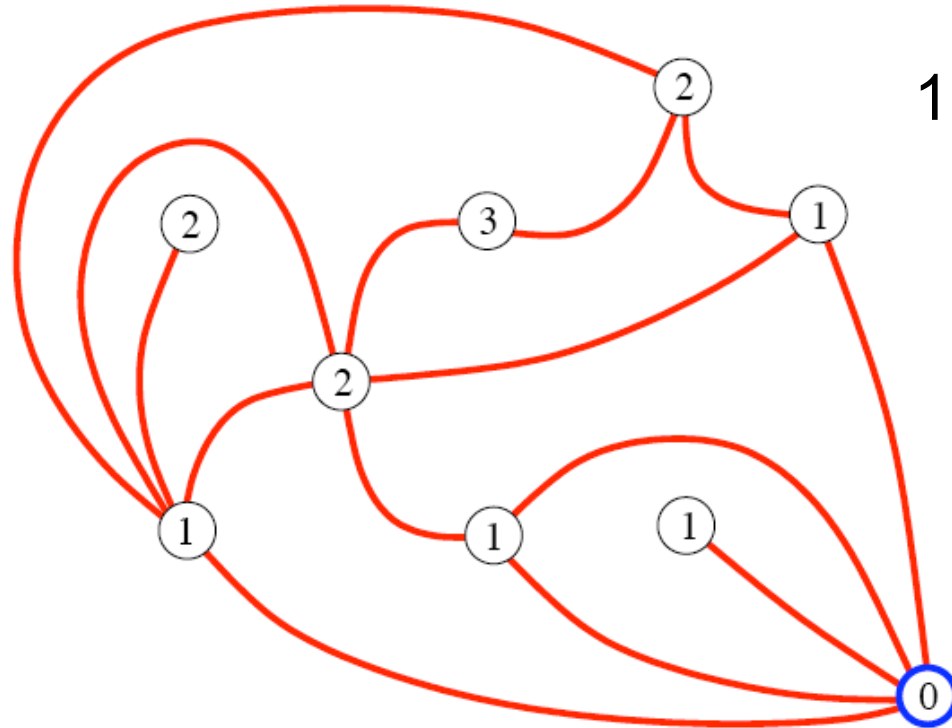
- There are several methods to count (rooted) maps:
 - **recursive** [Tutte'63]
 - **bijective**: either directly or **via quadrangulations**
- Bijection between maps and quadrangulations ([Tutte]):
 - 1) **Insert a star** inside each face
 - 2) Delete the edges of the original map



\Rightarrow #rooted maps n edges = #rooted quadrangulations n faces

Quad. \leftrightarrow well-labelled trees

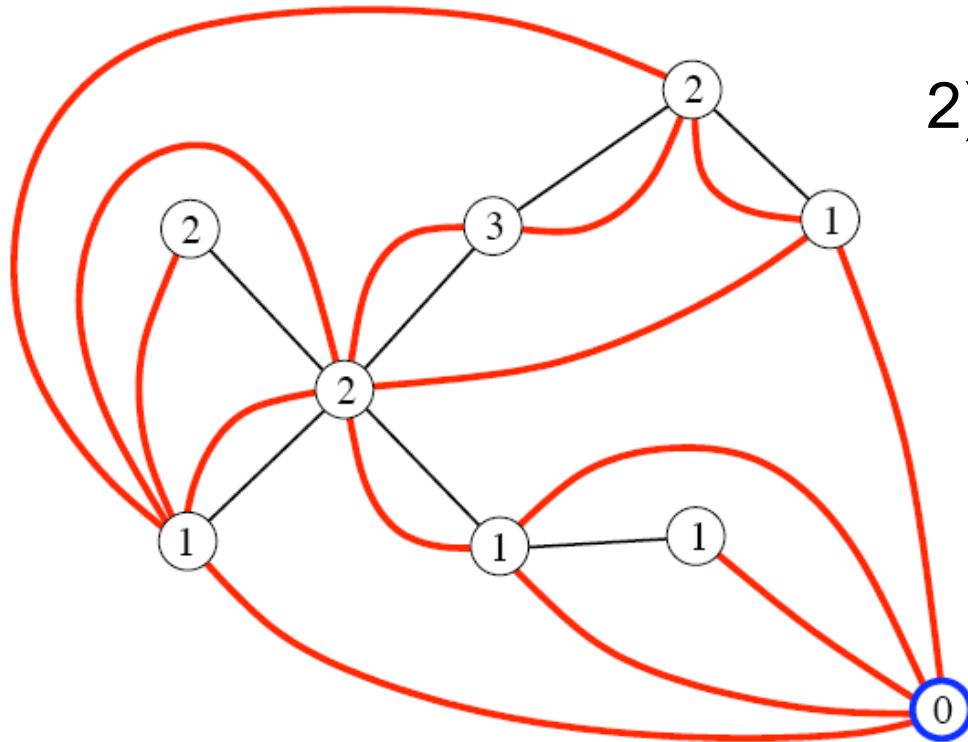
Bijection [Schaeffer'98] (also [Cori, Vauquelin'84]):



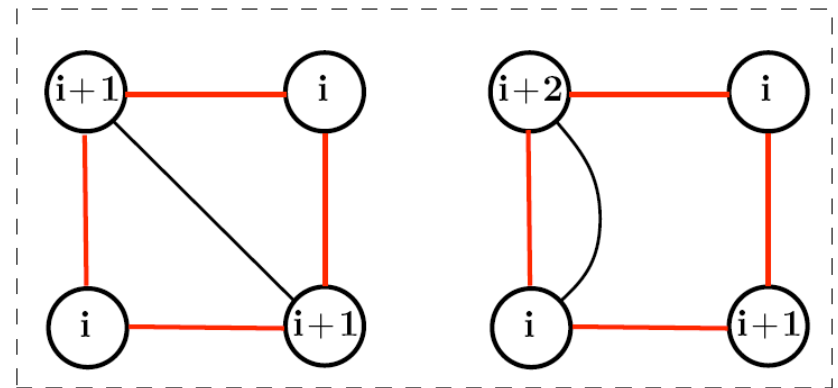
- 1) Start from a vertex-pointed quadrangulation with **vertices labelled by the distance** from the pointed vertex

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Bijection [Schaeffer'98] (also [Cori, Vauquelin'84]):

