

# Directed Algebraic Topology and Concurrency

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- From now on  $\mathcal{C}$  denotes a one-way category

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- Then  $\sigma$  is a **potential weak isomorphism** when it preserves both future cones and past cones. Potential weak isomorphisms compose.

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- Then  $\sigma$  is a **potential weak isomorphism** when it preserves both future cones and past cones. Potential weak isomorphisms compose.
- If  $\mathcal{C}(x, y)$  contains a potential weak isomorphism, then it is a singleton  
Requires the assumption that  $\mathcal{C}$  is one-way

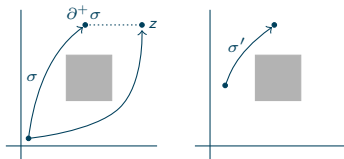
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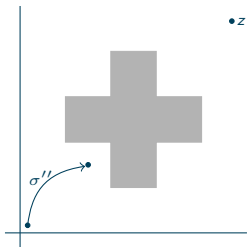


Due to the lower dipath, the  $\sigma, z$ -precomposition is not bijective; yet  $\sigma'$  is a potential weak isomorphism.

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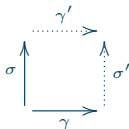


Note that  $\sigma''$  is a potential weak isomorphism though there exists a morphism from  $\partial^+ \sigma''$  to  $z$  but none from  $\partial^+ \sigma''$  to  $z$ .

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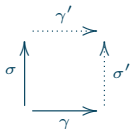
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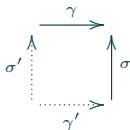


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- The collection  $\Sigma_\infty$  is stable under the action of  $\text{Aut}(\mathcal{C})$

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- Hence we suppose the systems of weak isomorphisms are closed under composition

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- The inverse image (resp. the direct image) of a system of weak isomorphisms by an equivalence of categories is a system of weak isomorphisms.



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- Given a one-way category  $\mathcal{C}$  we have:

All the systems of weak isomorphisms of  $\mathcal{C}$  are pure

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  - The greatest swi is invariant under the action of  $\text{Aut}(\mathcal{C})$

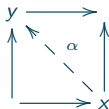
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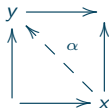
- By definition, a **filling square** category  $\mathcal{C}$  is such that for all commutative squares which are both pushout and pullback (see below), if  $\mathcal{C}(x, y) \neq \emptyset$  then there exists  $\alpha \in \mathcal{C}(x, y)$  that makes both triangles commute.



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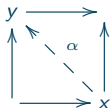


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- If  $\mathcal{C}$  satisfies the filling square property, then any collection of morphisms of  $\mathcal{C}$  that is stable under pushout and pullback is a system of weak isomorphisms.
- A conjecture:

For all loop-free isothetic region  $X$ ,  $\overrightarrow{\pi_1} X$  satisfies the square filling property



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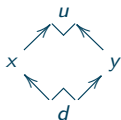


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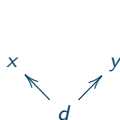


Diagram 2

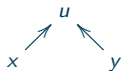


Diagram 3



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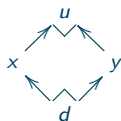


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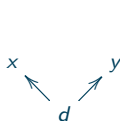


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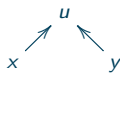


