

# Acceleration Methods based on Interval Arithmetic in Deterministic Global Optimization

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# Introduction

⇒ Solve NLP

## Problem

$$\left\{ \begin{array}{l} \min_{x \in X \subset \mathbb{R}^n} f(x) \\ s.t. \\ g_i(x) \leq 0, \forall i \in \{1, \dots, n_g\} \\ h_j(x) = 0, \forall j \in \{1, \dots, n_h\} \end{array} \right.$$

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- Problem: exponential complexity of memory and of time

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⇒ acceleration techniques

# Outline

- 1 IBBA
  - Principe
  - Interval Arithmetic
- 2 Constraints Propagation Techniques
  - Principe
  - Example
- 3 Automatic Relaxation Techniques
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- **STOP**  $\implies \max_{(Z,z) \in \mathcal{L}} \text{Wid}(Z) \leq \epsilon_L$   
 $\implies \widetilde{f(x)} - \min_{(Z,z) \in \mathcal{L}} z \leq \epsilon_f$

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Example :  $1/3 \rightarrow [0.33333331, 0.33333335]$   
 $\rightarrow [0.33333333333333331, 0.33333333333333338]$

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Every real number is represented by an interval of two floating point numbers which encloses it.

⇒ take care about numerical errors during calculations

$$f = 33.75y^6 + x^2(11x^2y^2 - y^6 - 121y^4 - 2) + 5.5y^8 + \frac{x}{2y}$$

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But the real value is :  $f = -\frac{54767}{66192} = -0.8273960599$

# Interval Arithmetic

$$\left\{ \begin{array}{l} [a, b] + [c, d] = [a + c, b + d] \\ [a, b] - [c, d] = [a - d, b - c] \\ [a, b] \times [c, d] = [\min\{a \times c, a \times d, b \times c, b \times d\}, \\ \quad \max\{a \times c, a \times d, b \times c, b \times d\}] \\ [a, b] \div [c, d] = [a, b] \times \left[\frac{1}{d}, \frac{1}{c}\right] \text{ si } 0 \notin [c, d] \end{array} \right.$$

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• Example :

$$\begin{aligned} [1; 2] + [3; 4] &= [4; 6] \\ [4; 6] - [3; 4] &= [0; 3] \\ [1; 2] \times [3; 4] &= [3; 8] \\ [3; 8] \div [3; 4] &= [3/4; 8/3] \\ [-1; 1] \times ([1; 2] + [3; 4]) &= [-6; 4] \\ [-1; 1] \times [1; 2] + [-1; 1] \times [3; 4] &= [-6; 6] \end{aligned}$$

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 Exact Bounds,  $\forall x \in [1; 2] \times [3; 4], f(x) \in [6.75; 13]$

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Idea: use some deduction steps for reducing the box  $X$ .

Linear case: if  $c(x) = \sum_{i=1}^n a_i x_i = [a, b]$  then:

$$X_k := \left( \frac{[a, b] - \sum_{i=1, i \neq k}^n a_i X_i}{a_k} \right) \cap X_k, \text{ si } a_k \neq 0. \quad (1)$$

where  $k$  is in  $\{1, \dots, n\}$  and  $X_i$  is the  $i^{th}$  component of  $X$ .

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**Other Idea**: construction of the calculus tree and propagation.

# Example of Propagation Technique based on the Calculus Tree

Let  $c(x) = 2x_3x_2 + x_1$  and

$$c(x) = 3$$

where  $x_i \in [1, 3]$  for all  $i \in \{1, 2, 3\}$ .

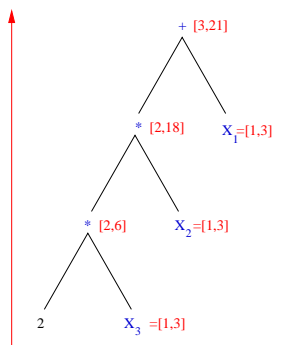


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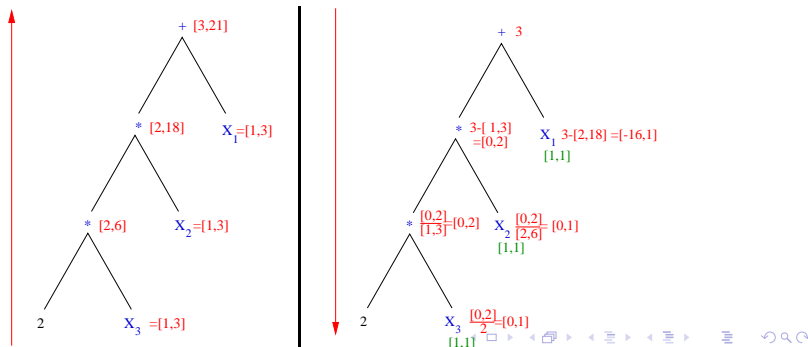