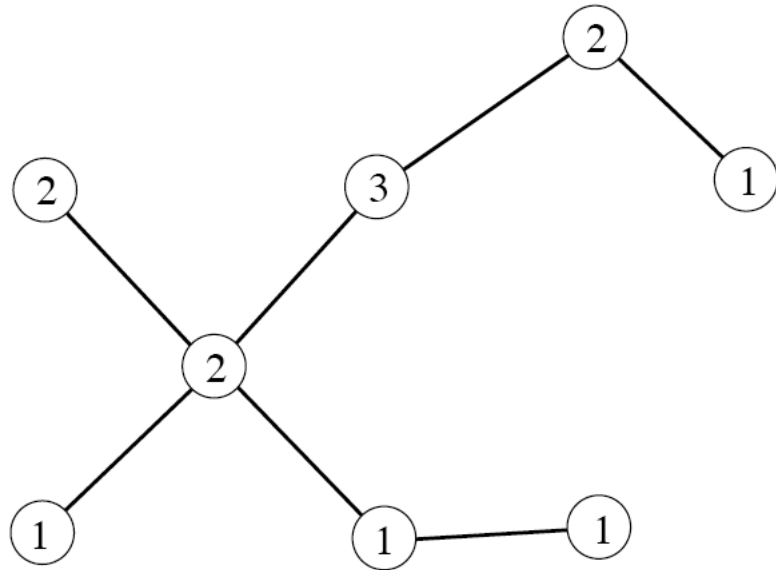


2) Draw a **black edge** in each **face** according to the rule:



Quad. \leftrightarrow well-labelled trees

Bijection [Schaeffer'98] (also [Cori, Vauquelin'84]):

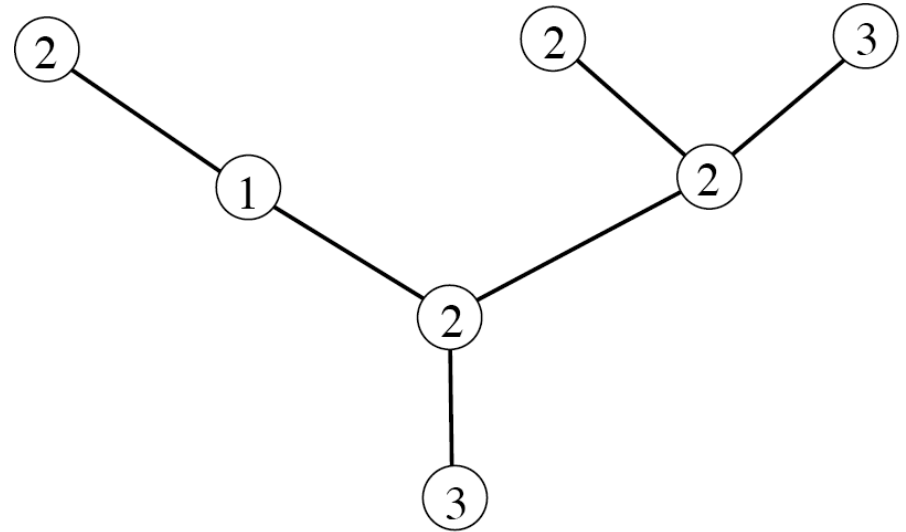


3) **Delete** the red edges and the pointed vertex

Quad. \leftrightarrow well-labelled trees

- A well-labelled tree is a plane tree where:

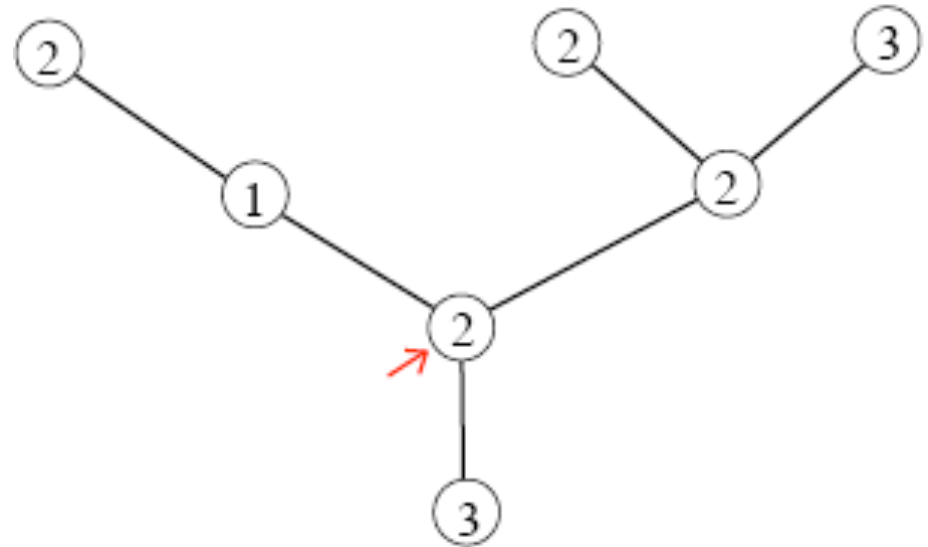
- each vertex v has a **positive** label
- the labels at each edge (v,v') differ by at most 1
- **at least one** vertex has **label 1**



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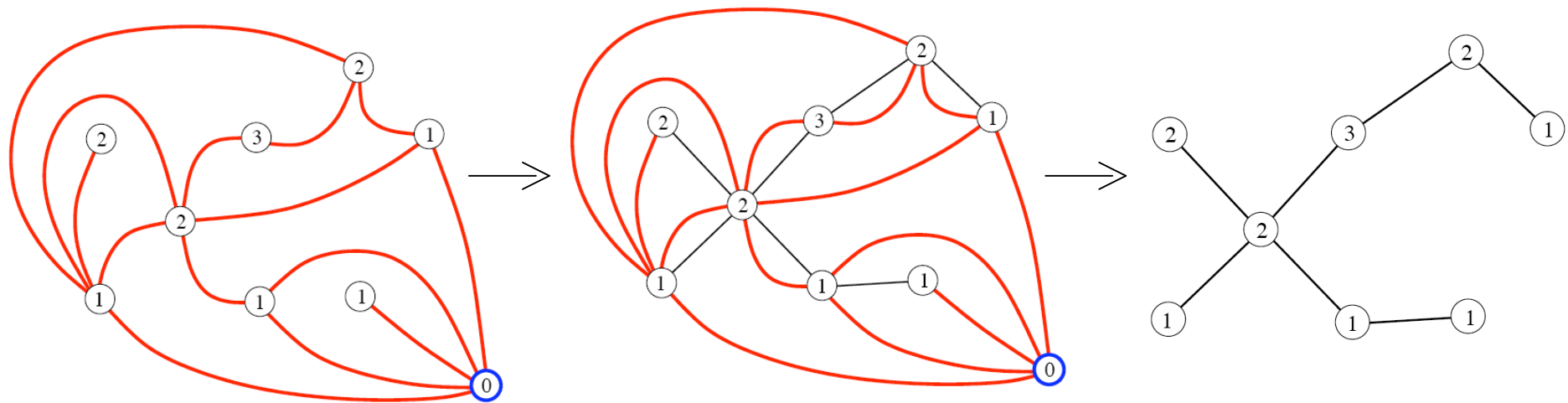