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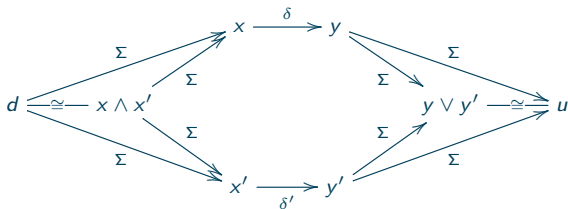
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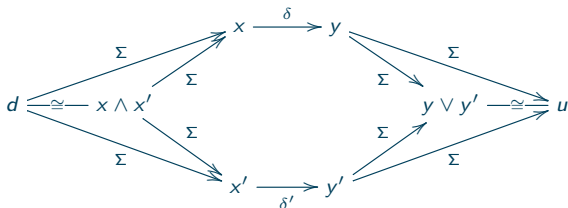
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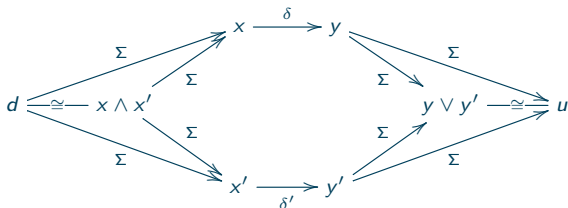


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- If  $\gamma \sim \delta$  then  $\partial^+ \gamma \sim \partial^+ \delta$  and  $\partial^- \gamma \sim \partial^- \delta$

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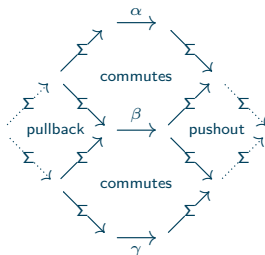
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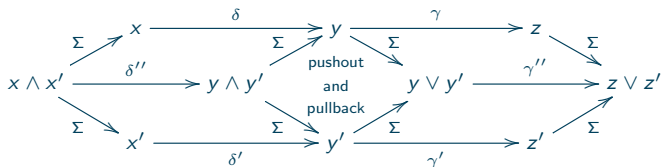
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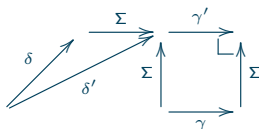
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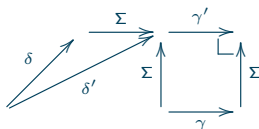
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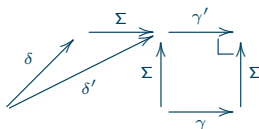


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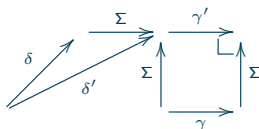


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there exists a unique  $G : \mathcal{C}/\Sigma \rightarrow \mathcal{D}$  s.t.  $F = G \circ Q$

A commutative triangle diagram illustrating the universal property of the quotient functor  $Q$ . The vertices are  $\mathcal{C}$  (bottom-left),  $\mathcal{C}/\Sigma$  (top-right), and  $\mathcal{D}$  (bottom-right). The edges are:  $Q : \mathcal{C} \rightarrow \mathcal{C}/\Sigma$  (diagonal arrow),  $F : \mathcal{C} \rightarrow \mathcal{D}$  (horizontal arrow), and  $G : \mathcal{C}/\Sigma \rightarrow \mathcal{D}$  (vertical arrow). The equation  $F = G \circ Q$  is satisfied.

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- If  $\mathcal{C}/\Sigma(Q(x), Q(y)) \neq \emptyset$  then there exist  $x'$  and  $y'$  such that  $\Sigma(x', x)$ ,  $\Sigma(y, y')$ ,  $\mathcal{C}(x', y)$ , and  $\mathcal{C}(x, y')$  are nonempty.

$$\begin{array}{ccccc}
 x & \xrightarrow{\quad} & c & \xrightarrow{\Sigma} & y' = y \vee c \\
 \Sigma \uparrow & & \lrcorner \uparrow \Sigma & & \uparrow \Sigma \\
 x \wedge a & \xrightarrow{\Sigma} & a & \xrightarrow{\alpha} & b \\
 \Sigma \uparrow & & \uparrow \Sigma & & \\
 x' & \xrightarrow{\quad} & b \wedge y & \xrightarrow{\Sigma} & y
 \end{array}$$

- The quotient functor  $Q$  preserves and reflects potential weak isomorphisms
- If  $\mathcal{C}$  is finite then so is the quotient  $\mathcal{C}/\Sigma$
- $\mathcal{C}$  is a preorder iff  $\mathcal{C}/\Sigma$  is a poset