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# Affine Arithmetic: MVA Andrade, J Comba, J Stolfi (1994)

## Definition

*All real numbers are represented by an affine form  $\hat{x}$*

$$\hat{x} = x_0 + \sum_{i=1}^n x_i \epsilon_i$$

*with  $\forall i \in [1; n], x_i \in \mathbb{R}$  and  $\epsilon_i = [-1; 1]$*

- Example : with  $n = 1$

$$\begin{aligned} 1/3 &\rightarrow 0.333 + 0.001 * [-1; 1] \\ \log(2) &\rightarrow 0.693 + 0.001 * [-1; 1] \end{aligned}$$

# Affine Operator

$$\begin{aligned}\hat{x} \pm \hat{y} &= (x_0 \pm y_0) + \sum_{i=1}^n (x_i \pm y_i) \epsilon_i \\ a \pm \hat{x} &= (a \pm x_0) + \sum_{i=1}^n x_i \epsilon_i \\ a \times \hat{x} &= ax_0 + \sum_{i=1}^n ax_i \epsilon_i\end{aligned}$$

- Example :  $A = [1; 3]$  and  $B = [-2; 0]$

$$\begin{aligned}\hat{A} &\rightarrow 2 + \epsilon_1 \\ \hat{B} &\rightarrow -1 + \epsilon_2 \\ \hat{A} + \hat{B} &= 1 + \epsilon_1 + \epsilon_2\end{aligned}$$

## Non-Affine Operator

$$\begin{aligned}x \times y &= \left(x_0 + \sum_{i=1}^n x_i \epsilon_i\right) \times \left(y_0 + \sum_{i=1}^n y_i \epsilon_i\right) \\&= x_0 y_0 + \sum_{i=1}^n (x_0 y_i + x_i y_0) \epsilon_i + \left(\sum_{i=1}^n |x_i| \times \sum_{i=1}^n |y_i|\right) \epsilon_{n+1}\end{aligned}$$

$$\hat{f}(\hat{x}) = \zeta + \alpha \hat{x} + \delta \epsilon_{n+1}$$

$$\text{with } \alpha, \delta, \zeta \in \mathbb{R} \text{ and } \hat{x} = x_0 + \sum_{i=1}^n x_i \epsilon_i$$

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$\implies$  All non-affine operations add a new variable

## Natural Extension

$$\forall x \in [1; 2] \times [3; 4], f(x) = x_1^2 + x_2^2 - x_1 x_2$$



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In real,  $f(x) \in [6.75; 13]$

## New Affine Forms : F. Messine (2002)

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- AF2 :  $\hat{x} = x_0 + \sum_{i=1}^n x_i \epsilon_i + e \epsilon_{n+1} + e_+ \epsilon_{n+2} + e_- \epsilon_{n+3}$   
with  $\epsilon_{n+1} = [-1; 1]$  ,  $\epsilon_{n+2} = [0; 1]$  and  $\epsilon_{n+3} = [-1; 0]$



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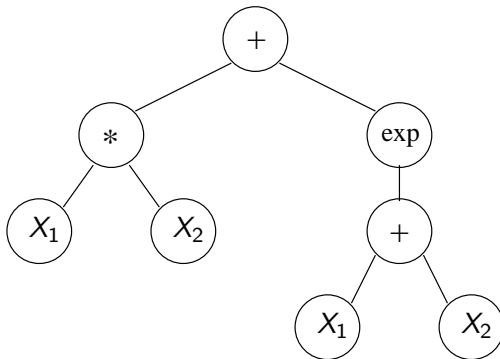
## Visualization using a tree

$$f(x_1, x_2) = x_1 * x_2 + \exp(x_1 + x_2) , x_1 = [-3, -1] \text{ and } x_2 = [3, 5]$$