- **Rooted** well-labelled tree = well-labelled tree + **marked corner**

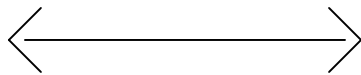
(there are $3^n \frac{(2n)!}{n!(n+1)!}$ such trees with n edges)

The mapping is a bijection

Theorem [Schaeffer'98]: The mapping is a **bijection** between **pointed quadrangulations** and **well-labelled trees**

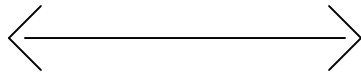


vertex label i



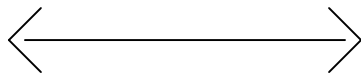
vertex at distance i

corner label i



edge at level i

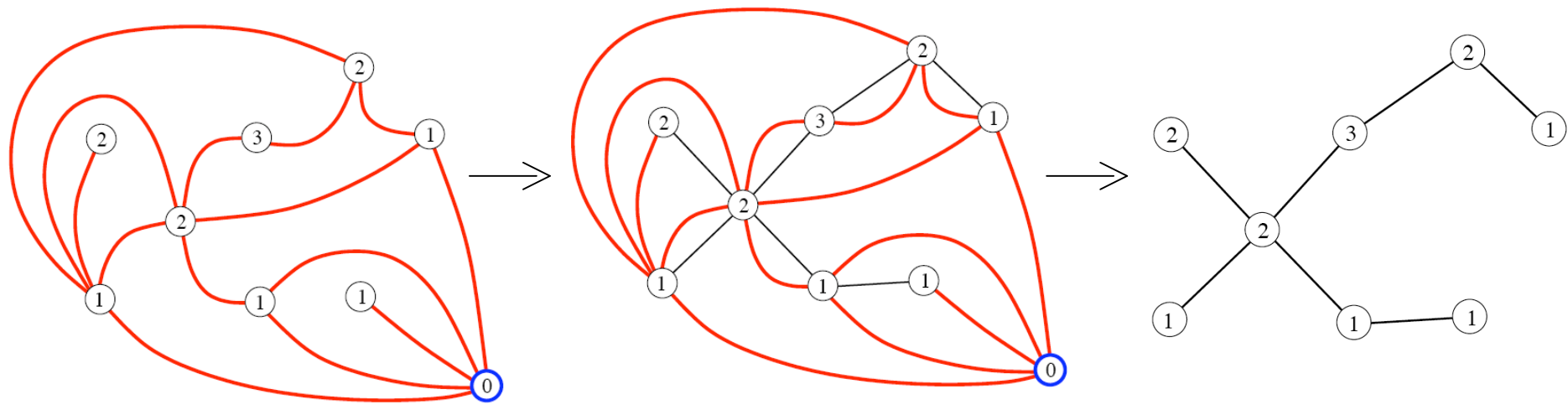
edge



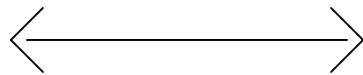
face

The mapping is a bijection

Theorem [Schaeffer'98]: The mapping is a bijection from well-labelled trees to pointed quadrangulations