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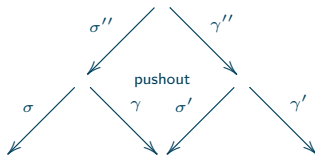
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  - In the diagram below we have  $Q(\gamma' \circ \gamma'') = Q(\gamma') \circ Q(\gamma'') = Q(\gamma') \circ Q(\gamma)$  hence the composite  $(\gamma' \circ \gamma'', \sigma \circ \sigma'')$  neither depend on the choice of the pushout nor on the representatives  $(\gamma, \sigma)$  and  $(\gamma', \sigma')$ .



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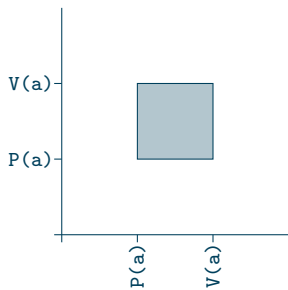
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# Plane without a square

$$X = \mathbb{R}_+^2 \setminus ]0, 1[{}^2$$

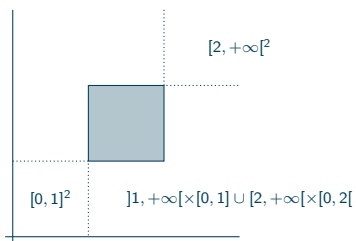
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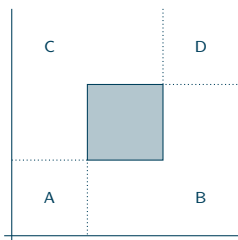
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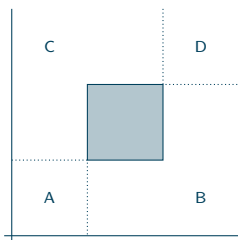
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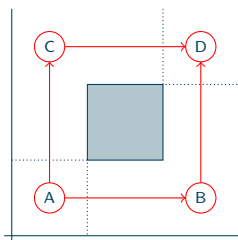


Let  $x, y$  such that  $x \leq^2 y$ , then  $\vec{\pi}_1 X(x, y)$  only depends on which elements of the partition  $x$  and  $y$  belong to

| $\rightarrow$ | $A$      | $B$      | $C$      | $D$  |
|---------------|----------|----------|----------|--|
| $A$           | $\sigma$ | $\beta$  | $\gamma$ | $\beta' \circ \beta$<br>$\alpha' \circ \alpha$ |
| $B$           |          | $\sigma$ |          | $\beta'$                                       |
| $C$           |          |          | $\sigma$ | $\gamma'$                                      |
| $D$           |          |          |          | $\sigma$                                       |

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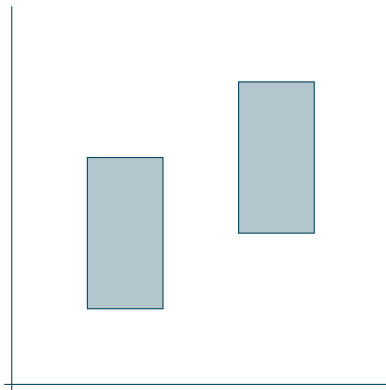


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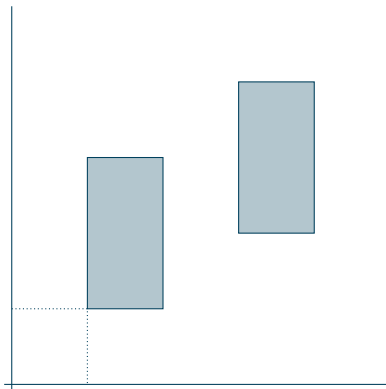
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| $C$           |          |          | $\sigma$ | $\gamma'$                                      |
| $D$           |          |          |          | $\sigma$                                       |