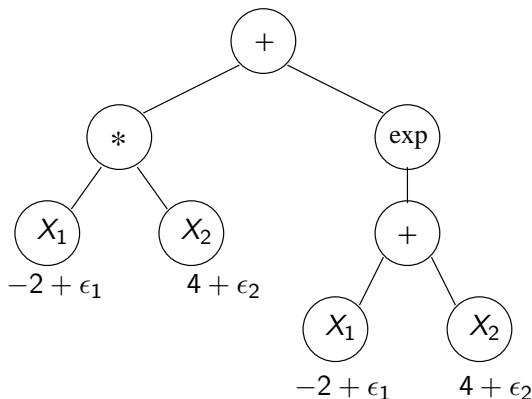
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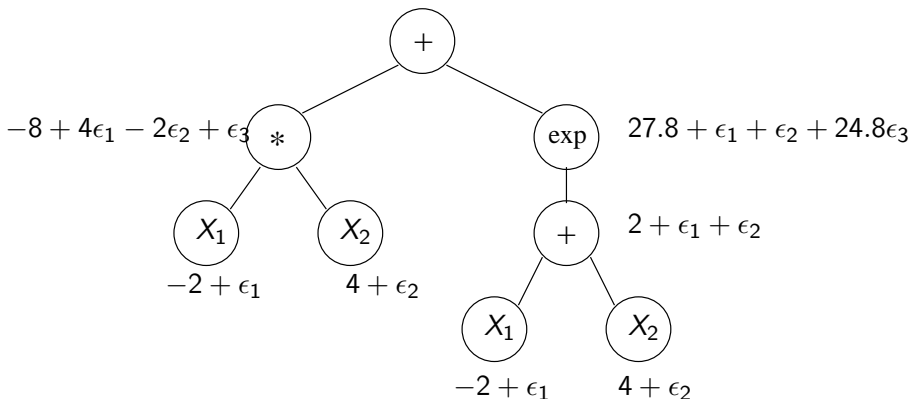
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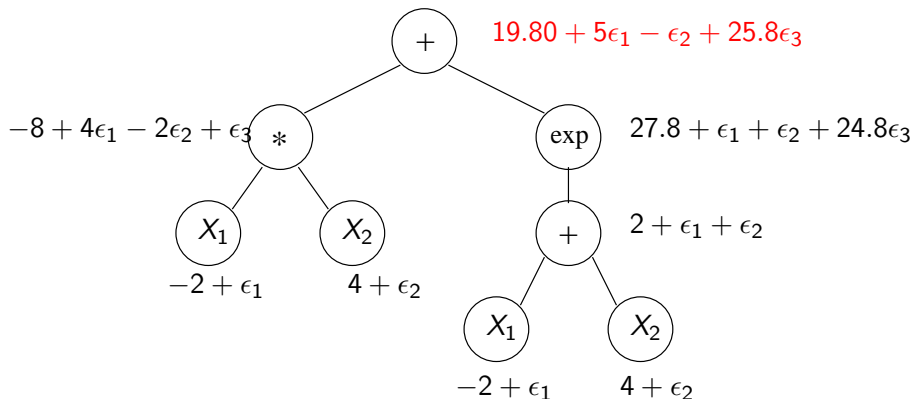
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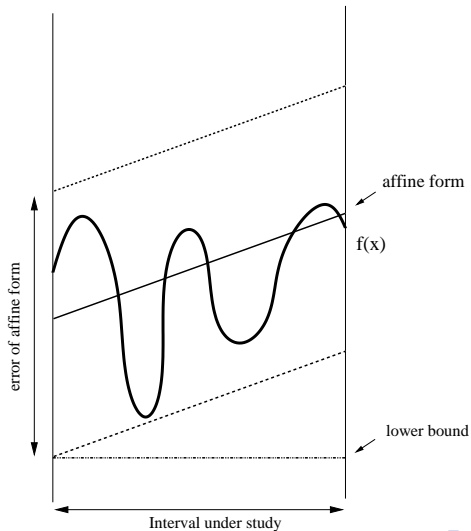


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# Graphical view



# Notation

$$\forall x \in X \subset \mathbb{R}^n, f(x) - L_f(T(x)) \in E_f$$

where  $X$  is the domain under study,  
 $T$  is the affine transformation of  $X$  to  $[-1; 1]^n$ ,  
 $L_f$  is a linear function of  $[-1; 1]^n$  to  $\mathbb{R}^n$ ,  
 $E_f = [\underline{E}_f; \overline{E}_f]$  is the interval corresponding to  
 the error generated.