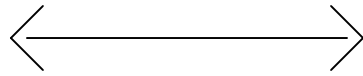


vertex label i



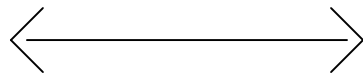
vertex at distance i

corner label i



edge at level i

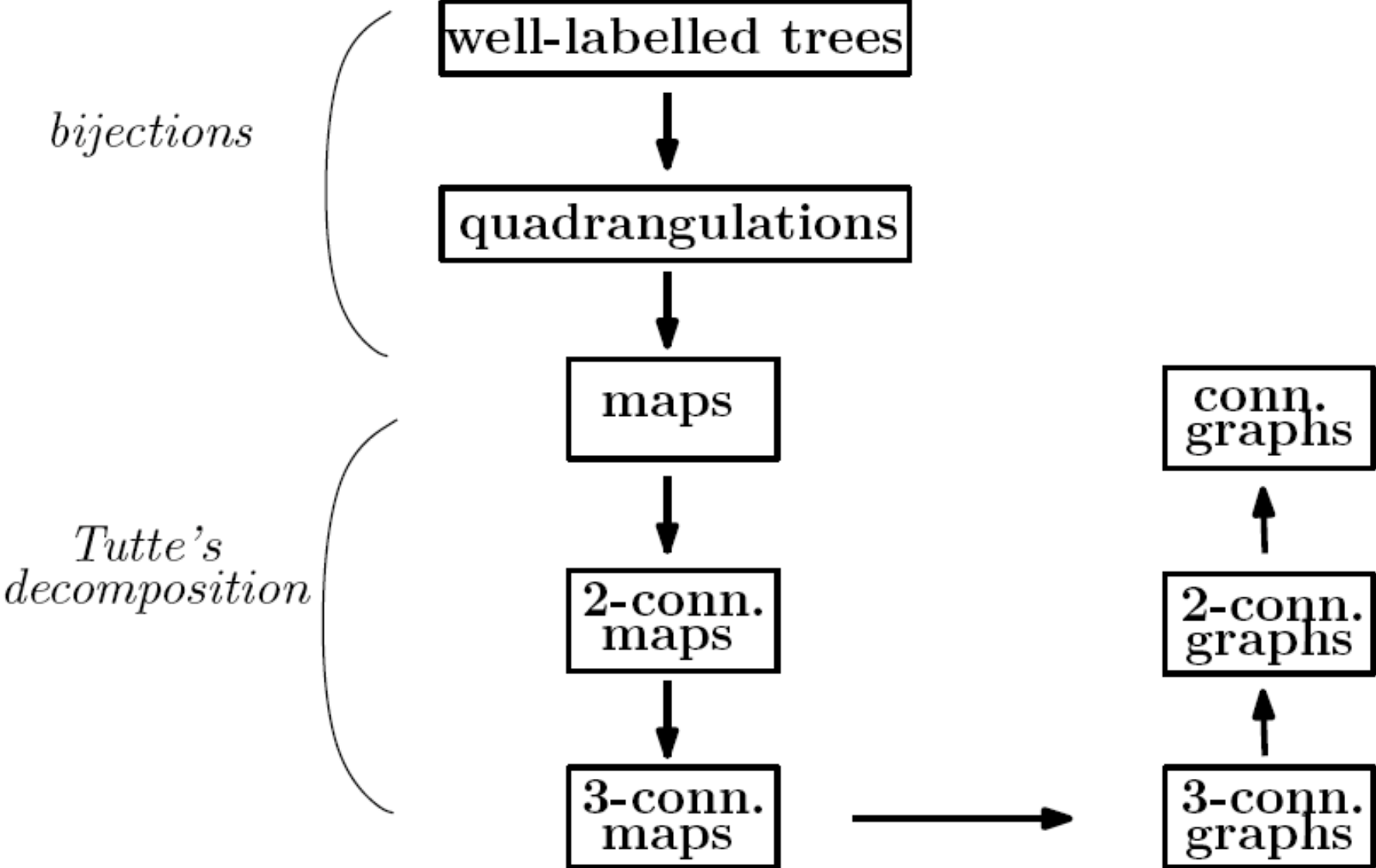
edge



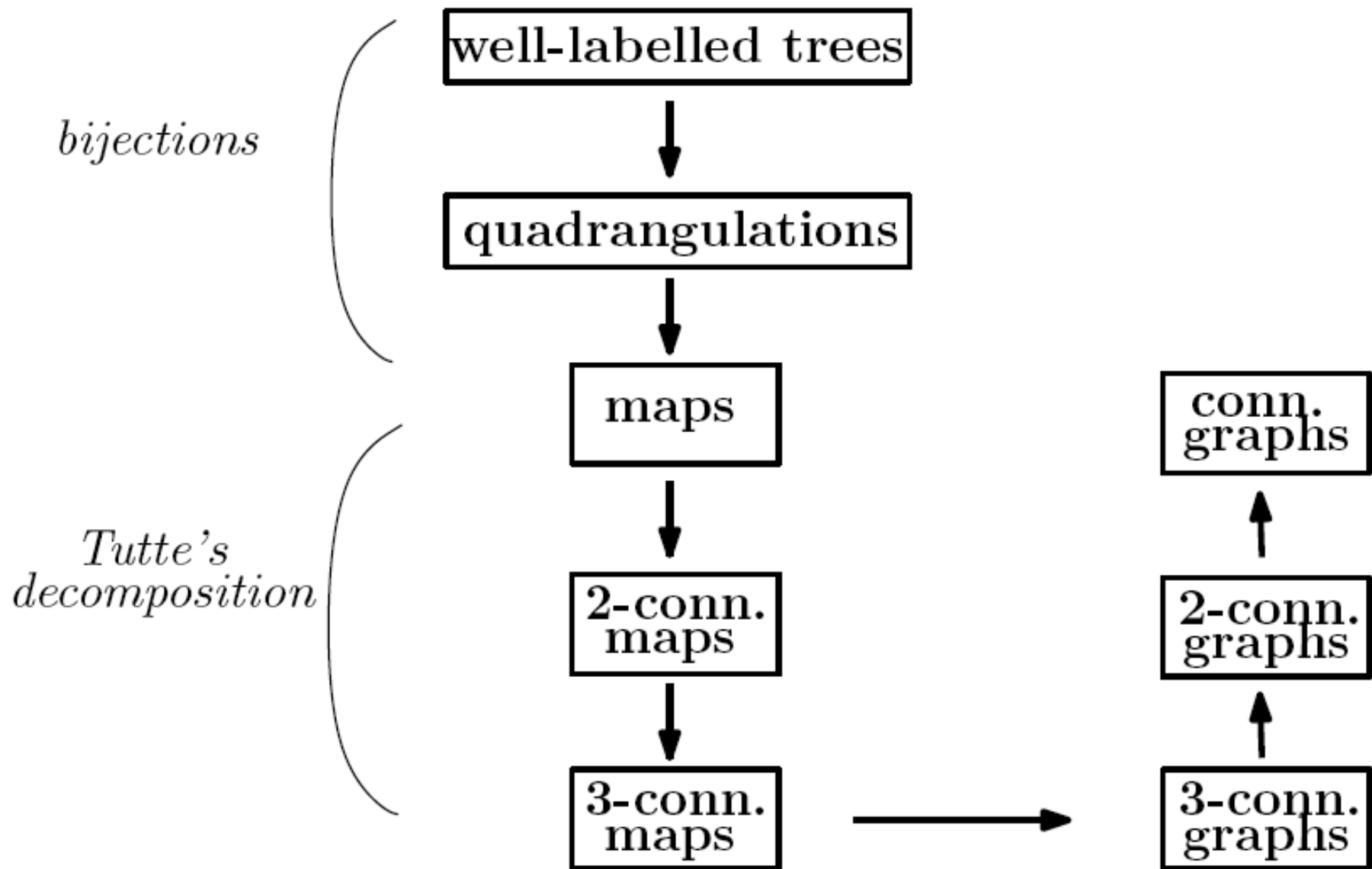
face

Corollary: there are $3^n \frac{(2n)!}{n!(n+1)!}$ quadrangulations with n faces, a marked vertex, and a marked edge