# Two rectangles

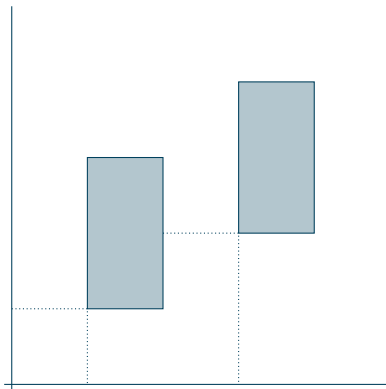
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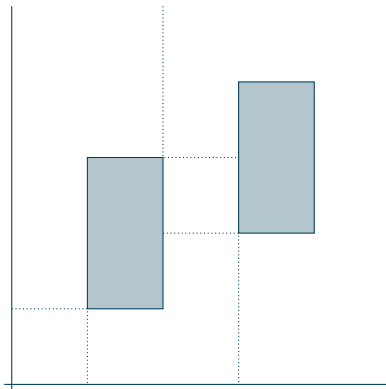
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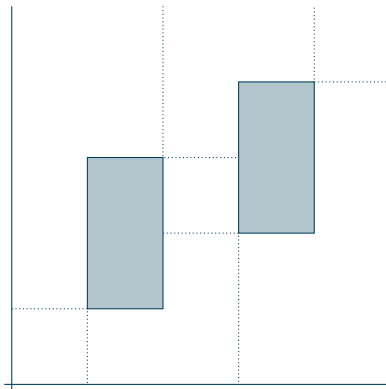


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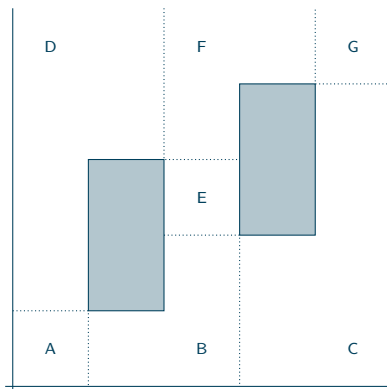




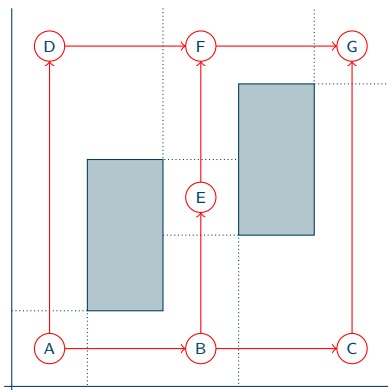
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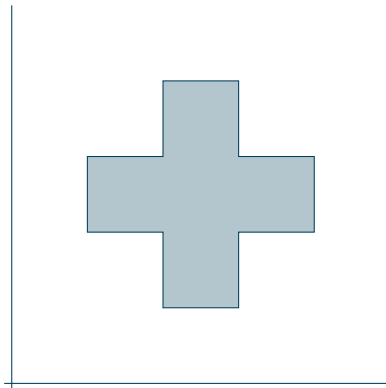


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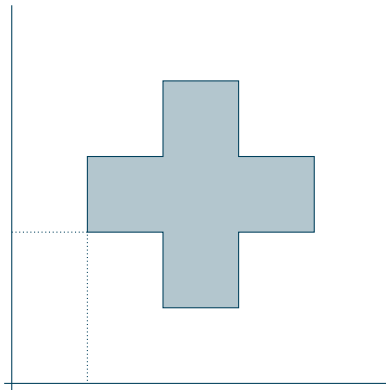