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With AF1,  $L_f(\epsilon) = \sum_{i=1}^n x_i \epsilon_i$  and  $E_f = x_0 + x_{n+1} \epsilon_{n+1}$

With AF2,  $L_f(\epsilon) = \sum_{i=1}^n x_i \epsilon_i$  and  $E_f = x_0 + e \epsilon_{n+1} + e_+ \epsilon_{n+2} + e_- \epsilon_{n+3}$

# Reformulation

$$\left\{ \begin{array}{l} \min_{x \in X^n} f(x) \\ \text{s. t.} \\ g_i(x) \leq 0, \forall i \in \{1, \dots, n_g\} \\ h_j(x) = 0, \forall j \in \{1, \dots, n_h\} \end{array} \right.$$

$$\left\{ \begin{array}{l} \min_{\epsilon \in [-1;1]^n} L_f(\epsilon) \\ \text{s. t.} \\ L_{g_i}(\epsilon) \leq -\underline{E}_{g_i}, \quad \forall i \in \{1, \dots, n_g\} \\ L_{h_j}(\epsilon) \leq -\underline{E}_{h_j}, \quad \forall j \in \{1, \dots, n_h\} \\ -L_{h_j}(\epsilon) \leq \overline{E}_{h_j}, \quad \forall j \in \{1, \dots, n_h\} \end{array} \right.$$



$$\left\{ \begin{array}{l} \min \\ x \in [1;1.5] \times [4.5;5] \times [3.5;4] \times [1;1.5] \end{array} \right. \quad \begin{array}{l} x_3 + (x_1 + x_2 + x_3)x_1x_4 \\ x_1x_2x_3x_4 \geq 25 \\ x_1^2 + x_2^2 + x_3^2 + x_4^2 = 40 \end{array}$$

$$\left\{ \begin{array}{l} \min_{x \in [1;1.5] \times [4.5;5] \times [3.5;4] \times [1;1.5]} \quad x_3 + (x_1 + x_2 + x_3)x_1x_4 \\ x_1x_2x_3x_4 \geq 25 \\ x_1^2 + x_2^2 + x_3^2 + x_4^2 = 40 \end{array} \right.$$

$$x_3 + (x_1 + x_2 + x_3)x_1x_4 \Leftrightarrow 18.98 + 3.44\epsilon_1 + 0.39\epsilon_2 + 0.64\epsilon_3 + 3.05\epsilon_4 + 1.13\epsilon_5$$

$$\left\{ \begin{array}{l} \min_{x \in [1;1.5] \times [4.5;5] \times [3.5;4] \times [1;1.5]} \quad x_3 + (x_1 + x_2 + x_3)x_1x_4 \\ x_1x_2x_3x_4 \geq 25 \\ x_1^2 + x_2^2 + x_3^2 + x_4^2 = 40 \end{array} \right.$$

$$\begin{aligned} x_3 + (x_1 + x_2 + x_3)x_1x_4 &\Leftrightarrow 18.98 + 3.44\epsilon_1 + 0.39\epsilon_2 + 0.64\epsilon_3 + 3.05\epsilon_4 + 1.13\epsilon_5 \\ 25 - x_1x_2x_3x_4 &\Leftrightarrow -2.83 - 5.57\epsilon_1 - 1.46\epsilon_2 - 1.86\epsilon_3 - 5.57\epsilon_4 + 2.71\epsilon_5 \end{aligned}$$

$$\left\{ \begin{array}{l} \min_{x \in [1;1.5] \times [4.5;5] \times [3.5;4] \times [1;1.5]} \quad x_3 + (x_1 + x_2 + x_3)x_1x_4 \\ x_1x_2x_3x_4 \geq 25 \\ x_1^2 + x_2^2 + x_3^2 + x_4^2 = 40 \end{array} \right.$$

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$$\left\{ \begin{array}{l} \min_{x \in [1;1.5] \times [4.5;5] \times [3.5;4] \times [1;1.5]} \quad x_3 + (x_1 + x_2 + x_3)x_1x_4 \\ x_1x_2x_3x_4 \geq 25 \\ x_1^2 + x_2^2 + x_3^2 + x_4^2 = 40 \end{array} \right.$$