Summary of counting scheme

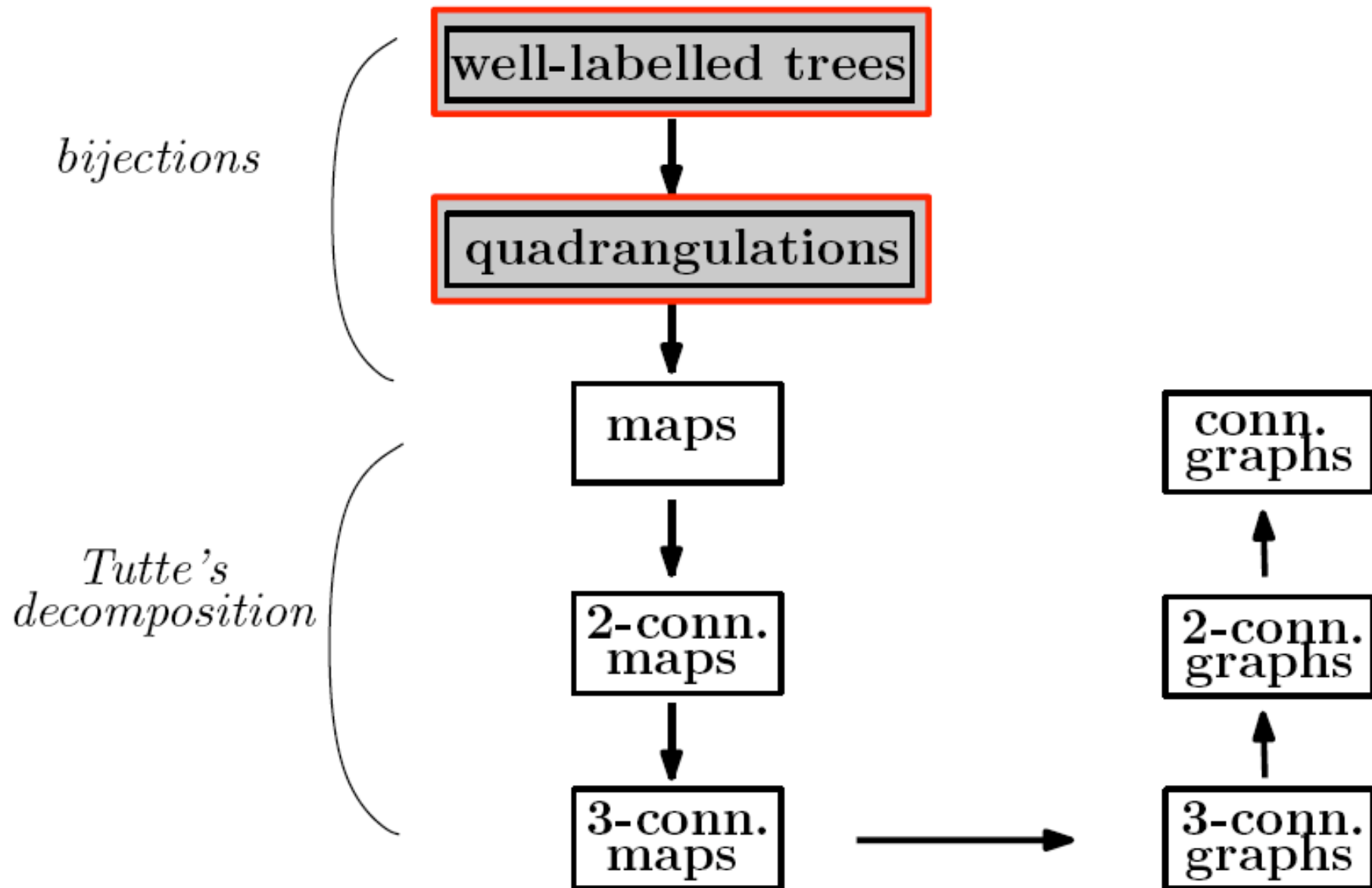


**Part II: carry large deviation
estimates for the diameter
along the counting scheme**



- We carry **large deviation** statements of the form:

$$\mathbb{P}(\text{Diam}(G_n) \notin [n^{1/4-\epsilon}, n^{1/4+\epsilon}]) = O(\exp(-n^{\Theta(\epsilon)}))$$



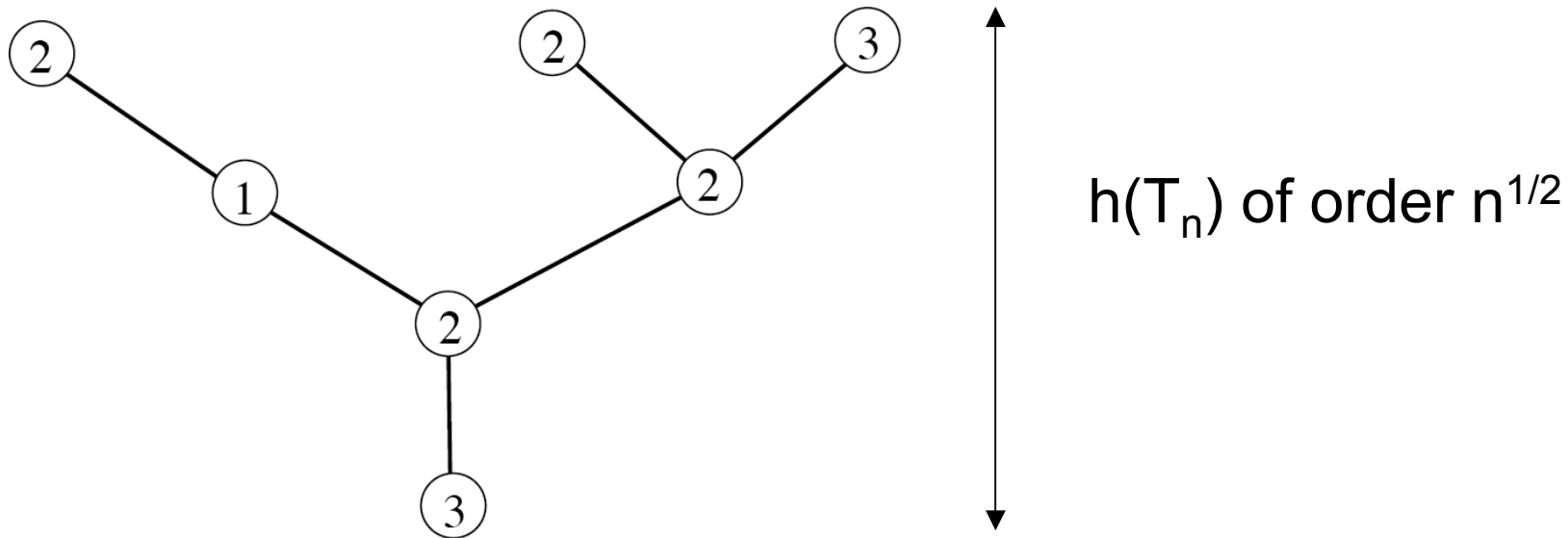
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Quad. via well-labelled trees

Let Q_n random pointed quadrangulation with $n+1$ vertices

$Q_n \simeq T_n$ the random well-labelled tree with n vertices



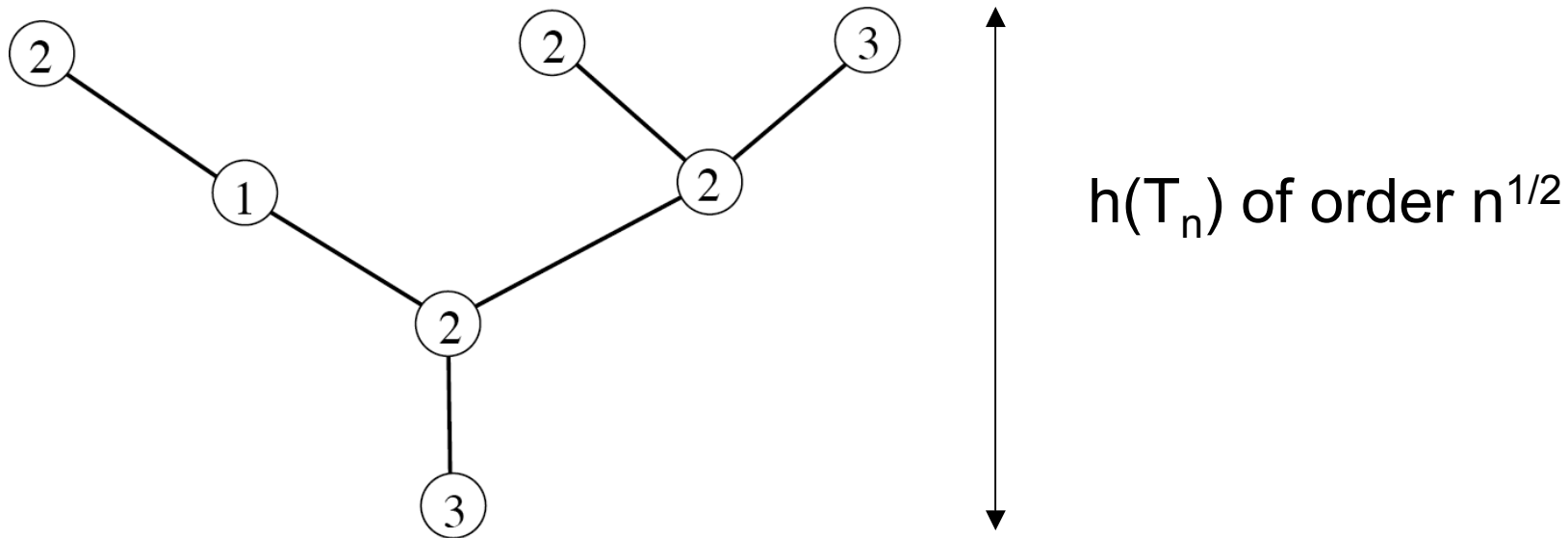
+ discrepancy of labels along k -path is of order $k^{1/2}$