# Swiss Flag

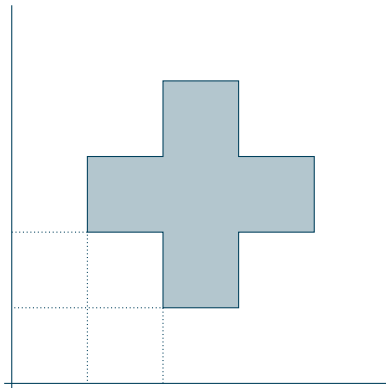
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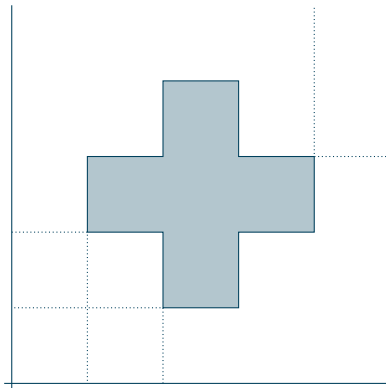
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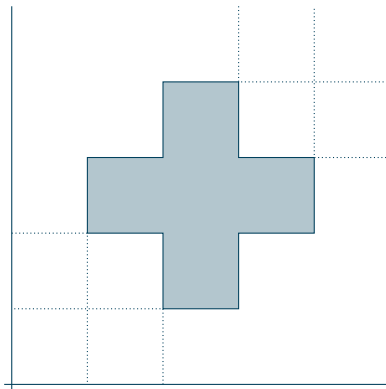


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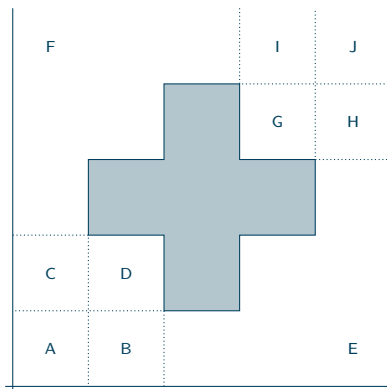




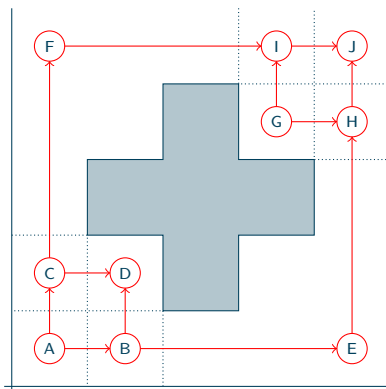
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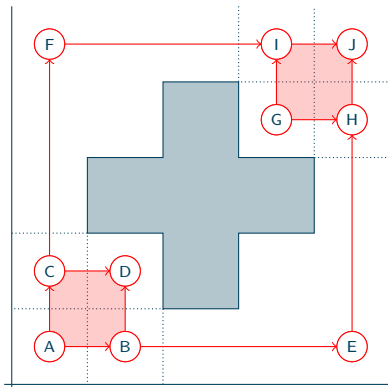
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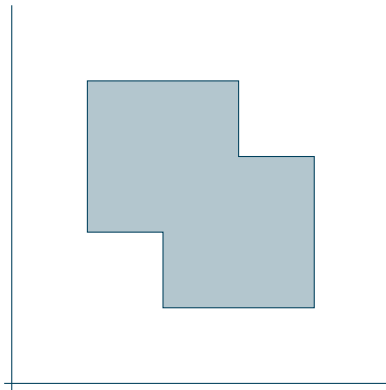


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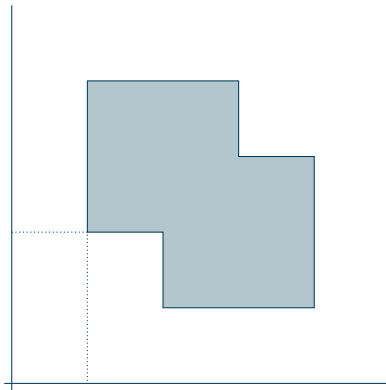