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$$\left\{ \begin{array}{l} \min_{\epsilon \in [-1;1]^4} \quad 3.44\epsilon_1 + 0.39\epsilon_2 + 0.64\epsilon_3 + 3.05\epsilon_4 \\ -5.57\epsilon_1 - 1.46\epsilon_2 - 1.86\epsilon_3 - 5.57\epsilon_4 \leq 5.54 \\ 0.62\epsilon_1 + 2.38\epsilon_2 + 1.88\epsilon_3 + 0.63\epsilon_4 \leq 0.5 \\ -0.62\epsilon_1 - 2.38\epsilon_2 - 1.88\epsilon_3 - 0.63\epsilon_4 \leq 0 \end{array} \right.$$

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Solution of linear program :  $\epsilon_R = (-1; -0.24; 1; -0.26)$   
 $L_f(\epsilon_R) = -3.70$



$$\left\{ \begin{array}{l} \min_{x \in [1;1.5] \times [4.5;5] \times [3.5;4] \times [1;1.5]} \quad x_3 + (x_1 + x_2 + x_3)x_1x_4 \\ x_1x_2x_3x_4 \geq 25 \\ x_1^2 + x_2^2 + x_3^2 + x_4^2 = 40 \end{array} \right.$$

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lower bound =  $L_f(\epsilon_R) + E_f^L$

$$\left\{ \begin{array}{l} \min_{x \in [1;1.5] \times [4.5;5] \times [3.5;4] \times [1;1.5]} \quad x_3 + (x_1 + x_2 + x_3)x_1x_4 \\ x_1x_2x_3x_4 \geq 25 \\ x_1^2 + x_2^2 + x_3^2 + x_4^2 = 40 \end{array} \right.$$

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Solution of linear program :  $\epsilon_R = (-1; -0.24; 1; -0.26)$   
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Lower bound by IA = 12.5 ; by AF1 = 10.34375

$$\left\{ \begin{array}{l} \min_{x \in [1;1.5] \times [4.5;5] \times [3.5;4] \times [1;1.5]} \quad x_3 + (x_1 + x_2 + x_3)x_1x_4 \\ x_1x_2x_3x_4 \geq 25 \\ x_1^2 + x_2^2 + x_3^2 + x_4^2 = 40 \end{array} \right.$$

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Solution of linear program :  $\epsilon_R = (-1; -0.24; 1; -0.26)$   
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Lower bound by IA = 12.5 ; by AF1 = 10.34375

New lower bound = 14.15

# Numerical Test

name	<i>n</i>	<i>m</i>	CP		AFFINE2+CP		AFFINE2	
			T(s)	Iter	T(s)	Iter	T(s)	Iter
TUY	4	2			4.95	<b>58 022</b>	18.40	134 219
2_1_1	5	1	0.02	26 309	0.00	<b>149</b>	0.00	149
2_1_2	6	2	0.03	118 209	1.08	113 545	0.90	<b>31 610</b>
2_1_3	13	9			1.25	114 439	0.43	<b>41 434</b>
2_1_4	6	5	0.00	9 881	0.10	<b>2 407</b>	0.15	2 598
2_1_5	10	11	1.83	1 449 757	37.45	<b>354 376</b>	45.25	392 234
2_1_6	10	5			0.03	<b>659</b>	0.02	660
3_1_2	5	6	0.00	11 330	0.03	<b>464</b>	0.12	595
3_1_3	6	6	0.02	34 148	0.25	<b>18 356</b>	0.18	3 794
3_1_4	3	3	0.02	35 902	0.02	<b>249</b>	0.03	291
7_3_1	4	7	0.05	30799	0.08	3592	0.12	<b>3444</b>
7_3_2	4	7	0.00	768	0.00	<b>155</b>	0.03	587
7_3_3	5	8	0.05	86282	0.03	<b>550</b>	0.12	1535
14_2_1	5	7	32.40	14 659 121	44.38	<b>944 184</b>	>	
14_2_2	4	5	0.02	11 708	0.10	<b>2 719</b>	0.18	6 378
14_2_3	6	9			110.03	2 177 710	37.25	<b>661 868</b>
14_2_4	5	7	11.92	2 284 630	1.10	<b>31 337</b>	2.58	54 418
14_2_5	4	5	0.03	20 656	0.07	<b>2 193</b>	0.25	5 129

# Conclusion

Acceleration Methods based on expression trees:

- Constraints Propagation Techniques
- Automatic Relaxation Techniques