Radius(Q_n) "of order" $n^{1/4}$

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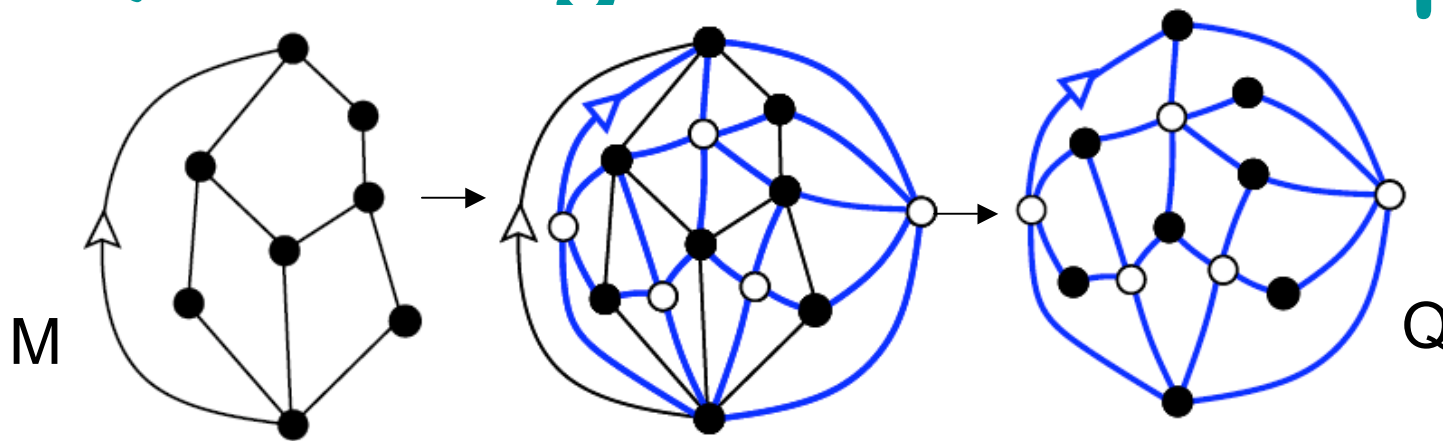


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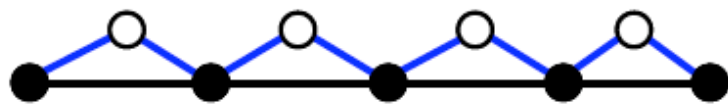
easily: $\text{Radius}(Q_n) \in [n^{1/4-\epsilon}, n^{1/4+\epsilon}]$ a.a.s.

(Precise convergence results in [Chassaing Schaeffer'04])

Quadrangulations -> Maps



$\text{dist}_Q \approx \text{dist}_M$ modulo max face-degree Δ :



$$\text{dist}_Q(v, v') \leq 2 \cdot \text{dist}_M(v, v')$$



$$\text{dist}_M(v, v') \leq \Delta \cdot \text{dist}_Q(v, v')$$

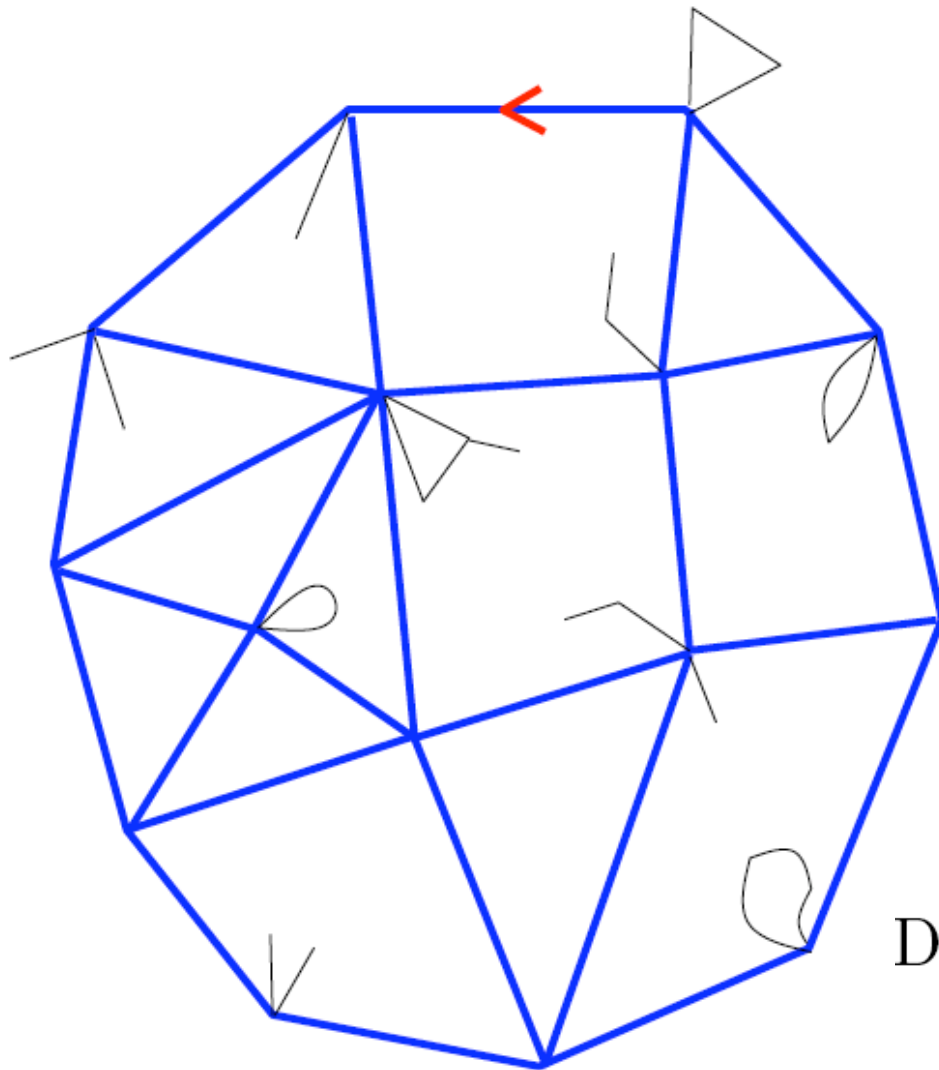
Δ is small: $\mathbf{P}_n(\Delta \geq k) = O(\exp(-ck))$

$\Rightarrow \text{Diam}(M_n) \in [n^{1/4-\epsilon}, n^{1/4+\epsilon}]$ a.a.s.