# Achronal overlapping square

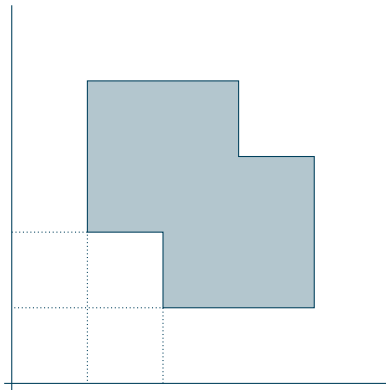
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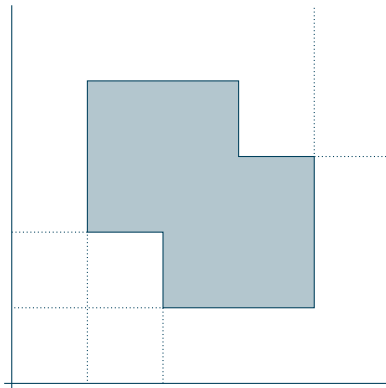


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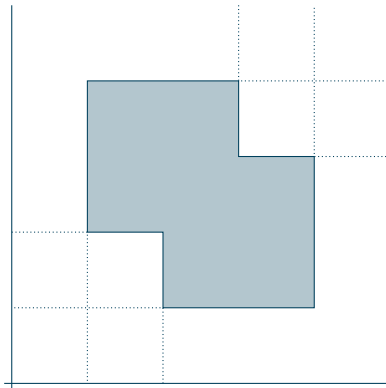




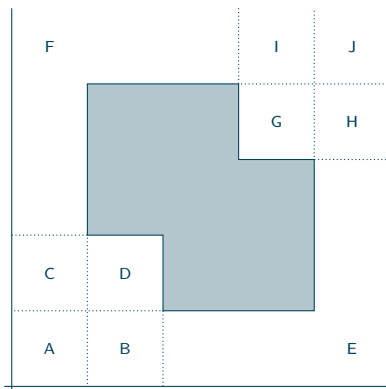
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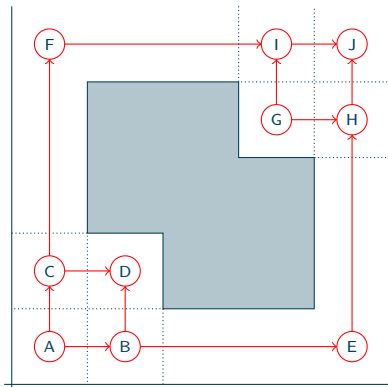
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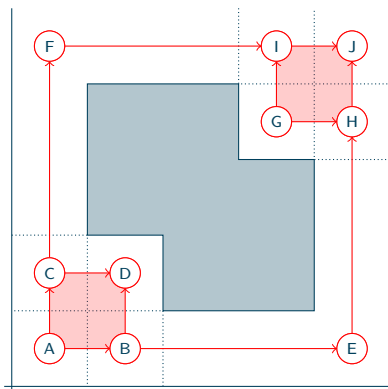
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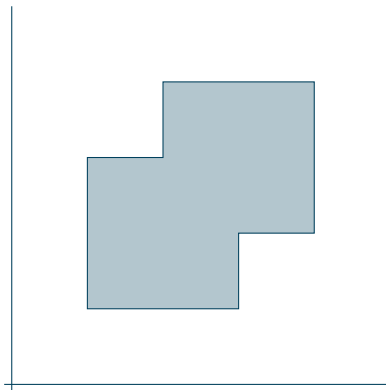


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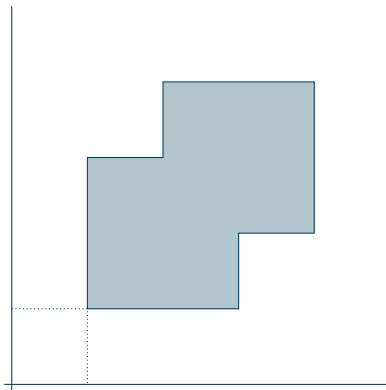