Maps \rightarrow 2-connected maps

[Gao, Wormald'99], [Banderier, Flajolet, Schaeffer, Soria'01]:

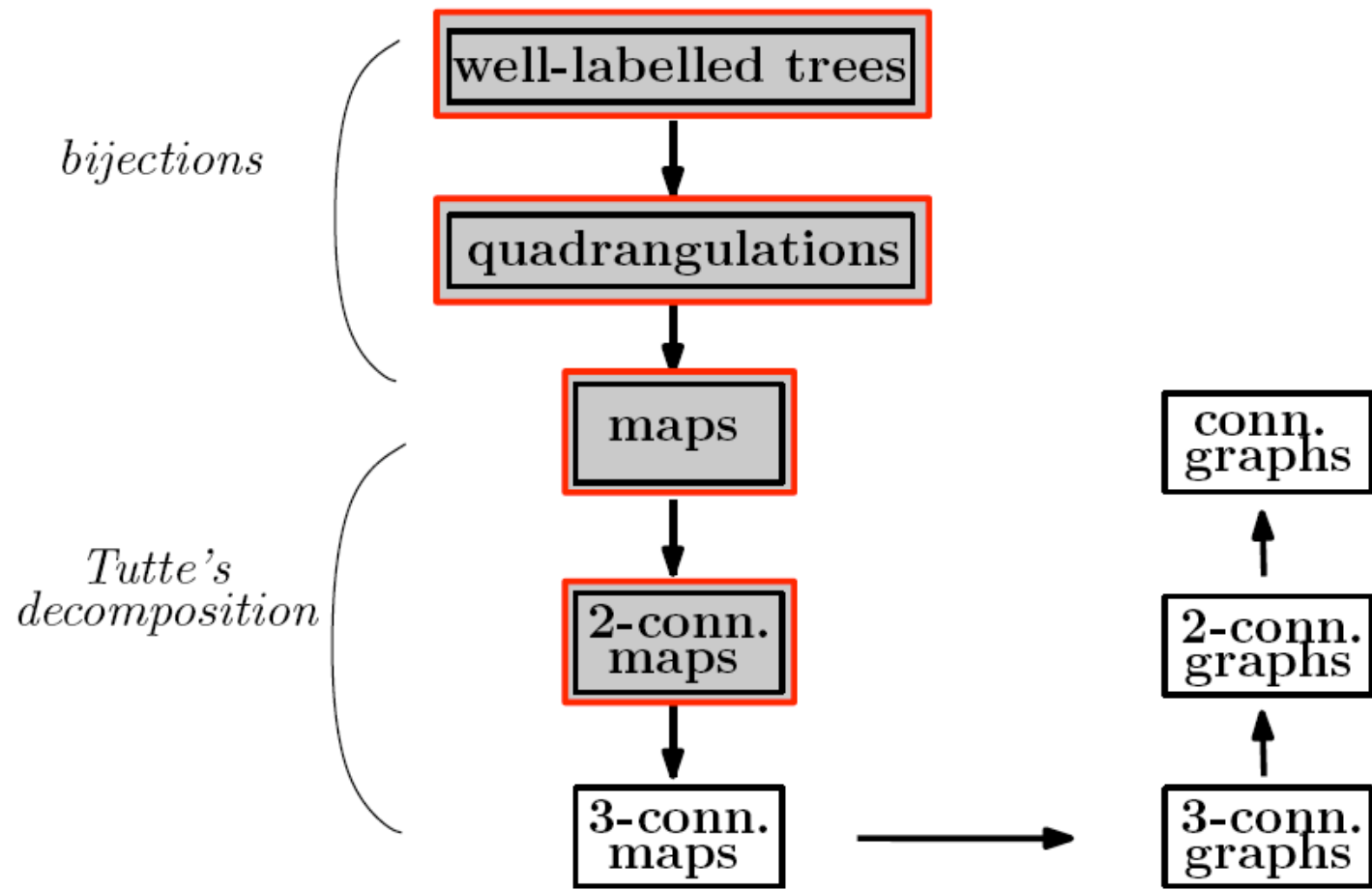
- The random map M_n has typically a **core of size $n/3 + n^{2/3} \theta$**
- The **other components are small** (size $O(n^{2/3})$)



Airy law

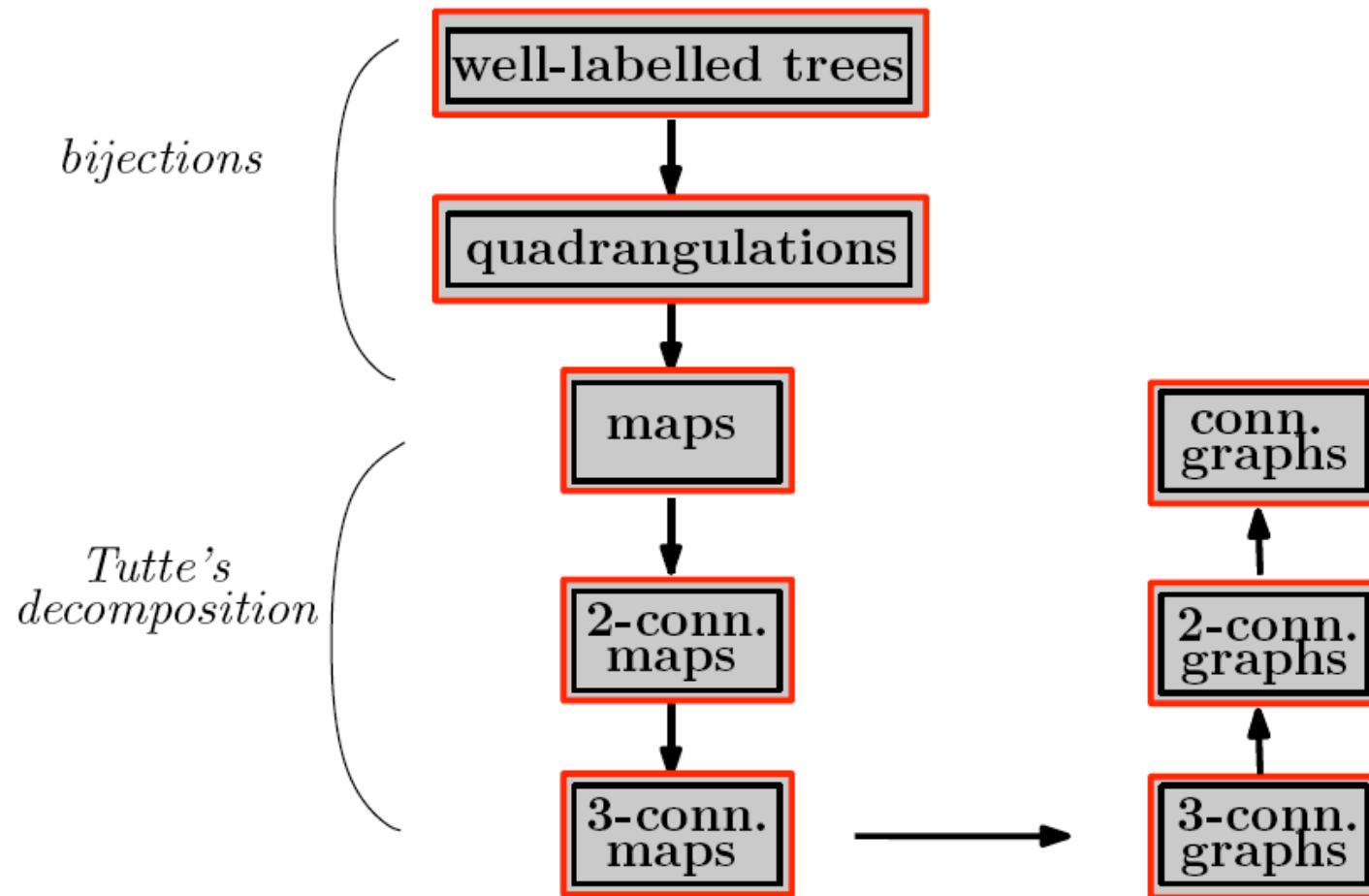
$$\text{Diam}(M_n) \text{ “} \approx \text{” } \text{Diam}(B_{n/3})$$

$$\text{Diam}(B_n) \in [n^{1/4-\epsilon}, n^{1/4+\epsilon}] \text{ a.a.s.}$$



- For each of the highlighted families:

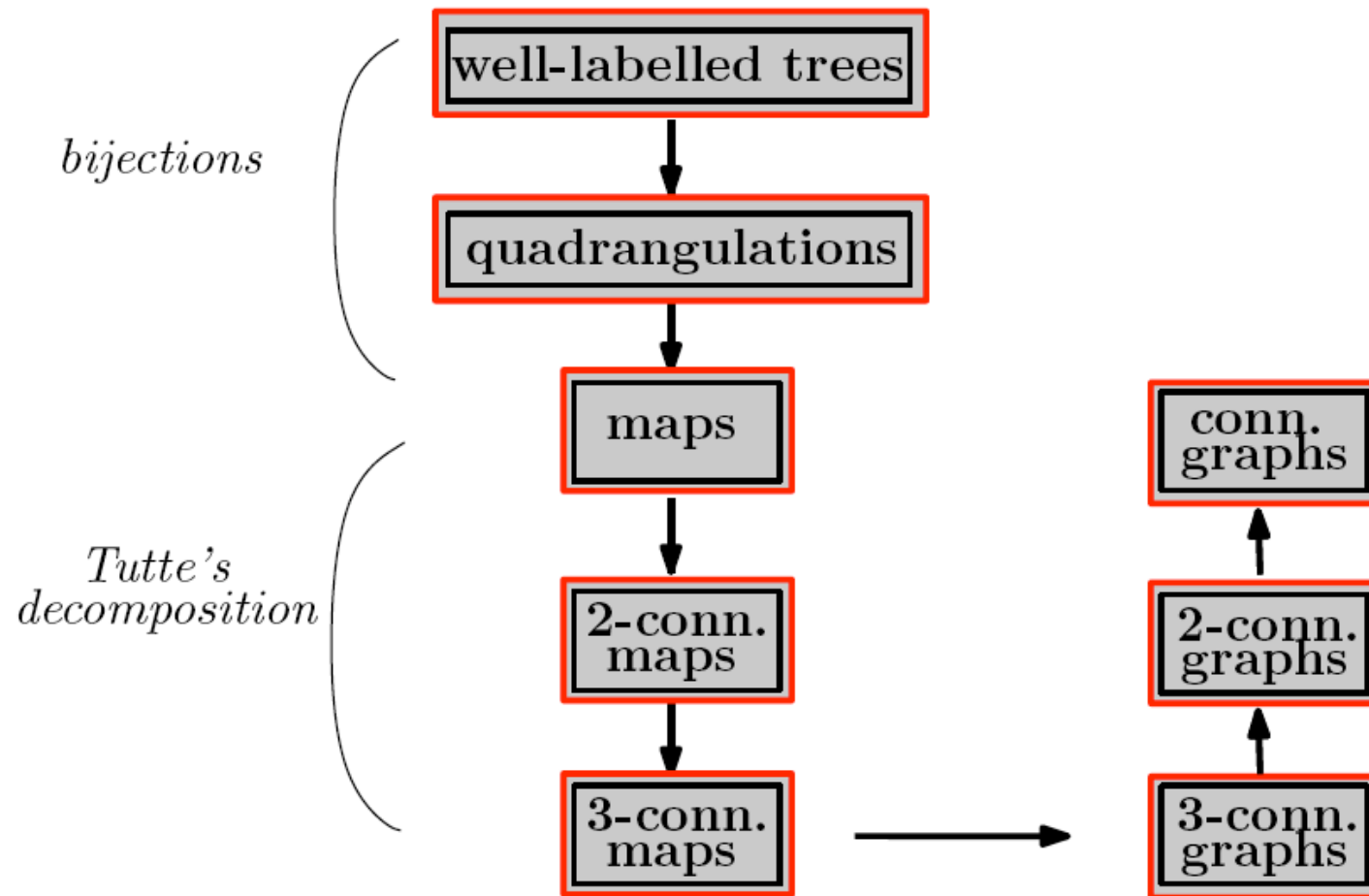
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- Can carry it **until connected planar graphs**



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$$\mathbb{P}(\text{Diam}(G_n) \notin [n^{1/4-\epsilon}, n^{1/4+\epsilon}]) = O(\exp(-n^{\Theta(\epsilon)}))$$

- Can carry it **until connected planar graphs**
- Is it possible to carry a statement such as $\mathbf{E}(\text{Diam}(G_n)) = \Theta(n^{1/4})$?