# Diagonal overlapping squares

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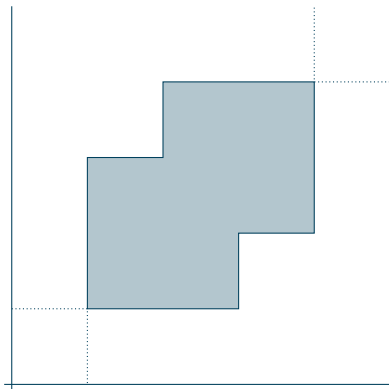


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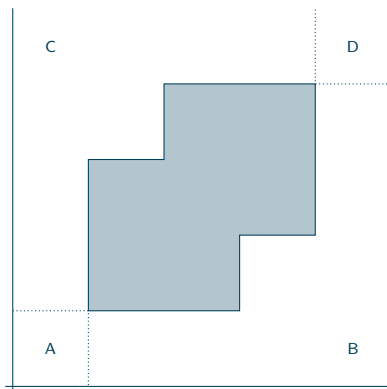




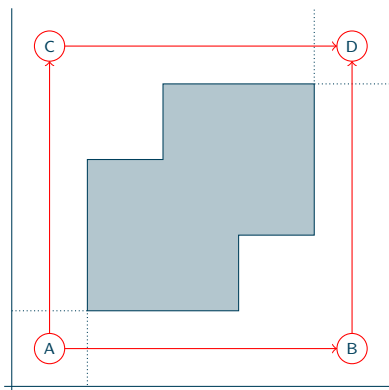
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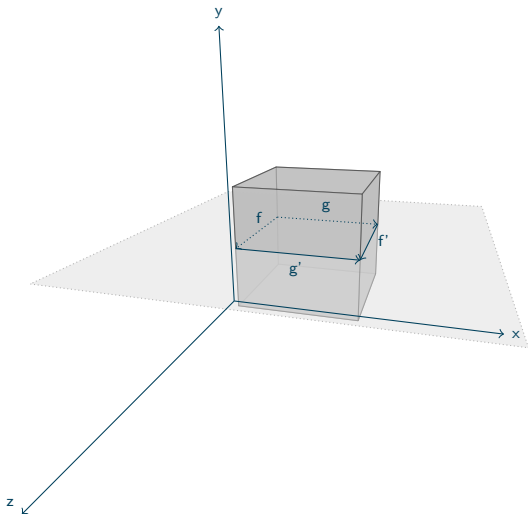


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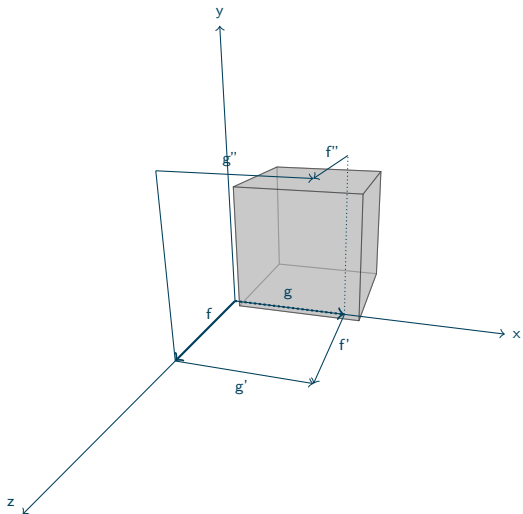


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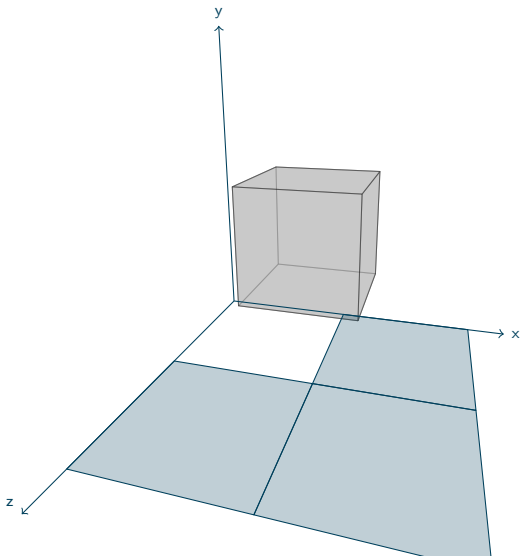
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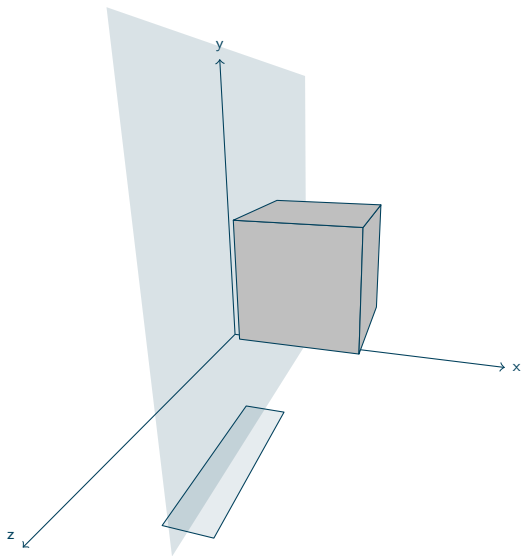
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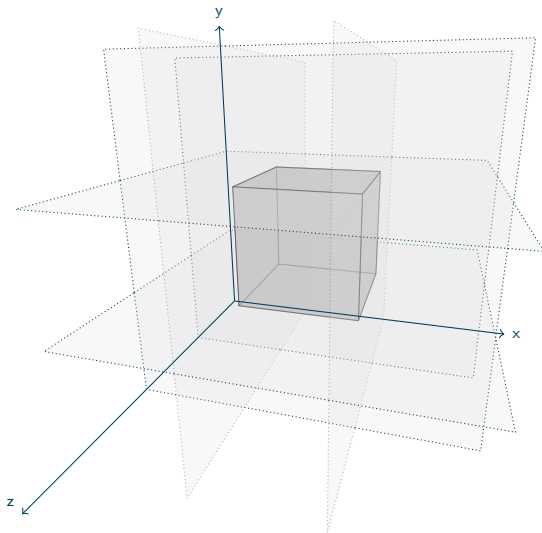
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$$\eta' \cdot \eta : X \times [0, r + r'] \rightarrow Y$$

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- Writing  $f \sim g$  when there is an  $A$ -homotopy from  $f$  to  $g$ , we define an equivalence relation over the mappings from  $X$  to  $Y$ .



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  - for  $n \geq 2$ , the group  $\pi_n(Y, p)$  is abelian



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i.e. for  $n \geq 2$ , the  $n^{\text{th}}$  homotopy group of a space is commutative

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- the function sending  $(n, d)$  to  $\pi_n(\mathbb{S}^d)$  is computable  
Finite Computability of Postnikov Complexes. E. H. Brown, Jr. Ann. of Math. 65(1). 1957

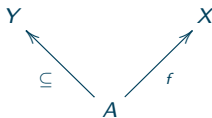
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- The CW-complexes arises in this way.

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- For  $x \in \mathcal{I}_n$ , an open  $n$ -cell is the image of  $\{x\} \times ]0, 1[^n$  under  $\Phi_n$ .  
It is a homeomorphic image.



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  - $F$  is fully faithful and any object of its codomain is isomorphic with an object of its image

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- For all  $n \geq 1$ ,  $\mathbb{R}^{n+1} \setminus \{0\}$  is homotopically equivalent to  $\mathbb{S}^n$

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- Whitehead theorem:

If  $X$  and  $Y$  are CW-complexes and  $f : X \rightarrow Y$  induces isomorphisms of  $n^{\text{th}}$  homotopy groups for all  $n$ , then  $f$  is a homotopy equivalence

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- The homotopy category is defined as the localization of  $\mathcal{T}op$  (or  $\mathcal{CGH}$  etc) with respect to the class of homotopy equivalences

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